

Reuse

Small example:

Cache: - direct mapped (1-way)
- cacheline = 2 doubles
- 2 cachelines



Array $x = x_{i0}, \dots, x_{i7}$ accessed twice, cold cache

1.) for $j = 0:1$
for $i = 0:7$
access($x[i]$)

8 misses, 8 hits
neighbor use ✓
data reuse 2

access pattern: 012...7012...7

2.) for $j = 0:1$
for $i = 0:2:7$
access($x[i]$)
for $i = 1:2:7$
access($x[i]$)

16 misses, 0 hits
neighbor use 2
data reuse 2

access pattern: 0246135702461357

3.) for $j = 0:1$
for $k = 0:1$
for $i = 0:3$
access($x[i+4j]$)

4 misses, 12 hits
(optimal)
neighbor use ✓
data reuse ✓

access pattern: 0123012345674567

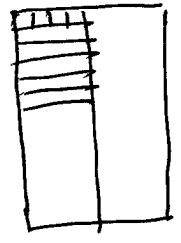
Types of reuse

Data reuse: reuse of data brought into cache before it is evicted

Neighbor use: use of data in the same cacheline of the requested data before it is evicted

Accessing a vector

Cache: - 2-way
- cacheline = 4 doubles

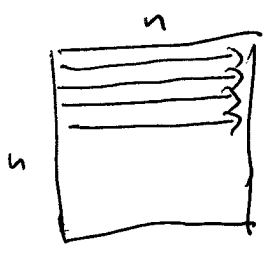


Array: $x_{i0}, \dots, x_{i(n-1)}$ $n \gg$ cache size
accessed once (\Rightarrow no data reuse)

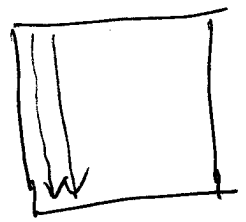
- 1.) sequential (stride = 1) = 0 1 2 ...
- neighbor use, $\frac{n}{4}$ misses
- 2.) stride = 2: 0 2 4 ... (n-2) 1 3 5 ... (n-1)
- some neighbor use, $\frac{n}{2}$ misses
- 3.) stride ≥ 4 : 0 4 8 ... (n-4) 1 5 9 ... (n-3) ...
- no neighbor use, n misses

Where does this occur?

2-D arrays (e.g. images) are stored in row-major order (in C):



Access to column-wise + n divisible by power of 2
 \Rightarrow no neighbor use



Typical pattern in image processing ($n = 256, 512, \dots$)