

Processing Data Where It Makes Sense in Modern Computing Systems: Enabling In-Memory Computation

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19 November 2019

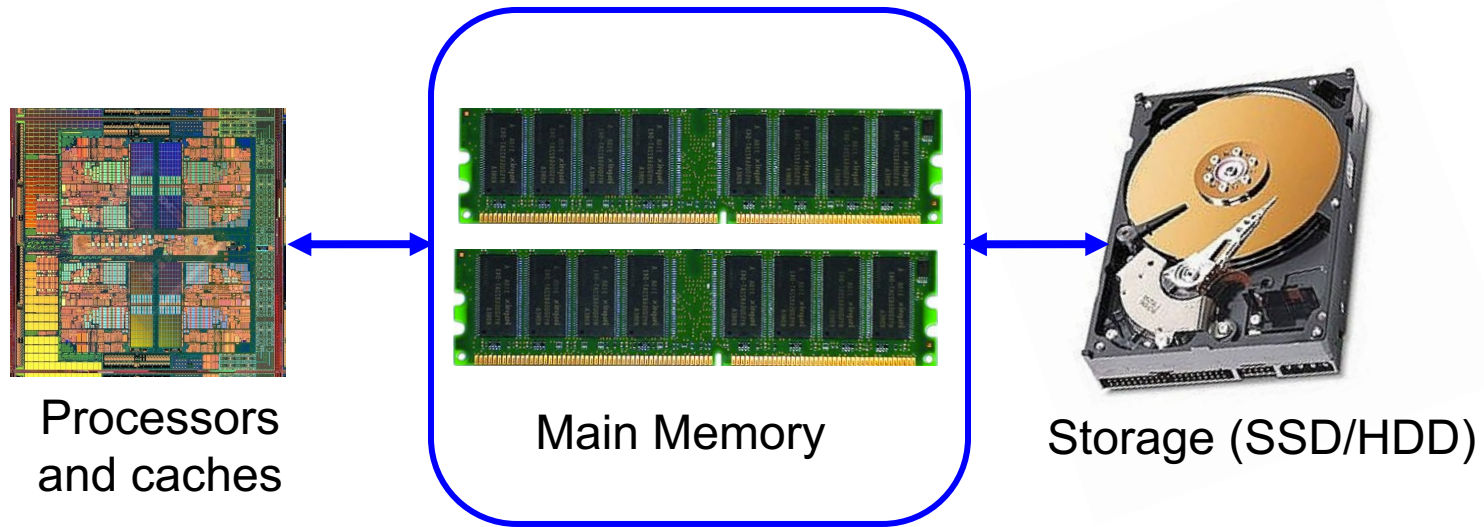
ICCD Keynote

SAFARI

ETH zürich

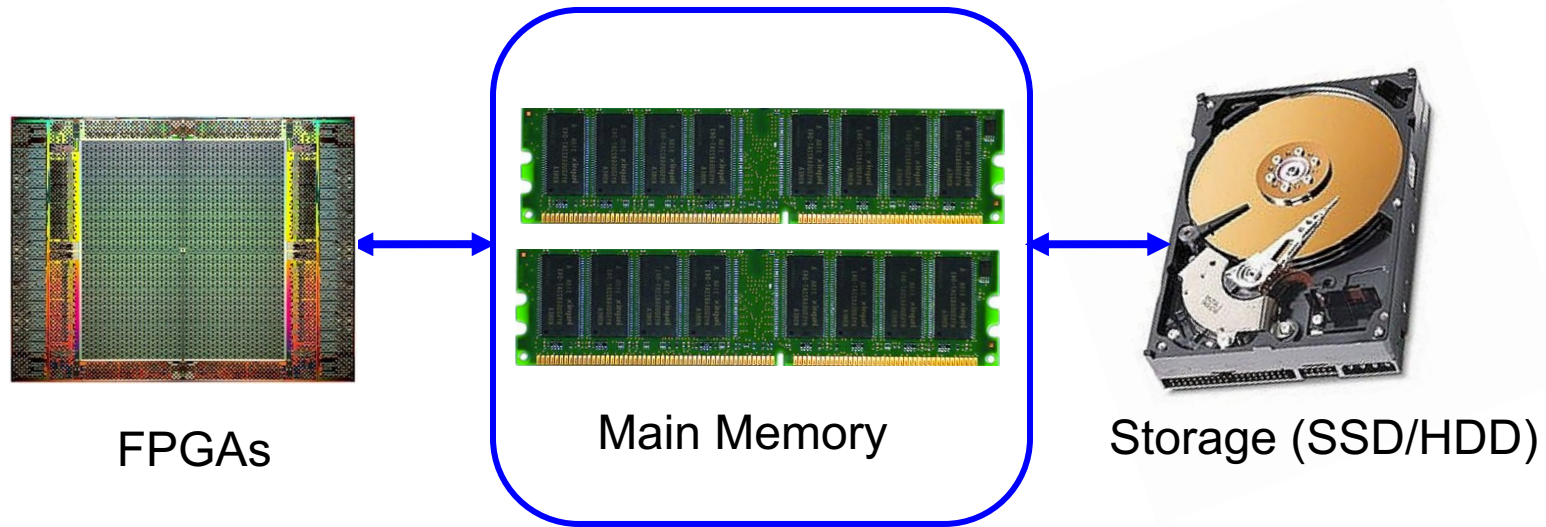
Carnegie Mellon

The Main Memory System



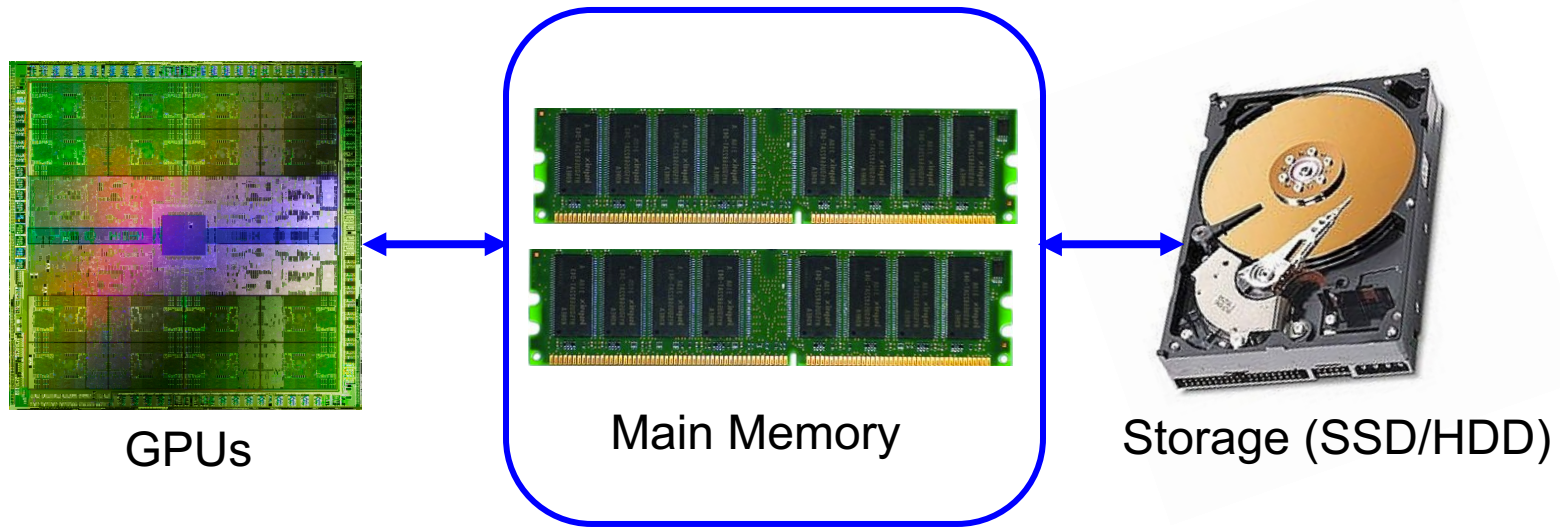
- Main memory is a critical component of all computing systems: server, mobile, embedded, desktop, sensor
- Main memory system must scale (in *size, technology, efficiency, cost, and management algorithms*) to maintain performance growth and technology scaling benefits

The Main Memory System



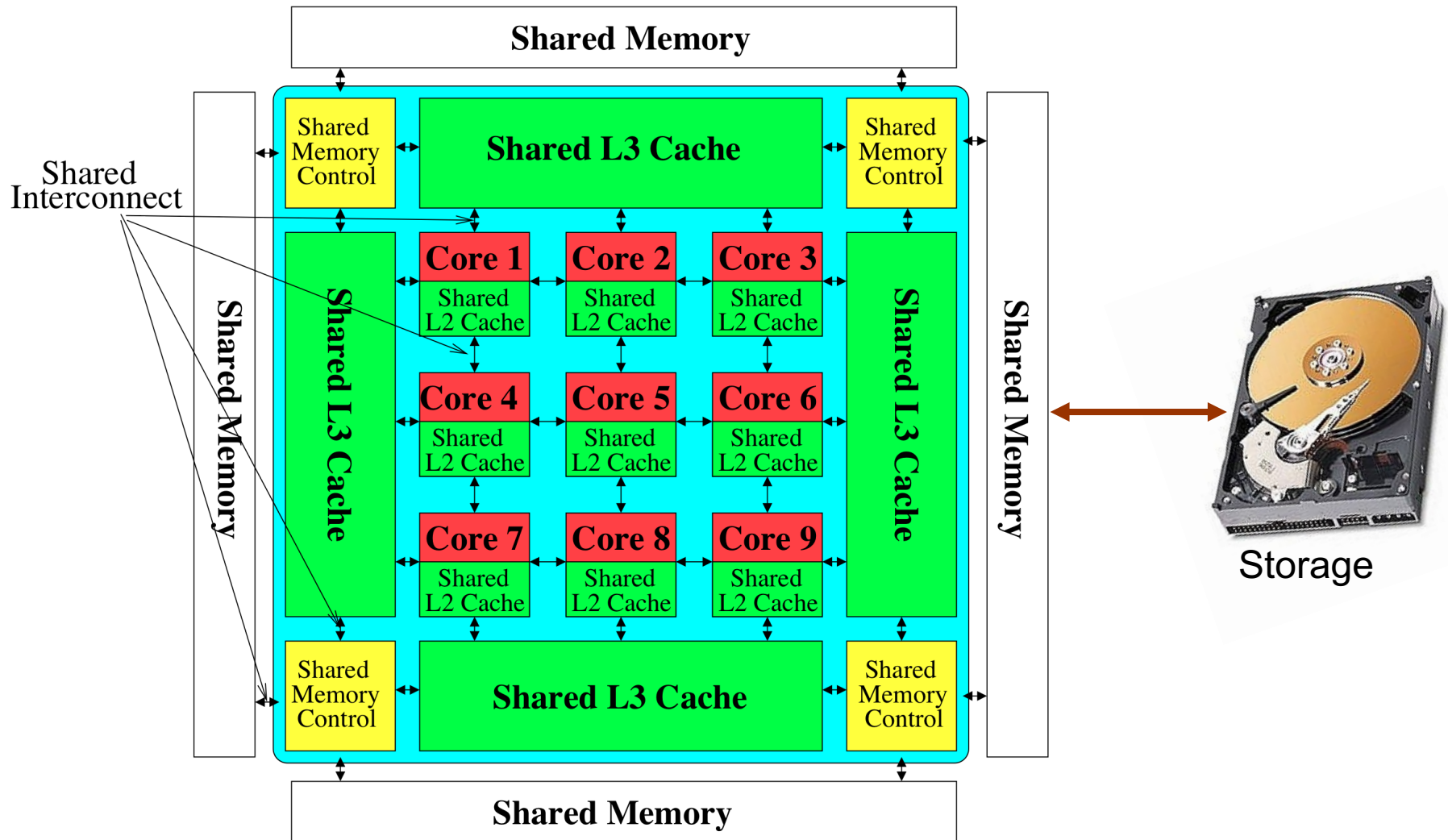
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Memory System: A *Shared Resource* View



