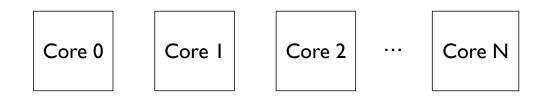
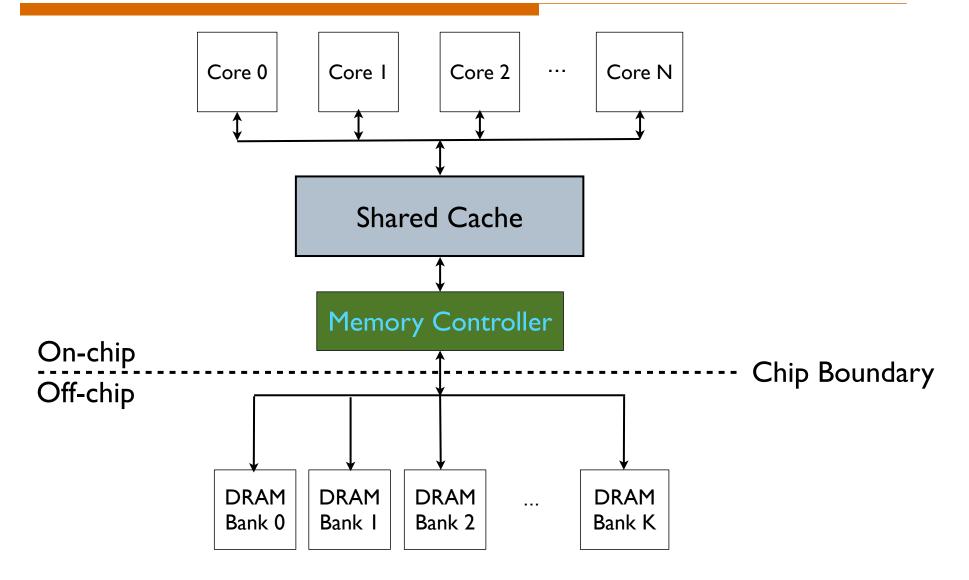
#### Fairness via Source Throttling:

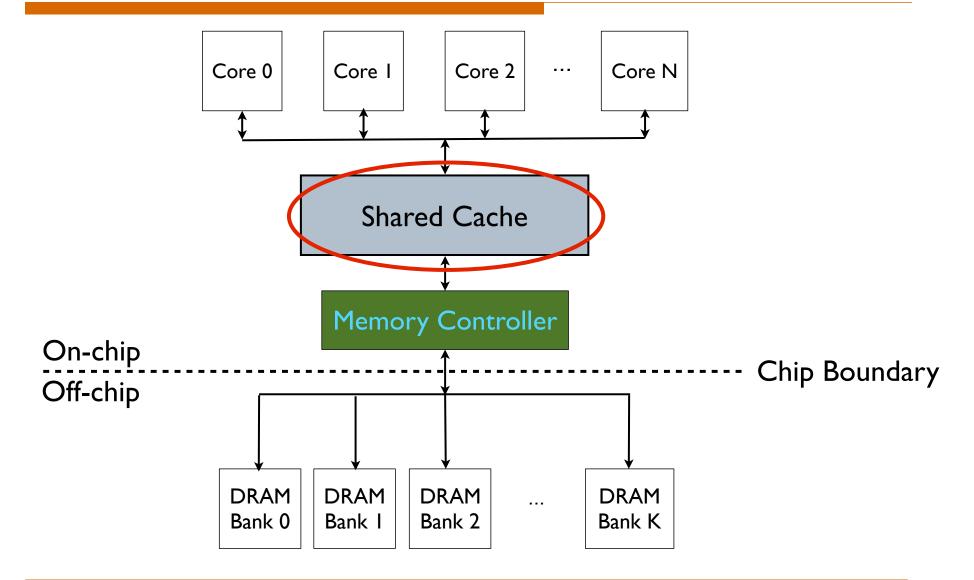
A configurable and high-performance fairness substrate for multi-core memory systems

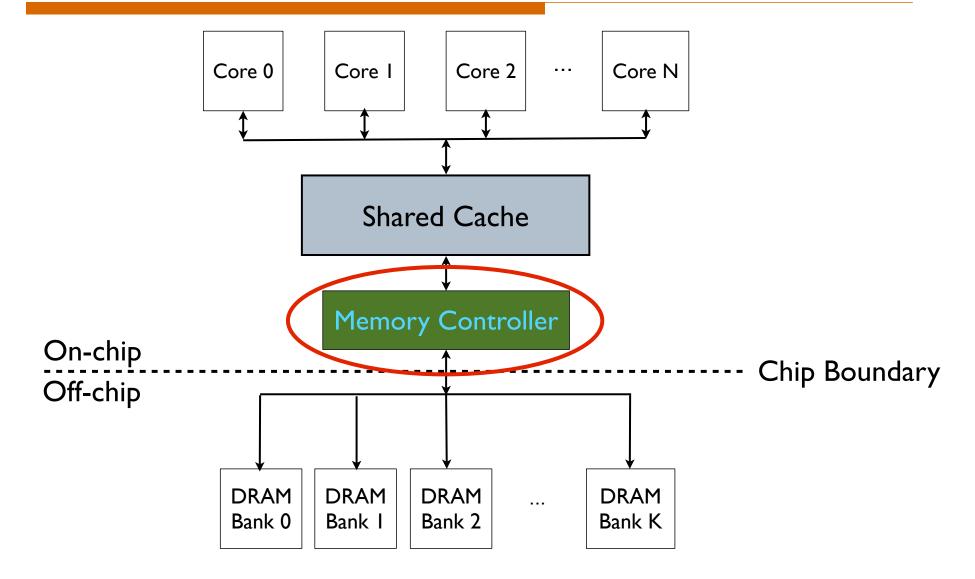
> Eiman Ebrahimi\* Chang Joo Lee\* Onur Mutlu<sup>‡</sup> Yale N. Patt\*

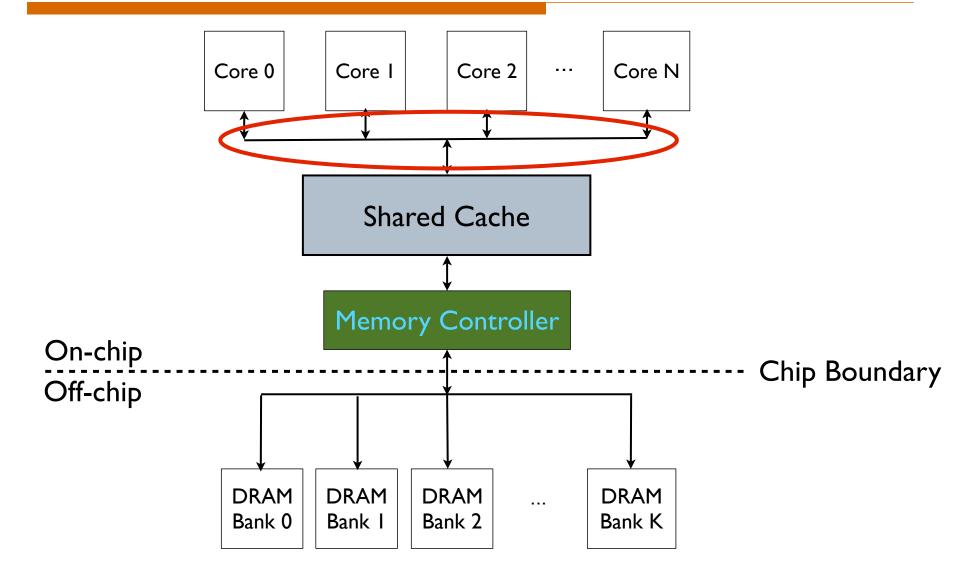
\* HPS Research Group The University of Texas at Austin **‡** Computer Architecture Laboratory Carnegie Mellon University

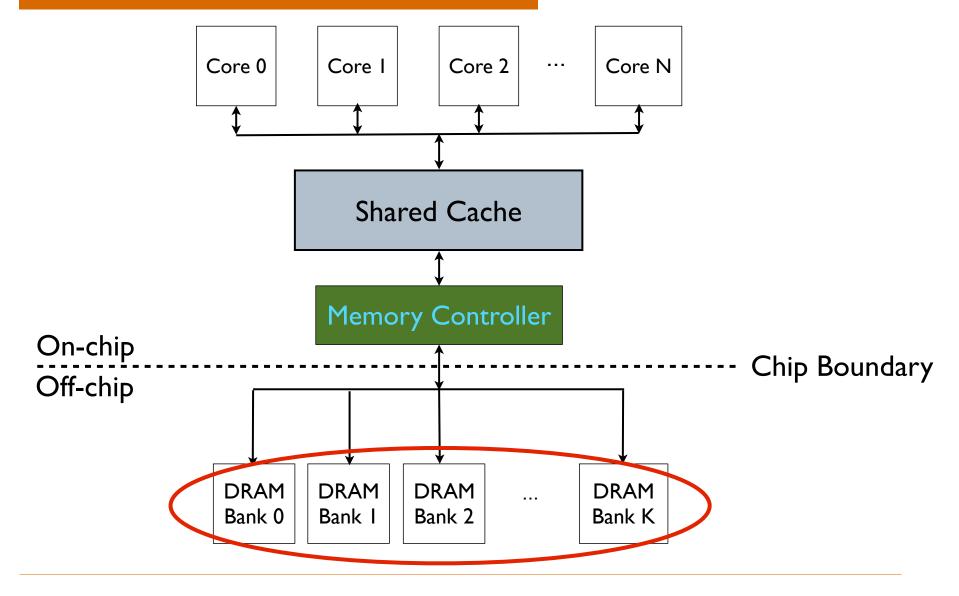


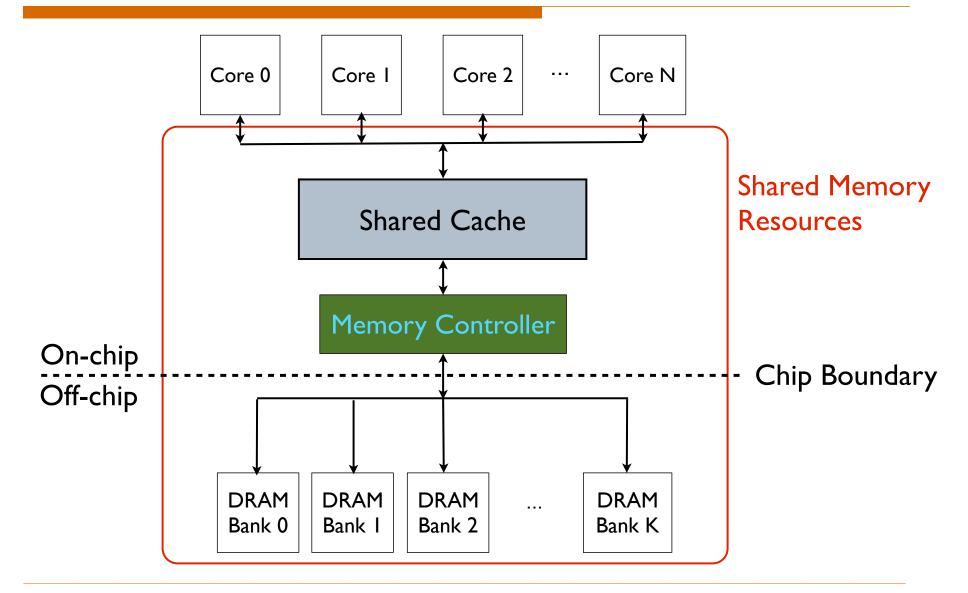












 Applications slow down due to interference from memory requests of other applications

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- A memory system is fair if slowdowns of same-priority applications are equal (MICRO '06, MICRO '07, ISCA '08)

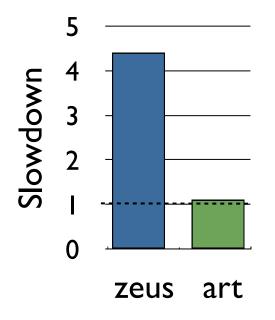
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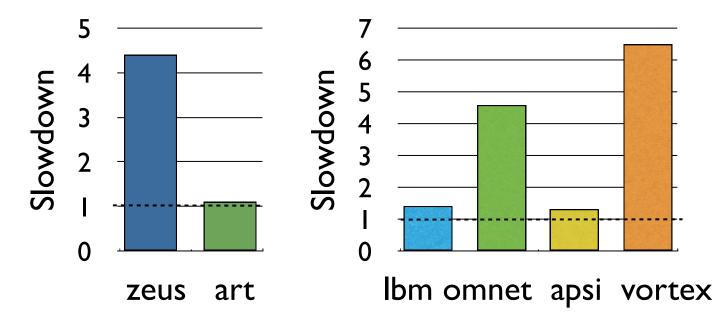
• Slowdown of application 
$$i = \frac{T_i^{\text{Shared}}}{T_i^{\text{Alone}}}$$

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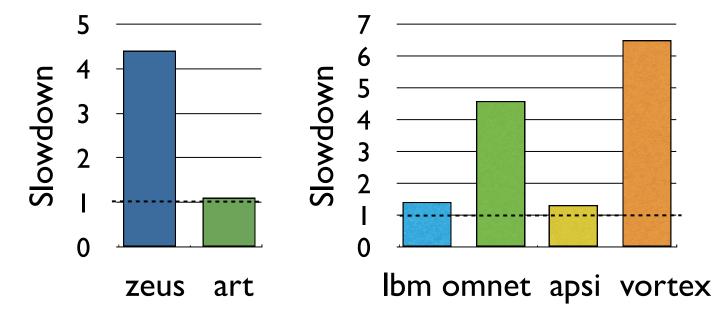
• Slowdown of application 
$$i = \frac{T_i^{\text{Shared}}}{T_i^{\text{Alone}}}$$

Unfairness = Max{Slowdown i} over all applications i
(MICRO '07)

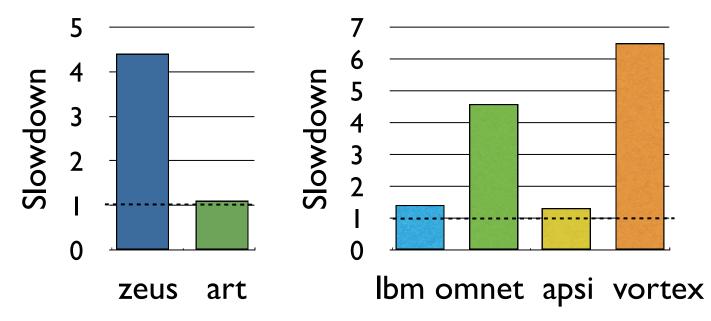




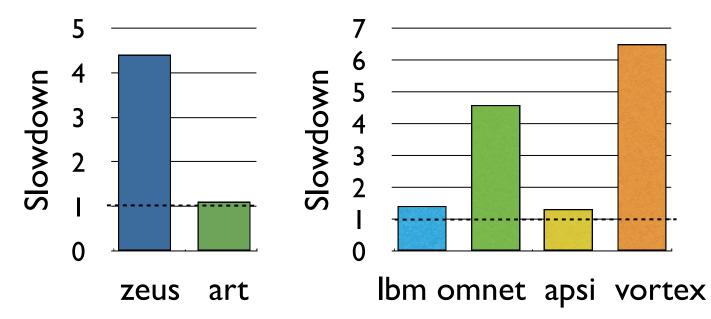
 Magnitude of each application's slowdown depends on concurrently running applications' memory behavior



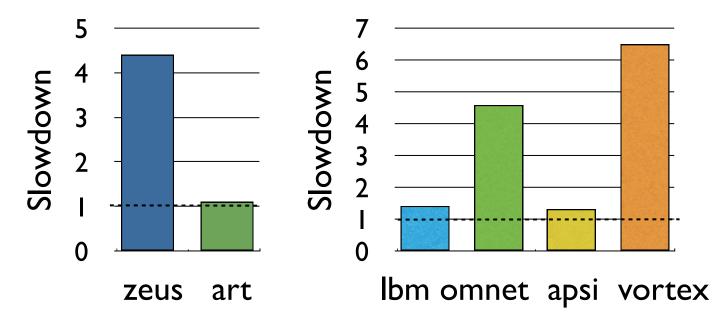
• Large disparities in slowdowns are unacceptable



- Large disparities in slowdowns are unacceptable
  - Low system performance



- Large disparities in slowdowns are unacceptable
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  - Vulnerability to denial of service attacks



- Large disparities in slowdowns are unacceptable
  - Low system performance
  - Vulnerability to denial of service attacks
  - Difficult for system software to enforce priorities

## Outline

- Background and Problem
- Motivation for Source Throttling
- Fairness via Source Throttling (FST)
- Evaluation
- Conclusion

- Primarily manage inter-application interference in only one particular resource
  - Shared Cache, Memory Controller, Interconnect, etc.

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Our Goal: Enable fair sharing of the entire memory system by dynamically detecting and controlling interference in a coordinated manner

 Manage inter-application interference at the cores, not at the shared resources

- Manage inter-application interference at the cores, not at the shared resources
- Dynamically estimate unfairness in the memory system

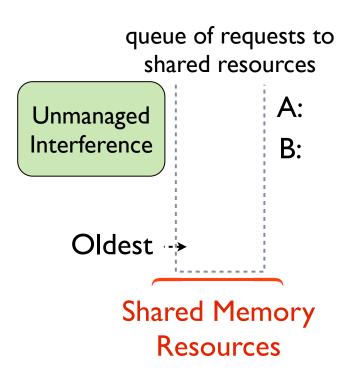
- Manage inter-application interference at the cores, not at the shared resources
- Dynamically estimate unfairness in the memory system
- If unfairness > system-software-specified target then throttle down core causing unfairness & throttle up core that was unfairly treated



A: B:

Fair Source Throttling A: B:

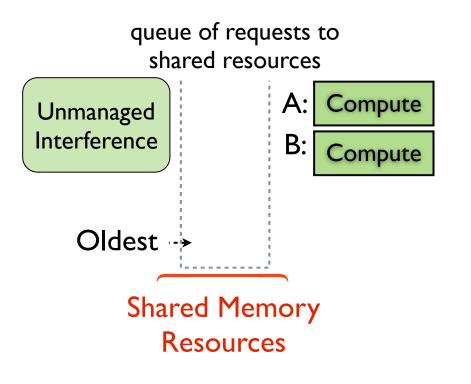
Wednesday, March 17, 2010



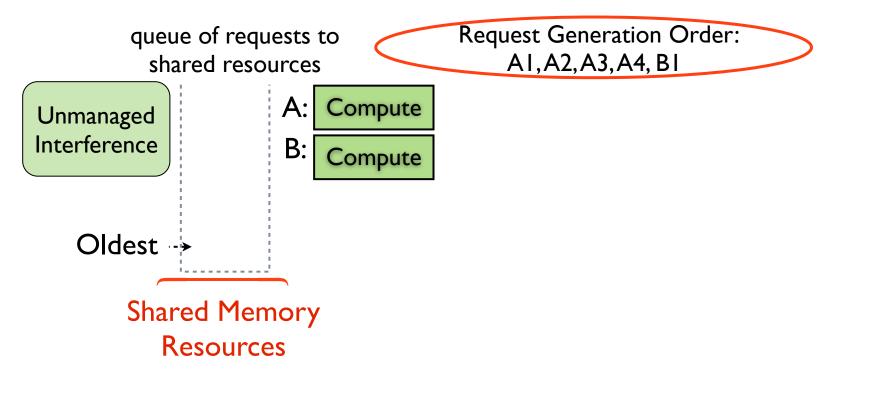


B:

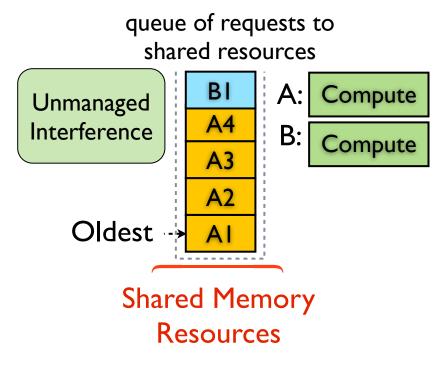
A:











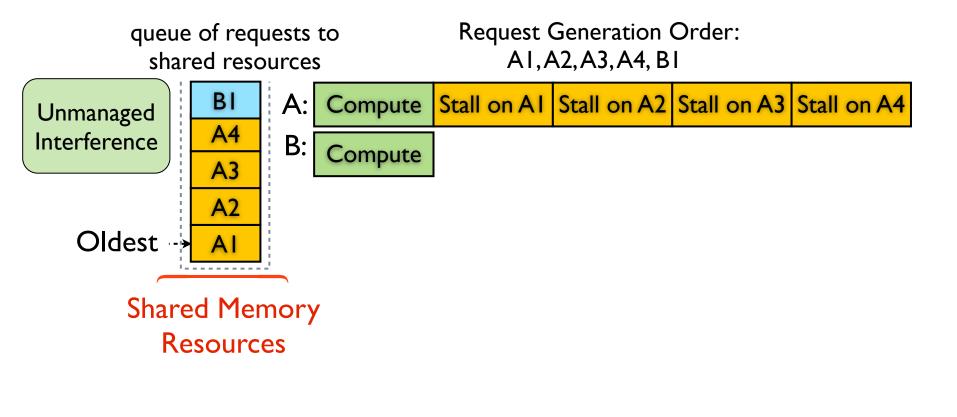
Fair Source Throttling



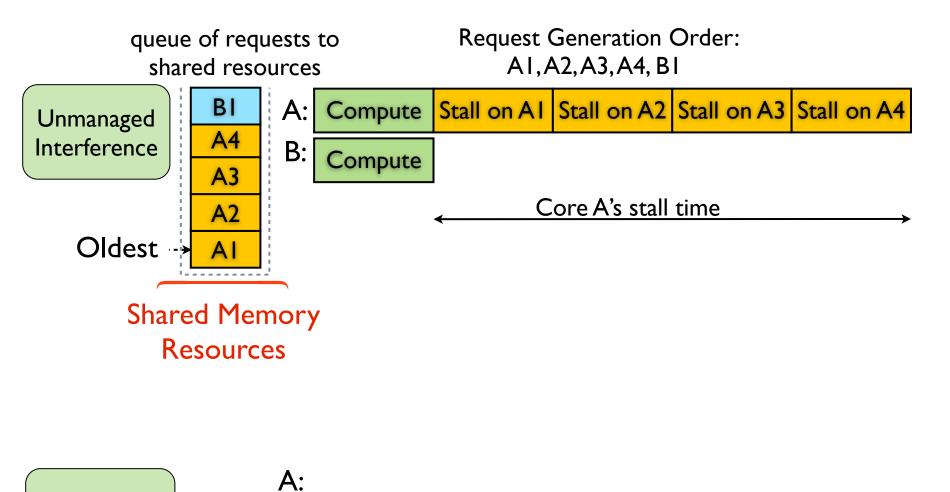
Request Generation Order:

