



Solving large scale eigenvalue problems

Lecture 3, March 7, 2018: Newton methods

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Survey of today's lecture

- ▶ Linear and nonlinear eigenvalue problems
- ▶ Eigenvalues as zeros of the determinant function
- ▶ Hyman's method for Hessenberg matrices
- ▶ Algorithmic differentiation
- ▶ Newton iterations
- ▶ Successive linear approximations

Linear and nonlinear eigenvalue problems

- ▶ Linear eigenvalue problems

Find values $\lambda \in \mathbb{C}$ such that $A - \lambda I$ is **singular**.

Or equivalently:

Find values $\lambda \in \mathbb{C}$ such that there is a nonzero (nontrivial) \mathbf{x} such that

$$(A - \lambda I)\mathbf{x} = \mathbf{0} \iff A\mathbf{x} = \lambda\mathbf{x}.$$

Linear and nonlinear eigenvalue problems (cont.)

► Nonlinear eigenvalue problems

More general: Find $\lambda \in \mathbb{C}$ such that $A(\lambda)\mathbf{x} = \mathbf{0}$ where $A(\lambda)$ is a matrix the elements of which depend on λ .

Examples: $A(\lambda) = \sum_{k=0}^d \lambda^k A_k$;

$d = 1$: $A(\lambda) = A_0 - \lambda A_1$, $A_0 = A$, $A_1 = I$.

Linear and nonlinear eigenvalue problems (cont.)

▶ Matrix polynomials

Matrix polynomials can be **linearized**.

Example: $A\mathbf{x} + \lambda K\mathbf{x} + \lambda^2 M\mathbf{x}$.

We can generate equivalent eigenvalue problems that are linear but have the size doubled: With $\mathbf{y} = \lambda\mathbf{x}$ we get

$$\begin{pmatrix} A & O \\ O & I \end{pmatrix} \begin{pmatrix} \mathbf{x} \\ \mathbf{y} \end{pmatrix} = \lambda \begin{pmatrix} -K & -M \\ I & O \end{pmatrix} \begin{pmatrix} \mathbf{x} \\ \mathbf{y} \end{pmatrix}$$

or

$$\begin{pmatrix} A & K \\ O & I \end{pmatrix} \begin{pmatrix} \mathbf{x} \\ \mathbf{y} \end{pmatrix} = \lambda \begin{pmatrix} O & -M \\ I & O \end{pmatrix} \begin{pmatrix} \mathbf{x} \\ \mathbf{y} \end{pmatrix}.$$

Many other linearizations exist.

(C.f. transformation of high order to first order ODE's.)

Numerical example

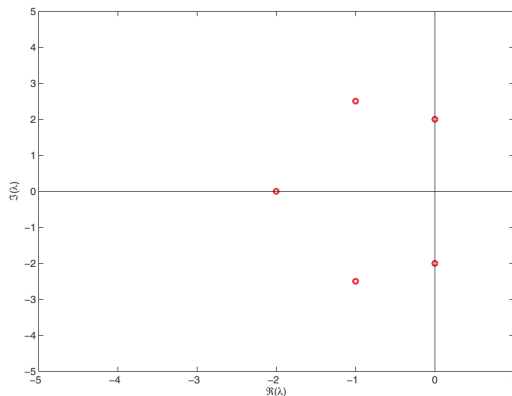
- ▶ Example: The matrix

$$A = \begin{pmatrix} -0.9880 & 1.8000 & -0.8793 & -0.5977 & -0.7819 \\ -1.9417 & -0.5835 & -0.1846 & -0.7250 & 1.0422 \\ 0.6003 & -0.0287 & -0.5446 & -2.0667 & -0.3961 \\ 0.8222 & 1.4453 & 1.3369 & -0.6069 & 0.8043 \\ -0.4187 & -0.2939 & 1.4814 & -0.2119 & -1.2771 \end{pmatrix}$$

has eigenvalues given approximately by $\lambda_1 = -2$,
 $\lambda_2 = -1 + 2.5i$, $\lambda_3 = -1 - 2.5i$, $\lambda_4 = 2i$, and $\lambda_5 = -2i$.

It is known that closed form formulas for the roots of a polynomial do not generally exist if the polynomial is of degree 5 or higher. Thus we cannot expect to be able to solve the eigenvalue problem in a finite procedure.

Numerical example (cont.)