Most of the system is dedicated to storing and moving data

State of the Main Memory System

- Recent technology, architecture, and application trends
 - lead to new requirements
 - exacerbate old requirements
- DRAM and memory controllers, as we know them today, are (will be) unlikely to satisfy all requirements
- Some emerging non-volatile memory technologies (e.g., PCM) enable new opportunities: memory+storage merging
- We need to rethink the main memory system
 - to fix DRAM issues and enable emerging technologies
 - to satisfy all requirements

Major Trends Affecting Main Memory (I)

- Need for main memory capacity, bandwidth, QoS increasing
- Main memory energy/power is a key system design concern
- DRAM technology scaling is ending

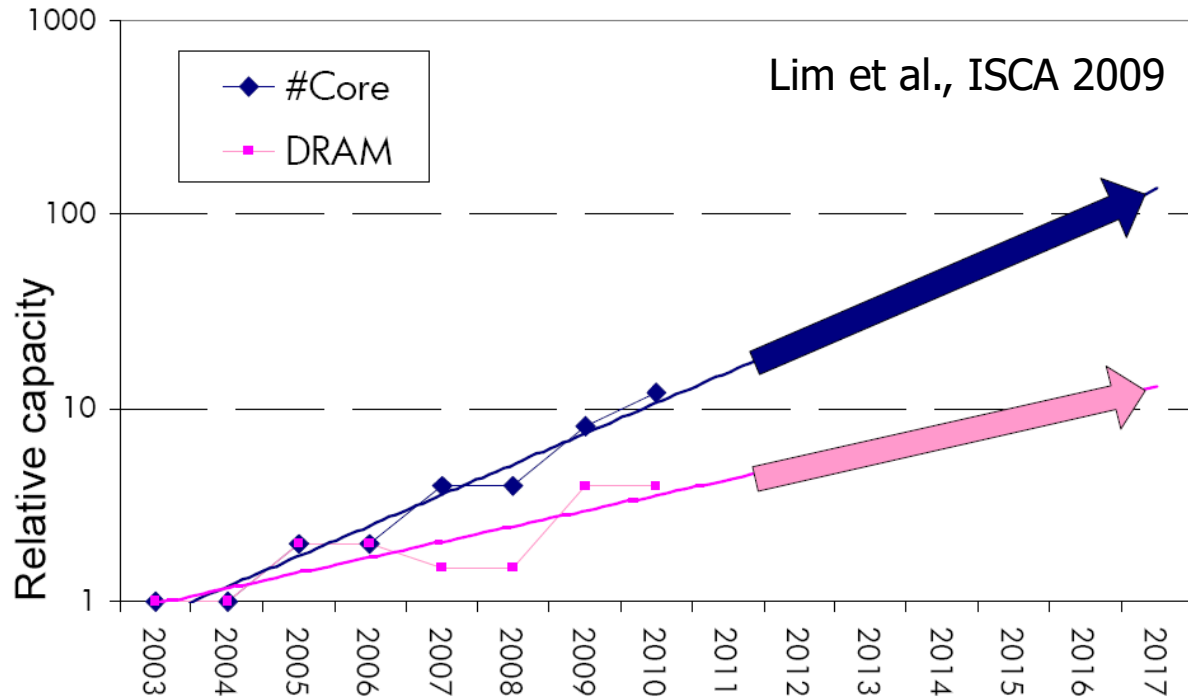
Major Trends Affecting Main Memory (II)

- Need for main memory capacity, bandwidth, QoS increasing
 - **Multi-core**: increasing number of cores/agents
 - **Data-intensive applications**: increasing demand/hunger for data
 - **Consolidation**: cloud computing, GPUs, mobile, heterogeneity
- Main memory energy/power is a key system design concern
- DRAM technology scaling is ending

Example: The Memory Capacity Gap

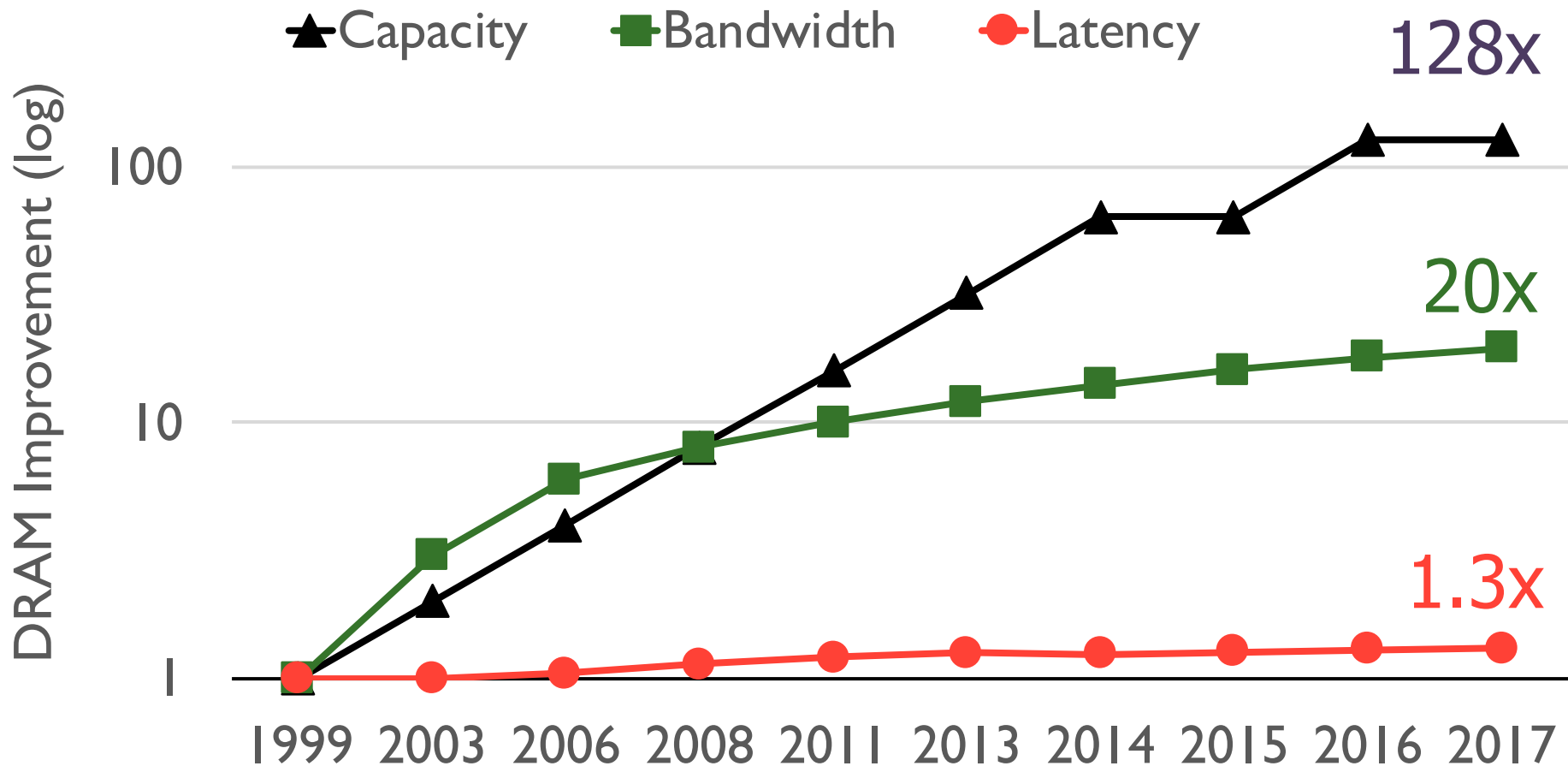
Core count doubling ~ every 2 years

DRAM DIMM capacity doubling ~ every 3 years

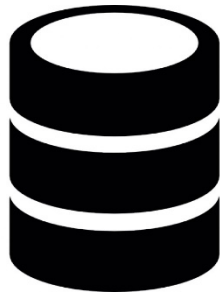


- *Memory capacity per core* expected to drop by 30% every two years
- Trends worse for *memory bandwidth per core*!

DRAM Capacity, Bandwidth & Latency



DRAM Is Critical for Performance



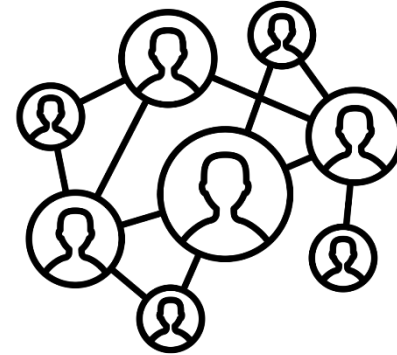
In-memory Databases

[Mao+, EuroSys'12;
Clapp+ (Intel), IISWC'15]



In-Memory Data Analytics

[Clapp+ (Intel), IISWC'15;
Awan+, BDCloud'15]



Graph/Tree Processing

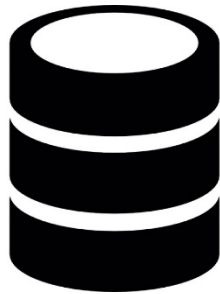
[Xu+, IISWC'12; Umuroglu+, FPL'15]



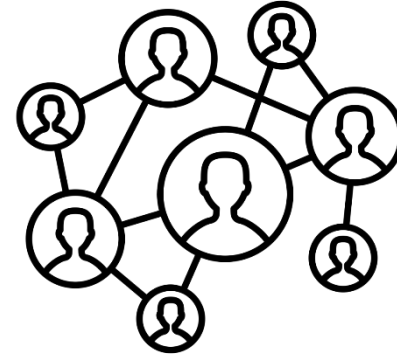
Datacenter Workloads

[Kanev+ (Google), ISCA'15]

DRAM Is Critical for Performance



In-memory Databases



Graph/Tree Processing

Memory → performance bottleneck



In-Memory Data Analytics

[Clapp+ (Intel), IISWC'15;
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Datacenter Workloads

[Kanev+ (Google), ISCA'15]

DRAM Is Critical for Performance



Chrome

Google's web browser



TensorFlow Mobile

Google's machine learning
framework

VP9



Video Playback

Google's **video codec**

VP9



Video Capture

Google's **video codec**

DRAM Is Critical for Performance



Chrome



TensorFlow Mobile

Memory → performance bottleneck

VP9



Video Playback

Google's **video codec**

VP9



Video Capture

Google's **video codec**

Major Trends Affecting Main Memory (III)

- Need for main memory capacity, bandwidth, QoS increasing
- Main memory energy/power is a key system design concern
 - ~40-50% energy spent in off-chip memory hierarchy [Lefurgy, IEEE Computer'03] >40% power in DRAM [Ware, HPCA'10][Paul, ISCA'15]
 - DRAM consumes power even when not used (periodic refresh)
- DRAM technology scaling is ending

Energy Waste in Mobile Devices

- Amirali Boroumand, Saugata Ghose, Youngsok Kim, Rachata Ausavarungnirun, Eric Shiu, Rahul Thakur, Daehyun Kim, Aki Kuusela, Allan Knies, Parthasarathy Ranganathan, and Onur Mutlu, **"Google Workloads for Consumer Devices: Mitigating Data Movement Bottlenecks"** *Proceedings of the 23rd International Conference on Architectural Support for Programming Languages and Operating Systems (ASPLOS)*, Williamsburg, VA, USA, March 2018.