AI, A2, A3, A4, BI

Wednesday, March 17, 2010

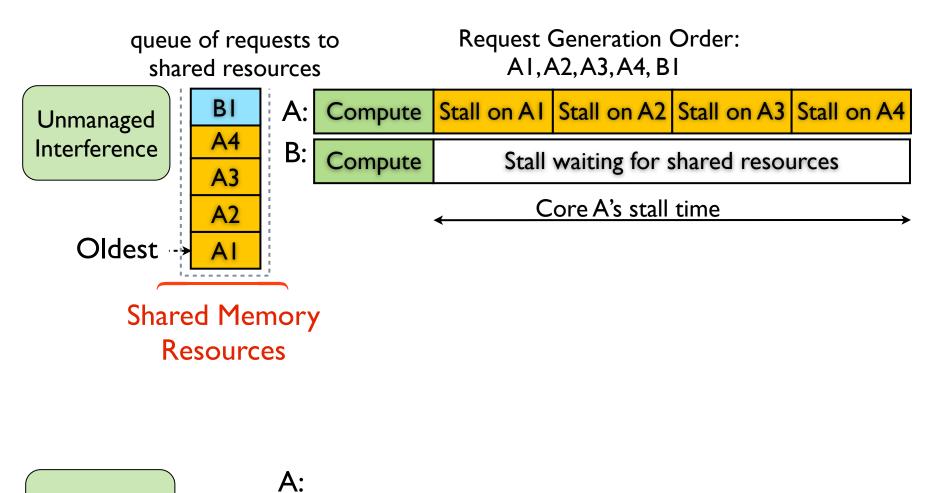






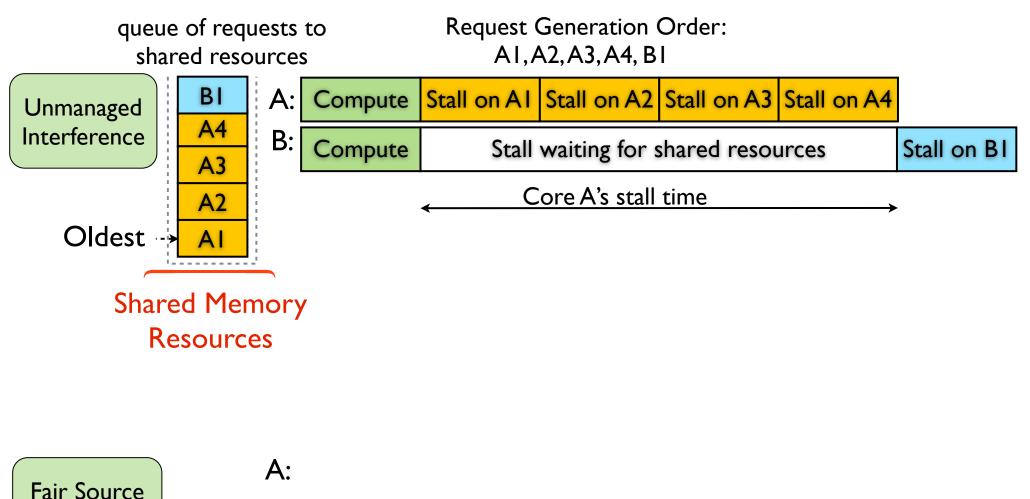






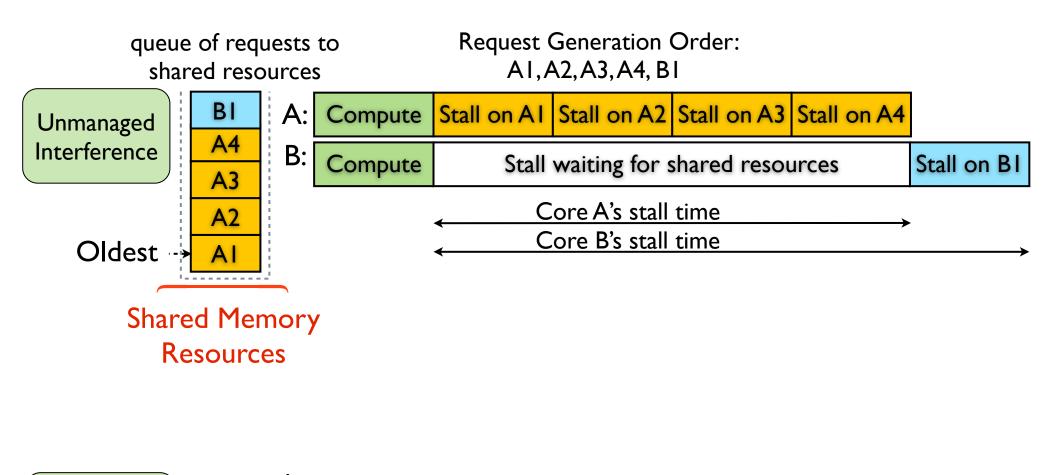
Fair Source Throttling





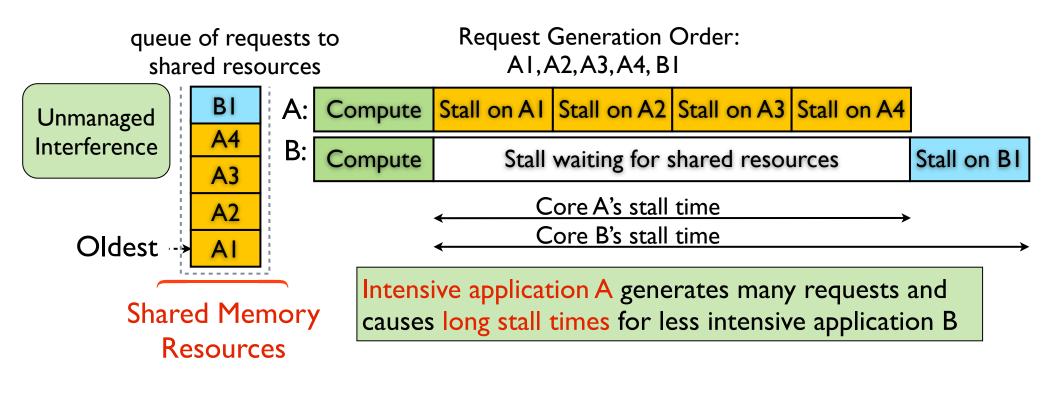
Throttling

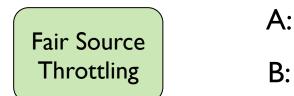


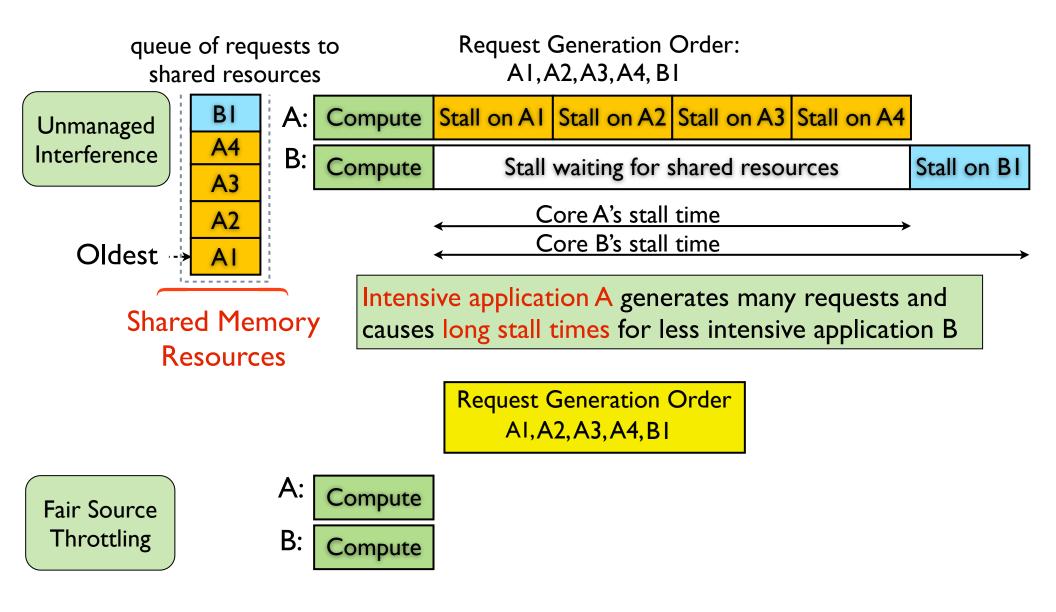


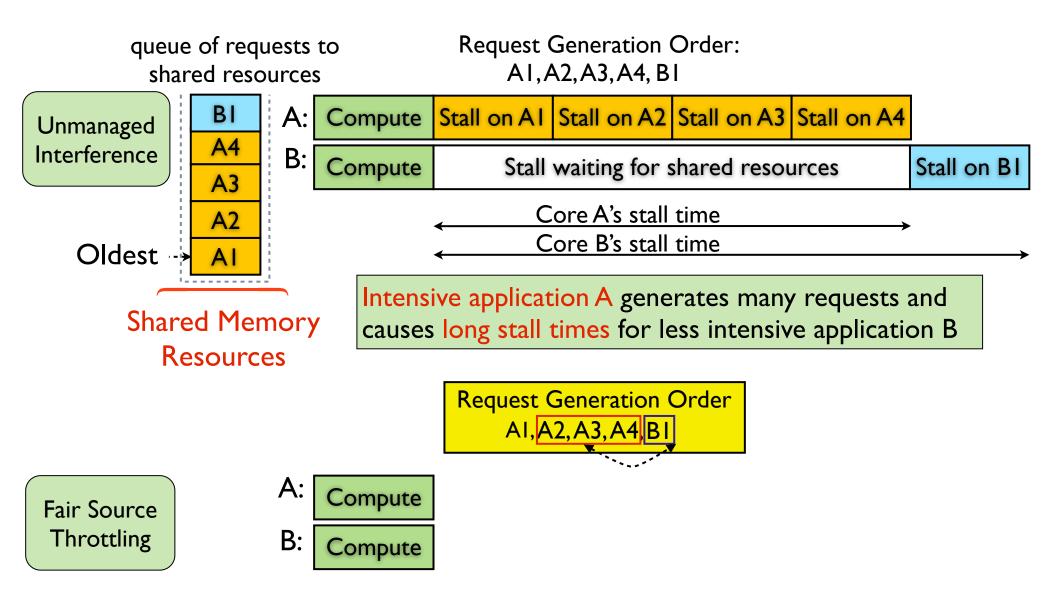


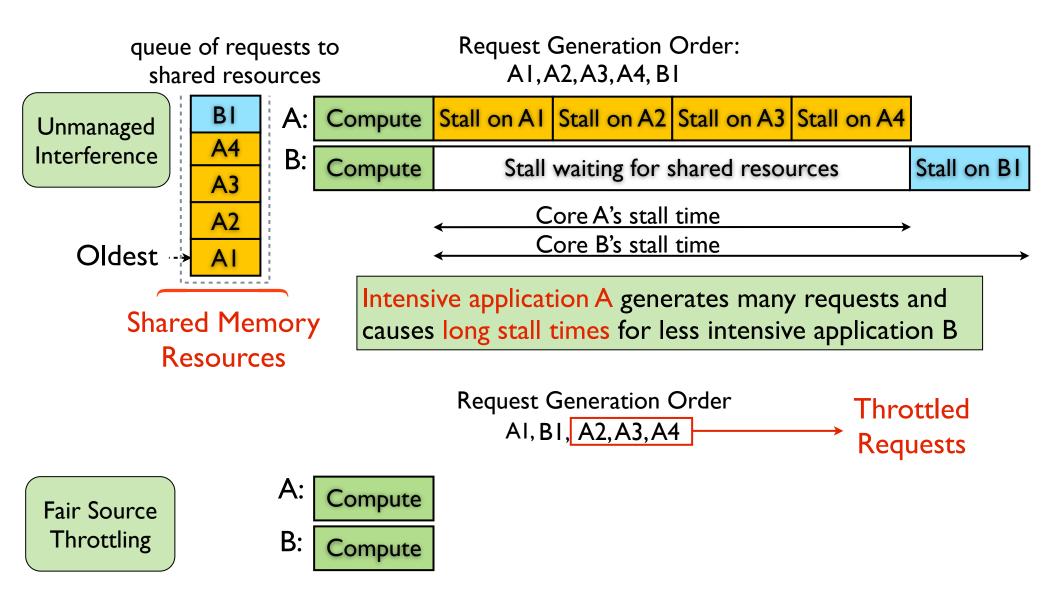
**B**:

