

How to Write Fast Numerical Code

Spring 2011

Lecture 19

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Schedule

May 2011

Sunday	Monday	Tuesday	Wednesday	Thursday	Friday	Saturday
24	25	26	27	28	29	30
1	2	3	4	5	6	7
8	9	10	11	12	13	14
15	16	17	18	19	20	21
22	23	24	25	26	27	28
29	30	31	1	2	3	4



Today



Lecture



Project meetings



Project presentations

- 10 minutes each
- random order
- random speaker

**Final project paper and code due:
A week or so (exact date still TBD) after semester end**

SIMD Extensions and SSE

- Overview
- SSE family, floating point, and x87
- SSE intrinsics
- *Compiler vectorization*

References:

Intel icc manual (currently 12.0) → Creating parallel applications

→ *Automatic vectorization*

<http://software.intel.com/sites/products/documentation/hpc/composerxe/en-us/cpp/lin/index.htm>

Compiler Vectorization

- Compiler flags
- Aliasing
- Proper code style
- Alignment

Compiler Flags (icc 12.0)

Linux* OS and Mac OS* X	Windows* OS	Description
-vec -no-vec	/Qvec /Qvec-	Enables or disables vectorization and transformations enabled for vectorization. Vectorization is enabled by default. To disable, use -no-vec (Linux* and MacOS* X) or /Qvec- (Windows*) option. Supported on IA-32 and Intel® 64 architectures only.
-vec-report	/Qvec-report	Controls the diagnostic messages from the vectorizer. See Vectorization Report .
-simd -no-simd	/Qsimd /Qsimd-	Controls user-mandated (SIMD) vectorization. User-mandated (SIMD) vectorization is enabled by default. Use the -no-simd (Linux* or MacOS* X) or /Qsimd- (Windows*) option to disable SIMD transformations for vectorization.

Architecture flags:

Linux: -xHost ⊃ -mHost

Windows: /QxHost ⊃ /Qarch:Host

Host in {SSE2, SSE3, SSSE3, SSE4.1, SSE4.2}

Default: -mSSE2, /Qarch:SSE2

How Do I Know the Compiler Vectorized?

- vec-report (previous slide)
- Look at assembly: mulps, addps, xxxps
- Generate assembly with source code annotation:
 - Visual Studio + icc: /Fas
 - icc on Linux/Mac: -S

```

void myadd(float *a, float *b, const int n) {
    for (int i = 0; i < n; i++)
        a[i] = a[i] + b[i];
}

```

Example

unvectorized: /Qvec-

```

<more>
;;;     a[i] = a[i] + b[i];
movss     xmm0, DWORD PTR [rcx+rax*4]
addss     xmm0, DWORD PTR [rdx+rax*4]
movss     DWORD PTR [rcx+rax*4], xmm0
<more>

```

vectorized:

```

<more>
;;;     a[i] = a[i] + b[i];
movss     xmm0, DWORD PTR [rcx+r11*4]
addss     xmm0, DWORD PTR [rdx+r11*4]
movss     DWORD PTR [rcx+r11*4], xmm0
...
movups    xmm0, XMMWORD PTR [rdx+r10*4]
movups    xmm1, XMMWORD PTR [16+rdx+r10*4]
addps     xmm0, XMMWORD PTR [rcx+r10*4]
addps     xmm1, XMMWORD PTR [16+rcx+r10*4]
movaps    XMMWORD PTR [rcx+r10*4], xmm0
movaps    XMMWORD PTR [16+rcx+r10*4], xmm1
<more>

```



why this?



why everything twice?
why movups and movaps?

↑
unaligned

↑
aligned

Aliasing

```
for (i = 0; i < n; i++)
    a[i] = a[i] + b[i];
```

Cannot be vectorized in a straightforward way due to potential aliasing.

However, in this case compiler can insert runtime check:

```
if (a + n < b || b + n < a)
    /* vectorized loop */
    ...
else
    /* serial loop */
    ...
```

Removing Aliasing

- Globally with compiler flag:
 - `-fno-alias`, `/Oa`
 - `-fargument-noalias`, `/Qalias-args-` (function arguments only)
- For one loop: pragma

```
void add(float *a, float *b, int n) {  
    #pragma ivdep  
    for (i = 0; i < n; i++)  
        a[i] = a[i] + b[i];  
}
```

- For specific arrays: restrict (needs compiler flag `-restrict`, `/Qrestrict`)

```
void add(float *restrict a, float *restrict b, int n) {  
    for (i = 0; i < n; i++)  
        a[i] = a[i] + b[i];  
}
```

Proper Code Style

- Use countable loops = number of iterations known at runtime
 - *Number of iterations is a:*
 - constant
 - loop invariant term
 - linear function of outermost loop indices
- Countable or not?

```
for (i = 0; i < n; i++)
    a[i] = a[i] + b[i];
```

```
void vsum(float *a, float *b, float *c) {
    int i = 0;

    while (a[i] > 0.0) {
        a[i] = b[i] * c[i];
        i++;
    }
}
```

