Flash Correct-and-Refresh Retention-Aware Error Management for Increased Flash Memory Lifetime

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Executive Summary

- NAND flash memory has low endurance: a flash cell dies after 3k P/E cycles vs. 50k desired → Major scaling challenge for flash memory
- Flash error rate increases exponentially over flash lifetime
- Problem: Stronger error correction codes (ECC) are ineffective and undesirable for improving flash lifetime due to
 - diminishing returns on lifetime with increased correction strength
 - prohibitively high power, area, latency overheads
- Our Goal: Develop techniques to tolerate high error rates w/o strong ECC
- Observation: Retention errors are the dominant errors in MLC NAND flash
 - flash cell loses charge over time; retention errors increase as cell gets worn out
- Solution: Flash Correct-and-Refresh (FCR)
 - Periodically read, correct, and reprogram (in place) or remap each flash page before it accumulates more errors than can be corrected by simple ECC
 - Adapt "refresh" rate to the severity of retention errors (i.e., # of P/E cycles)
- Results: FCR improves flash memory lifetime by 46X with no hardware changes and low energy overhead; outperforms strong ECCs

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Outline

- Executive Summary
- The Problem: Limited Flash Memory Endurance/Lifetime
- Error and ECC Analysis for Flash Memory
- Flash Correct and Refresh Techniques (FCR)
- Evaluation
- Conclusions

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Problem: Limited Endurance of Flash Memory

- NAND flash has limited endurance
 - □ A cell can tolerate a small number of Program/Erase (P/E) cycles
 - □ 3x-nm flash with 2 bits/cell \rightarrow 3K P/E cycles
- Enterprise data storage requirements demand very high endurance
 - □ >50K P/E cycles (10 full disk writes per day for 3-5 years)
- Continued process scaling and more bits per cell will reduce flash endurance
- One potential solution: stronger error correction codes (ECC)
 Stronger ECC not effective enough and inefficient

Decreasing Endurance with Flash Scaling



Ariel Maislos, "A New Era in Embedded Flash Memory", Flash Summit 2011 (Anobit)

- Endurance of flash memory decreasing with scaling and multi-level cells
- Error correction capability required to guarantee storage-class reliability (UBER < 10⁻¹⁵) is increasing exponentially to reach *less* endurance

UBER: Uncorrectable bit error rate. Fraction of erroneous bits after error correction. SAFARI
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The Problem with Stronger Error Correction

- Stronger ECC detects and corrects more raw bit errors → increases P/E cycles endured
- Two shortcomings of stronger ECC:
 - 1. High implementation complexity
 - → Power and area overheads increase super-linearly, but correction capability increases sub-linearly with ECC strength
 - 2. Diminishing returns on flash lifetime improvement
 - → Raw bit error rate increases exponentially with P/E cycles, but correction capability increases sub-linearly with ECC strength

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Methodology: Error and ECC Analysis

- Characterized errors and error rates of 3x-nm MLC NAND flash using an experimental FPGA-based flash platform
 - Cai et al., "Error Patterns in MLC NAND Flash Memory: Measurement, Characterization, and Analysis," DATE 2012.
- Quantified Raw Bit Error Rate (RBER) at a given P/E cycle
 - Raw Bit Error Rate: Fraction of erroneous bits without any correction

- Quantified error correction capability (and area and power consumption) of various BCH-code implementations
 - Identified how much RBER each code can tolerate
 - \rightarrow how many P/E cycles (flash lifetime) each code can sustain

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NAND Flash Error Types

- Four types of errors [Cai+, DATE 2012]
- Caused by common flash operations
 - Read errors
 - Erase errors
 - Program (interference) errors
- Caused by flash cell losing charge over time
 - Retention errors
 - Whether an error happens depends on required retention time
 - Especially problematic in MLC flash because voltage threshold window to determine stored value is smaller

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Observations: Flash Error Analysis