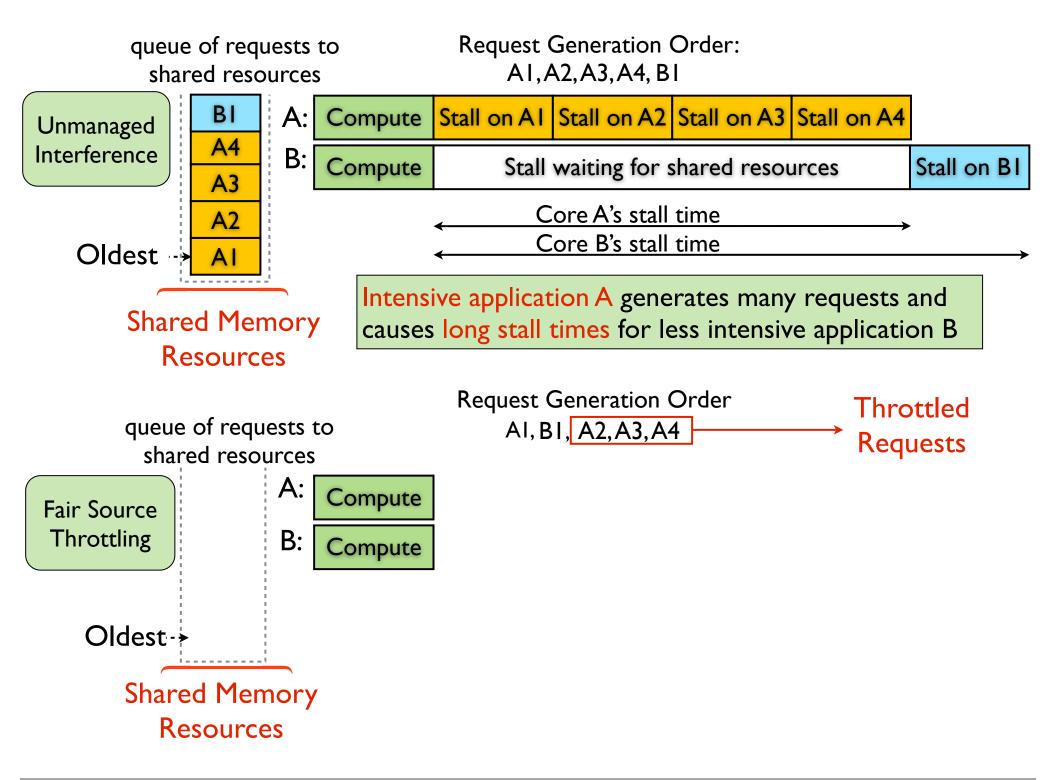


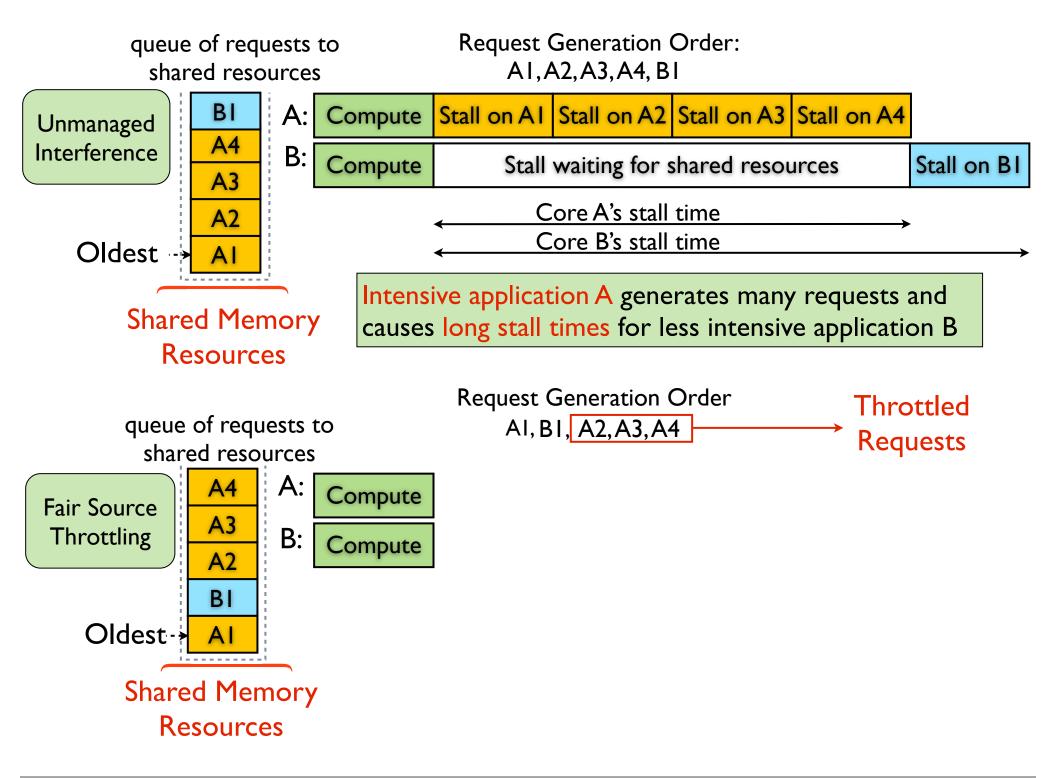


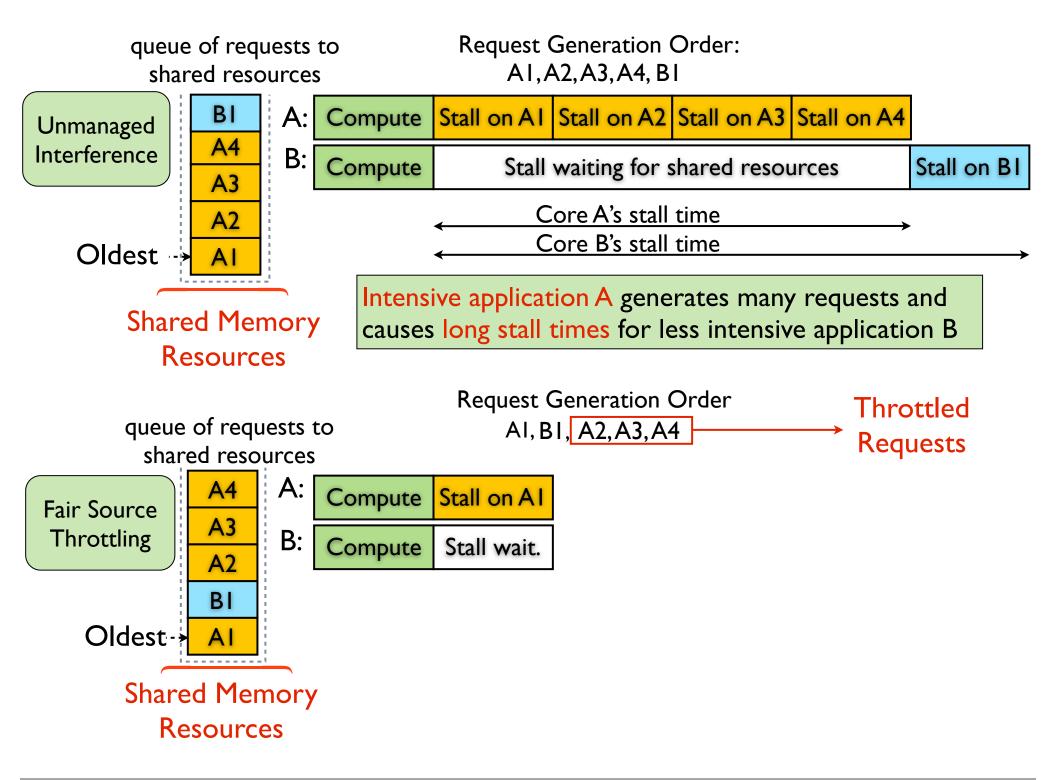


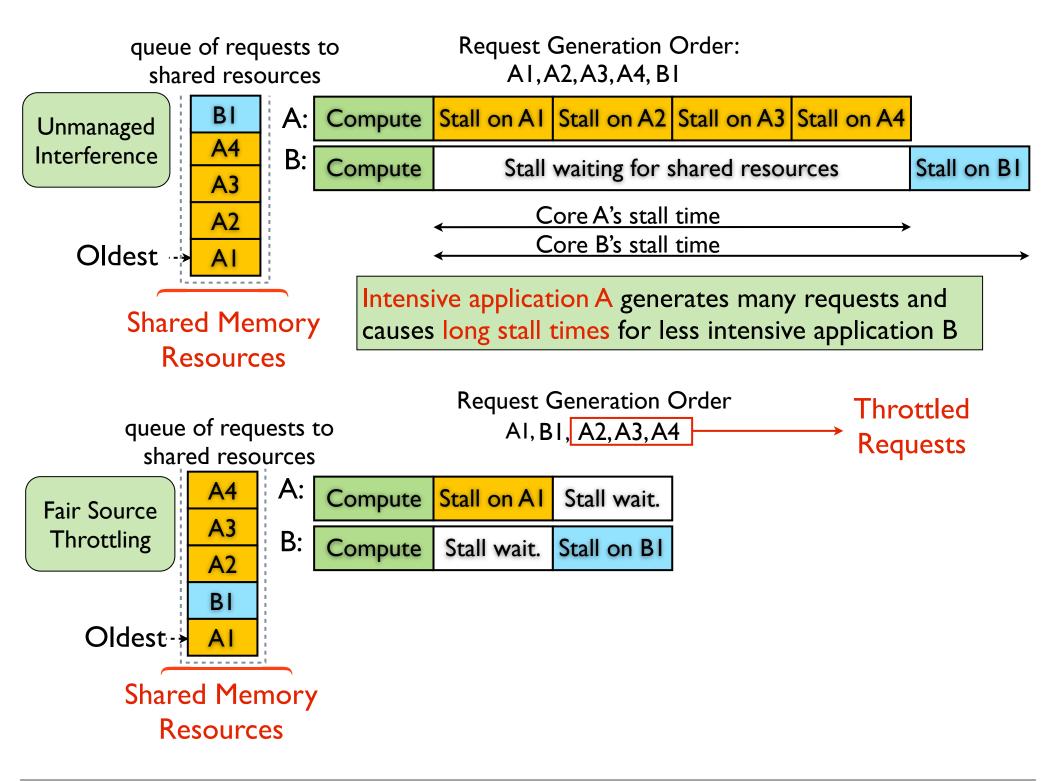


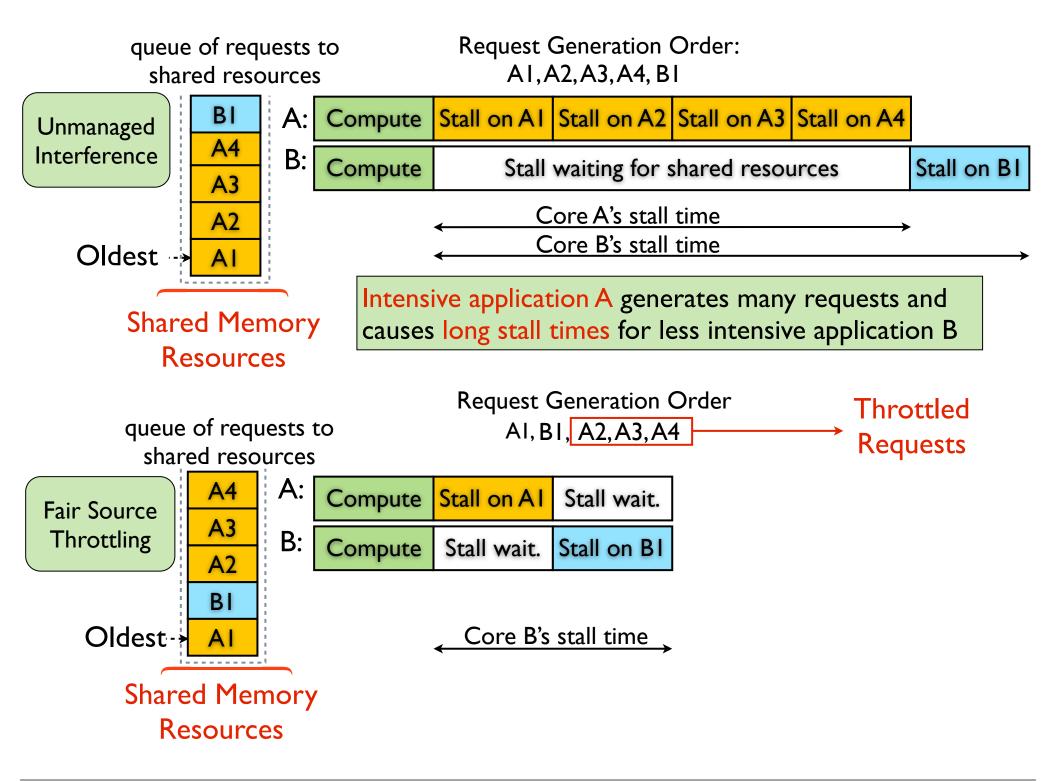


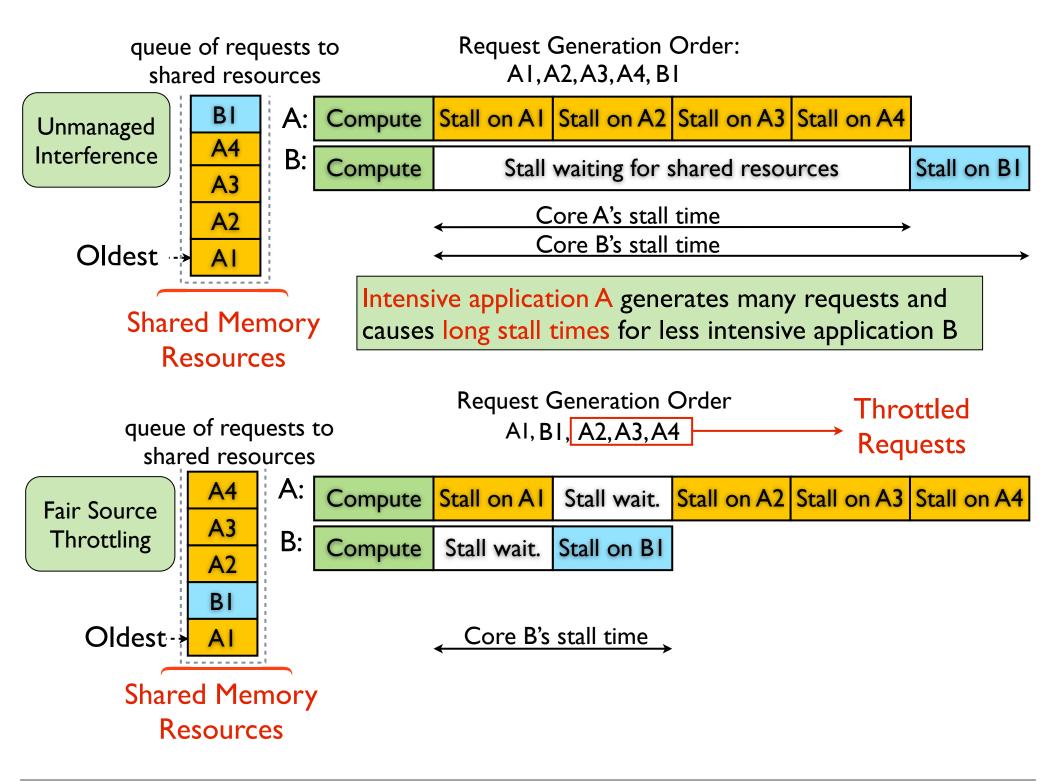


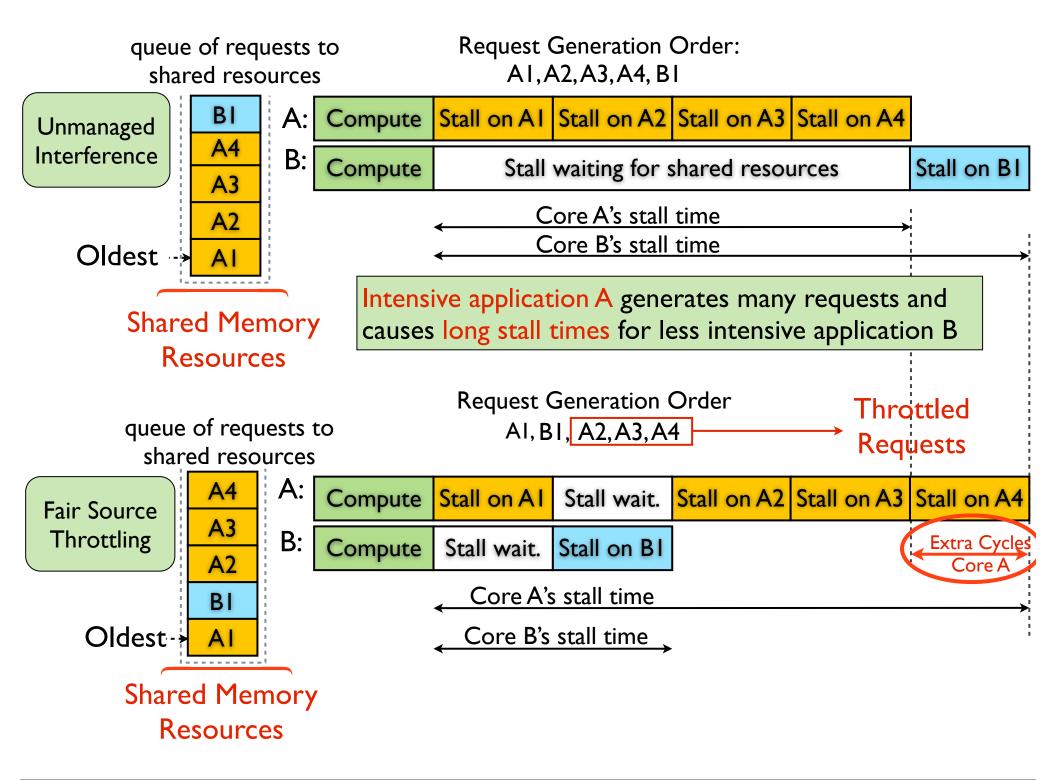


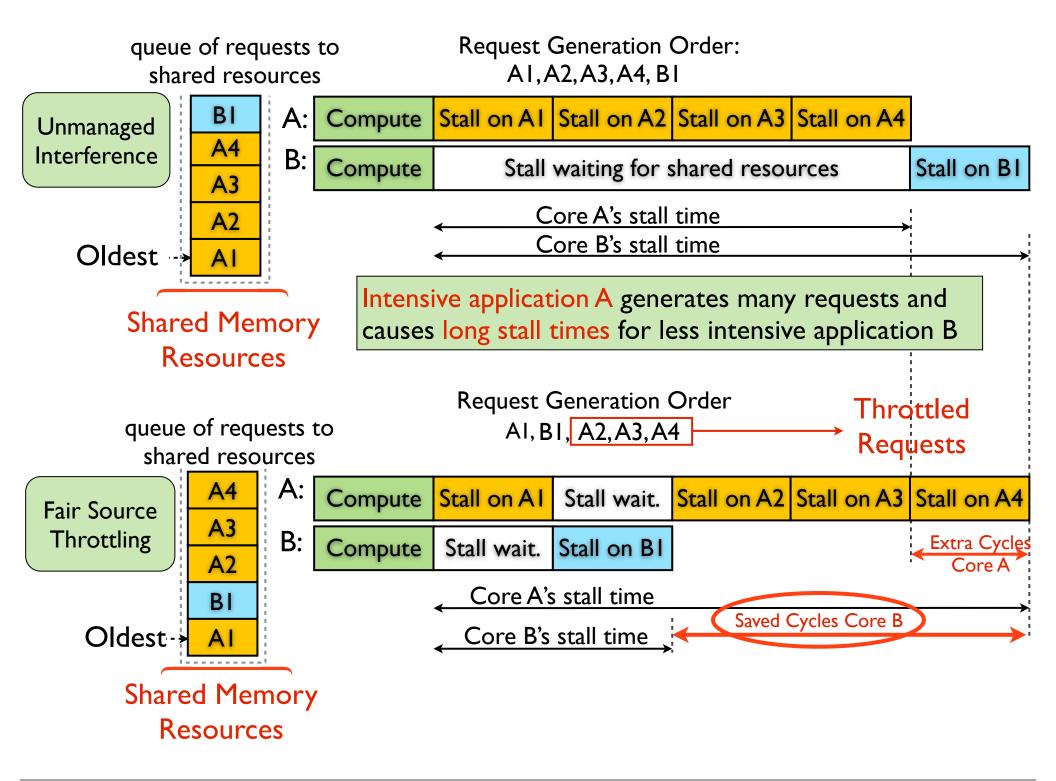


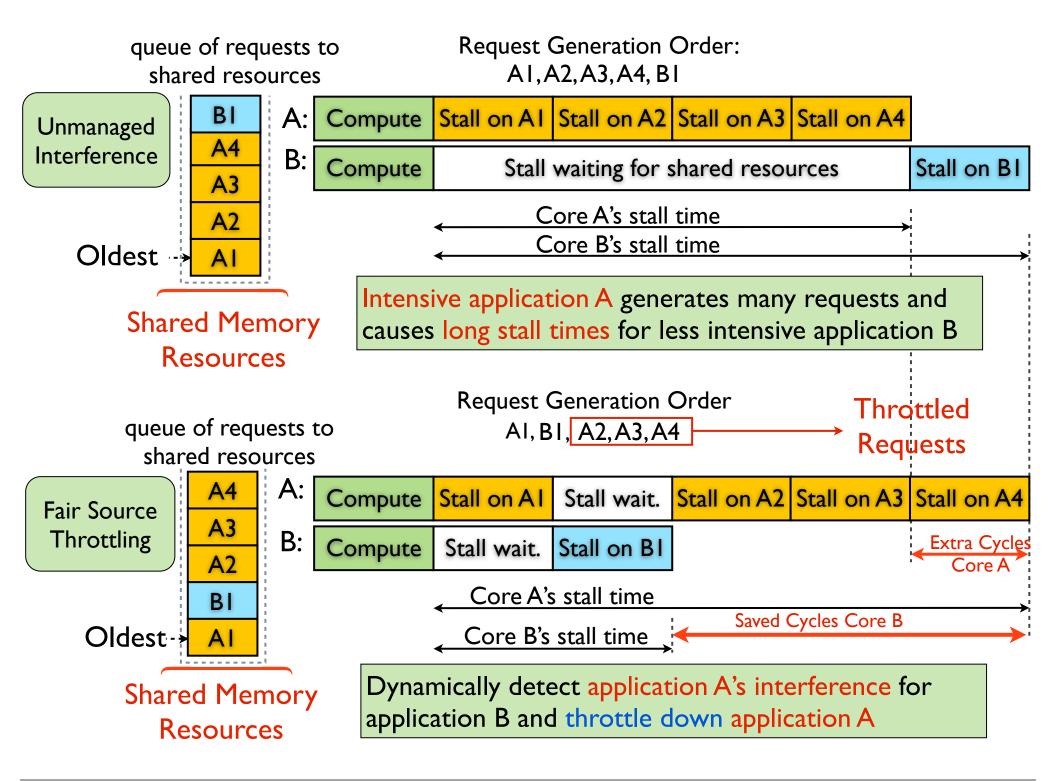










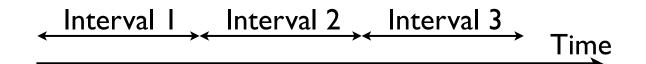


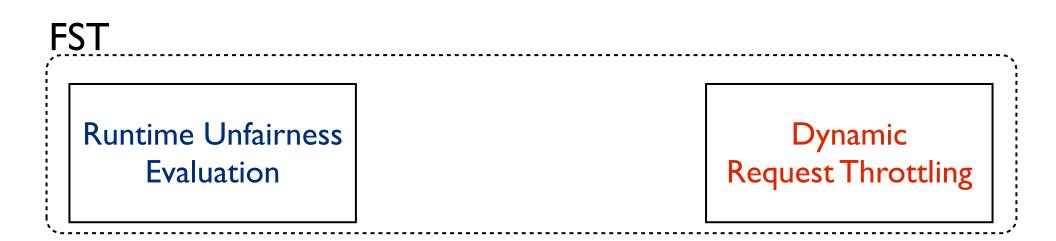
## Outline

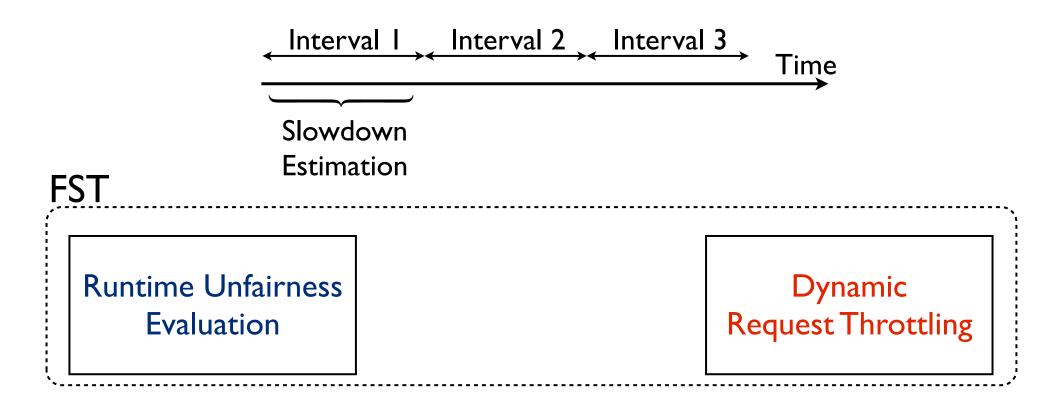
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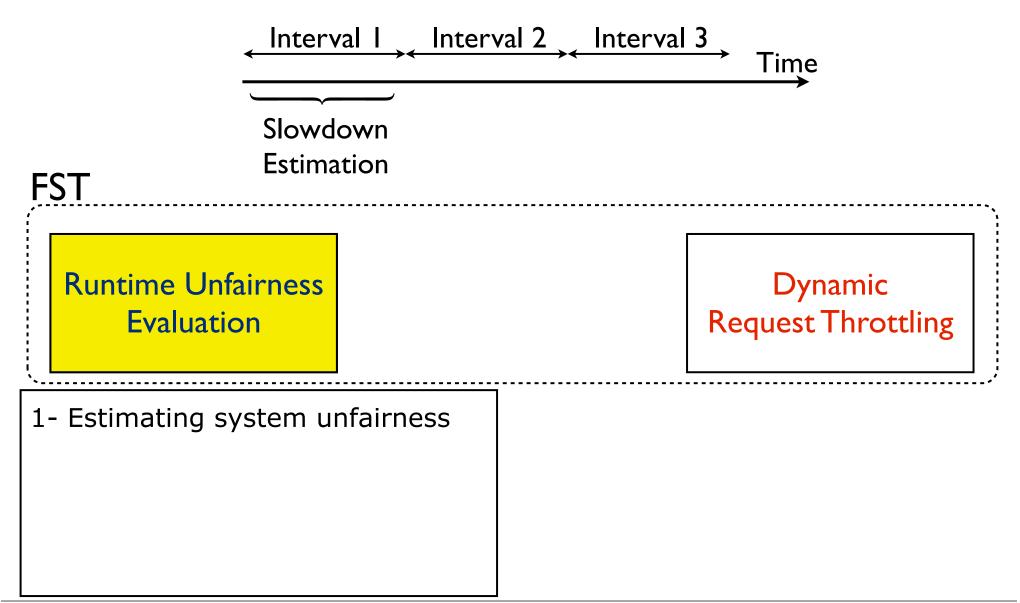
- Runtime Unfairness Evaluation
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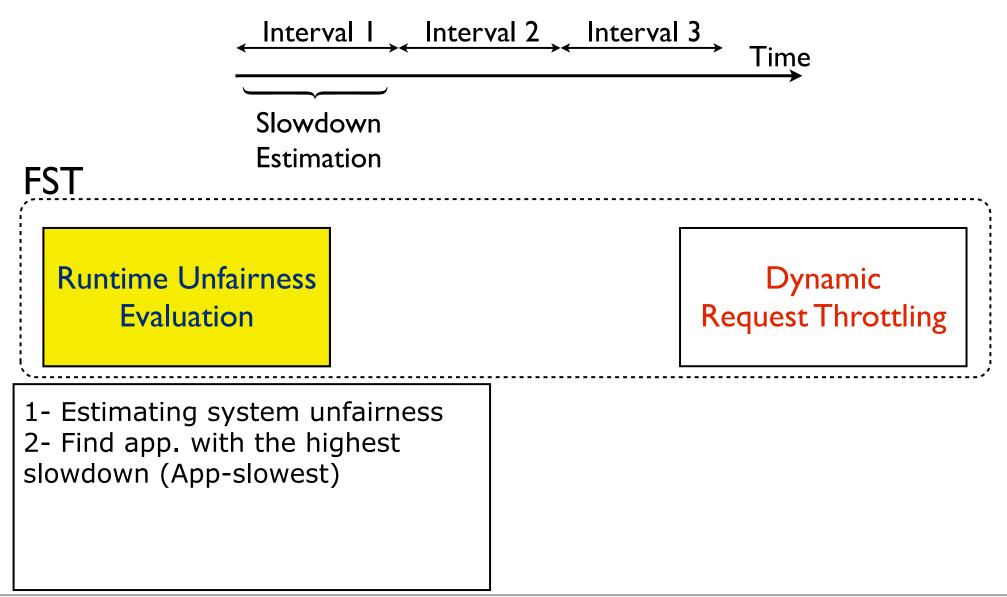
- Runtime Unfairness Evaluation
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- Dynamic Request Throttling
  - Adjusts how aggressively each core makes requests to the shared resources