Eigenvalues in \mathbb{C} . For real matrices, the complex eigenvalues come in pairs. If λ is an eigenvalue, then so is $\bar{\lambda}$.

Zeros of determinant

Find values $\lambda \in \mathbb{C}$ such that $A - \lambda I$ is singular.

Equivalent:

Find values $\lambda \in \mathbb{C}$ such that

$$\det A(\lambda) = 0. \quad (1)$$

Apply zero finder to eq. (1).

Questions:

1. What zero finder?
2. How to compute $f(\lambda) = \det A(\lambda)$?
3. How to compute $f'(\lambda) = \frac{d}{d\lambda} \det A(\lambda)$?

Gaussian elimination with partial pivoting (GEPP)

Let the factorization

$$P(\lambda)A(\lambda) = L(\lambda)U(\lambda)$$

be obtained by GEPP.

P : permutation matrix,

L : lower unit triangular matrix,

U : upper triangular matrix.

$$\det P(\lambda) \cdot \det A(\lambda) = \det L(\lambda) \cdot \det U(\lambda).$$

$$\pm 1 \cdot \det A(\lambda) = 1 \cdot \prod_{i=1}^n u_{ii}(\lambda).$$

Newton iteration

Need the derivative $f'(\lambda)$ of $f(\lambda) = \det A(\lambda)$.

$$\begin{aligned} f'(\lambda) &= \pm 1 \cdot \sum_{i=1}^n u'_{ii}(\lambda) \prod_{j \neq i}^n u_{jj}(\lambda) \\ &= \pm 1 \cdot \sum_{i=1}^n \frac{u'_{ii}(\lambda)}{u_{ii}(\lambda)} \prod_{j=1}^n u_{jj}(\lambda) = \sum_{i=1}^n \frac{u'_{ii}(\lambda)}{u_{ii}(\lambda)} f(\lambda). \end{aligned}$$

How do we compute the u'_{ii} ?

Possibility: **algorithmic differentiation**

See: Arbenz & Gander: Solving Nonlinear Eigenvalue Problems by Algorithmic Differentiation. Computing 36, 205 – 215 (1986).

Algorithmic differentiation

Example: **Horner scheme** to evaluate polynomial

$$f(z) = \sum_{i=1}^n c_i z^i.$$

$$p_0(z) = c_0 + z (c_1 + z (c_2 + \cdots + z (c_n)))$$

by the recurrence

$$p_n := c_n,$$

$$p_i := z p_{i+1} + c_i, \quad i = n-1, n-2, \dots, 0$$

$$f(z) := p_0.$$

Consider the p_i as functions (polynomials) in z .

Algorithmic differentiation (cont.)

$$dp_n := 0, \quad p_n := c_n,$$

$$dp_i := p_{i+1} + z dp_{i+1}, \quad p_i := z p_{i+1} + c_i, \quad i = n-1, n-2, \dots, 0,$$

$$f'(z) := dp_0, \quad f(z) := p_0.$$

Can proceed in a similar fashion for computing $\det A(\lambda)$.

Need to be able to compute derivatives a'_{ij} . Then, derive each single assignment in the algorithm of Gaussian elimination.

Discussion

We restrict ourselves to the standard eigenvalue problem $A\mathbf{x} = \lambda\mathbf{x}$, i.e., $A(\lambda) = A - \lambda I$.

Then $A'(\lambda) = -I$.

In the Newton method we have to compute the determinant for possibly many values λ .

Computing the determinant costs $\frac{2}{3}n^3$ flops (floating point operations).

Can we do better?

Idea: Transform A by a similarity transformation to **Hessenberg form**.

Hessenberg matrices

Definition

A matrix H is a **Hessenberg matrix** if its elements below the lower off-diagonal are zero,

$$h_{ij} = 0, \quad i > j + 1.$$

Any matrix A can be transformed into a Hessenberg matrix by a sequence of elementary Householder transformations, for details see QR algorithm.

Let $S^*AS = H$, where S is **unitary**. Then

$$A\mathbf{x} = \lambda\mathbf{x} \iff H\mathbf{y} = \lambda\mathbf{y}, \quad \mathbf{x} = S\mathbf{y}.$$

We assume that H is **unreduced**, i.e., $h_{i+1,i} \neq 0$ for all i .

Hessenberg matrices (cont.)

Let λ be an eigenvalue of H and

$$(H - \lambda I)\mathbf{x} = \mathbf{0}, \quad (2)$$

i.e., \mathbf{x} is an eigenvector of H associated with the eigenvalue λ .