**62.7% of the total system energy
is spent on data movement**

Google Workloads for Consumer Devices: Mitigating Data Movement Bottlenecks

Amirali Boroumand¹

Saugata Ghose¹

Youngsok Kim²

Rachata Ausavarungnirun¹

Eric Shiu³

Rahul Thakur³

Daehyun Kim^{4,3}

Aki Kuusela³

Allan Knies³

Parthasarathy Ranganathan³

Onur Mutlu^{5,1}

Major Trends Affecting Main Memory (IV)

- Need for main memory capacity, bandwidth, QoS increasing
- Main memory energy/power is a key system design concern
- DRAM technology scaling is ending
 - ITRS projects DRAM will not scale easily below X nm
 - Scaling has provided many benefits:
 - higher capacity (density), lower cost, lower energy

Major Trends Affecting Main Memory (V)

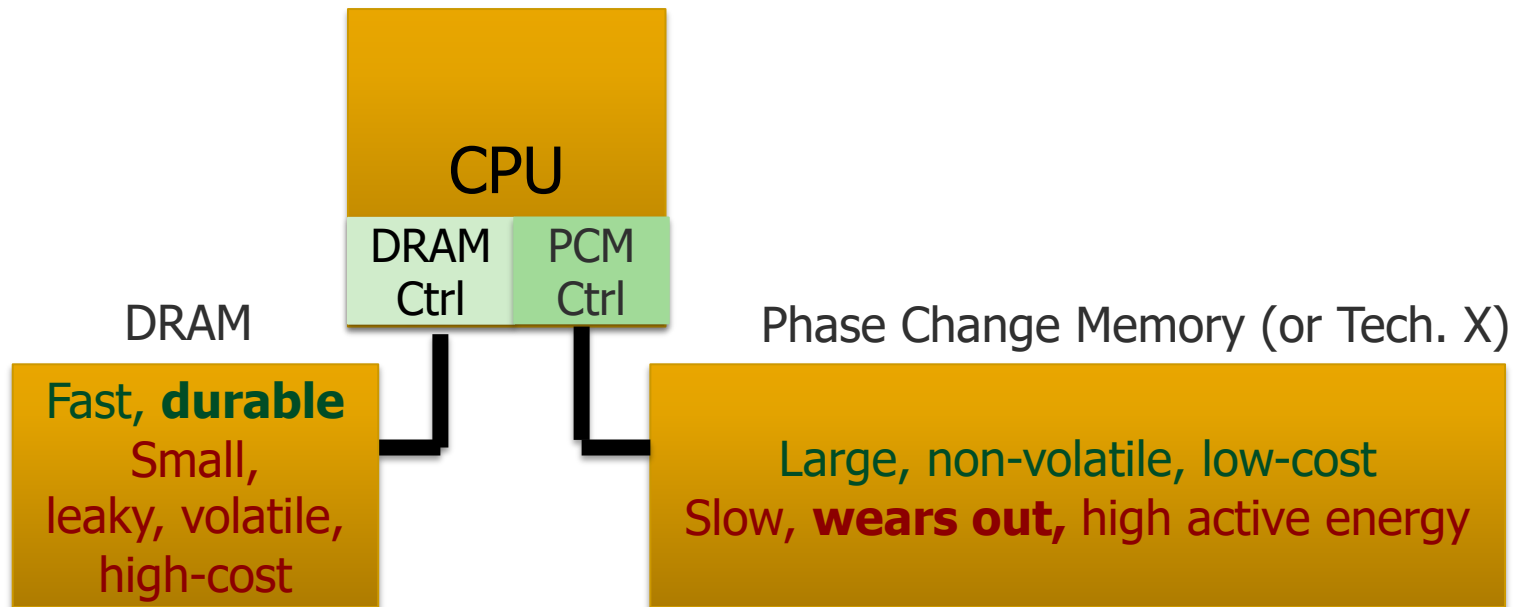
- DRAM scaling has already become increasingly difficult
 - Increasing cell leakage current, reduced cell reliability, increasing manufacturing difficulties [Kim+ ISCA 2014], [Liu+ ISCA 2013], [Mutlu IMW 2013], [Mutlu DATE 2017]
 - **Difficult to significantly improve capacity, energy**
- **Emerging memory technologies** are promising

Major Trends Affecting Main Memory (V)

- DRAM scaling has already become increasingly difficult
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 - **Difficult to significantly improve capacity, energy**
- **Emerging memory technologies** are promising

3D-Stacked DRAM	higher bandwidth	smaller capacity
Reduced-Latency DRAM (e.g., RL/TL-DRAM, FLY-RAM)	lower latency	higher cost
Low-Power DRAM (e.g., LPDDR3, LPDDR4, Voltron)	lower power	higher latency higher cost
Non-Volatile Memory (NVM) (e.g., PCM, STTRAM, ReRAM, 3D Xpoint)	larger capacity	higher latency higher dynamic power lower endurance

Major Trend: Hybrid Main Memory



Hardware/software manage data allocation and movement
to achieve the best of multiple technologies

Meza+, "[Enabling Efficient and Scalable Hybrid Memories](#)," IEEE Comp. Arch. Letters, 2012.
Yoon+, "[Row Buffer Locality Aware Caching Policies for Hybrid Memories](#)," ICCD 2012 Best Paper Award.

Main Memory Needs Intelligent Controllers

Industry Is Writing Papers About It, Too

DRAM Process Scaling Challenges

❖ Refresh

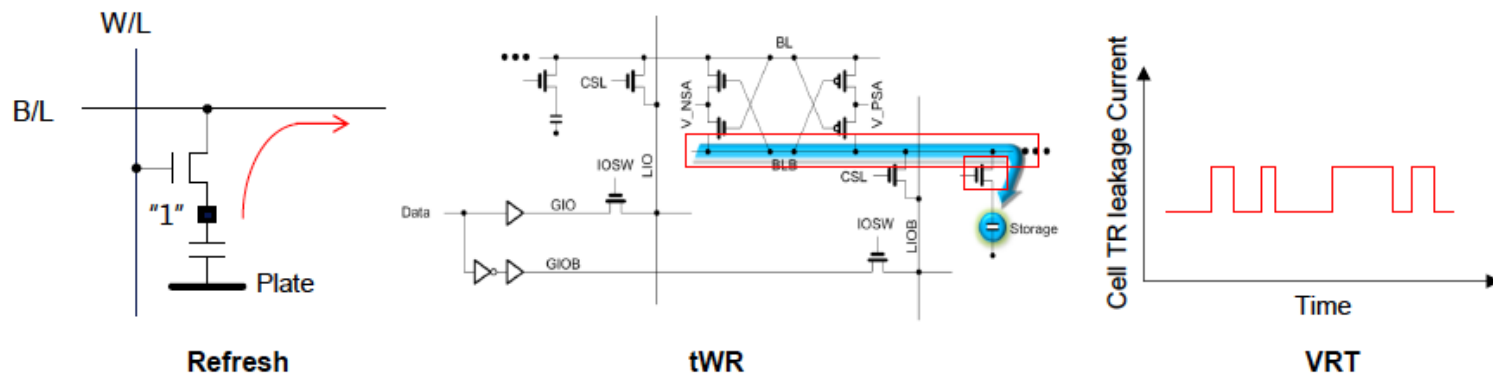
- Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance
- Leakage current of cell access transistors increasing

❖ tWR

- Contact resistance between the cell capacitor and access transistor increasing
- On-current of the cell access transistor decreasing
- Bit-line resistance increasing

❖ VRT

- Occurring more frequently with cell capacitance decreasing



Call for Intelligent Memory Controllers

DRAM Process Scaling Challenges

❖ Refresh

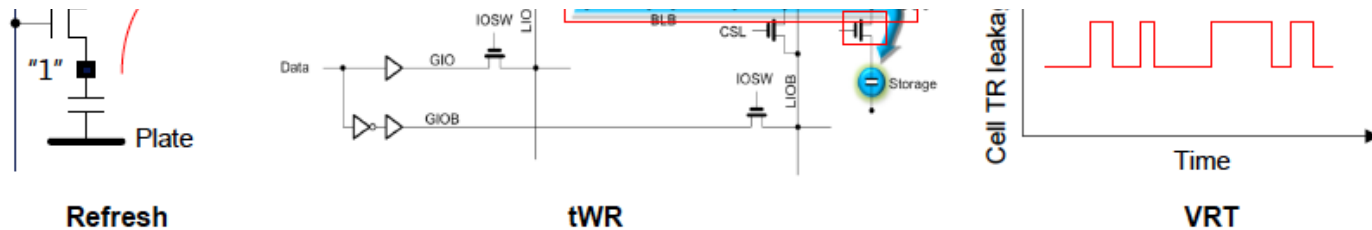
- Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance

THE MEMORY FORUM 2014

Co-Architecting Controllers and DRAM to Enhance DRAM Process Scaling

Uksong Kang, Hak-soo Yu, Churoo Park, *Hongzhong Zheng,
**John Halbert, **Kuljit Bains, SeongJin Jang, and Joo Sun Choi