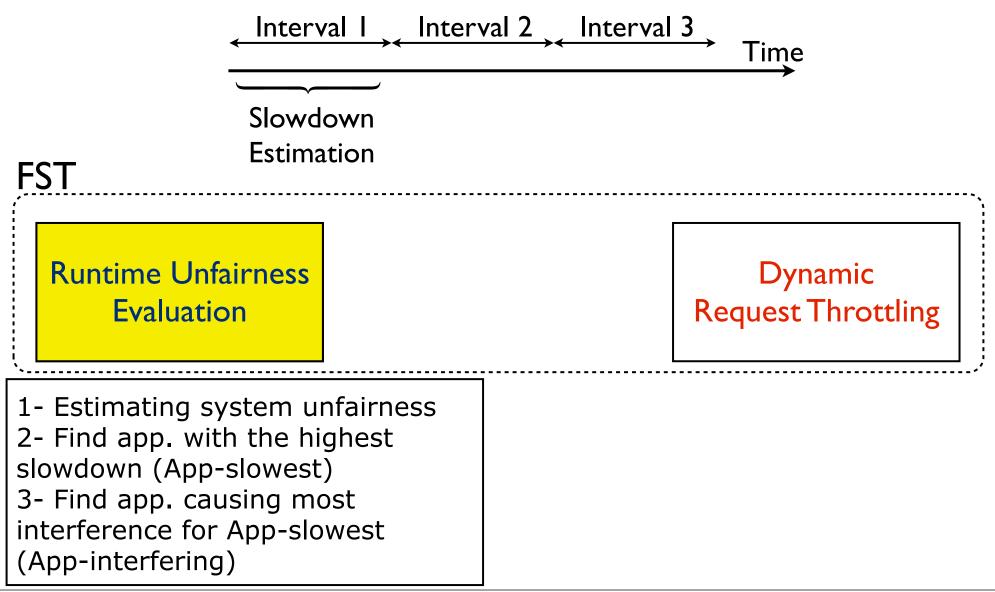


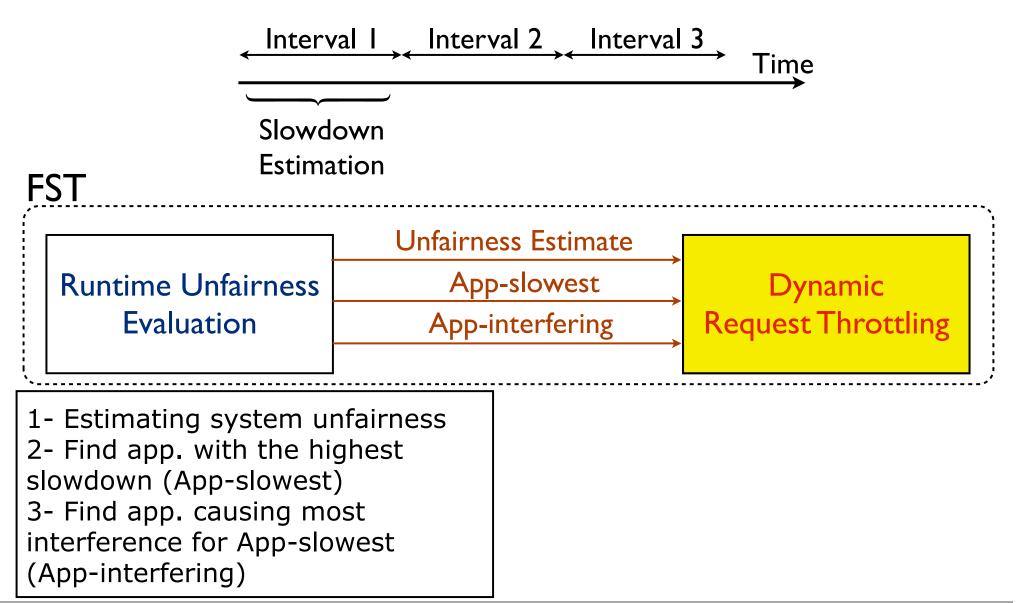


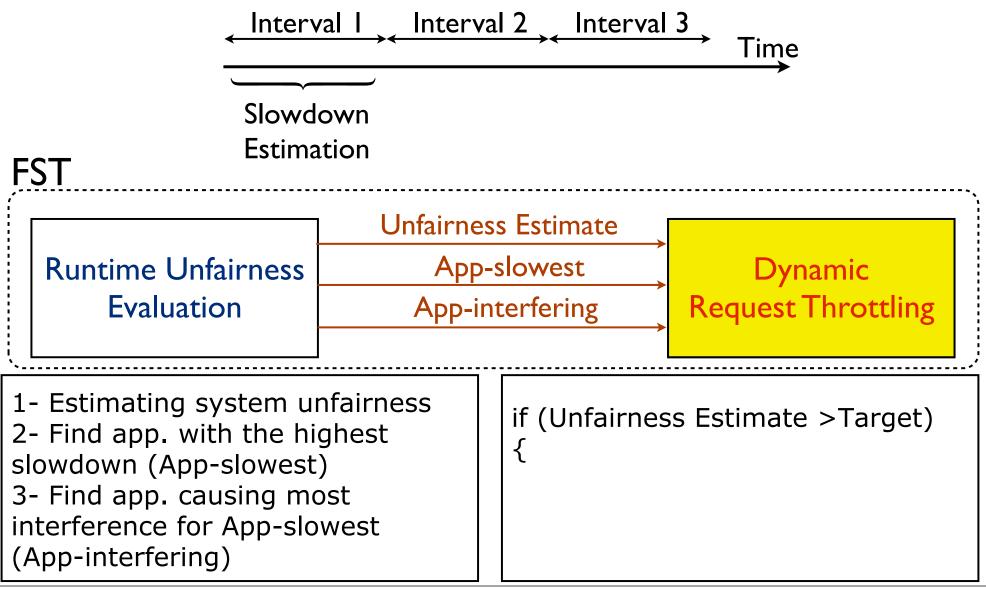




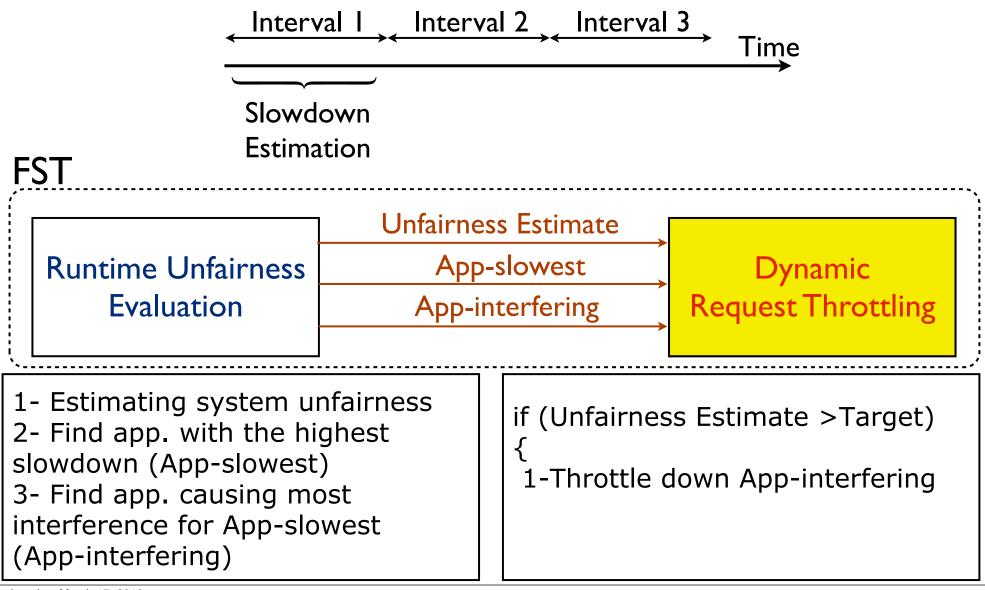


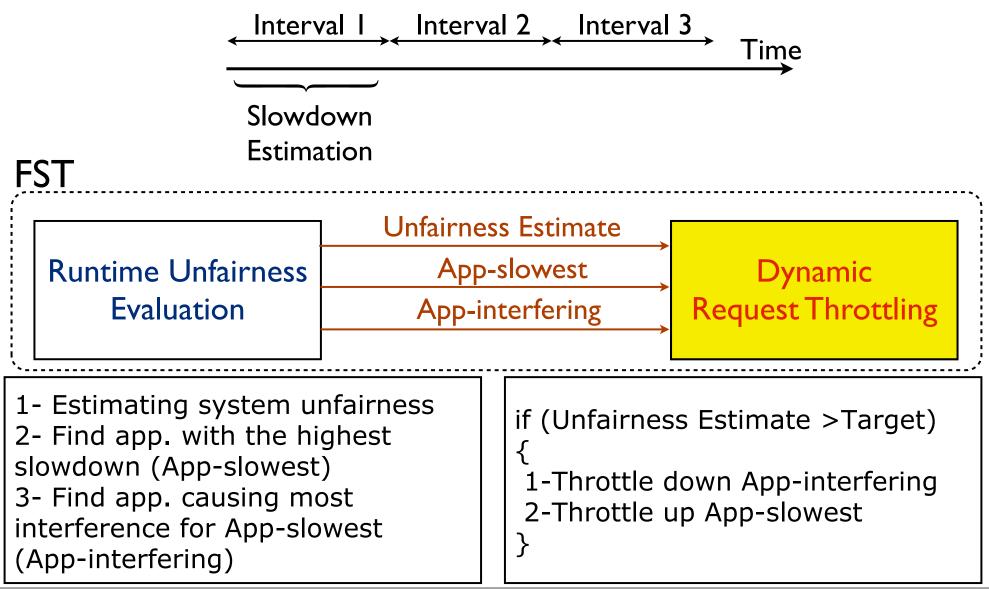


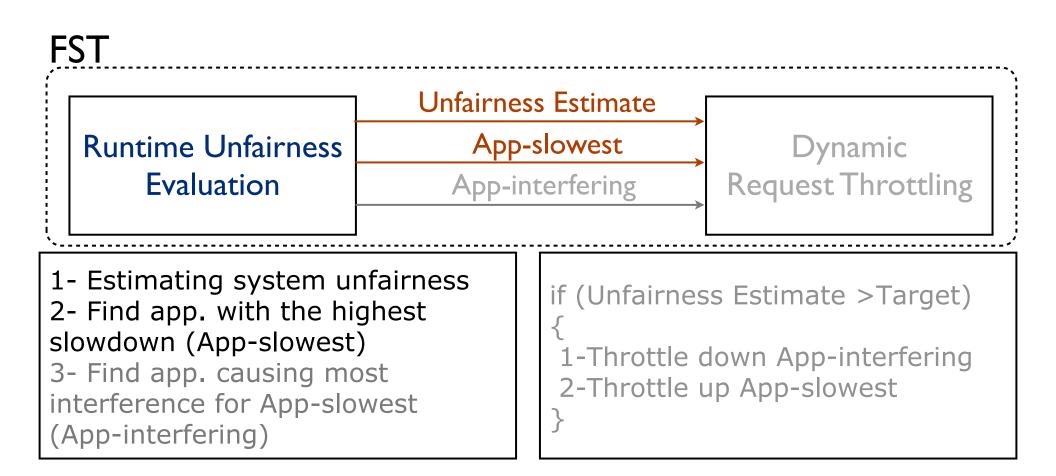


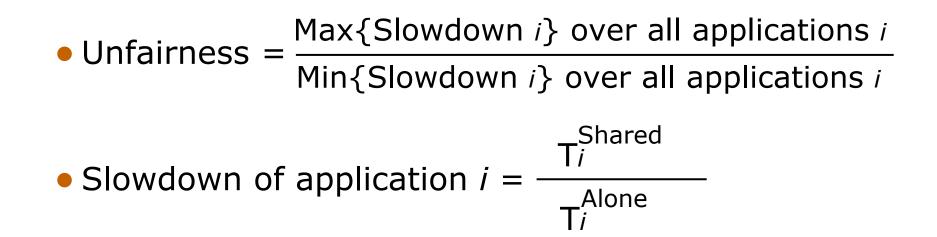


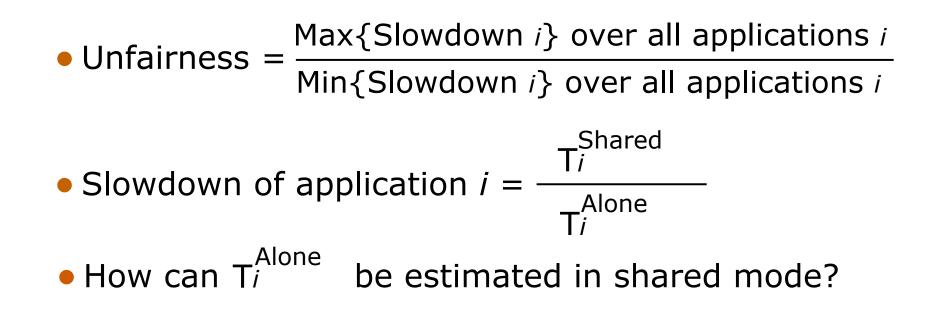
Wednesday, March 17, 2010





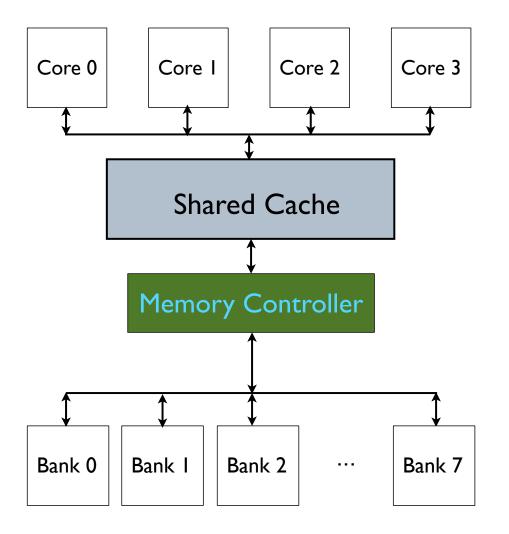


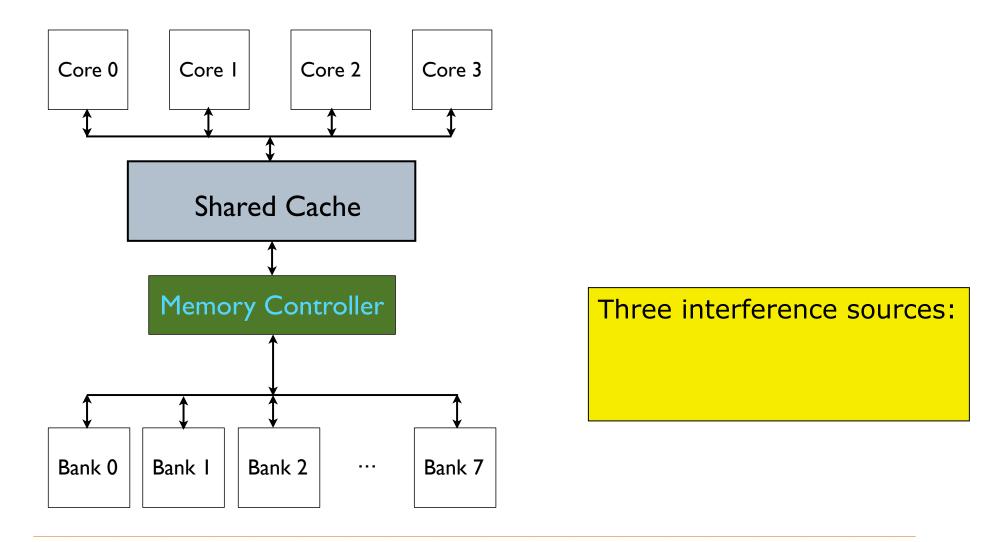


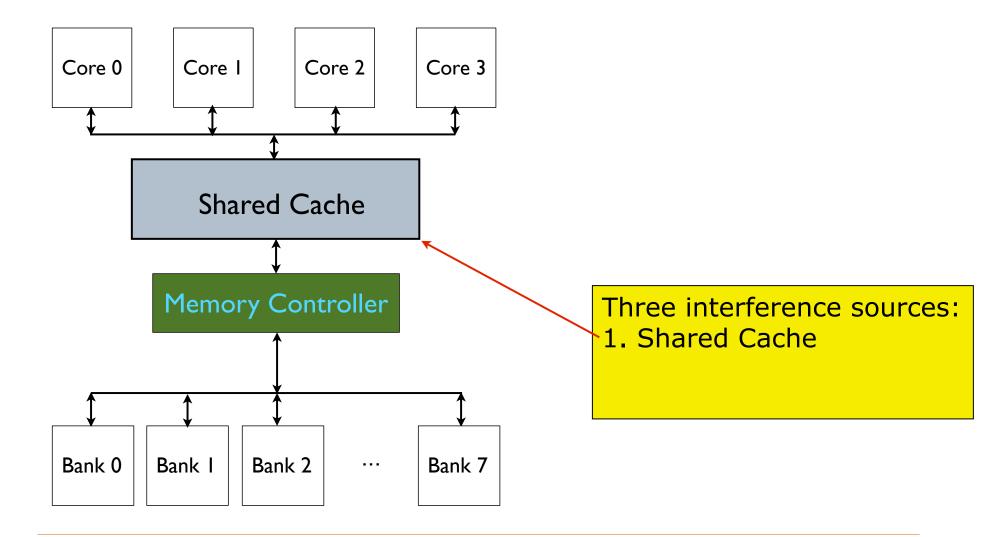


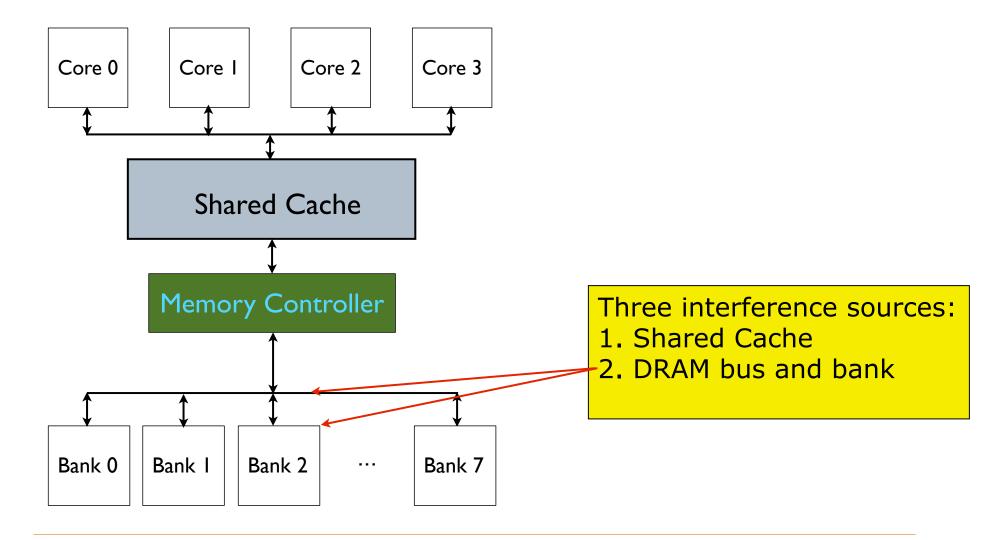
- Unfairness =  $\frac{\text{Max}\{\text{Slowdown }i\}\text{ over all applications }i}{\text{Min}\{\text{Slowdown }i\}\text{ over all applications }i}$ • Slowdown of application  $i = \frac{T_i^{\text{Shared}}}{T_i^{\text{Alone}}}$ • How can  $T_i^{\text{Alone}}$  be estimated in shared mode?
  - T<sup>Excess</sup> is the number of extra cycles it takes application *i* to execute due to interference

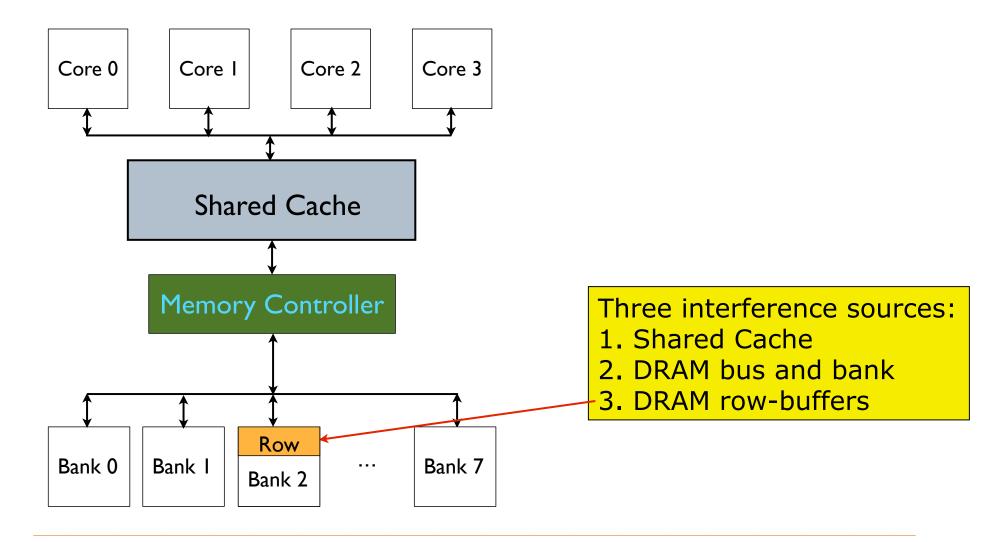
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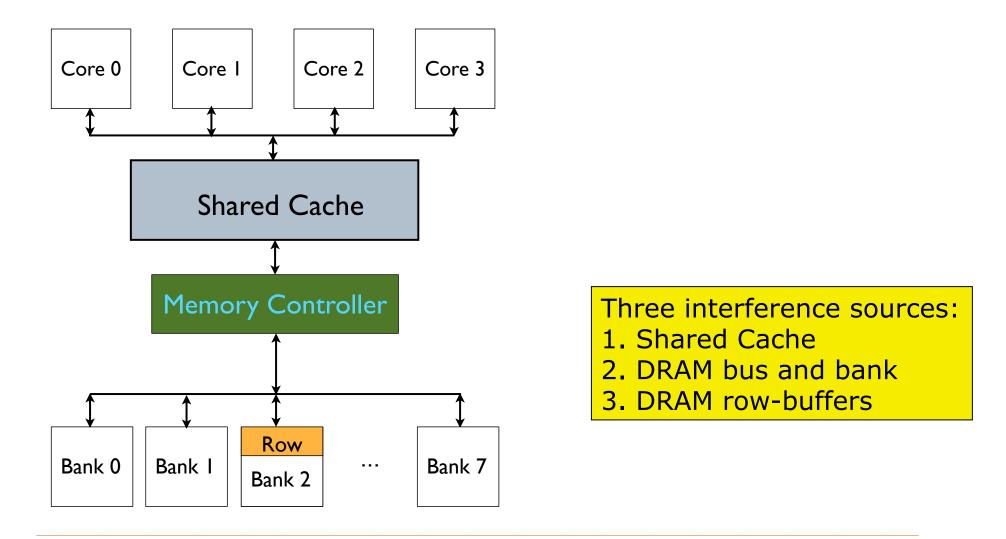