- Raw bit error rate increases exponentially with P/E cycles
- Retention errors are dominant (>99% for 1-year ret. time)
- Retention errors increase with retention time requirement

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ECC Strength Analysis

Error correction capability increases sub-linearly

Power and area overheads increase super-linearly

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Resulting Flash Lifetime with Strong ECC

Lifetime improvement comparison of various BCH codes



Strong ECC is very inefficient at improving lifetime

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Develop new techniques to improve flash lifetime without relying on stronger ECC

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Flash Correct-and-Refresh (FCR)

- Key Observations:
 - Retention errors are the dominant source of errors in flash memory [Cai+ DATE 2012][Tanakamaru+ ISSCC 2011]
 → limit flash lifetime as they increase over time
 - Retention errors can be corrected by "refreshing" each flash page periodically
- Key Idea:
 - Periodically read each flash page,
 - Correct its errors using "weak" ECC, and
 - Either remap it to a new physical page or reprogram it in-place,
 - Before the page accumulates more errors than ECC-correctable
 - Optimization: Adapt refresh rate to endured P/E cycles

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FCR Intuition

	Errors with No refresh	Errors with Periodic refresh
Program Page	×	×
After time T	× × ×	××××
After time 2T	× × × × ×	× × × ×
After time 3T	$\times \times \times \times \times \times \times$	××××

× Retention Error × Program Error

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FCR: Two Key Questions

- How to refresh?
 - Remap a page to another one
 - Reprogram a page (in-place)
 - Hybrid of remap and reprogram
- When to refresh?
 - Fixed period
 - Adapt the period to retention error severity

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 - 3. Adaptive-Rate FCR
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Remapping Based FCR

- Idea: Periodically remap each page to a different physical page (after correcting errors)
 Select next Block
 - Also [Pan et al., HPCA 2012]
 - □ FTL already has support for changing logical → physical flash block/page mappings
 - Deallocated block is erased by garbage collector



• Problem: Causes additional erase operations \rightarrow more wearout

- Bad for read-intensive workloads (few erases really needed)
- Lifetime degrades for such workloads (see paper)

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In-Place Reprogramming Based FCR

- Idea: Periodically reprogram (in-place) each physical page (after correcting errors)
 - Flash programming techniques (ISPP) can correct retention errors in-place by recharging flash cells



■ Problem: Program errors accumulate on the same page → may not be correctable by ECC after some time



Pro: No remapping needed \rightarrow no additional erase operations

Con: Increases the occurrence of program errors

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Program Errors in Flash Memory

- When a cell is being programmed, voltage level of a neighboring cell changes (unintentionally) due to parasitic capacitance coupling
 - \rightarrow can change the data value stored
- Also called program interference error
- Program interference causes neighboring cell voltage to shift to the right

Problem with In-Place Reprogramming



Problem: Program errors can accumulate over time

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Hybrid Reprogramming/Remapping Based FCR

Idea:

- Monitor the count of right-shift errors (after error correction)
- □ If count < threshold, in-place reprogram the page
- Else, remap the page to a new page
- Observation:
 - □ Program errors much less frequent than retention errors → Remapping happens only infrequently
- Benefit:
 - Hybrid FCR greatly reduces erase operations due to remapping

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Adaptive-Rate FCR

Observation:

- Retention error rate strongly depends on the P/E cycles a flash page endured so far
- □ No need to refresh frequently (at all) early in flash lifetime

Idea:

- □ Adapt the refresh rate to the P/E cycles endured by each page
- Increase refresh rate gradually with increasing P/E cycles
- Benefits:
 - Reduces overhead of refresh operations
 - Can use existing FTL mechanisms that keep track of P/E cycles

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Adaptive-Rate FCR (Example)