*Samsung Electronics, Hwasung, Korea / *Samsung Electronics, San Jose / **Intel*



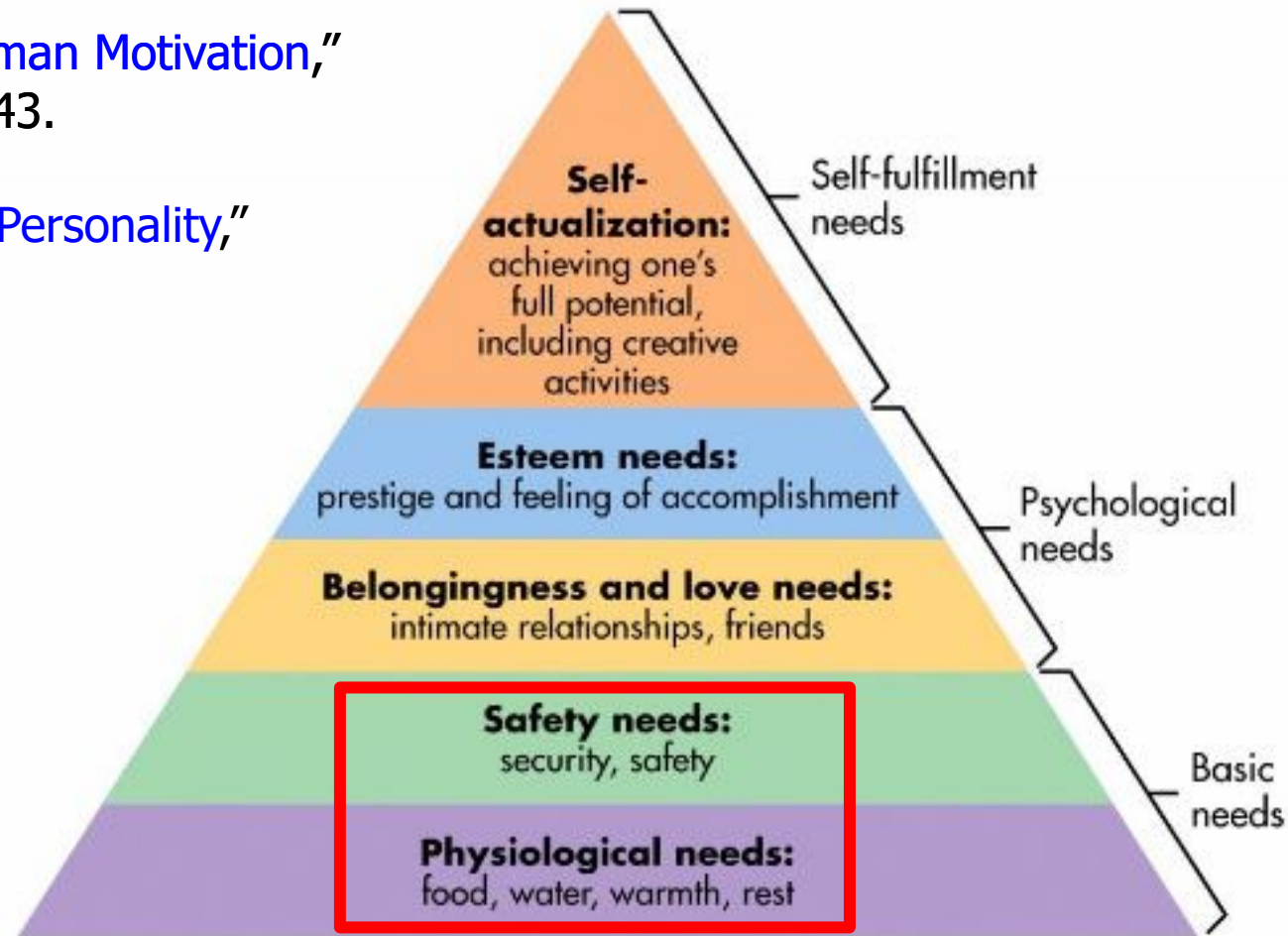
Agenda

- Major Trends Affecting Main Memory
- The Need for Intelligent Memory Controllers
 - Bottom Up: Push from Circuits and Devices
 - Top Down: Pull from Systems and Applications
- Processing in Memory: Two Directions
 - Minimally Changing Memory Chips
 - Exploiting 3D-Stacked Memory
- How to Enable Adoption of Processing in Memory
- Conclusion

Maslow's (Human) Hierarchy of Needs

Maslow, "A Theory of Human Motivation,"
Psychological Review, 1943.

Maslow, "Motivation and Personality,"
Book, 1954-1970.



- We need to start with reliability and security...

How Reliable/Secure/Safe is This Bridge?



Collapse of the “Galloping Gertie”



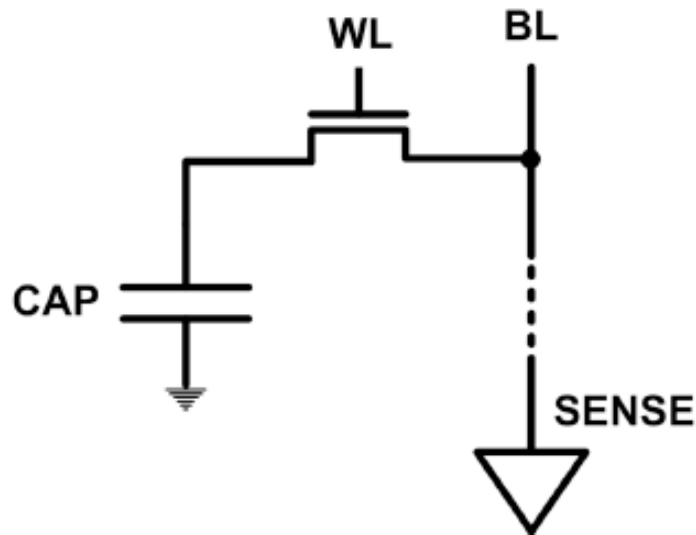
How Secure Are These People?



Security is about preventing unforeseen consequences

The DRAM Scaling Problem

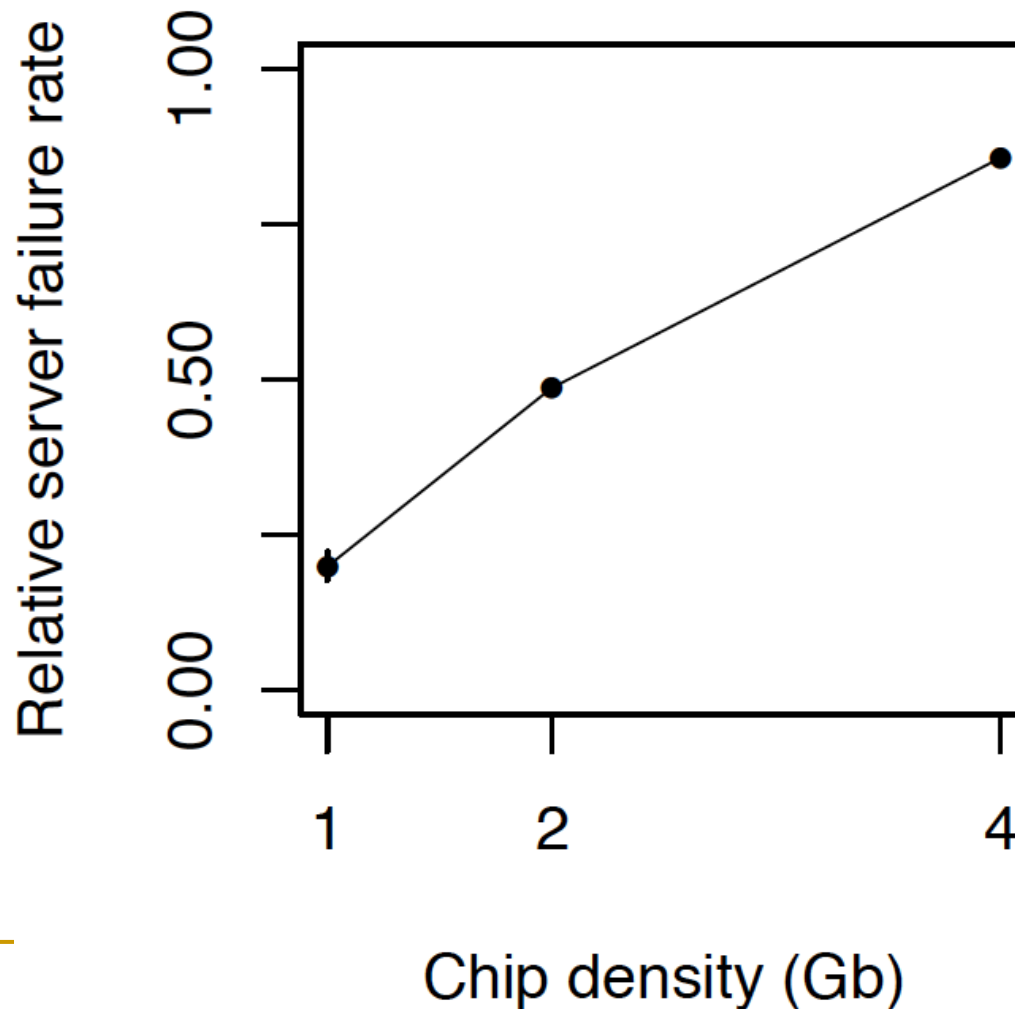
- DRAM stores charge in a capacitor (charge-based memory)
 - Capacitor must be large enough for reliable sensing
 - Access transistor should be large enough for low leakage and high retention time
 - Scaling beyond 40-35nm (2013) is challenging [ITRS, 2009]



- DRAM capacity, cost, and energy/power hard to scale

As Memory Scales, It Becomes Unreliable

- Data from all of Facebook's servers worldwide
- Meza+, "Revisiting Memory Errors in Large-Scale Production Data Centers," DSN'15.



*Intuition:
quadratic
increase
in
capacity*

Large-Scale Failure Analysis of DRAM Chips

- Analysis and modeling of memory errors found in all of Facebook's server fleet
- Justin Meza, Qiang Wu, Sanjeev Kumar, and Onur Mutlu,
"Revisiting Memory Errors in Large-Scale Production Data Centers: Analysis and Modeling of New Trends from the Field"
Proceedings of the 45th Annual IEEE/IFIP International Conference on Dependable Systems and Networks (DSN), Rio de Janeiro, Brazil, June 2015.
[[Slides \(pptx\)](#)] [[pdf](#)] [[DRAM Error Model](#)]

Revisiting Memory Errors in Large-Scale Production Data Centers: Analysis and Modeling of New Trends from the Field

Justin Meza Qiang Wu* Sanjeev Kumar* Onur Mutlu
Carnegie Mellon University * Facebook, Inc.