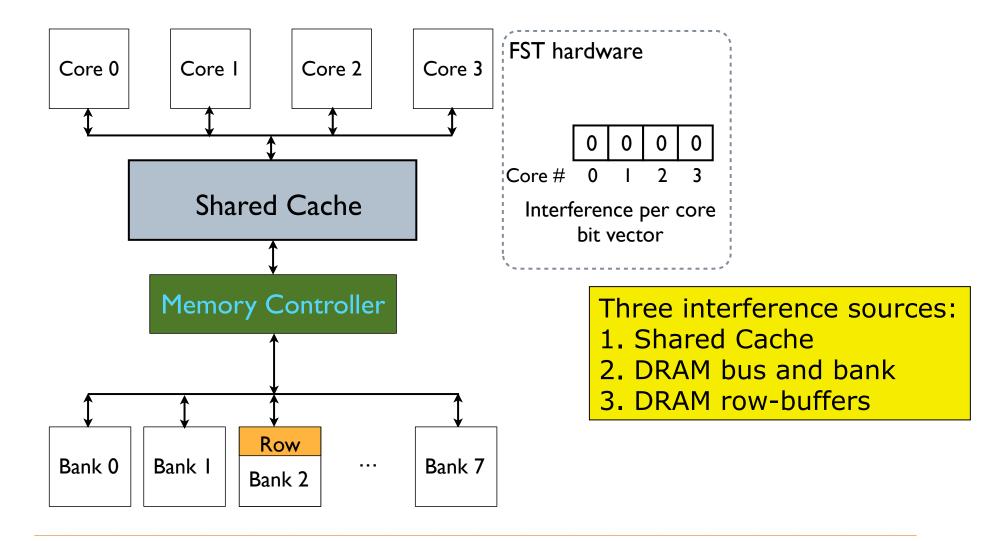


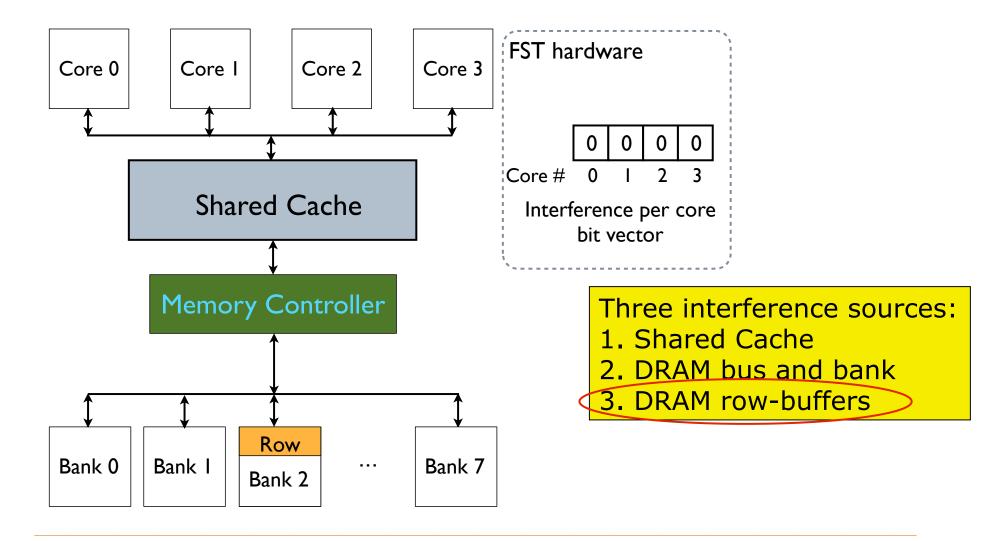








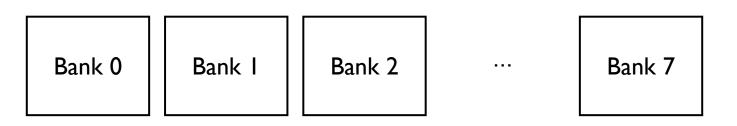


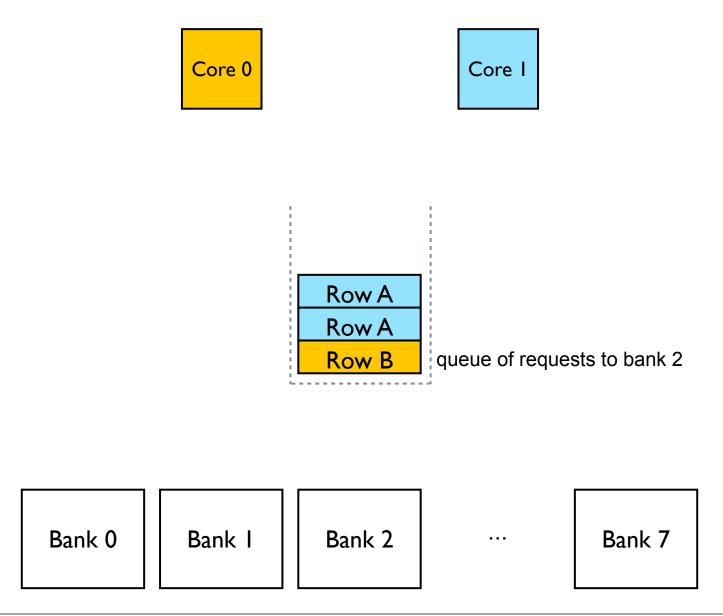


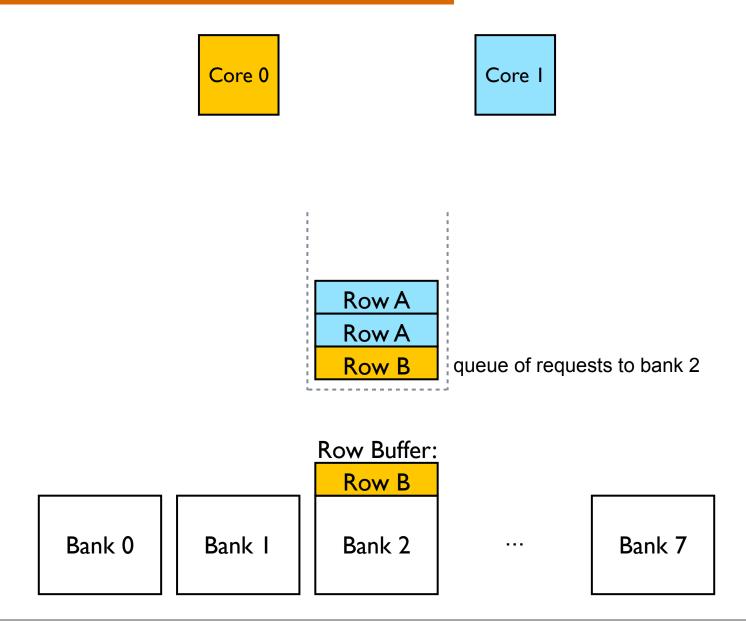
Core 0

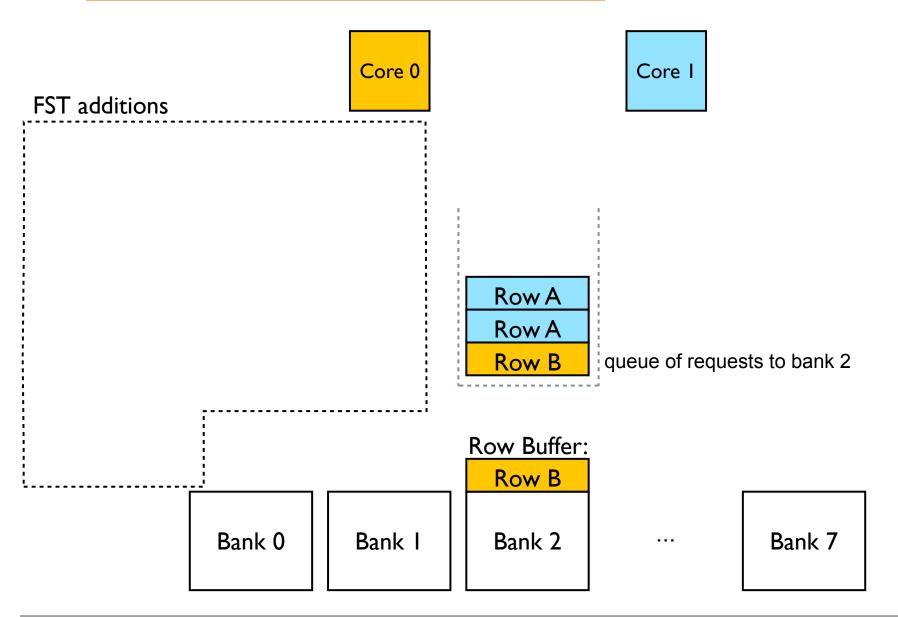


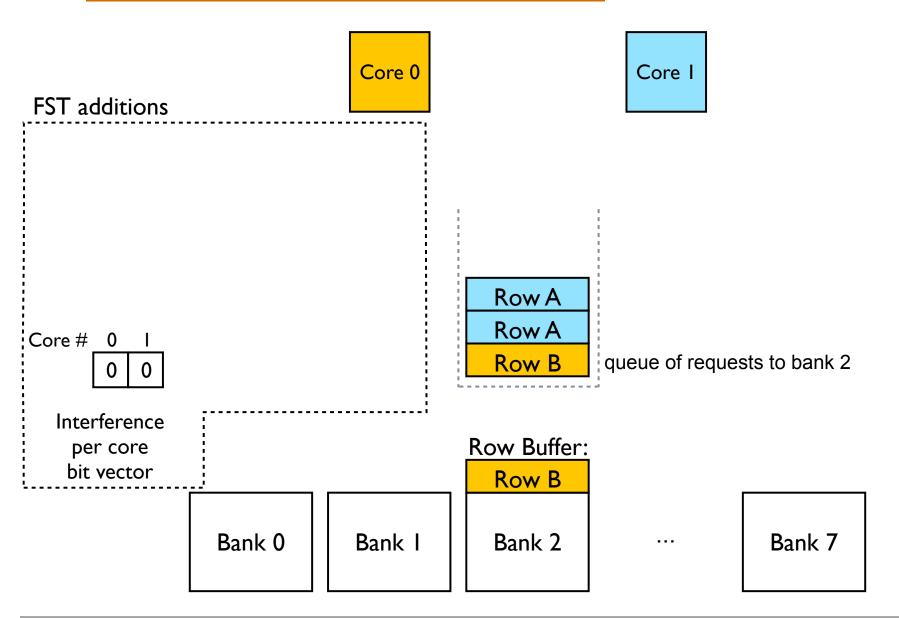
Core I

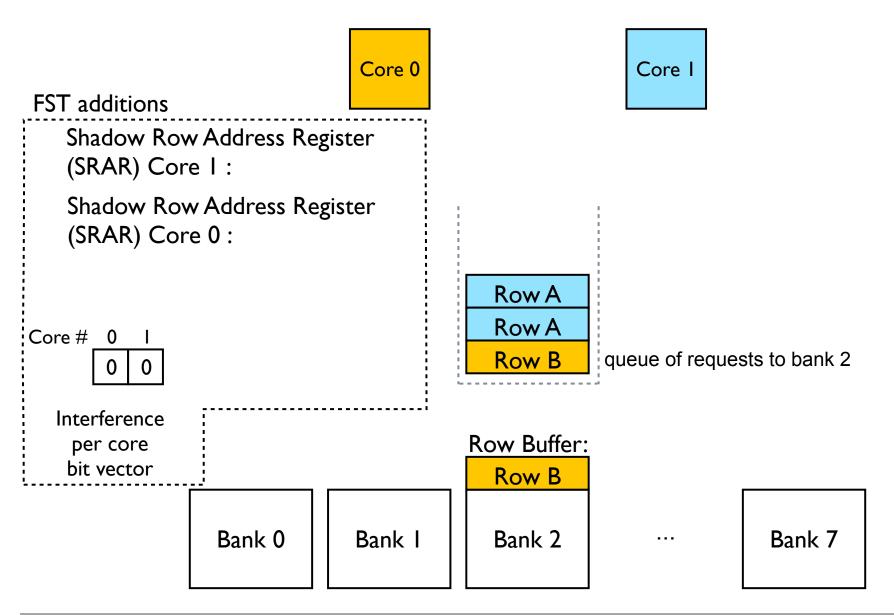


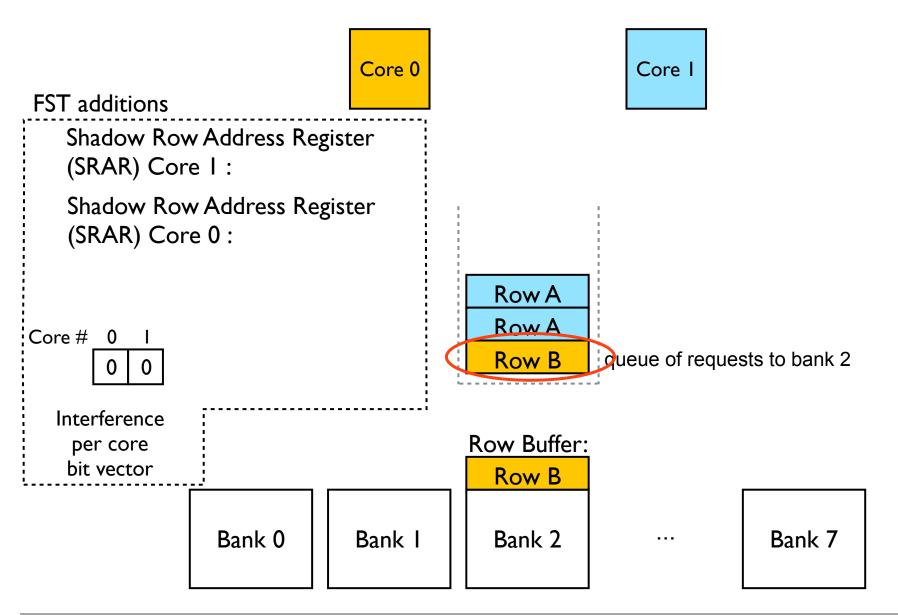


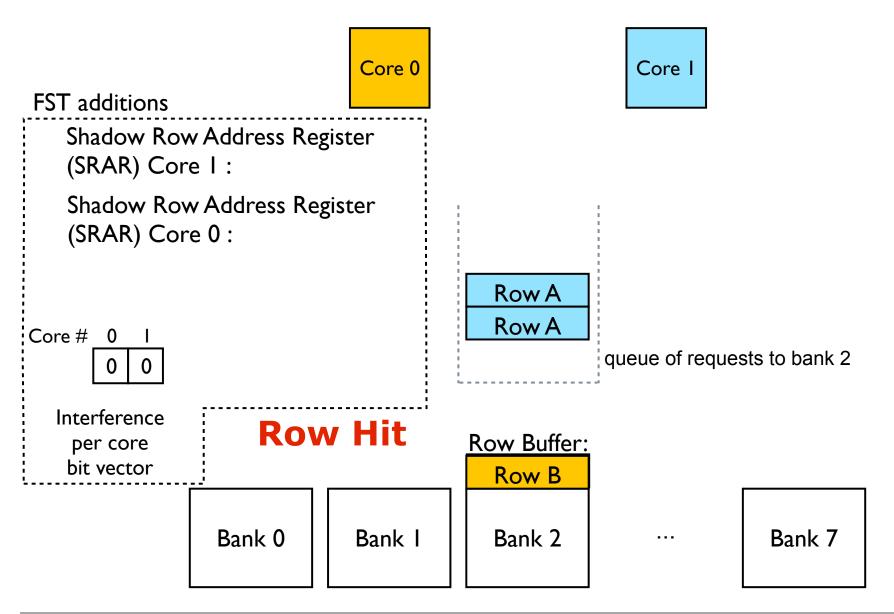


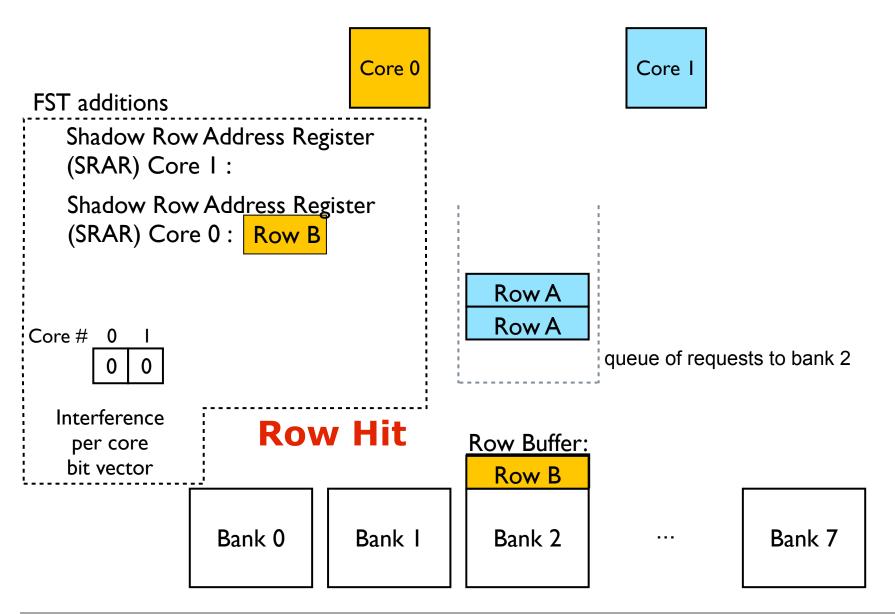


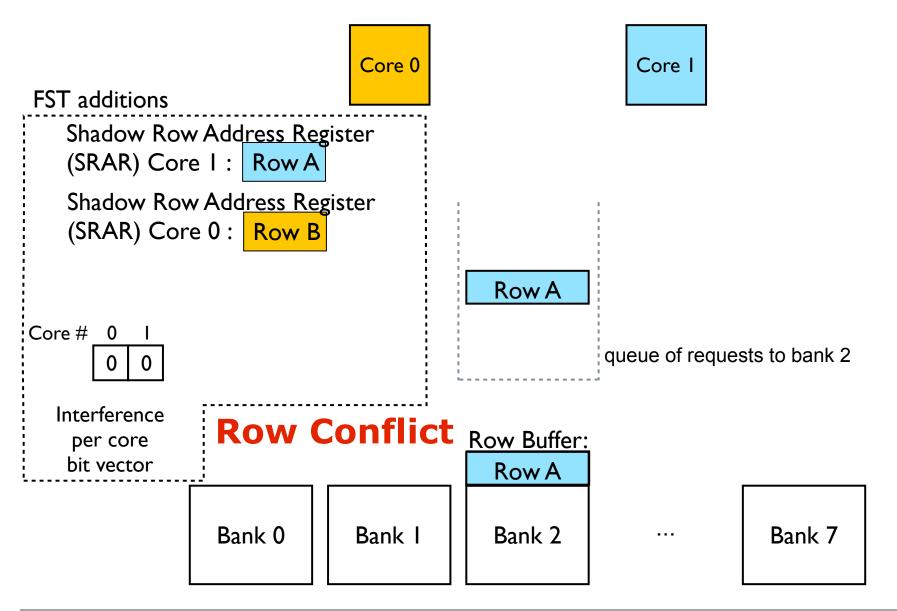


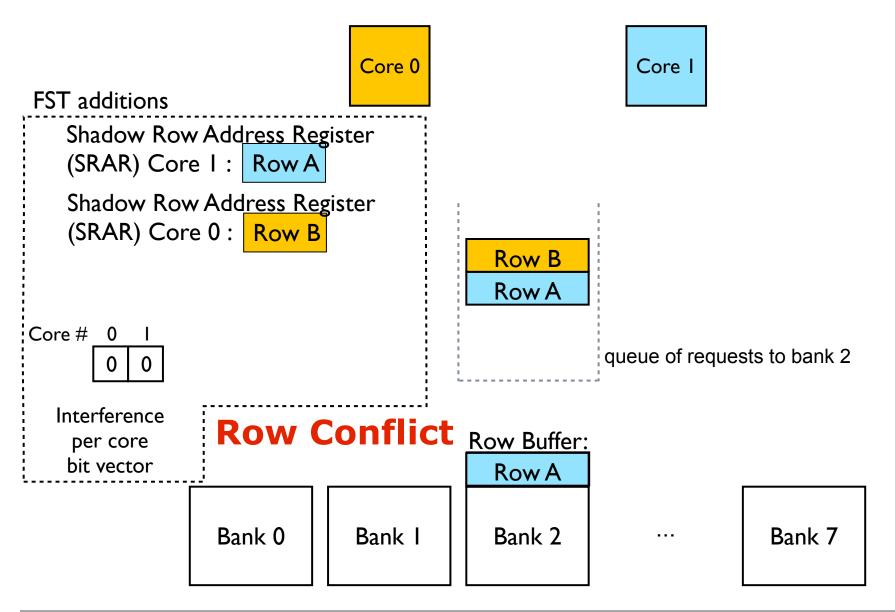


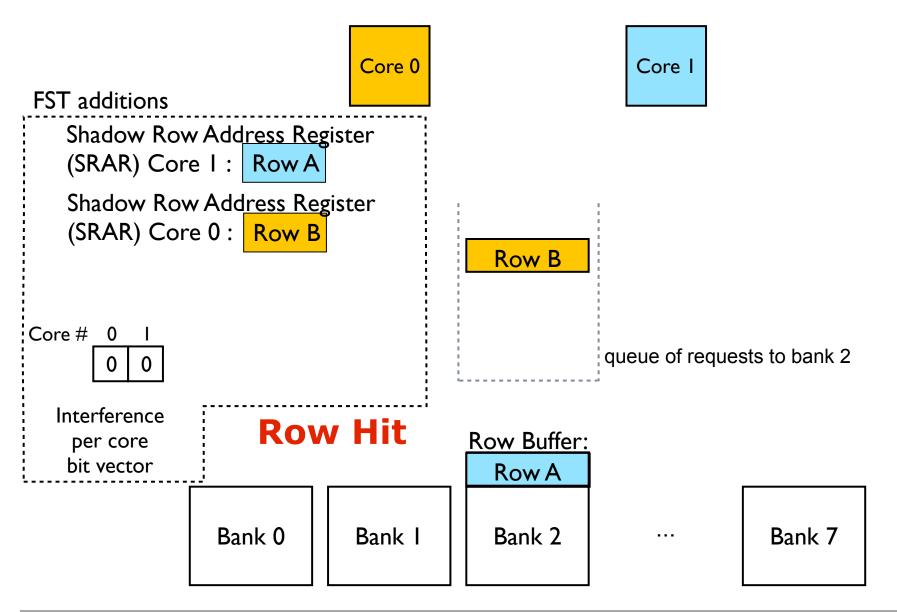


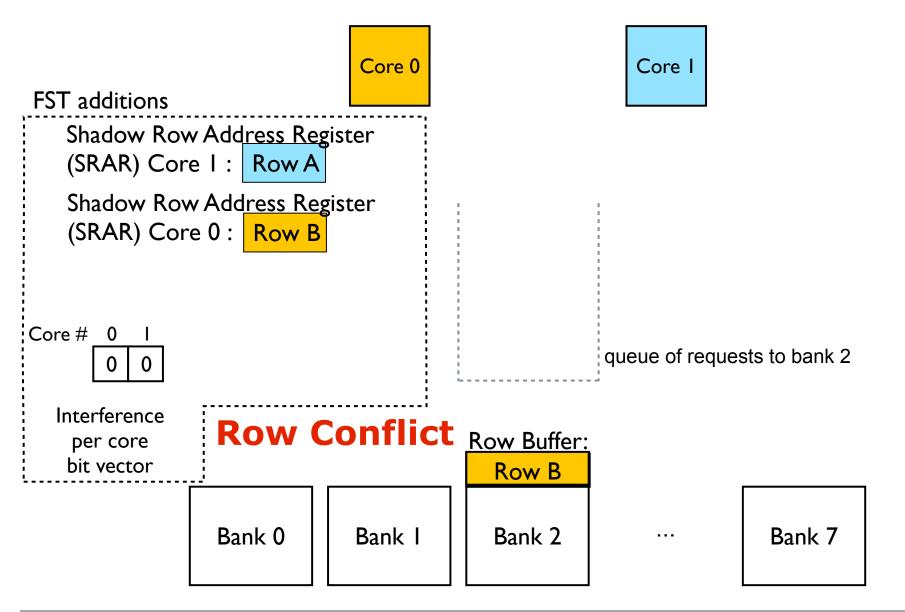


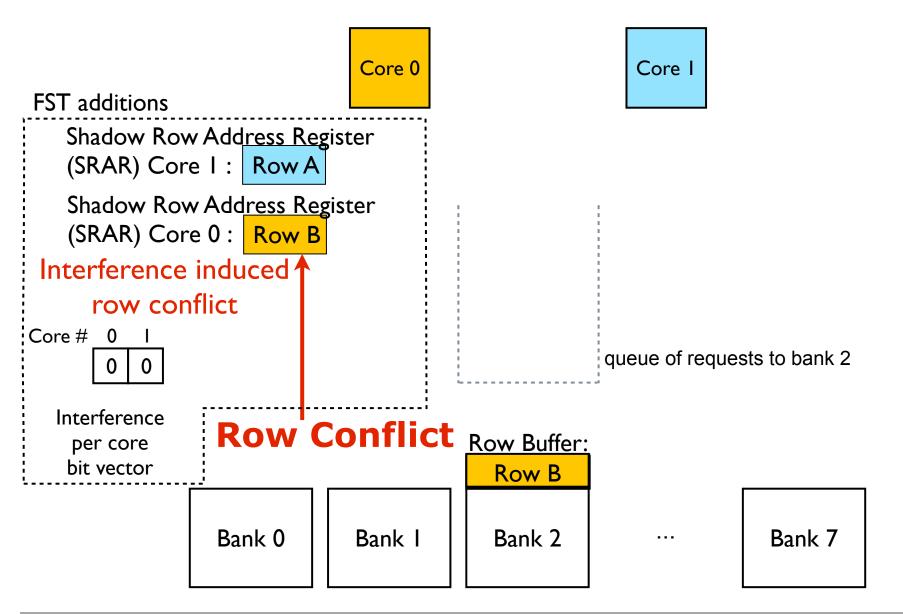


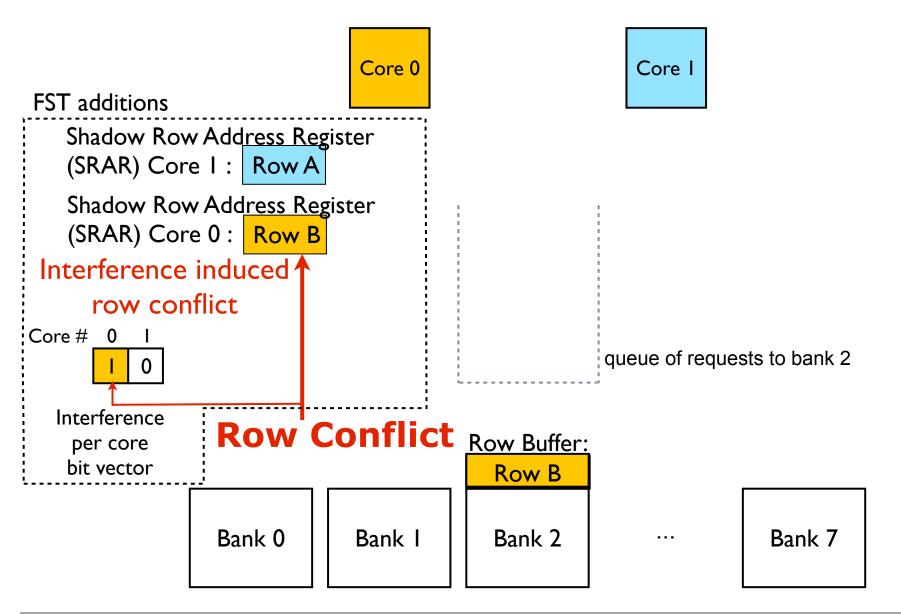


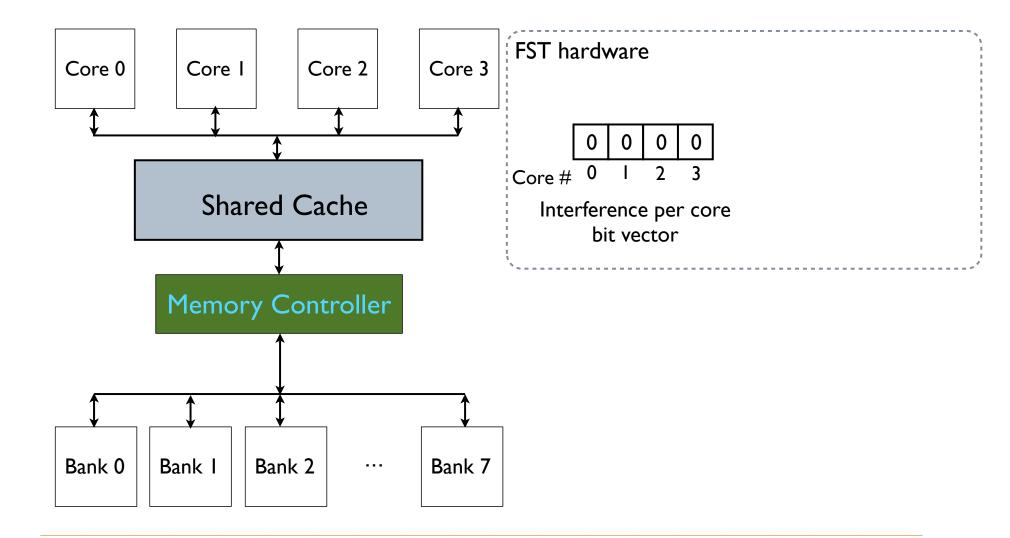




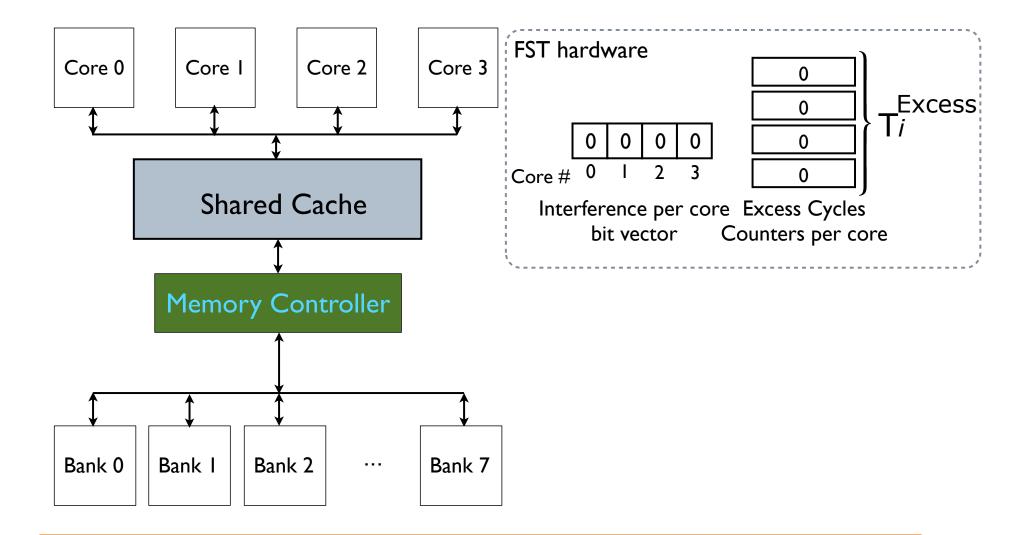


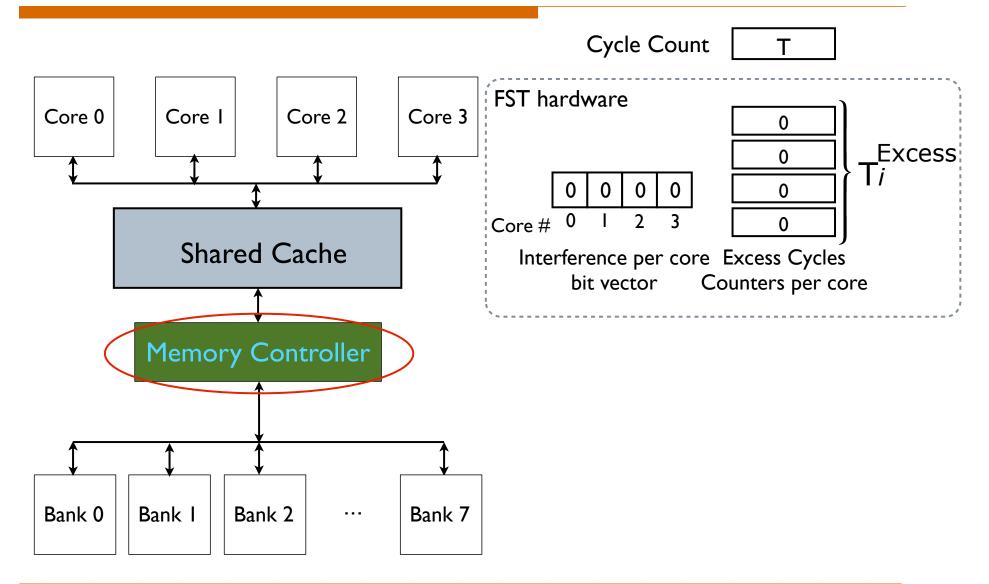


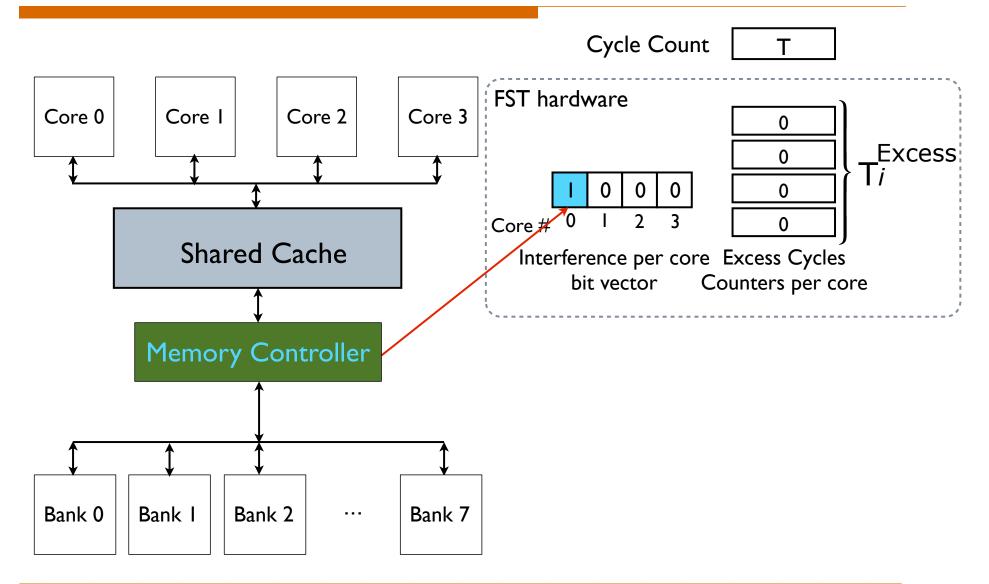


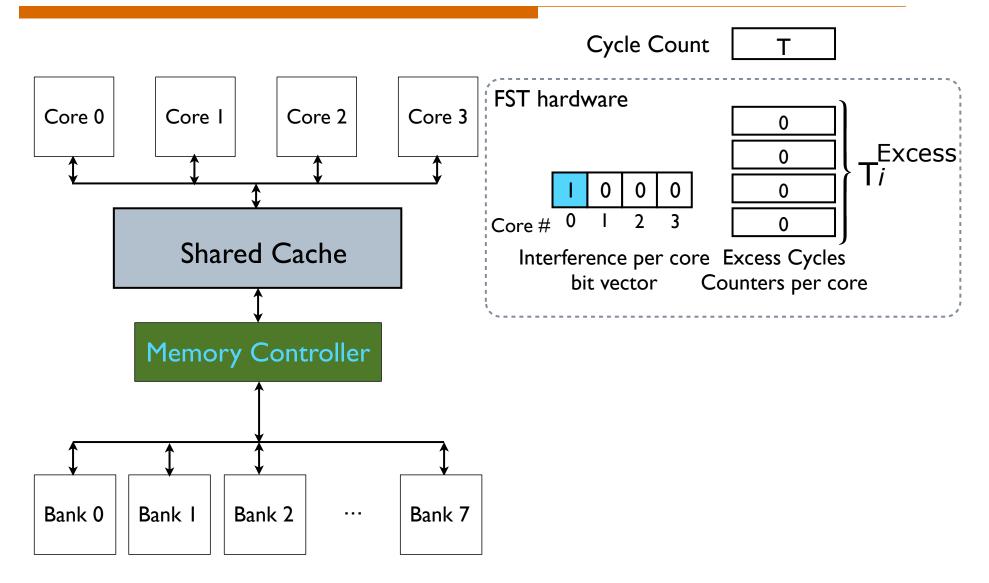


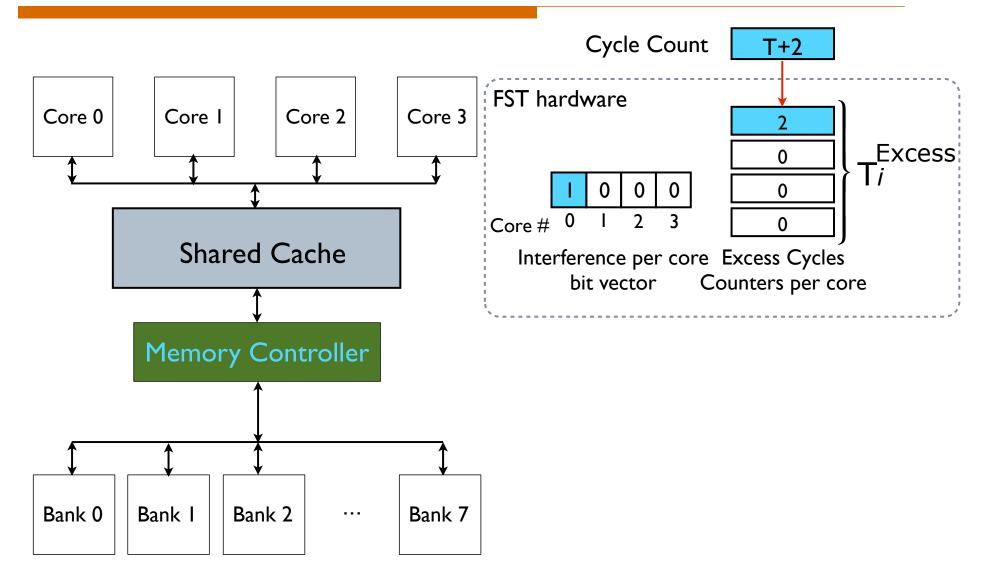
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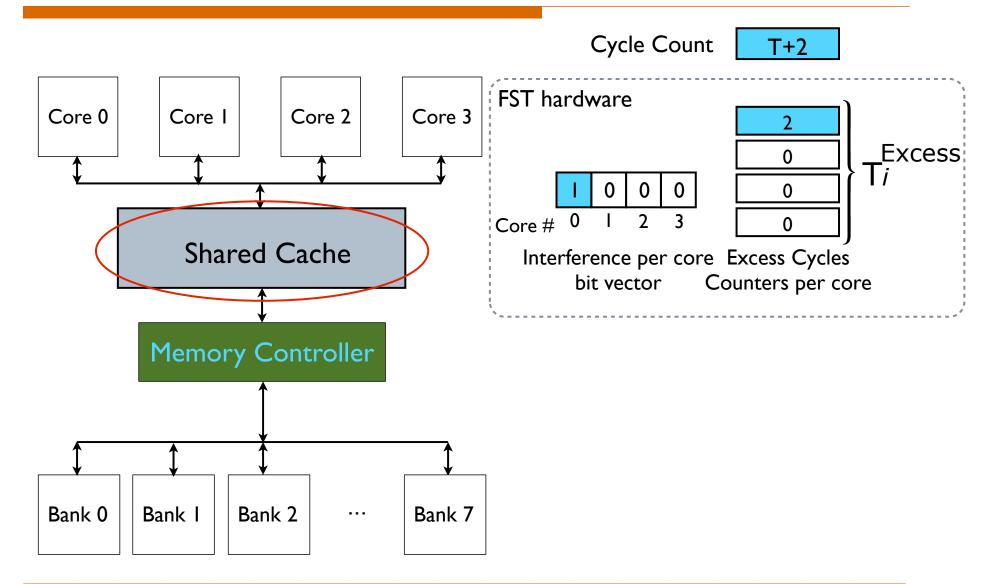


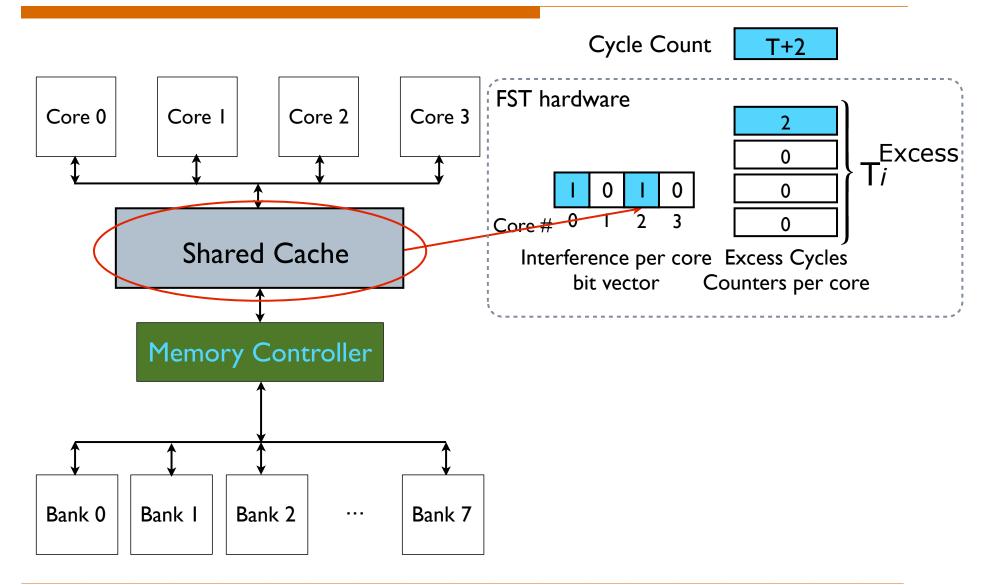


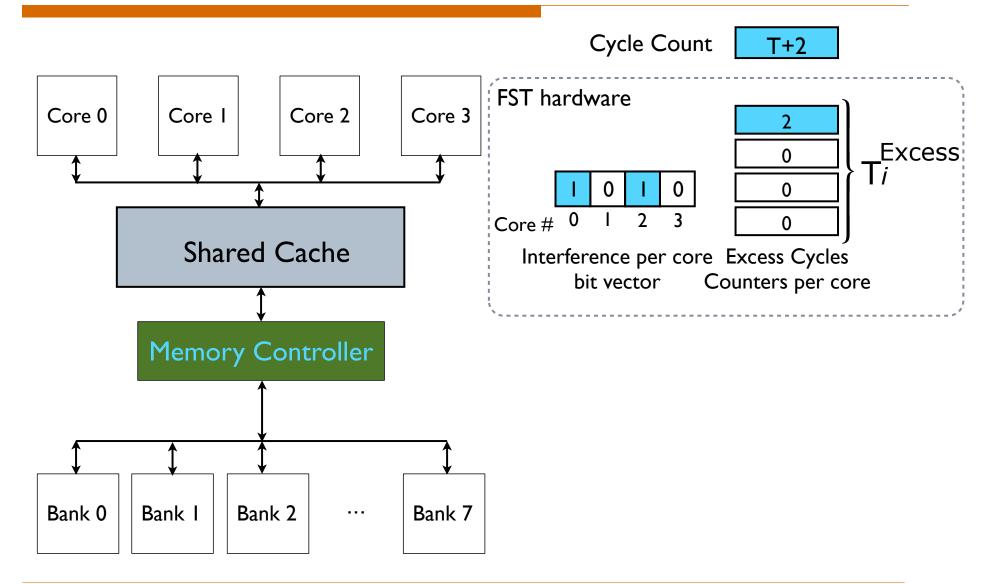


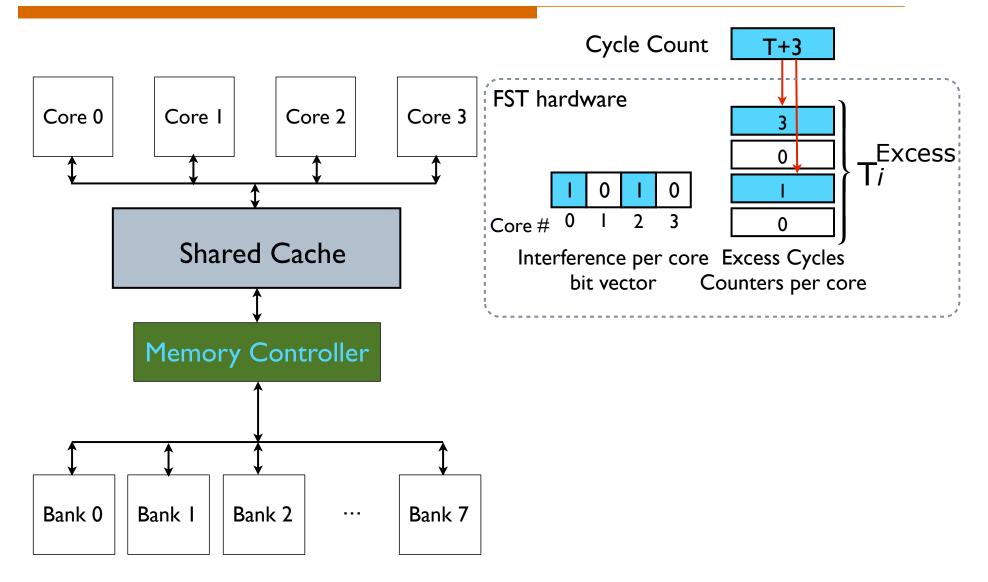




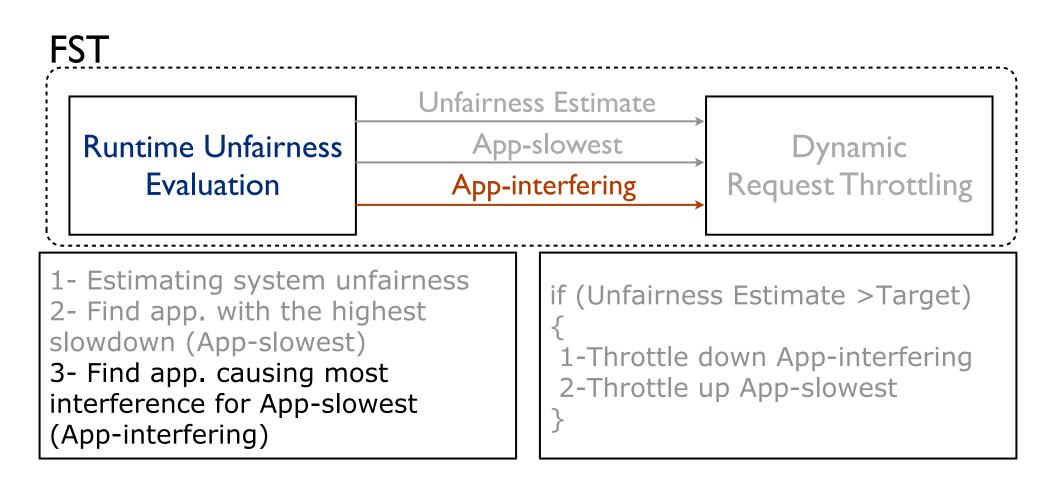








### Fairness via Source Throttling (FST)



• To identify App-interfering, for each core *i* 

- To identify App-interfering, for each core *i* 
  - FST separately tracks interference caused by each core j ( j ! i )

To identify App-interfering, for each core *i*FST separately tracks interference

caused by each core j ( j ! i )

