

# How to Write Fast Code

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# Complexity of the DFT

- **Measure:  $L_c$ ,  $2 \leq c$** 
  - Complex adds count 1
  - Complex mults by a constant  $a$  with  $|a| < c$  counts 1
  - $L_2$  is strictest,  $L_{\text{infinity}}$  the loosest (and most natural)
- $n = 2^k$ :  $L_2(\text{DFT}_n) \leq 3/2 n \log_2(n)$  (using Cooley-Tukey FFT)
- General  $n$ :  $L_2(\text{DFT}_n) \leq 8 n \log_2(n)$  (needs Bluestein FFT)
- **Theorem (Morgenstern):  $c < \text{infinity}$**   
 $L_c(\text{DFT}_n) \geq \frac{1}{2} n \log_c(n)$   
Implies: in the measure  $L_c$ , the DFT is  $\Theta(n \log(n))$
- More details: [Algebraic Complexity Theory](#)

# The History of Fast Transform Algorithms

- ... starts with the FFT
- The advent of digital signal processing is often attributed to the FFT (Cooley-Tukey 1965)
- **History:**
  - Around 1805: FFT discovered by Gauss [1]  
(Fourier publishes the concept of Fourier analysis in 1807!)
  - 1965: Rediscovered by Cooley-Tukey
  - 2002: James W. Cooley receives the IEEE Jack S. Kilby Signal Processing Medal "For pioneering the Fast Fourier Transform (FFT) algorithm."

# Carl-Friedrich Gauss



1777 - 1855

- **Contender for the greatest mathematician of all times**
- **Some contributions:** Modular arithmetic, least square analysis, normal distribution, fundamental theorem of algebra, Gauss elimination, Gauss quadrature, Gauss-Seidel, non-euclidean geometry, ...