Infrastructures to Understand Such Issues



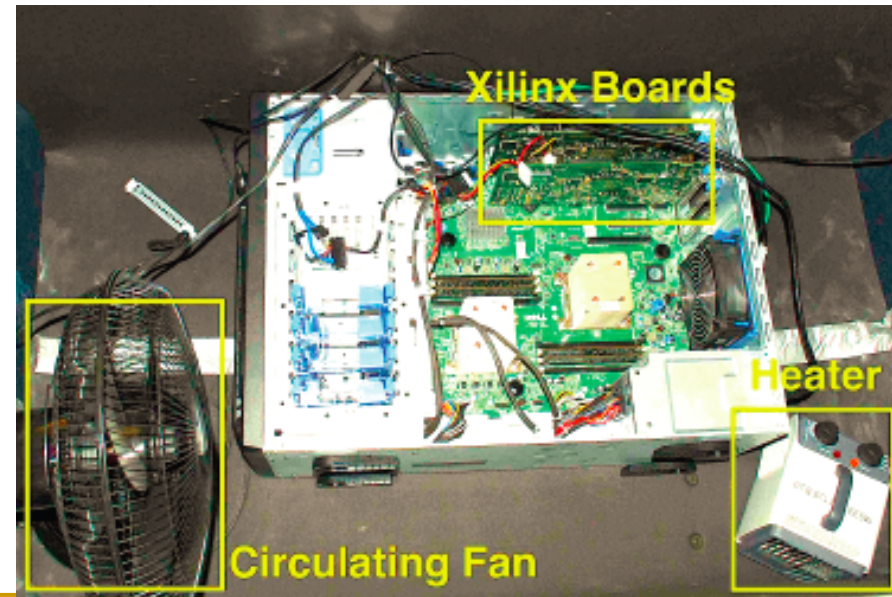
An Experimental Study of Data Retention Behavior in Modern DRAM Devices: Implications for Retention Time Profiling Mechanisms (Liu et al., ISCA 2013)

The Efficacy of Error Mitigation Techniques for DRAM Retention Failures: A Comparative Experimental Study (Khan et al., SIGMETRICS 2014)

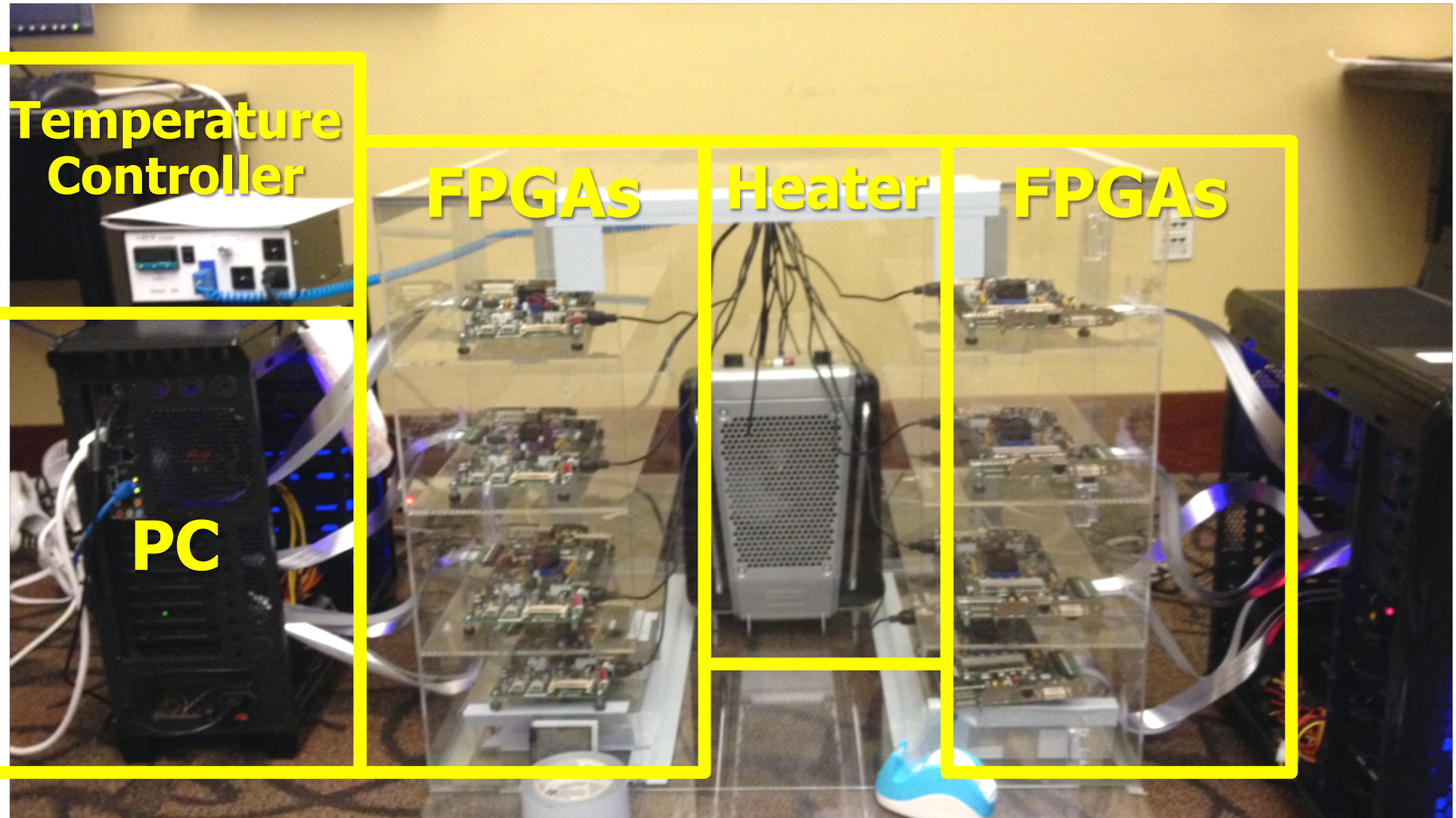
Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors (Kim et al., ISCA 2014)

Adaptive-Latency DRAM: Optimizing DRAM Timing for the Common-Case (Lee et al., HPCA 2015)

AVATAR: A Variable-Retention-Time (VRT) Aware Refresh for DRAM Systems (Qureshi et al., DSN 2015)

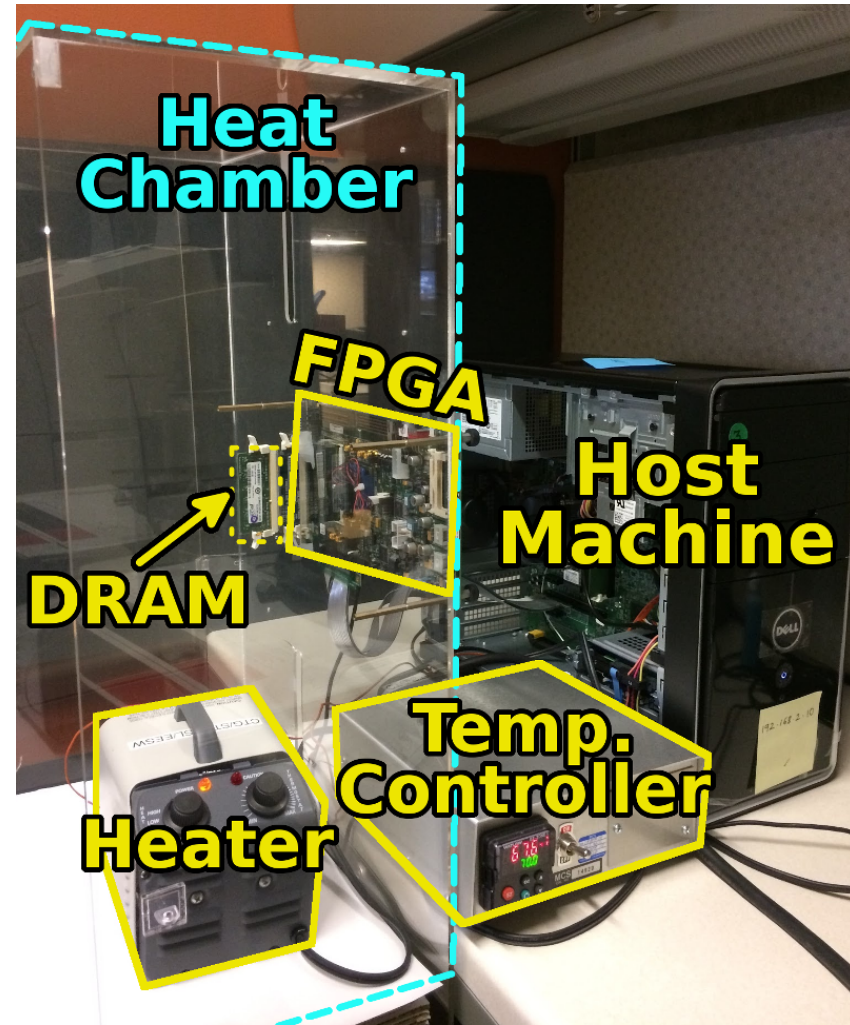


Infrastructures to Understand Such Issues



SoftMC: Open Source DRAM Infrastructure

- Hasan Hassan et al., “**SoftMC: A Flexible and Practical Open-Source Infrastructure for Enabling Experimental DRAM Studies**,” HPCA 2017.
- Flexible
- Easy to Use (C++ API)
- Open-source
github.com/CMU-SAFARI/SoftMC



- <https://github.com/CMU-SAFARI/SoftMC>

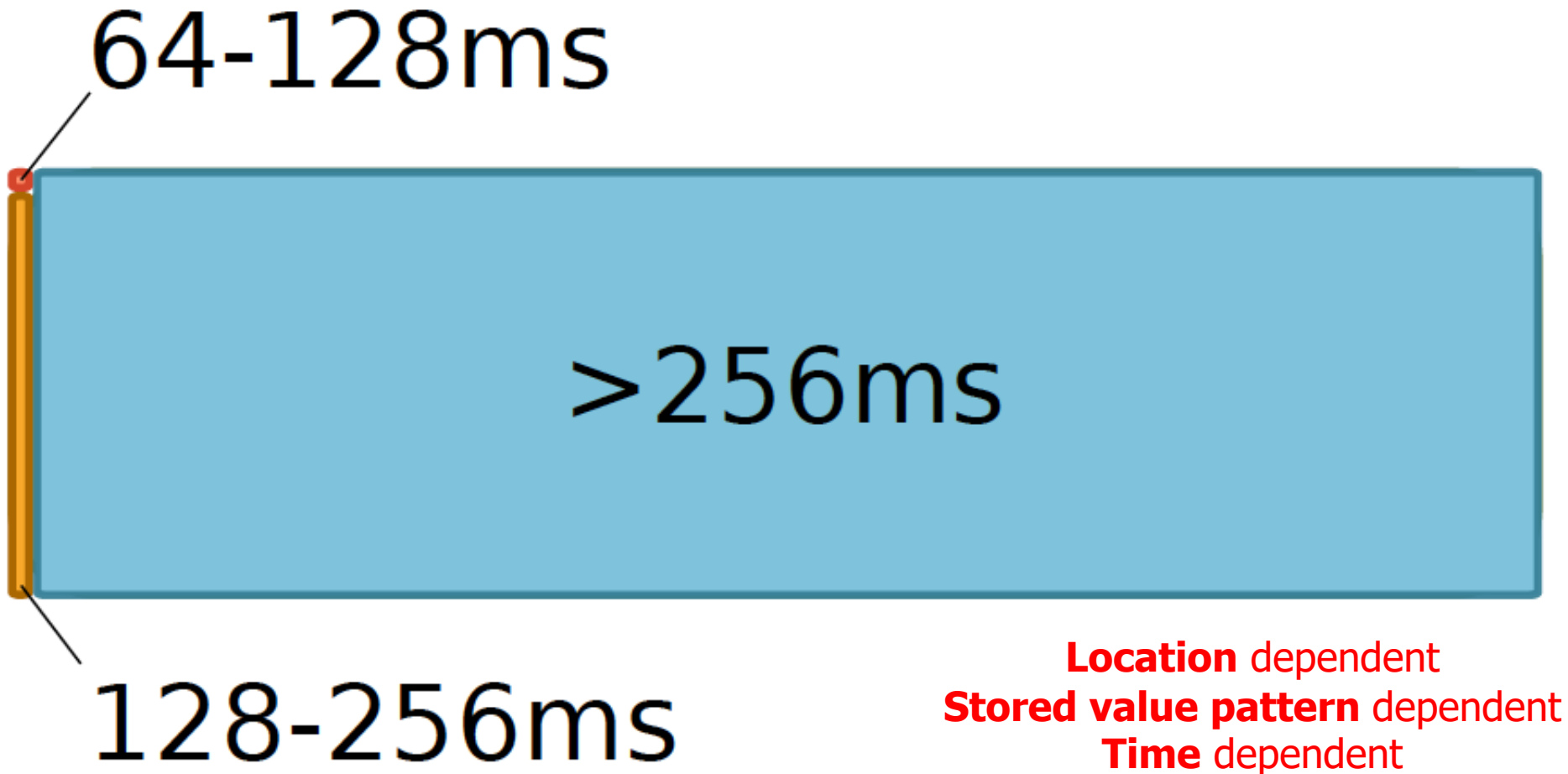
SoftMC: A Flexible and Practical Open-Source Infrastructure for Enabling Experimental DRAM Studies