Interference per core bit vector

Core #0 | 2 3

To identify App-interfering, for each core *i*FST separately tracks interference caused by each core *j* (*j* ! *i*)

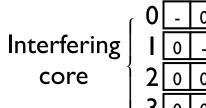
Interference per core bit vector

Core #	0		2	3
0	-	0	0	0
1	0	-	0	0
2	0	0	-	0
3	0	0	0	-

To identify App-interfering, for each core *i*FST separately tracks interference

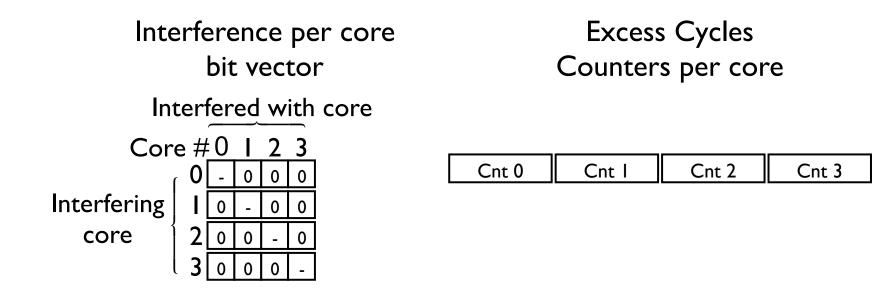
caused by each core j ( j ! i )

Interference per core bit vector Interfered with core Core #0 1 2 3



П	.0			<u> </u>	_
)	-	0	0	0	
	0	-	0	0	
2	0	0	-	0	
3	0	0	0	-	

To identify App-interfering, for each core *i*FST separately tracks interference caused by each core *j* (*j* ! *i*)



• To identify App-interfering, for each core *i*  FST separately tracks interference caused by each core j ( j ! i )

> Interference per core bit vector

**Excess Cycles** Counters per core



Core #0 | 2 3 Interfering core

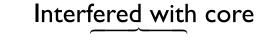
$\pi$	. 0			<u> </u>	-
)	-	0	0	0	
	0	-	0	0	
2	0	0	-	0	
3	0	0	0	-	

-	Cnt 0, I	Cnt 0,2	Cnt 0,3
Cnt I,0	-	Cnt I,2	Cnt I,3
Cnt 2,0	Cnt 2, I	-	Cnt 2,3
Cnt 3,0	Cnt 3, I	Cnt 3,2	-

• To identify App-interfering, for each core *i*  FST separately tracks interference caused by each core j ( j ! i )

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**Excess Cycles** Counters per core



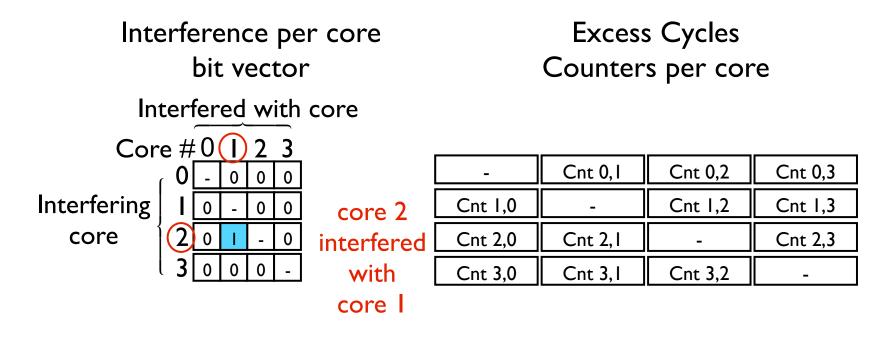
Interfering core

Cor	e #	0		2	3
	0	-	0	0	0
ring		0	-	0	0
e	2	0	I	-	0
	3	0	0	0	

-	Cnt 0, I	Cnt 0,2	Cnt 0,3
Cnt I,0	-	Cnt I,2	Cnt I,3
Cnt 2,0	Cnt 2, I	-	Cnt 2,3
Cnt 3,0	Cnt 3, I	Cnt 3,2	_

