

## Exercise 02

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## 1 Rounding for Bin Packing

For the construction of a PTAS for the Bin packing problem we needed a tricky way of rounding. Here we investigate why this is necessary.

Construct a set of items with sizes  $0 < x_1, \dots, x_n < 1$  with the following property: while it is possible to pack these into some number  $m$  of bins, if we for any  $\varepsilon > 0$  consider the set of items with sizes  $(1 + \varepsilon)x_1, \dots, (1 + \varepsilon)x_n$ , we need at least  $\alpha m$  bins for some  $\alpha > 1$ . What is the best  $\alpha$  you can get and how does it relate to the issue of designing a PTAS for Bin packing?

## 2 Target shooting

Suppose we have a ground set  $S$  and we wish to estimate the size of a subset  $T \subseteq S$ . We can do it by repeatedly sampling from  $S$ : we sample  $m$  times a uniform element from  $S$  and let  $X_i$  to be an indicator for the event of the  $i$ -th sample being an element of  $T$ . One can use the Chernoff bound to prove that the variable  $(X_1 + X_2 + \dots + X_m)/m$  is within  $(1 + \varepsilon)$  multiplicative error of  $|T|/|S|$  with probability at least  $1 - \delta$  if we set

$$m = \Theta\left(\frac{|S|}{|T|} \varepsilon^{-2} \log(1/\delta)\right).$$

1. Argue that the lemma really follows from the Chernoff bound.
2. Argue that if we sample only  $O(\frac{|S|}{|T|})$  elements, with constant probability, we never sample an element from  $T$  throughout the whole procedure and, hence, we cannot get a reasonable estimate.
3. Suppose that we have a FPRAS that returns value that is in the correct range  $[(1 - \varepsilon)OPT, (1 + \varepsilon)OPT]$  with probability  $2/3$ . Recall the median trick from the lecture and show how to amplify the probability to  $1 - \delta$  with  $O(\log 1/\delta)$  calls of the original algorithm.

## 3 Counting satisfying assignments

In the lecture we have seen a FPRAS that estimates the number of satisfying assignments of a given DNF formula. Suppose that we wish to solve this problem and we know that one of the clauses contains only 10 variables. Can you propose a faster algorithm than the one you know from the lecture?