Hasan Hassan^{1,2,3} Nandita Vijaykumar³ Samira Khan^{4,3} Saugata Ghose³ Kevin Chang³
Gennady Pekhimenko^{5,3} Donghyuk Lee^{6,3} Oguz Ergin² Onur Mutlu^{1,3}

¹*ETH Zürich* ²*TOBB University of Economics & Technology* ³*Carnegie Mellon University*
⁴*University of Virginia* ⁵*Microsoft Research* ⁶*NVIDIA Research*

Data Retention in Memory [Liu et al., ISCA 2013]

- Retention Time Profile of DRAM looks like this:



Main Memory Needs Intelligent Controllers

More on DRAM Refresh (I)

- Jamie Liu, Ben Jaiyen, Richard Veras, and Onur Mutlu,
"RAIDR: Retention-Aware Intelligent DRAM Refresh"
*Proceedings of the 39th International Symposium on
Computer Architecture (ISCA)*, Portland, OR, June 2012.
[Slides \(pdf\)](#)

RAIDR: Retention-Aware Intelligent DRAM Refresh

Jamie Liu Ben Jaiyen Richard Veras Onur Mutlu
Carnegie Mellon University

More on DRAM Refresh (II)

- Jamie Liu, Ben Jaiyen, Yoongu Kim, Chris Wilkerson, and Onur Mutlu,
"An Experimental Study of Data Retention Behavior in Modern DRAM Devices: Implications for Retention Time Profiling Mechanisms"
Proceedings of the 40th International Symposium on Computer Architecture (ISCA), Tel-Aviv, Israel, June 2013. [Slides \(ppt\)](#) [Slides \(pdf\)](#)

An Experimental Study of Data Retention Behavior in Modern DRAM Devices: Implications for Retention Time Profiling Mechanisms

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onur@cmu.edu

More on DRAM Refresh (III)

- Minesh Patel, Jeremie S. Kim, and Onur Mutlu,
"The Reach Profiler (REAPER): Enabling the Mitigation of DRAM Retention Failures via Profiling at Aggressive Conditions"
Proceedings of the 44th International Symposium on Computer Architecture (ISCA), Toronto, Canada, June 2017.
[[Slides \(pptx\)](#)] [[pdf](#)]
[[Lightning Session Slides \(pptx\)](#)] [[pdf](#)]
- First experimental analysis of (mobile) LPDDR4 chips
- Analyzes the complex tradeoff space of retention time profiling
- Idea: enable fast and robust profiling at higher refresh intervals & temperatures

The Reach Profiler (REAPER): Enabling the Mitigation of DRAM Retention Failures via Profiling at Aggressive Conditions

Minesh Patel^{§‡} Jeremie S. Kim^{‡§} Onur Mutlu^{§‡}
[§]ETH Zürich [‡]Carnegie Mellon University

A Curious Discovery [Kim et al., ISCA 2014]

One can
predictably induce errors
in most DRAM memory chips

DRAM RowHammer

A simple hardware failure mechanism
can create a widespread
system security vulnerability

WIRED

Forget Software—Now Hackers Are Exploiting Physics

BUSINESS	CULTURE	DESIGN	GEAR	SCIENCE
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ANDY GREENBERG SECURITY 08.31.16 7:00 AM

SHARE



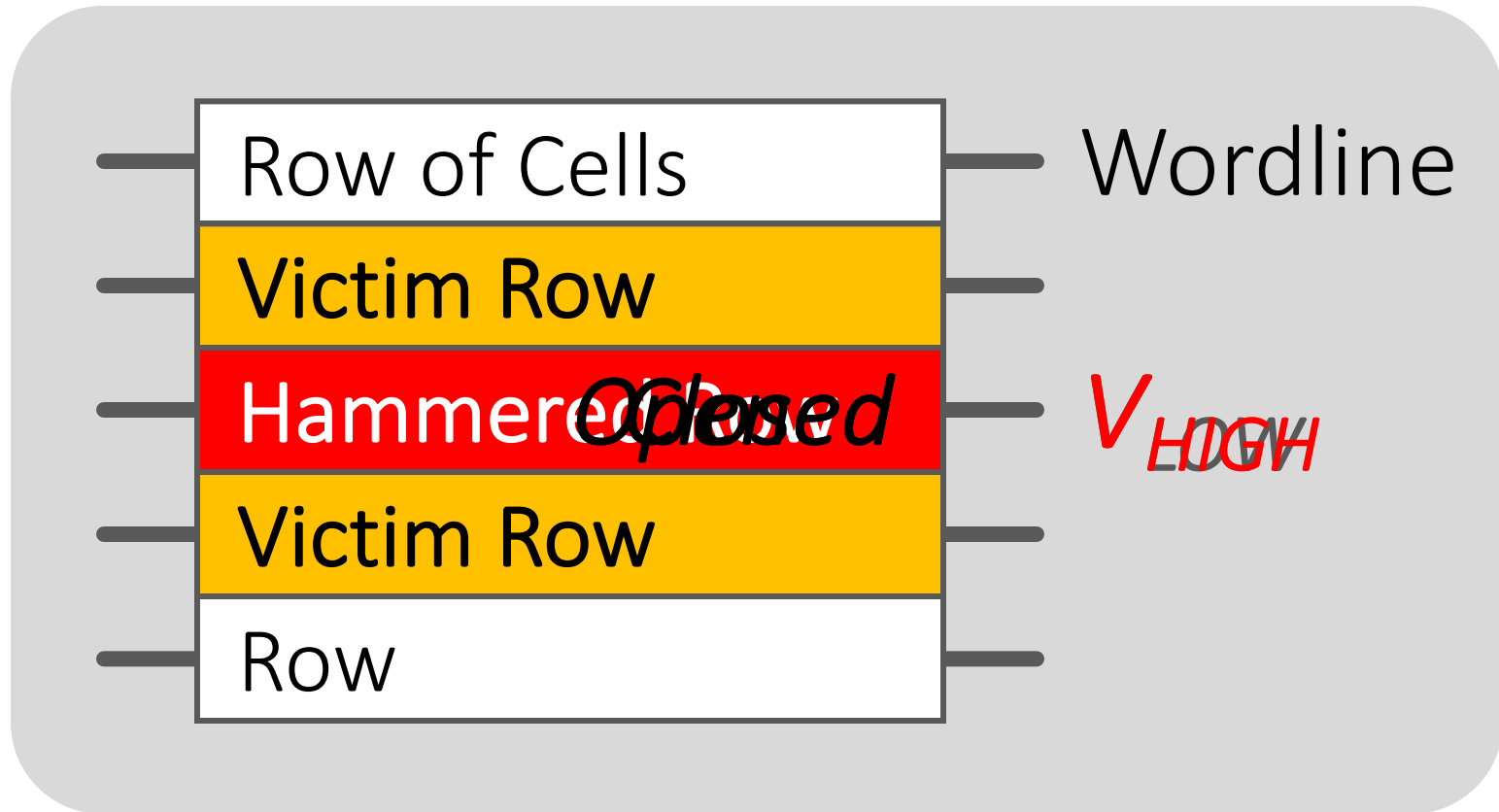
SHARE
18276



TWEET

FORGET SOFTWARE—NOW HACKERS ARE EXPLOITING PHYSICS

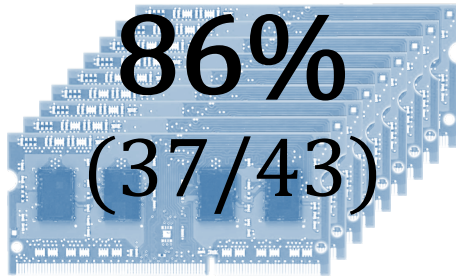
Modern DRAM is Prone to Disturbance Errors



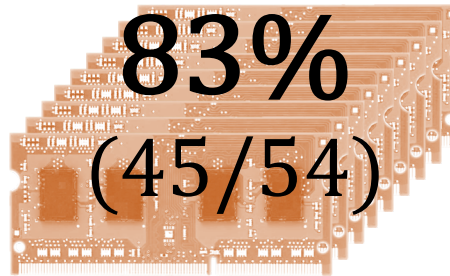
Repeatedly reading a row enough times (before memory gets refreshed) induces **disturbance errors** in adjacent rows in **most real DRAM chips you can buy today**