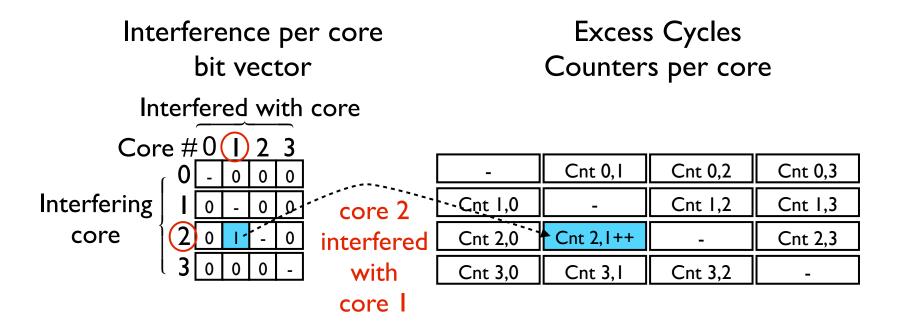
To identify App-interfering, for each core *i*FST separately tracks interference

caused by each core j ( j ! i )



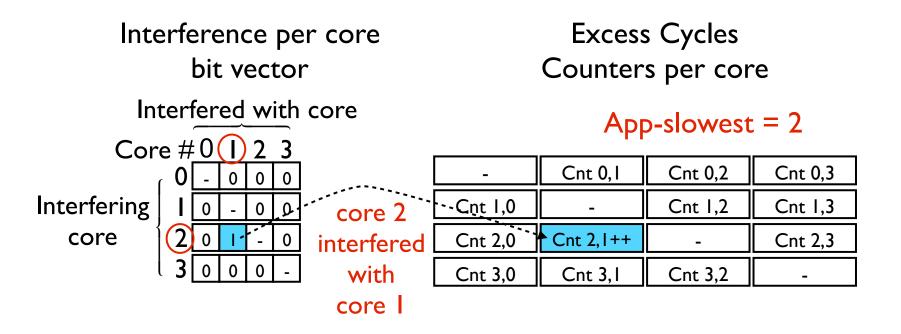
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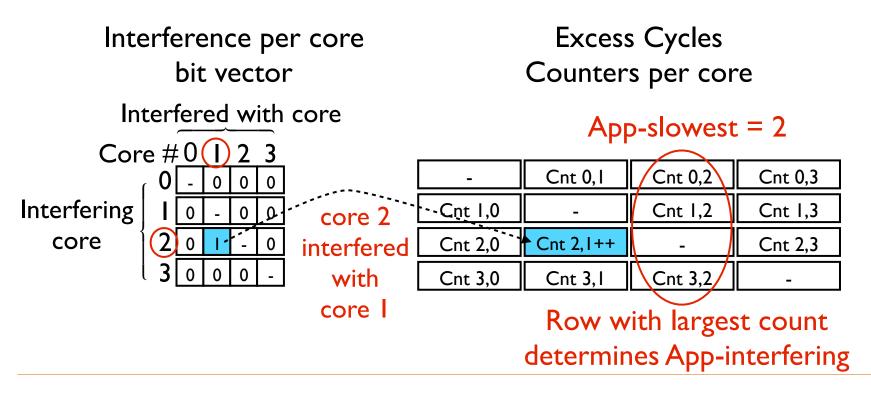


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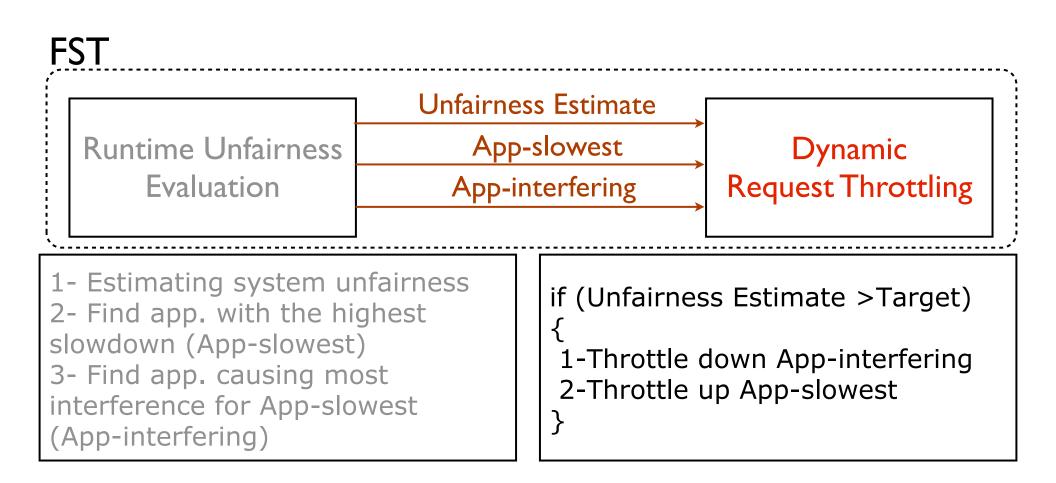
caused by each core j ( j ! i )



To identify App-interfering, for each core *i*FST separately tracks interference caused by each core *j* (*j* ! *i*)



## Fairness via Source Throttling (FST)



• Goal: Adjust how aggressively each core makes requests to the shared resources

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- Mechanisms:
  - Miss Status Holding Register (MSHR) quota

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- Miss Status Holding Register (MSHR) quota
  - Controls the number of concurrent requests accessing shared resources from each application
- Request injection frequency
  - Controls how often memory requests are issued to the last level cache from the MSHRs

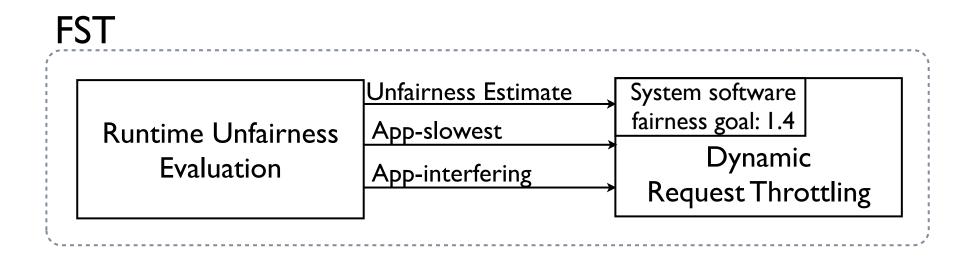
 Throttling level assigned to each core determines both MSHR quota and request injection rate

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	Throttling level	MSHR quota	Request Injection Rate
	100%	128	Every cycle
	50%	64	Every other cycle
	25%	32	Once every 4 cycles
	10%	12	Once every 10 cycles
	5%	6	Once every 20 cycles
	4%	5	Once every 25 cycles
	3%	3	Once every 30 cycles
Total # of MSHRs: 128	2%	2	Once every 50 cycles

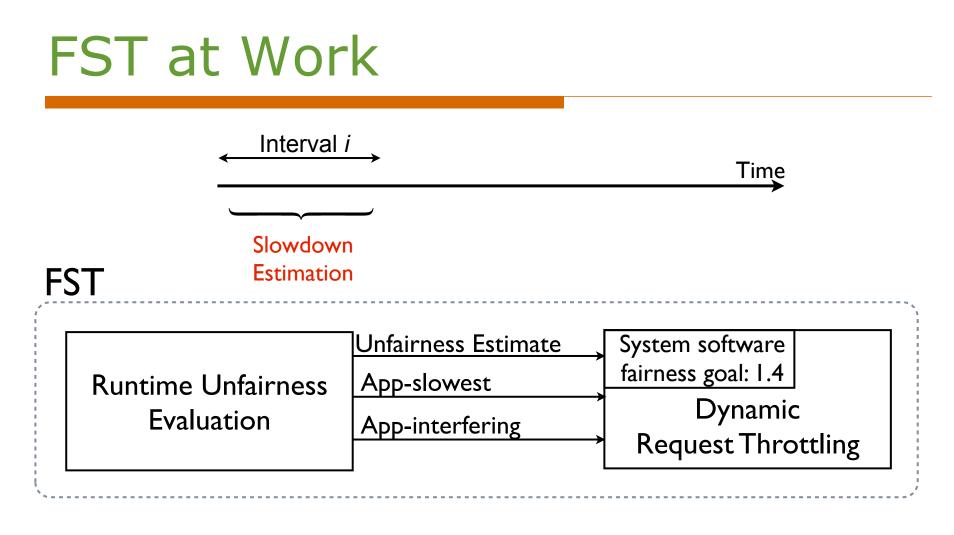
 Throttling level assigned to each core determines both MSHR quota and request injection rate

	Throttling level	MSHR quota	Request Injection Rate
	100%	128	Every cycle
	50%	64	Every other cycle
	25%	32	Once every 4 cycles
$\langle$	0%	12	Once every 10 cycles
	5%	6	Once every 20 cycles
	4%	5	Once every 25 cycles
<b>T</b>	3%	3	Once every 30 cycles
Total # of MSHRs: I 28	2%	2	Once every 50 cycles

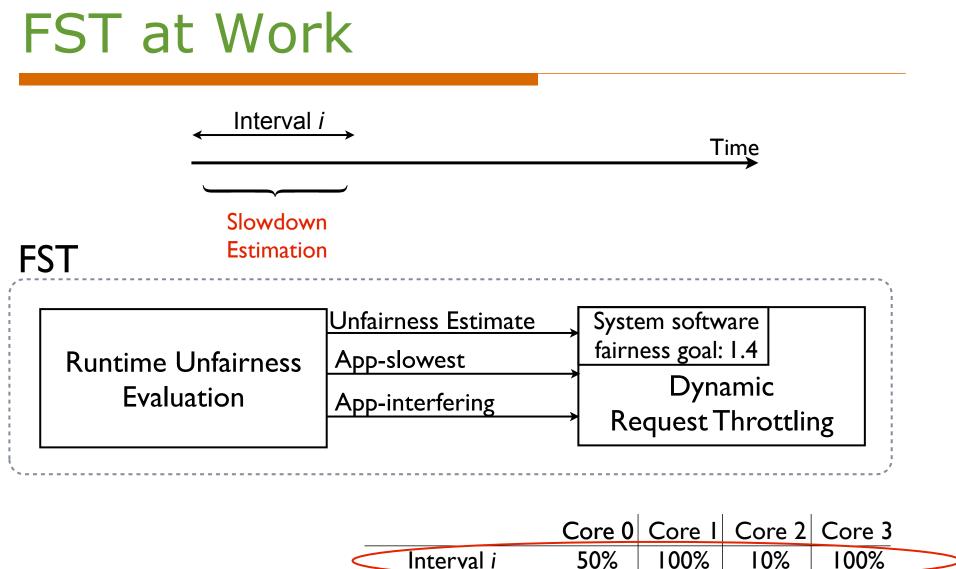


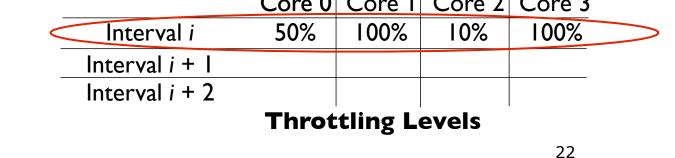
Core 0Core 1Core 2Core 3Interval iInterval i + 1Interval i + 2Interval i + 2Interval i + 2Interval i + 2Interval i + 2Throttling Levels

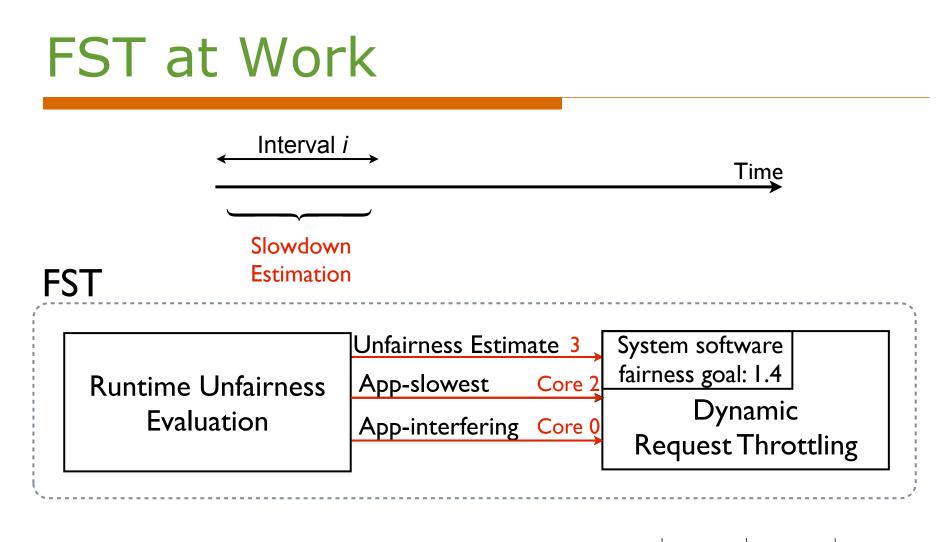
Time



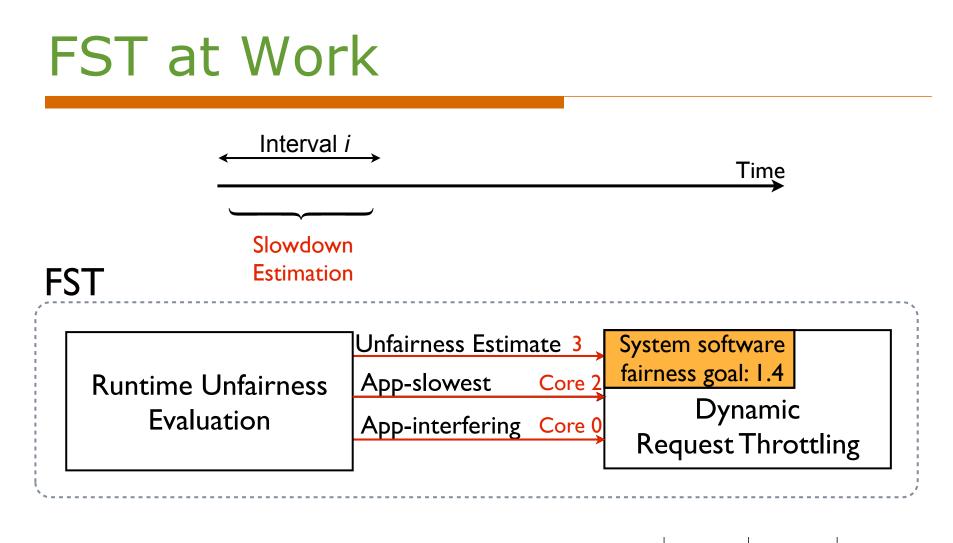
		Core 0	Core I	Core 2	Core 3	
Inter	val i	50%	100%	10%	100%	
Interva	<i>i</i> +					
Interva	l i + 2					
	Throttling Levels					



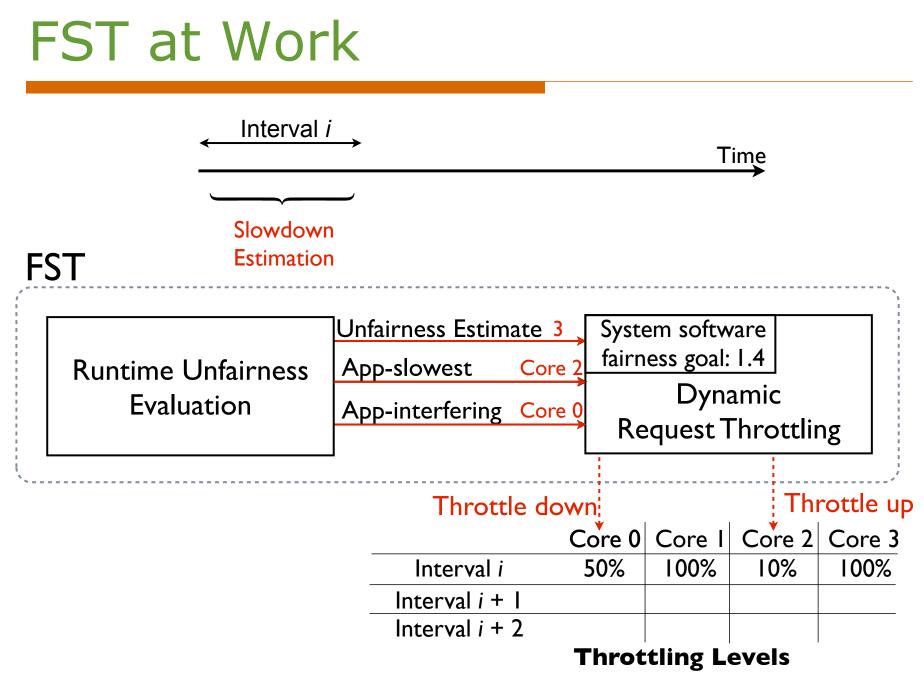




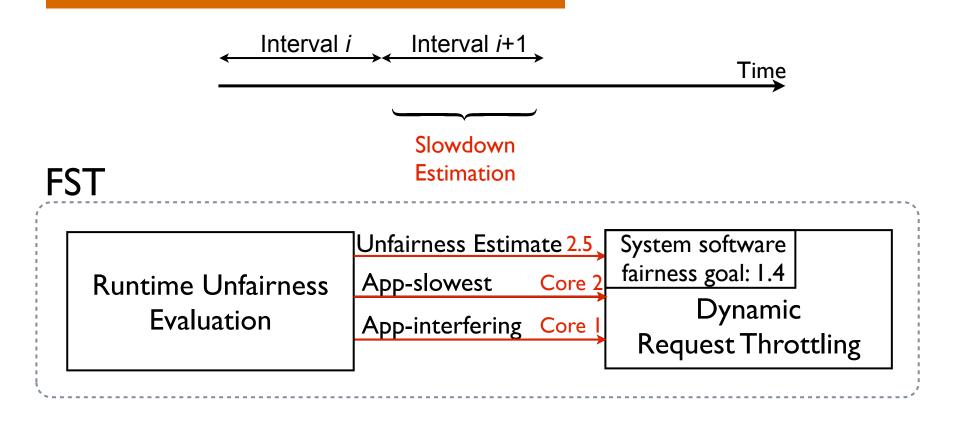
	Core 0	Core I	Core 2	Core 3		
Interval i	50%	100%	10%	100%		
Interval <i>i</i> + 1						
Interval <i>i</i> + 2						
	Throttling Levels					



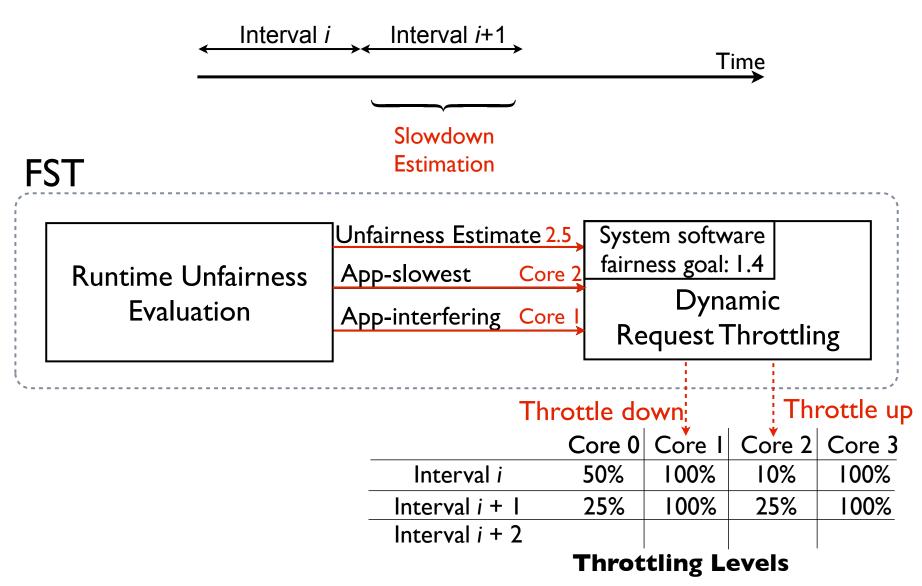
	Core 0	Core I	Core 2	Core 3		
Interval i	50%	100%	10%	100%		
Interval <i>i</i> + 1						
Interval <i>i</i> + 2						
	Throttling Levels					

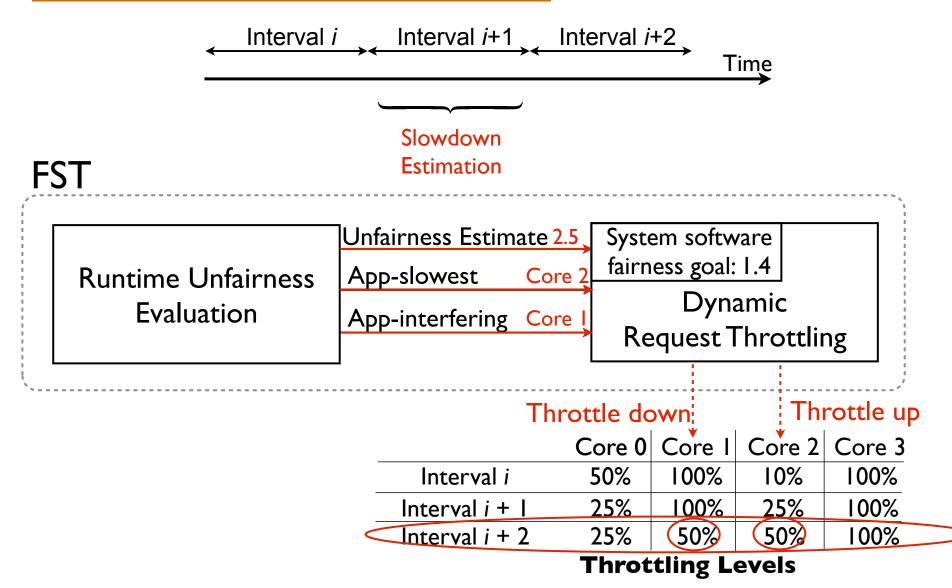


#### FST at Work Interval *i* \_ Interval *i*+1 Time Slowdown **Estimation** FST Unfairness Estimate 3 System software fairness goal: 1.4 Runtime Unfairness App-slowest Core 2 Dynamic Evaluation App-interfering Core Q **Request Throttling** Throttle up Throttle down Core 2 Core 3 Core 0 Core 1 Interval *i* 50% 100% 10% 100% 25% Interval i + | (25%)100% 100% Interval i + 2**Throttling Levels**



	Core 0	Core I	Core 2	Core 3	
Interval i	50%	100%	10%	100%	
Interval i + 1	25%	100%	25%	100%	
Interval <i>i</i> + 2					
	Throttling Levels				





## System Software Support

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 Different fairness objectives can be configured by system software

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- Different fairness objectives can be configured by system software
  - Estimated Unfairness > Target Unfairness

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  - Estimated Unfairness > Target Unfairness
  - Estimated Max Slowdown > Target Max Slowdown
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- Support for thread priorities

Weighted Slowdown(i) =
Estimated Slowdown(i) x Weight(i)