Select refresh frequency such that error rate is below acceptable rate

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FCR: Other Considerations

Implementation cost

- No hardware changes
- FTL software/firmware needs modification
- Response time impact
 - □ FCR not as frequent as DRAM refresh; low impact
- Adaptation to variations in retention error rate
 Adapt refresh rate based on, e.g., temperature [Liu+ ISCA 2012]
- FCR requires power
 - Enterprise storage systems typically powered on

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Evaluation Methodology

- Experimental flash platform to obtain error rates at different P/E cycles [Cai+ DATE 2012]
- Simulation framework to obtain P/E cycles of real workloads: DiskSim with SSD extensions
- Simulated system: 256GB flash, 4 channels, 8 chips/ channel, 8K blocks/chip, 128 pages/block, 8KB pages

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- Workloads
 - File system applications, databases, web search
 - Categories: Write-heavy, read-heavy, balanced
- Evaluation metrics
 - Lifetime (extrapolated)
 - Energy overhead, P/E cycle overhead

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Normalized Flash Memory Lifetime

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Lifetime Evaluation Takeaways

Significant average lifetime improvement over no refresh

- Adaptive-rate FCR: 46X
- Hybrid reprogramming/remapping based FCR: 31X
- Remapping based FCR: 9X
- FCR lifetime improvement larger than that of stronger ECC
 46X vs. 4X with 32-kbit ECC (over 512-bit ECC)
 FCR is less complex and less costly than stronger ECC
- Lifetime on all workloads improves with Hybrid FCR
 - Remapping based FCR can degrade lifetime on read-heavy WL
 - Lifetime improvement highest in write-heavy workloads

Energy Overhead

Remapping-based Refresh Hybrid Refresh

 Adaptive-rate refresh: <1.8% energy increase until daily refresh is triggered

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Overhead of Additional Erases

- Additional erases happen due to remapping of pages
- Low (2%-20%) for write intensive workloads
- High (up to 10X) for read-intensive workloads
- Improved P/E cycle lifetime of all workloads largely outweighs the additional P/E cycles due to remapping

More Results in the Paper

- Detailed workload analysis
- Effect of refresh rate

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Conclusion

- NAND flash memory lifetime is limited due to uncorrectable errors, which increase over lifetime (P/E cycles)
- Observation: Dominant source of errors in flash memory is retention errors → retention error rate limits lifetime
- Flash Correct-and-Refresh (FCR) techniques reduce retention error rate to improve flash lifetime
 - Periodically read, correct, and remap or reprogram each page before it accumulates more errors than can be corrected
 - Adapt refresh period to the severity of errors
- FCR improves flash lifetime by 46X at no hardware cost
 - More effective and efficient than stronger ECC
 - Can enable better flash memory scaling

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Thank You.

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Backup Slides

Effect of Refresh Rate on Lifetime

Lifetime: Remapping vs. Hybrid FCR

Energy Overhead

Average Lifetime Improvement

Individual Workloads: Remapping-Based FCR

Individual Workloads: Hybrid FCR

Individual Workloads: Adaptive-Rate FCR

P/E Cycle Overhead

- P/E cycle overhead of hybrid FCR is lower than that of remapping-based FCR
- P/E cycle overhead for write-intensive applications is low
 - Remapping-based FCR (20%), Hybrid FCR (2%)
- Read-intensive applications have higher P/E cycle overhead

Motivation for Refresh: A Different Way

- NAND flash endurance can be increased via
 - Stronger error correction codes (4x)
 - Tradeoff guaranteed storage time for one write for high endurance (> 50x)

FTL Implementation

□ FCR can be implemented just as a module in FTL software

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Flash Cells Can Be Reprogrammed In-Place

- Observations:
 - Retention errors occur due to loss of charge
 - Simply recharging the cells can correct the retention errors
 - Flash programming mechanisms can accomplish this recharging
- ISPP (Incremental Step Pulse Programming)
 - Iterative programming mechanism that increases the voltage level of a flash cell step by step
 - After each step, voltage level compared to desired voltage threshold
 - Can inject more electrons but cannot remove electrons