# A Simple Parallel/Distributed Sampling Technique: Local Glauber Dynamics

### Manuela Fischer

ETH Zurich, Switzerland

joint work with Mohsen Ghaffari Sampling Proper Colorings





Markov chain

- over set of proper colorings
- uniform distribution as unique stationary distribution
- rapidly mixing



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t = 2

Markov chain

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t = 3

Markov chain

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t = 4

Markov chain

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 $t \geq t_{mix}$ 

update color of a random node to a random color, if proper

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[Jerrum 1995] Single-Site Glauber

#### **O**(**n log n**) steps

 $q \ge 2\Delta + 1$ 

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#### $O(n\log n)$ steps

 $q \ge 2\Delta + 1$ 

**Decentralized** 

[Jerrum 1995] Single-Site Glauber

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 $q \geq 2\Delta + 1$ 

### **Decentralized**

# [Feng, Sun, Yi 2017] What can be sampled locally?

local/decentralized sampling techniques? local/decentralized transition rules for Markov chain?

[Jerrum 1995] Single-Site Glauber

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 $q \ge 2\Delta + 1$ 

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[Feng, Sun, Yi 2017] LubyGlauber

 $O(\Delta \log n)$  steps

 $q \geq lpha \Delta$  for lpha > 2

[Jerrum 1995] update a single node

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[Feng, Sun, Yi 2017] $O(\Delta \log n)$  steps $q \ge \alpha \Delta$  for  $\alpha > 2$ [Feng, Sun, Yi 2017] $O(\log n)$  steps $q \ge \alpha \Delta$  for  $\alpha > 2 + \sqrt{2}$ 

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[Feng, Sun, Yi 2017] update an independent set	$O(\Delta \log n)$ steps	$q\geq lpha \Delta$ for $lpha>2$
[Feng, Sun, Yi 2017] <b>update all nodes</b>	<b>O</b> ( <b>log n</b> ) steps	$q\geq lpha \Delta$ for $lpha>2+\sqrt{2}$
[F., Ghaffari 2018] <b>Local Glauber</b>	<b>O</b> ( <b>log n</b> ) steps	$q\geq lpha \Delta$ for $lpha>2$

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[Feng, Sun, Yi 2017] <b>update all nodes</b>	<b>O</b> (log n) steps	$q\geq lpha \Delta$ for $lpha>2+\sqrt{2}$
[F., Ghaffari 2018] <b>update an almost independent set</b>	<b>O</b> ( <b>log</b> <i>n</i> ) steps	$q\geq lpha \Delta$ for $lpha>2$

**Local Glauber Dynamics**
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[*F., Ghaffari 2018*] **Theorem** Local Glauber converges to uniform distribution over proper q-colorings in  $O(\log n)$  steps if  $q \ge \alpha \Delta$  for  $\alpha > 2$ . **Proof Sketch** 

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### Improved Path Coupling for Single-Site Glauber Dynamics

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to have different colors, node needs to be on a **red-blue**-path starting from  $v_0$ 





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[Feng, Sun, Yi 2017] **LubyGlauber** 

[Feng, Sun, Yi 2017] **LocalMetropolis** 

[F., Ghaffari 2018] Local Glauber  $O(\Delta \log n)$  steps  $q \ge \alpha \Delta$  for  $\alpha > 2$  $O(\log n)$  steps  $q \ge \alpha \Delta$  for  $\alpha > 2 + \sqrt{2}$ 

almost independent set suffices $O(\log n)$  steps $q \ge \alpha \Delta$  for  $\alpha > 2$ 

#### **Centralized**

[F., Ghaffari 2018] <b>Local Glauber</b>	almost independent set suffices $O(\log n)$ steps $q \ge \alpha \Delta$ for $\alpha > 2$	
[Feng, Sun, Yi 2017] <b>LocalMetropolis</b>	<b>O</b> ( <b>log n</b> ) steps	$q\geq lpha \Delta$ for $lpha>2+\sqrt{2}$
Decentralized [Feng, Sun, Yi 2017] LubyGlauber	<b>Ο</b> ( <b>Δ log n</b> ) steps	fewer colors? $q \ge \alpha \Delta$ for $\alpha > 2$
[Vigoda 2000], [Chen, Moitra 2018], [Delcourt, Perarnau, Postle 2018]	<b>poly</b> ( <i>n</i> ) steps	$q\geq lpha \Delta$ for $lpha>rac{11}{6}$
[Jerrum 1995] Single-Site Glauber	<b>O</b> ( <b>n log n</b> ) steps	$q \geq 2\Delta + 1$

### **<u>Centralized</u>**

Centralized		Thank you!
[Jerrum 1995]		
Single-Site Glauber	$O(n\log n)$ steps	$q \geq 2\Delta + 1$
[Vigoda 2000],		
[Chen, Moitra 2018],	<b>poly</b> ( <b>n</b> ) steps	$q\geq lpha\Delta$ for $lpha>rac{11}{6}$
[Delcourt, Perarnau, Postle 2018]		0
<b>Decentralized</b>		
[Fena, Sun, Yi 2017]		fewer colors?
LubyGlauber	$O(\Delta \log n)$ steps	$q\geq lpha\Delta$ for $lpha>2$
[Feng, Sun, Yi 2017]		
Local Metropolis	<b>O</b> (log n) steps	$q \geq lpha \Delta$ for $lpha > 2 + \sqrt{2}$
[F., Ghaffari 2018]	almost independent set suffices	
Local Glauber	<b>O</b> (log n) steps	$q\geq lpha\Delta$ for $lpha>2$