 Total storage cost required for 4 cores is ~ 12KB

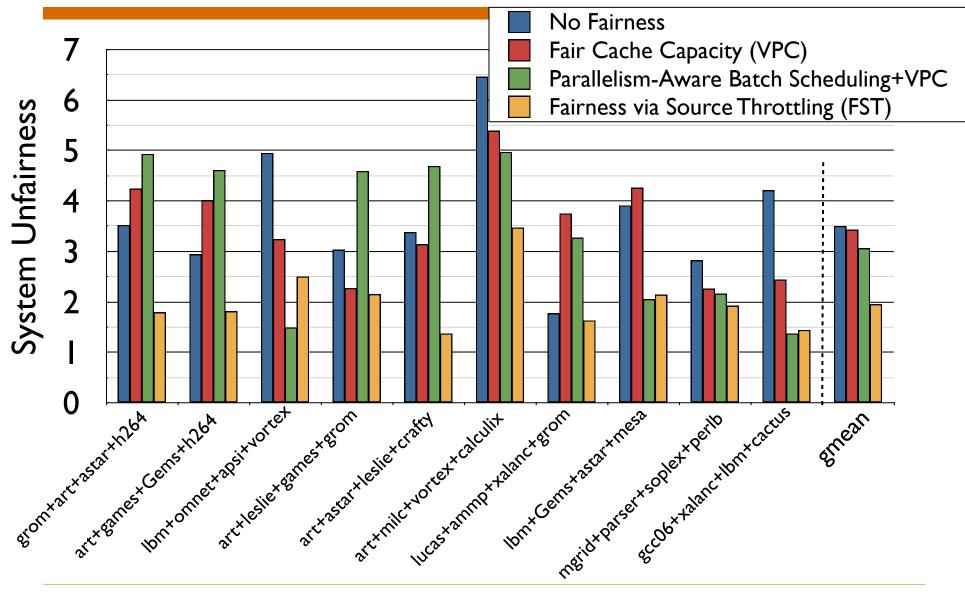
 FST does not require any structures or logic that are on the processor's critical path

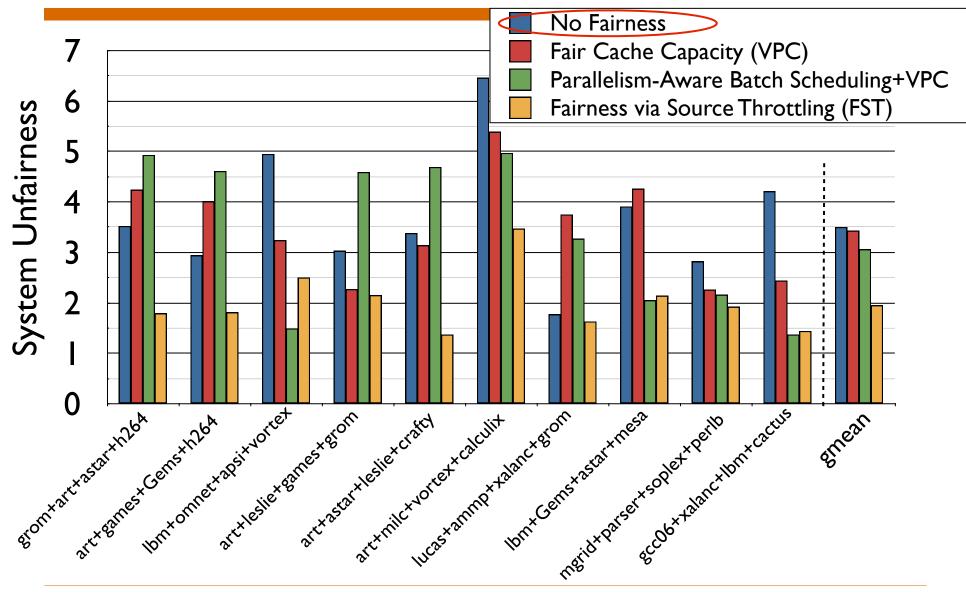
# Outline

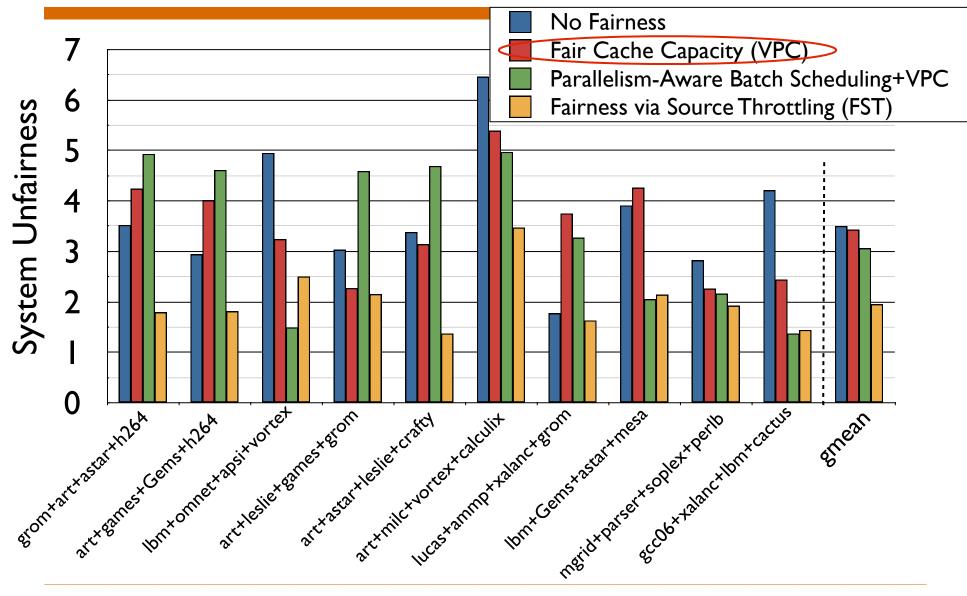
- Background and Problem
- Motivation for Source Throttling
- Fairness via Source Throttling (FST)
- Evaluation
- Conclusion

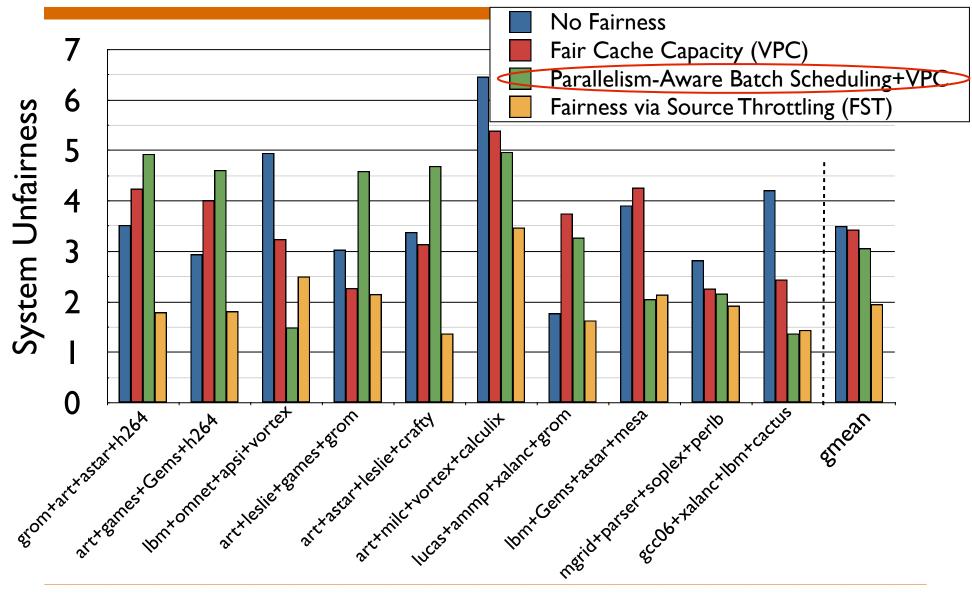
# **Evaluation Methodology**

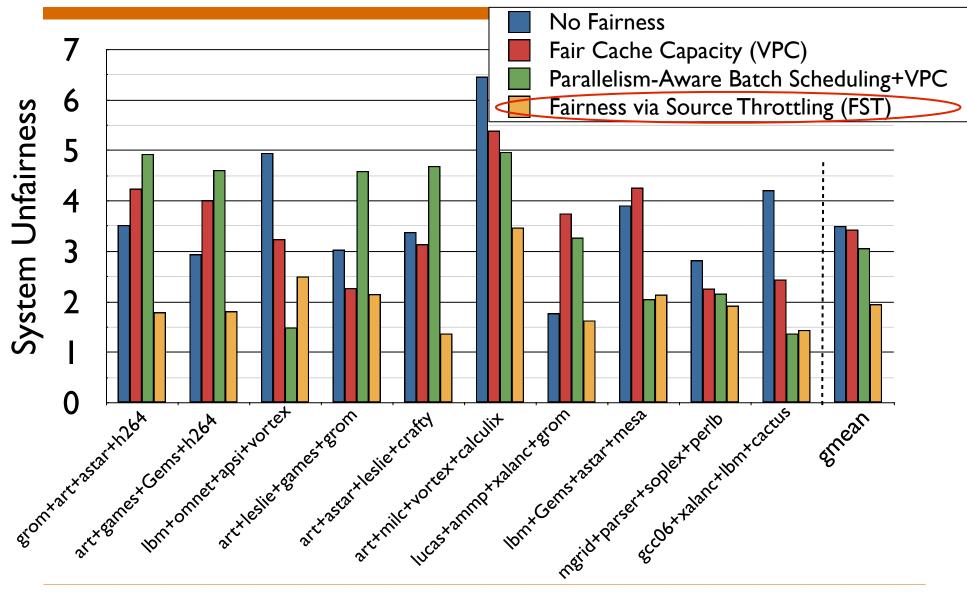
- x86 cycle accurate simulator
- Baseline processor configuration
  - Per-core
    - 4-wide issue, out-of-order, 256 entry ROB
  - Shared (4-core system)
    - 128 MSHRs
    - 2 MB, 16-way L2 cache
  - Main Memory
    - DDR3 1333 MHz
    - Latency of 15ns per command (*tRP*, *tRCD*, *CL*)
    - 8B wide core to memory bus

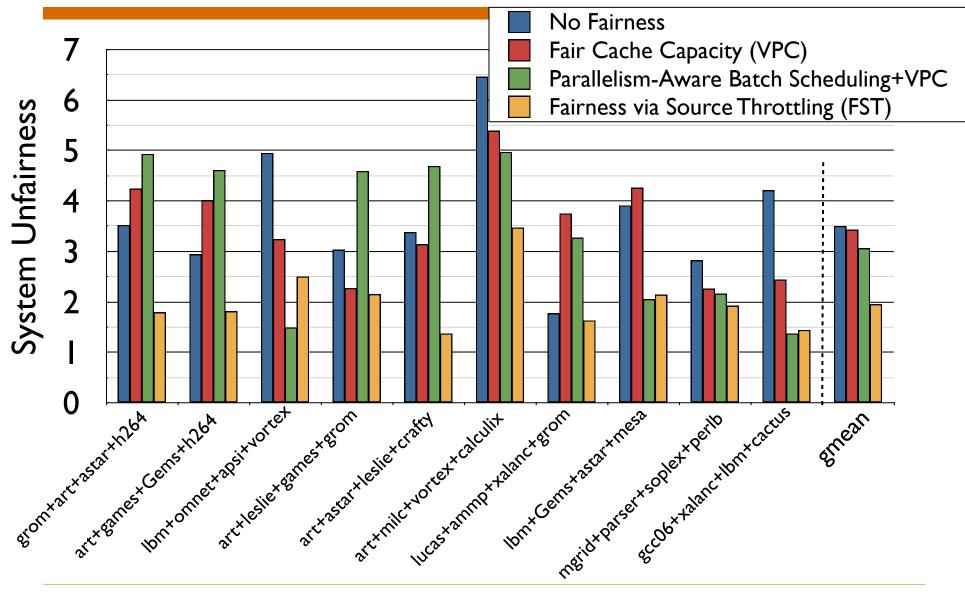


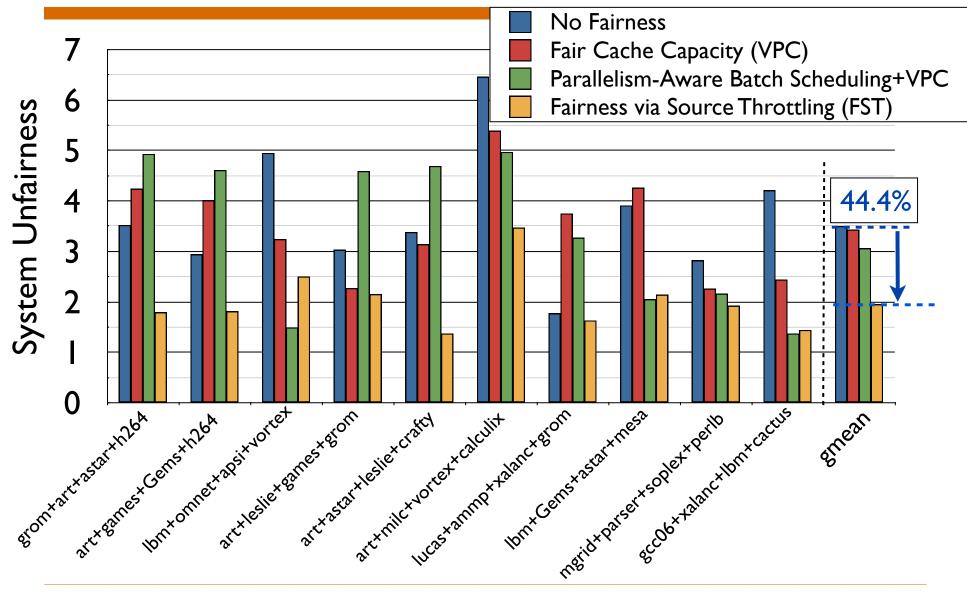


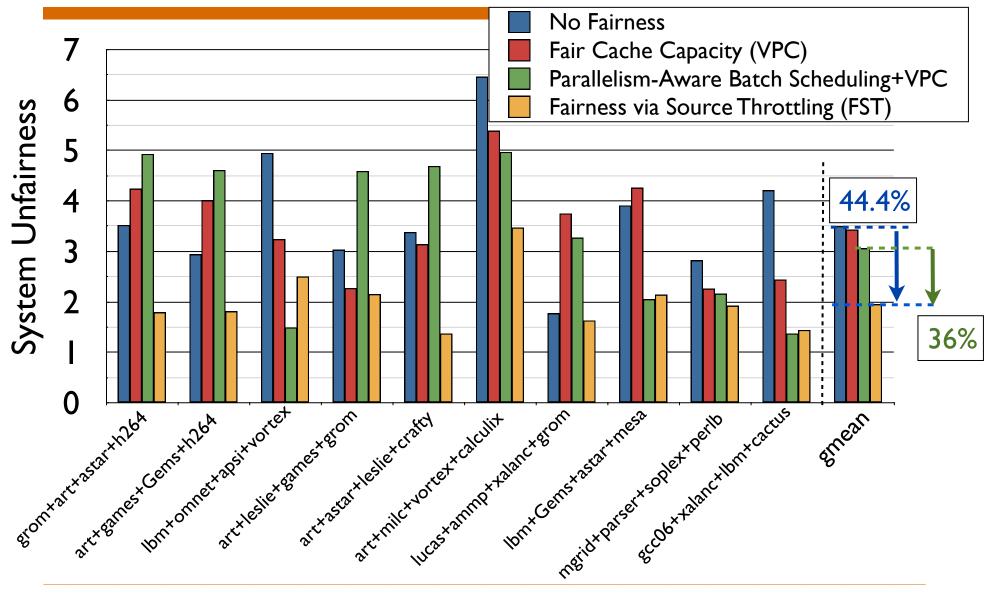


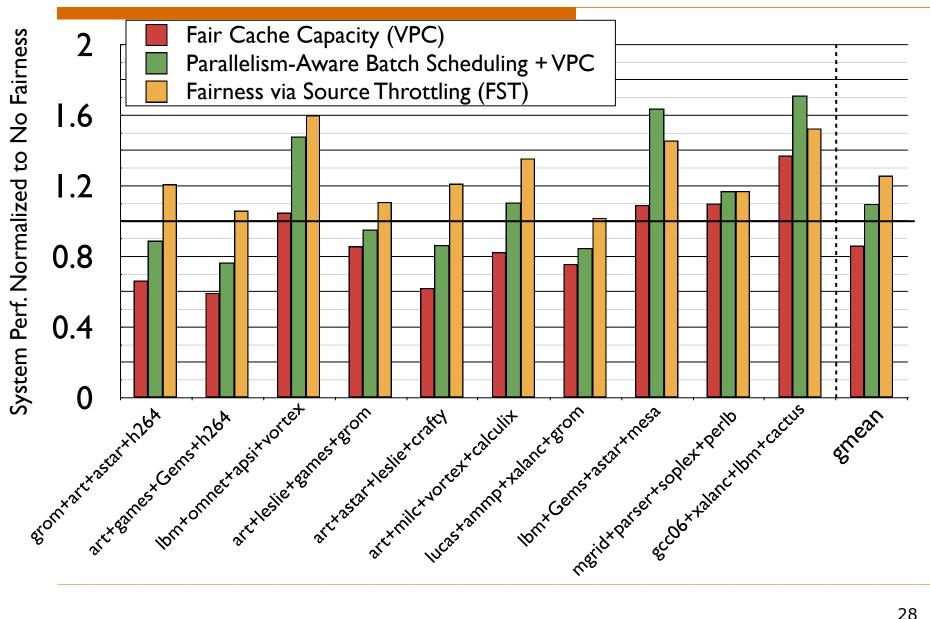


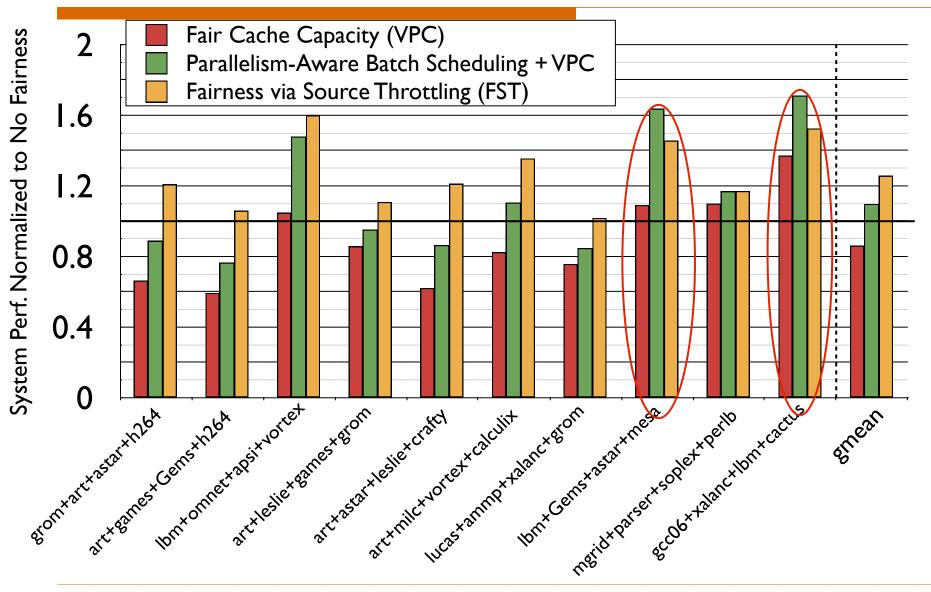


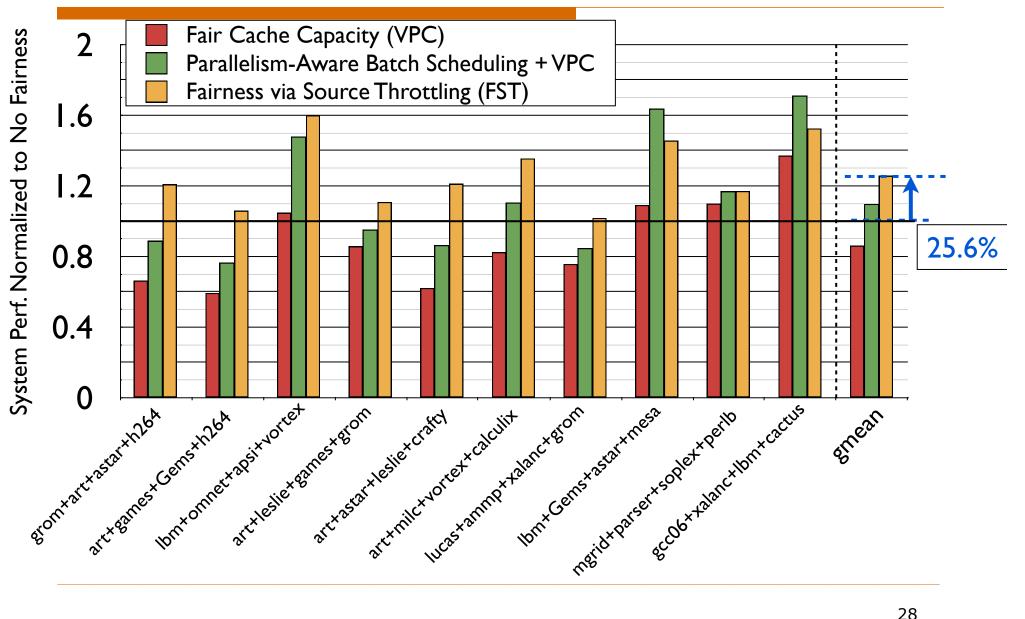


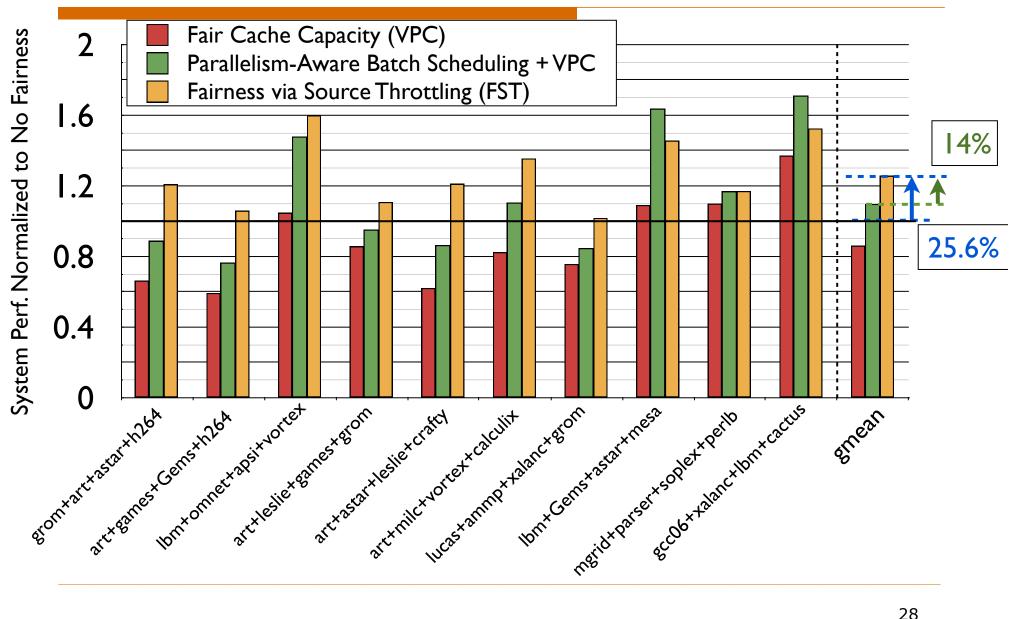












#### Conclusion

- Fairness via Source Throttling (FST) is a new fair and high-performance shared resource management approach for CMPs
- Dynamically monitors unfairness and throttles down sources of interfering memory requests
- Eliminates the need for and complexity of multiple per-resource fairness techniques
- Improves both system fairness and performance
- Incorporates thread weights and enables different fairness objectives

#### Fairness via Source Throttling:

A configurable and high-performance fairness substrate for multi-core memory systems

> Eiman Ebrahimi\* Chang Joo Lee\* Onur Mutlu<sup>‡</sup> Yale N. Patt\*

\* HPS Research Group The University of Texas at Austin **‡** Computer Architecture Laboratory Carnegie Mellon University