Most DRAM Modules Are Vulnerable

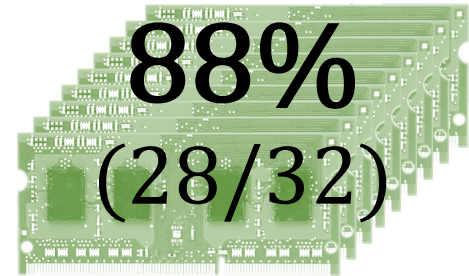
A company



B company



C company

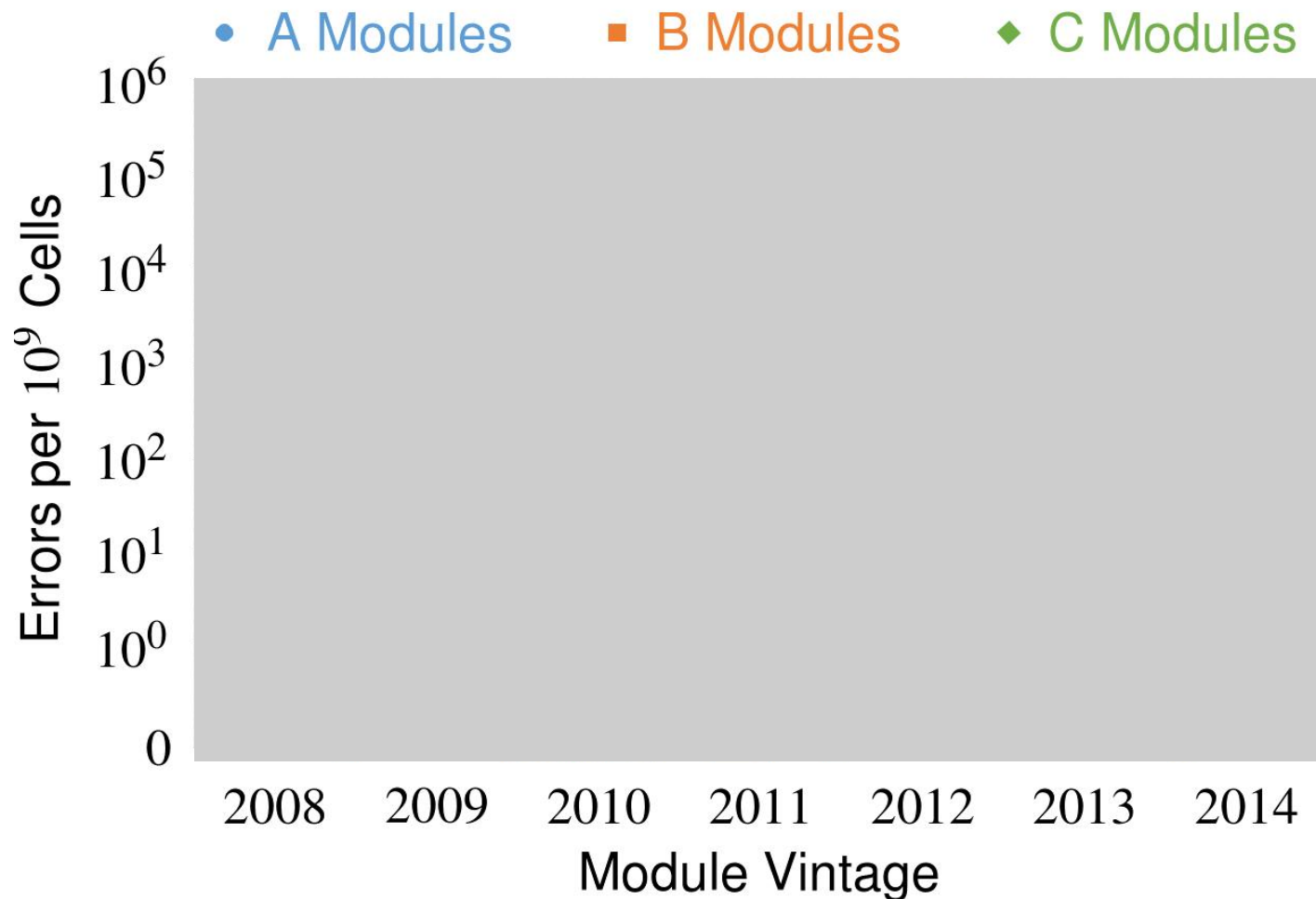


Up to
 1.0×10^7
errors

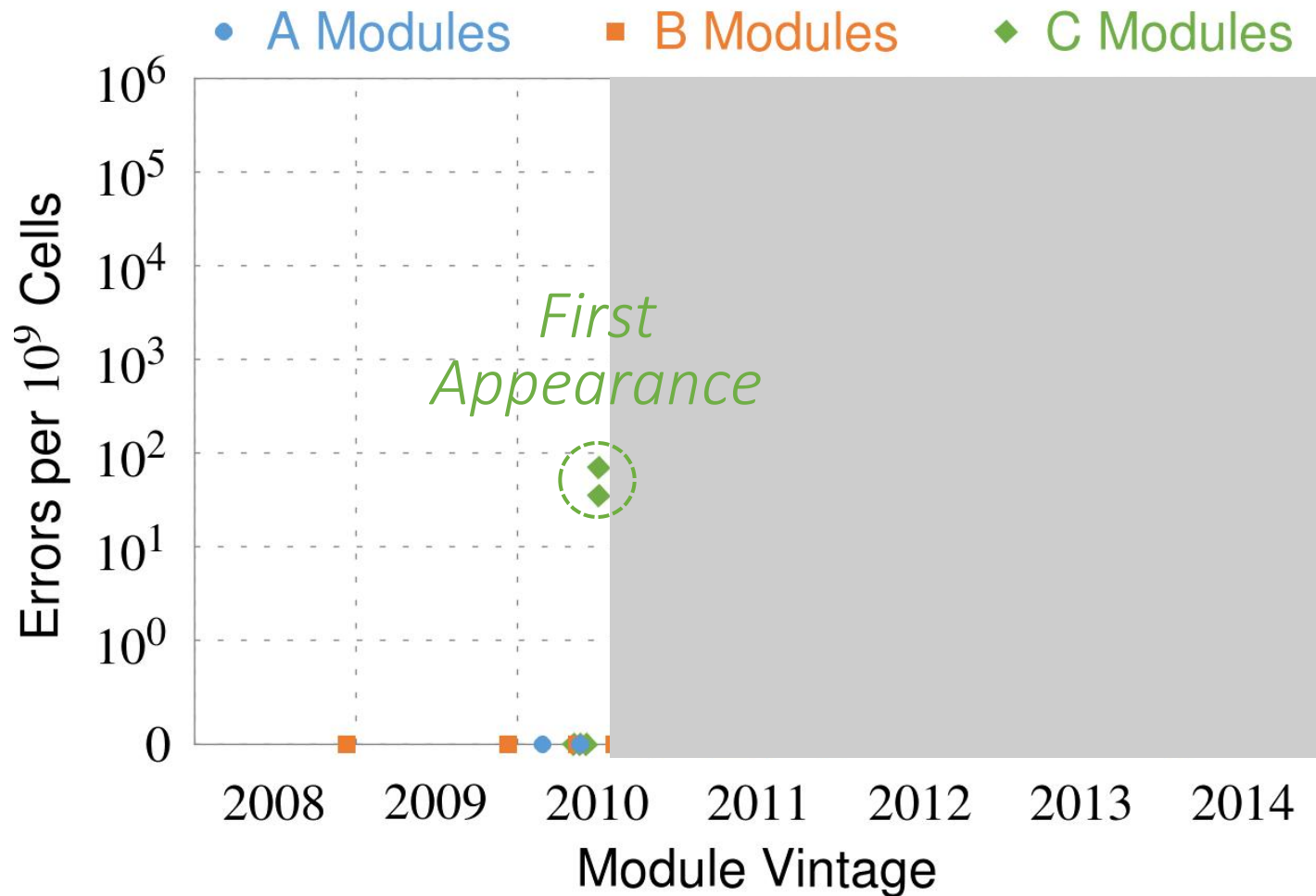
Up to
 2.7×10^6
errors

Up to
 3.3×10^5
errors

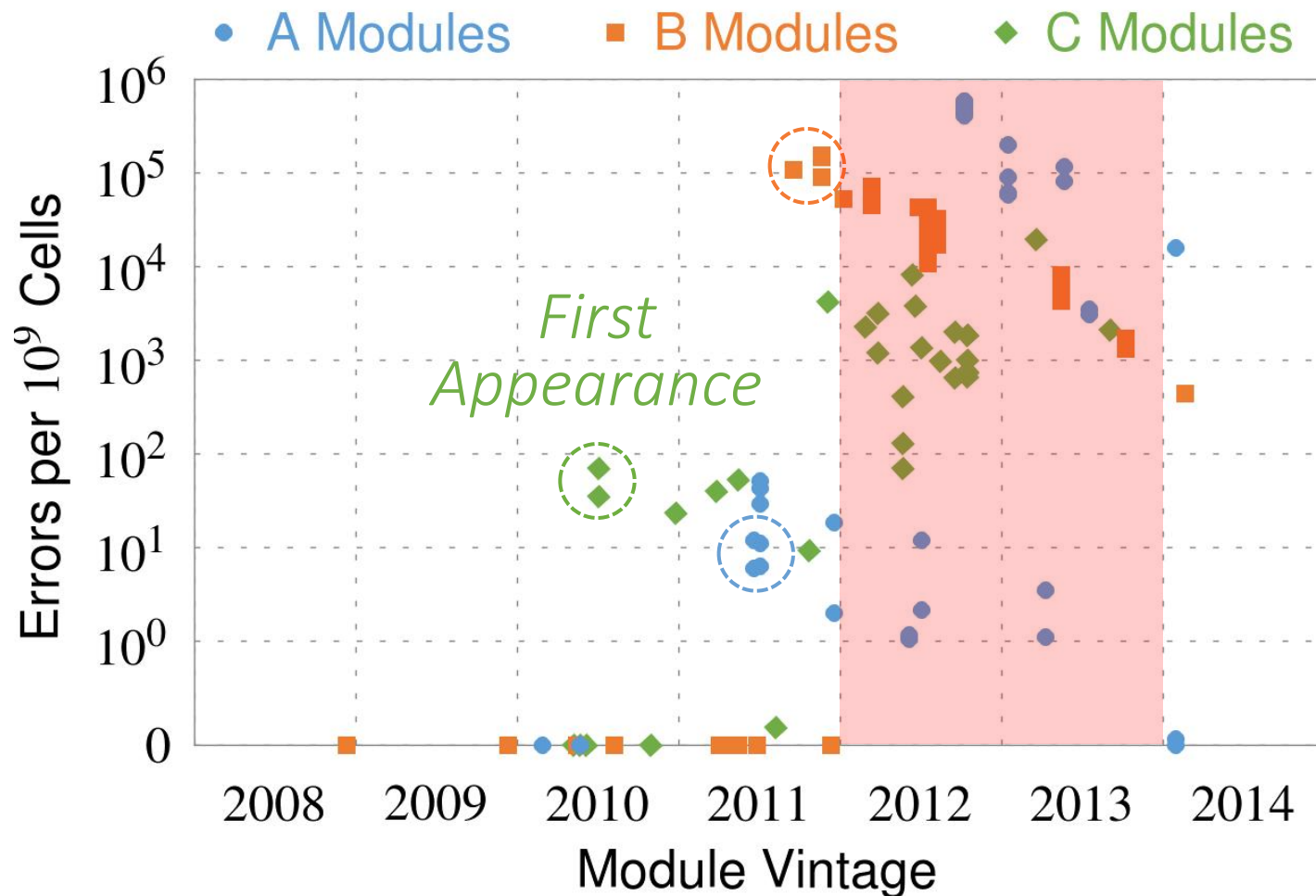
Recent DRAM Is More Vulnerable



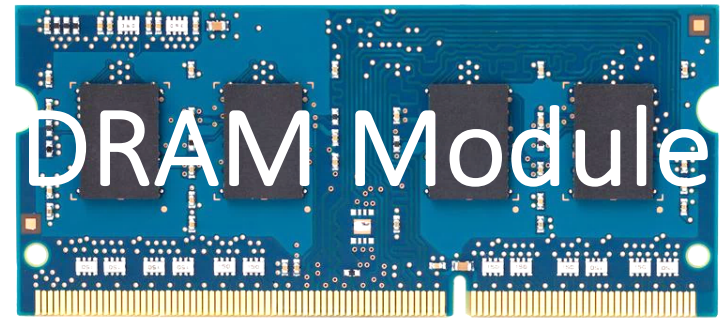
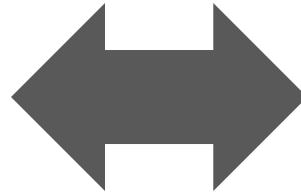
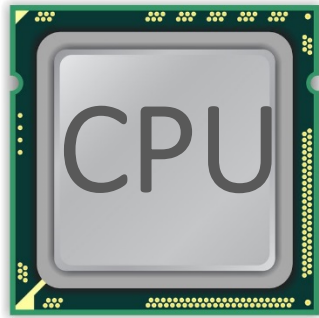
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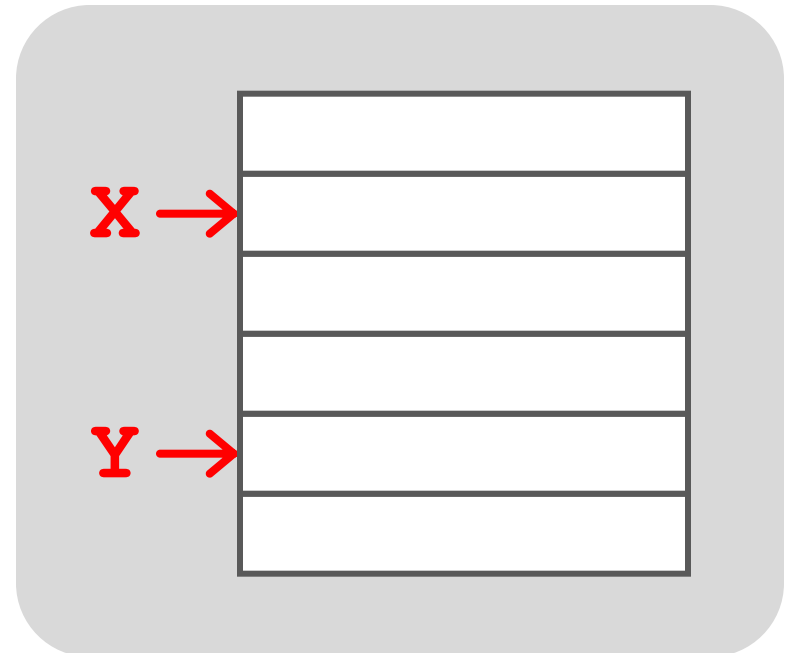
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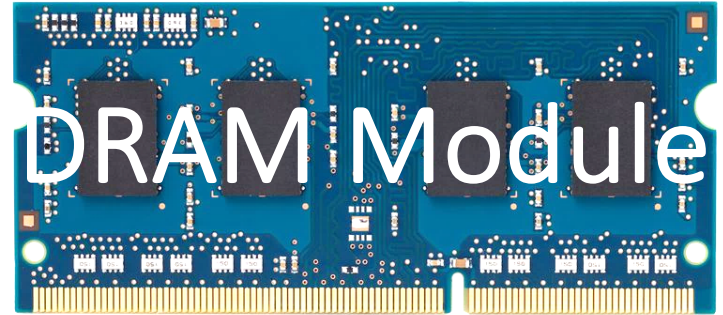
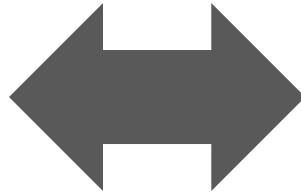
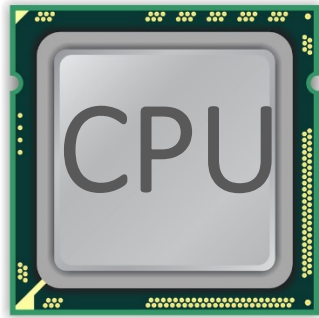