Then $x_n \neq 0$. (Proof by contradiction.)

W.l.o.g., we can set $x_n = 1$.

If λ is an eigenvalue then there are x_i , $1 \leq i < n$, such that

$$\begin{pmatrix} h_{11} - \lambda & h_{12} & h_{13} & h_{14} \\ h_{21} & h_{22} - \lambda & h_{23} & h_{24} \\ & h_{32} & h_{33} - \lambda & h_{34} \\ & & h_{43} & h_{44} - \lambda \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}.$$

Hessenberg matrices (cont.)

If λ is **not** an eigenvalue then we determine the x_i such that

$$\left(\begin{array}{ccc|c} h_{11} - \lambda & h_{12} & h_{13} & h_{14} \\ h_{21} & h_{22} - \lambda & h_{23} & h_{24} \\ & h_{32} & h_{33} - \lambda & h_{34} \\ & & h_{43} & h_{44} - \lambda \end{array} \right) \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ 1 \end{pmatrix} = \begin{pmatrix} * \\ 0 \\ 0 \\ 0 \end{pmatrix}. \quad (*)$$

Determine the $n - 1$ numbers $x_{n-1}, x_{n-2}, \dots, x_1$ by the equations n down to **2** of the equation above

$$x_i = \frac{-1}{h_{i+1,i}} \left((h_{i+1,i+1} - \lambda) x_{i+1} + h_{i+1,i+2} x_{i+2} + \dots + h_{i+1,n} \underbrace{x_n}_1 \right).$$

The first equation gives

$$(h_{1,1} - \lambda) x_1 + h_{1,2} x_2 + \dots + h_{1,n} x_n = c \cdot f(\lambda). \quad (3)$$

Hessenberg matrices (cont.)

We can consider the x_i as functions of λ , in fact, $x_i \in \mathbb{P}_{n-i}$.

Therefore, we can algorithmically differentiate the x'_i to get $f'(\lambda)$.

For $i = n - 1, \dots, 1$ we have

$$x'_i = \frac{-1}{h_{i+1,i}} \left(-x_{i+1} + (h_{i+1,i+1} - \lambda) x'_{i+1} + h_{i+1,i+2} x'_{i+2} + \dots + h_{i+1,n-1} x'_{n-1} \right).$$

Finally,

$$c \cdot f'(\lambda) = -x_1 + (h_{1,1} - \lambda) x'_1 + h_{1,2} x'_2 + \dots + h_{1,n-1} x'_{n-1}.$$

Hessenberg matrices (matrix form)

In matrix form (*) reads

$$(H - \lambda I) \begin{pmatrix} \mathbf{x}(\lambda) \\ 1 \end{pmatrix} = \begin{bmatrix} \mathbf{h}(\lambda) & h_{1n} \\ R(\lambda) & \mathbf{k}(\lambda) \end{bmatrix} \begin{pmatrix} \mathbf{x} \\ 1 \end{pmatrix} = \begin{pmatrix} p(\lambda) \\ \mathbf{0} \end{pmatrix}.$$

Computing p :

$$R(\lambda)\mathbf{x}(\lambda) + \mathbf{k}(\lambda) = \mathbf{0} \implies \mathbf{x}(\lambda) = -R(\lambda)^{-1}\mathbf{k}(\lambda),$$
$$p(\lambda) = \mathbf{h}(\lambda)\mathbf{x}(\lambda) + h_{1n}.$$

Hessenberg matrices (matrix form) (cont.)

Computing $q = p'$:

$$R'(\lambda)\mathbf{x}(\lambda) + R(\lambda)\mathbf{x}'(\lambda) = -\mathbf{k}'(\lambda) = \begin{pmatrix} 0 \\ \vdots \\ 0 \\ 1 \end{pmatrix}, \quad R'(\lambda) = \begin{bmatrix} 0 & -1 & & \\ & 0 & \ddots & \\ & & \ddots & -1 \\ & & & 0 \end{bmatrix}.$$

$$\mathbf{x}'(\lambda) = R(\lambda)^{-1}[-\mathbf{k}'(\lambda) - R'(\lambda)\mathbf{x}] = R(\lambda)^{-1} \begin{pmatrix} x_2 \\ \vdots \\ x_{n-1} \\ 1 \end{pmatrix}$$

$$q(\lambda) = \mathbf{h}'(\lambda)\mathbf{x}(\lambda) + \mathbf{h}(\lambda)\mathbf{x}'(\lambda), \quad \mathbf{h}'(\lambda) = [-1, 0, \dots, 0].$$

Hyman's algorithm

We have shown that we can compute $f(\lambda) = \det(H(\lambda))$ and its derivative $f'(\lambda)$ of a Hessenberg matrix H in $\mathcal{O}(n^2)$ operations.