Proper Code Style

- Use arrays, structs of arrays, not arrays of structs
- Ideally: unit stride access in innermost loop

```
void mmm1(float *a, float *b, float *c) {  
    int N = 100;  
    int i, j, k;  
  
    for (i = 0; i < N; i++)  
        for (j = 0; j < N; j++)  
            for (k = 0; k < N; k++)  
                c[i][j] = c[i][j] + a[i][k] * b[k][j];  
}
```

```
void mmm2(float *a, float *b, float *c) {  
    int N = 100;  
    int i, j, k;  
  
    for (i = 0; i < N; i++)  
        for (k = 0; k < N; k++)  
            for (j = 0; j < N; j++)  
                c[i][j] = c[i][j] + a[i][k] * b[k][j];  
}
```

Alignment

```
float x[1024];
int i;

for (i = 0; i < 1024; i++)
    x[i] = 1;
```

Cannot be vectorized in a straightforward way since x may not be aligned

However, the compiler can peel the loop to extract aligned part:

```
float x[1024];
int i;

peel = x & 0x0f; /* x mod 16 */
if (peel != 0) {
    peel = 16 - peel;
    /* initial segment */
    for (i = 0; i < peel; i++)
        x[i] = 1;
}
/* 16-byte aligned access */
for (i = peel; i < 1024; i++)
    x[i] = 1;
```

Ensuring Alignment

- Align arrays to 16-byte boundaries (see earlier discussion)
- If compiler cannot analyze:

- Use pragma for loops

```
float x[1024];
int i;

#pragma vector aligned
for (i = 0; i < 1024; i++)
    x[i] = 1;
```

- For specific arrays:
`__assume_aligned(a, 16);`

```

void myadd(float *a, float *b, const int n) {
    for (int i = 0; i < n; i++)
        a[i] = a[i] + b[i];
}

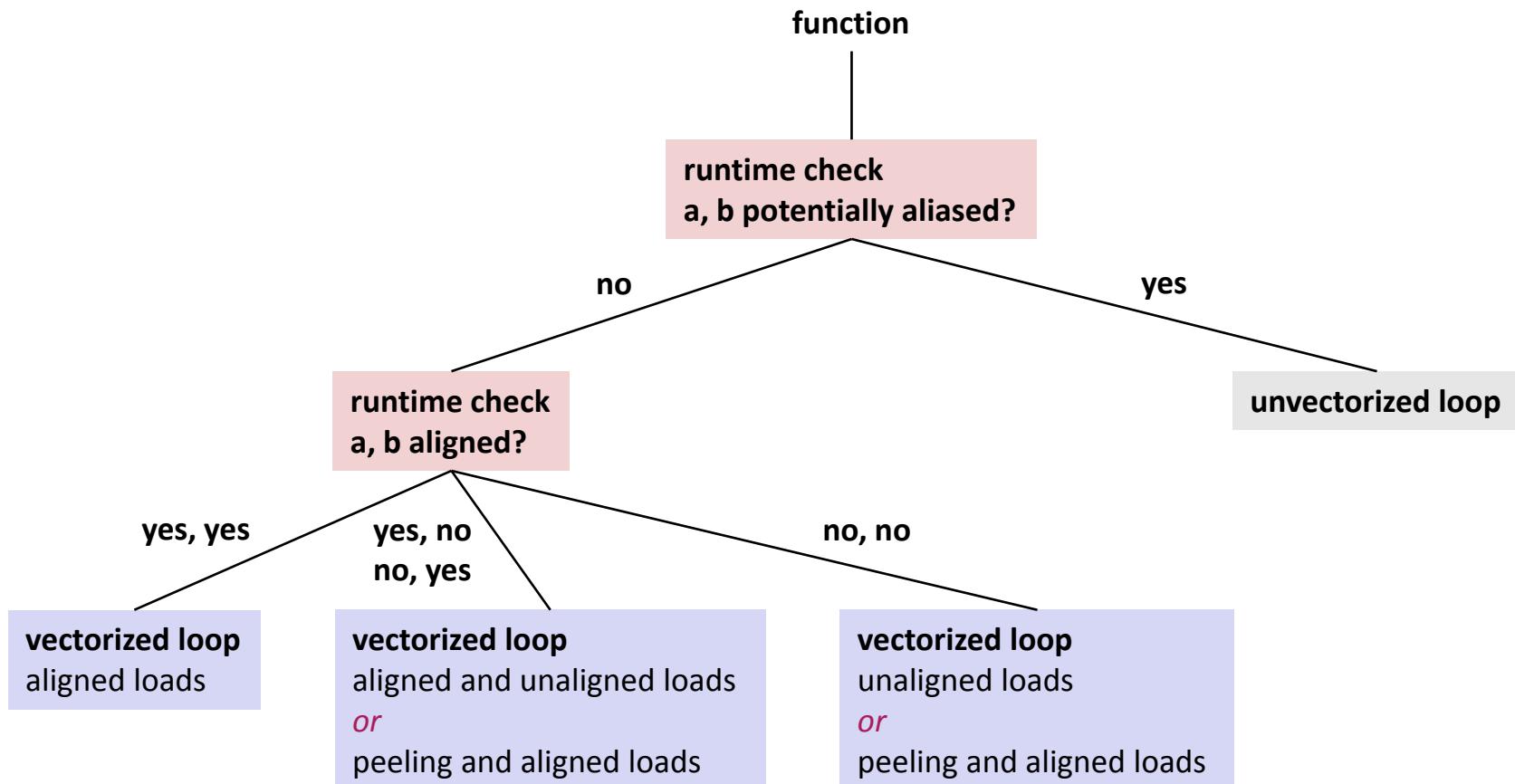
```