Apply **Newton iteration**:

Choose initial guess λ_0 .

While not converged,

$$\lambda_{k+1} = \lambda_k - \frac{f(\lambda_k)}{f'(\lambda_k)}, \quad k = 0, 1, \dots$$

Note: Higher order derivatives of f can be computed in an analogous fashion. Higher order zero finders are then applicable (e.g. Laguerre's zero finder).

Computing multiple zeros

If we have found a zero z of $f(x) = 0$ and want to compute another one, we want to avoid recomputing the already found z .

We can **explicitly deflate** the zero by defining a new function

$$f_1(x) := \frac{f(x)}{x - z},$$

and apply our method of choice to f_1 . This procedure can in particular be done with polynomials. The coefficients of f_1 are however very sensitive to inaccuracies in z .

We can proceed similarly for multiple zeros z_1, \dots, z_m .

Explicit deflation is not recommended and often not feasible since f is not given explicitly.

Computing multiple zeros (cont.)

For the reciprocal Newton correction for f_1 we get

$$\frac{f_1'(x)}{f_1(x)} = \frac{\frac{f'(x)}{x-z} - \frac{f(x)}{(x-z)^2}}{\frac{f(x)}{x-z}} = \frac{f'(x)}{f(x)} - \frac{1}{x-z}.$$

Then a Newton correction becomes

$$x^{(k+1)} = x_k - \frac{1}{\frac{f'(x_k)}{f(x_k)} - \frac{1}{x_k - z}}$$

and similarly for multiple zeros z_1, \dots, z_m .

The above procedure is called **implicit deflation**. f is not modified.

In this way errors in z are not propagated to f_1

Inverse Iteration

We consider again the nonlinear eigenvalue problem

$$\begin{aligned} A(\lambda) \mathbf{x} &= \mathbf{0}, \\ \mathbf{c}^T \mathbf{x} &= 1, \end{aligned} \tag{4}$$

where \mathbf{c} is some given vector.

For the *linear* eigenvalue problem we have $A(\lambda) = \lambda I - A$.

Solving (4) is equivalent with finding a zero of the nonlinear function $\mathbf{f}(\mathbf{x}, \lambda)$,

$$\mathbf{f}(\mathbf{x}, \lambda) = \begin{pmatrix} A(\lambda) \mathbf{x} \\ \mathbf{c}^T \mathbf{x} - 1 \end{pmatrix} = \begin{pmatrix} \mathbf{0} \\ 0 \end{pmatrix}. \tag{5}$$

Inverse Iteration (cont.)

To apply Newton's zero finding method we need the **Jacobian** of \mathbf{f} ,

$$J(\mathbf{x}, \lambda) \equiv \frac{\partial \mathbf{f}(\mathbf{x}, \lambda)}{\partial (\mathbf{x}, \lambda)} = \begin{pmatrix} A(\lambda) & A'(\lambda)\mathbf{x} \\ \mathbf{c}^T & 0 \end{pmatrix}. \quad (6)$$

Then a step of Newton's iteration is given by

$$\begin{pmatrix} \mathbf{x}_{k+1} \\ \lambda_{k+1} \end{pmatrix} = \begin{pmatrix} \mathbf{x}_k \\ \lambda_k \end{pmatrix} - J(\mathbf{x}_k, \lambda_k)^{-1} \mathbf{f}(\mathbf{x}_k, \lambda_k), \quad (7)$$

or, with the abbreviations $A_k := A(\lambda_k)$ and $A'_k := A'(\lambda_k)$,

$$\begin{pmatrix} A_k & A'_k \mathbf{x}_k \\ \mathbf{c}^T & 0 \end{pmatrix} \begin{pmatrix} \mathbf{x}_{k+1} - \mathbf{x}_k \\ \lambda_{k+1} - \lambda_k \end{pmatrix} = \begin{pmatrix} -A_k \mathbf{x}_k \\ 1 - \mathbf{c}^T \mathbf{x}_k \end{pmatrix}. \quad (8)$$

Inverse Iteration (cont.)

If \mathbf{x}_k is normalized, $\mathbf{c}^T \mathbf{x}_k = 1$, then second equation in (8) yields

$$\mathbf{c}^T (\mathbf{x}_{k+1} - \mathbf{x}_k) = 0 \iff \mathbf{c}^T \mathbf{x}_{k+1} = 1. \quad (9)$$

First equation in (8) gives

$$\begin{aligned} A_k (\mathbf{x}_{k+1} - \mathbf{x}_k) + (\lambda_{k+1} - \lambda_k) A'_k \mathbf{x}_k &= -A_k \mathbf{x}_k \\ \iff A_k \mathbf{x}_{k+1} &= -(\lambda_{k+1} - \lambda_k) A'_k \mathbf{x}_k. \end{aligned}$$

Introduce auxiliary vector \mathbf{u}_{k+1} :

$$A_k \mathbf{u}_{k+1} = A'_k \mathbf{x}_k. \quad (10)$$

\mathbf{u}_{k+1} points in the desired direction; it just needs to be normalized.

Inverse Iteration (cont.)

Normalizing \mathbf{u}_{k+1} gives

$$1 = \mathbf{c}^T \mathbf{x}_{k+1} = -(\lambda_{k+1} - \lambda_k) \mathbf{c}^T \mathbf{u}_{k+1}, \quad (11)$$

or

$$\lambda_{k+1} = \lambda_k - \frac{1}{\mathbf{c}^T \mathbf{u}_{k+1}}. \quad (12)$$

Algorithm: Newton iteration for solving (5)

1: Choose a starting vector $\mathbf{x}_0 \in \mathbb{R}^n$ with $\mathbf{c}^T \mathbf{x}_0 = 1$. $k := 0$.

2: **repeat**

3: Solve $A(\lambda_k) \mathbf{u}_{k+1} := A'(\lambda_k) \mathbf{x}_k$ for \mathbf{u}_{k+1} ; (10)

4: $\mu_k := \mathbf{c}^T \mathbf{u}_{k+1}$;

5: $\mathbf{x}_{k+1} := \mathbf{u}_{k+1} / \mu_k$; (Normalize \mathbf{u}_{k+1})

6: $\lambda_{k+1} := \lambda_k - 1 / \mu_k$; (12)

7: $k := k + 1$;

8: **until** some convergence criterion is satisfied

Note: • For linear eigenvalue problems we have $A'(\lambda)\mathbf{x} = \mathbf{x}$.

• In above algorithm: In each iteration step a linear system has to be solved.

Successive linear approximations

$$A(\lambda)\mathbf{x} \approx (A(\lambda_k) - \vartheta A'(\lambda_k))\mathbf{x} = \mathbf{0}, \quad \lambda = \lambda_k - \vartheta.$$

This suggests the method of **successive linear problems**.

- 1: Start with approximation λ_1 of an eigenvalue of $A(\lambda)$.
- 2: **for** $k = 1, 2, \dots$ **do**
- 3: Solve the linear eigenvalue problem $A(\lambda)\mathbf{u} = \vartheta A'(\lambda)\mathbf{u}$.
- 4: Choose an eigenvalue ϑ smallest in modulus.
- 5: $\lambda_{k+1} := \lambda_k - \vartheta$;
- 6: **end for**

Remark: If A is twice continuously differentiable, and λ is an eigenvalue of problem (1) such that $A'(\lambda)$ is singular and 0 is an algebraically simple eigenvalue of $A'(\lambda)^{-1}A(\lambda)$, then the method in Algorithm 3 converges quadratically towards λ .

Discussion

- ▶ Methods of today can be used to compute a few eigenvalues of small and/or dense matrices.
- ▶ Methods require a factorization of a matrix in each iteration step.
This may lead to excessive flop counts.
- ▶ Hyman's method is designed for Hessenberg matrices.
Transformation of large sparse matrices to Hessenberg form leads to dense matrices. So, it not suited for large sparse matrices.

References

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- [2] A. Ruhe: *Algorithms for the nonlinear eigenvalue problem*.
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- [3] F. Tisseur and K. Meerbergen: *The quadratic eigenvalue problem*.
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Exercise 3

<http://people.inf.ethz.ch/arbenz/ewp/Exercises/exercise03.pdf>