Assume:

- No aliasing information
- No alignment information

Can compiler vectorize?

Yes: Through versioning



Compiler Vectorization

- Read manual

Linear Transforms

- Overview
- Discrete Fourier transform
- Fast Fourier transforms
- Optimized implementation and autotuning (FFTW)
- Automatic implementation and optimization (Spiral)

Linear Transforms

- Very important class of functions: signal processing, scientific computing, ...
- *Mathematically:* Change of basis = Multiplication by a fixed matrix T

$$\begin{pmatrix} y_0 \\ y_1 \\ \vdots \\ y_{n-1} \end{pmatrix} = y = Tx \quad \xleftarrow{\hspace{1cm}} \quad T \cdot \quad \xleftarrow{\hspace{1cm}} \quad x = \begin{pmatrix} x_0 \\ x_1 \\ \vdots \\ x_{n-1} \end{pmatrix}$$

$T = [t_{k,\ell}]_{0 \leq k,\ell < n}$

Output

Input

- Equivalent definition: Summation form

$$y_k = \sum_{\ell=0}^{n-1} t_{k,\ell} x_\ell, \quad 0 \leq k < n$$

Smallest Relevant Example

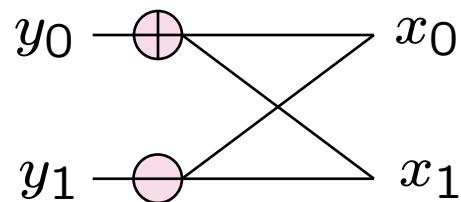
Transform (matrix): $T = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$

Computation: $y = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} x$

or

$$\begin{aligned} y_0 &= x_0 + x_1 \\ y_1 &= x_0 - x_1 \end{aligned}$$

As graph (direct acyclic graph or DAG):



called a butterfly



[http://charlottesmartypants.blogspot.com/
2011_02_01_archive.html](http://charlottesmartypants.blogspot.com/2011_02_01_archive.html)

Transforms: Examples

- A few dozen transforms are relevant
- Some examples

$$\mathbf{DFT}_n = [e^{-2k\ell\pi i/n}]_{0 \leq k, \ell < n}$$

$$\mathbf{RDFT}_n = [r_{k\ell}]_{0 \leq k, \ell < n}, \quad r_{k\ell} = \begin{cases} \cos \frac{2\pi k\ell}{n}, & k \leq \lfloor \frac{n}{2} \rfloor \\ -\sin \frac{2\pi k\ell}{n}, & k > \lfloor \frac{n}{2} \rfloor \end{cases}$$

universal tool

$$\mathbf{DHT} = [\cos(2k\ell\pi/n) + \sin(2k\ell\pi/n)]_{0 \leq k, \ell < n}$$

$$\mathbf{WHT}_n = \begin{bmatrix} \mathbf{WHT}_{n/2} & \mathbf{WHT}_{n/2} \\ \mathbf{WHT}_{n/2} & -\mathbf{WHT}_{n/2} \end{bmatrix}, \quad \mathbf{WHT}_2 = \mathbf{DFT}_2$$

$$\mathbf{IMDCT}_n = [\cos((2k+1)(2\ell+1+n)\pi/4n)]_{0 \leq k < 2n, 0 \leq \ell < n}$$

MPEG

$$\mathbf{DCT-2}_n = [\cos(k(2\ell+1)\pi/2n)]_{0 \leq k, \ell < n}$$

JPEG

$$\mathbf{DCT-3}_n = \mathbf{DCT-2}_n^T \quad (\text{transpose})$$

$$\mathbf{DCT-4}_n = [\cos((2k+1)(2\ell+1)\pi/4n)]_{0 \leq k, \ell < n}$$

Blackboard

- Discrete Fourier transform (DFT)
- Transform algorithms
- Fast Fourier transform, size 4