A Simple Program Can Induce Many Errors



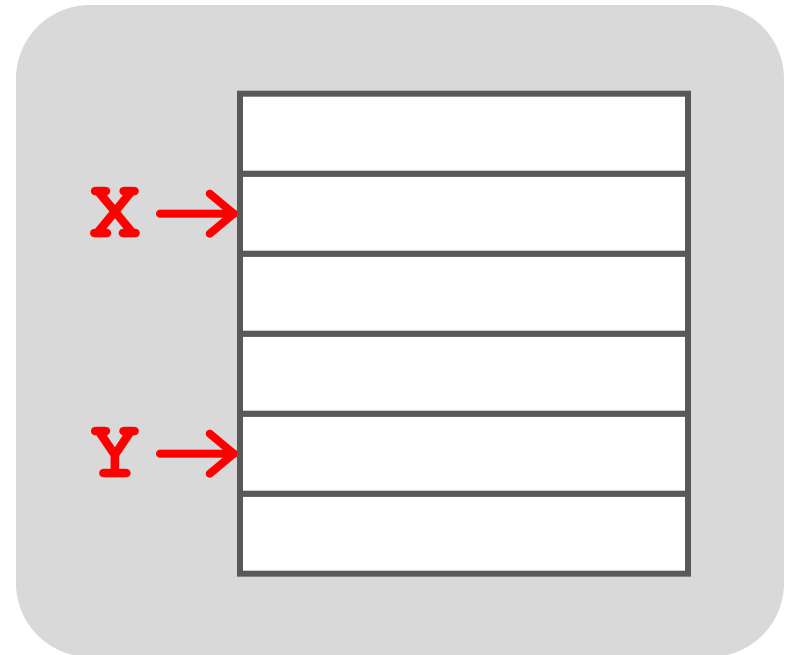
```
loop:  
  mov  (X), %eax  
  mov  (Y), %ebx  
  clflush (X)  
  clflush (Y)  
  mfence  
  jmp  loop
```



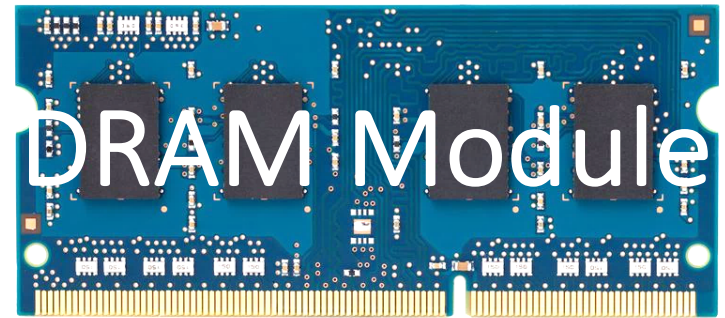
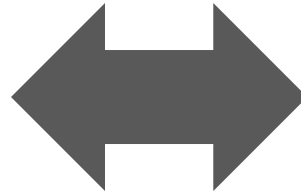
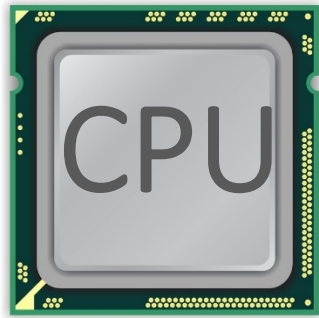
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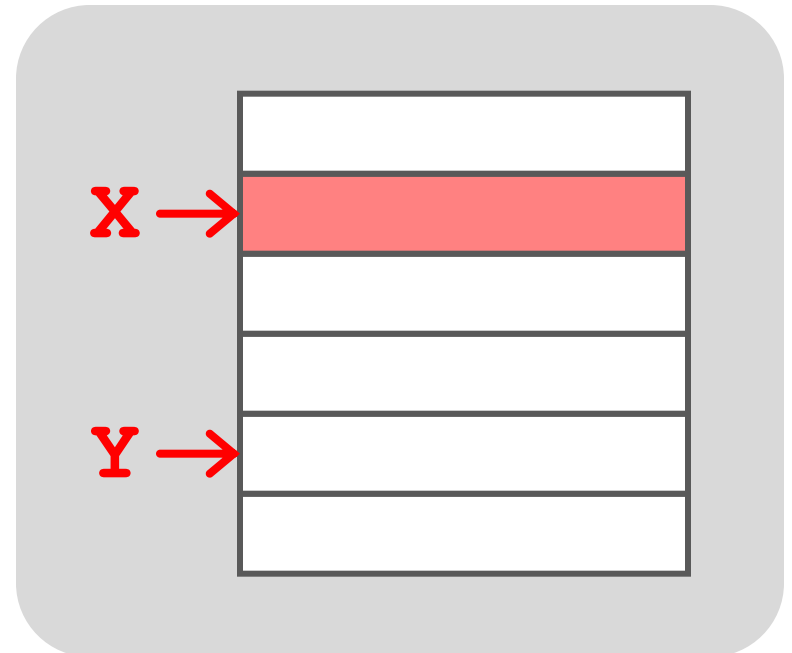
1. Avoid *cache hits*
 - Flush **X** from cache
2. Avoid *row hits* to **X**
 - Read **Y** in another row



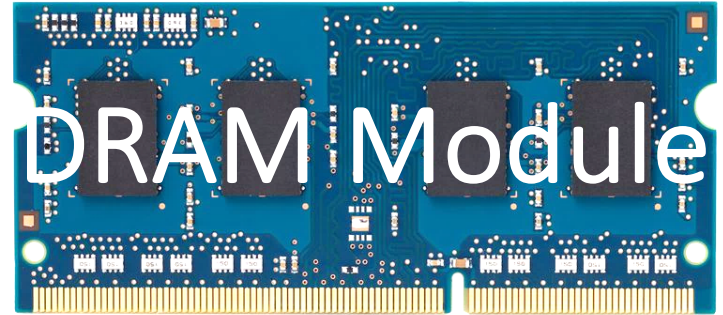
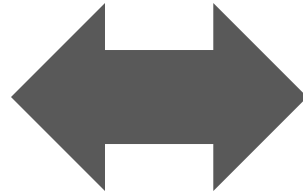
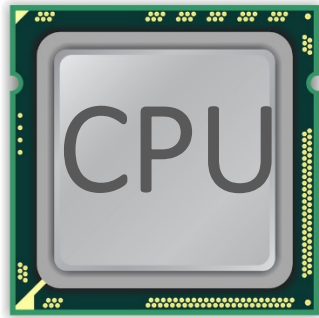
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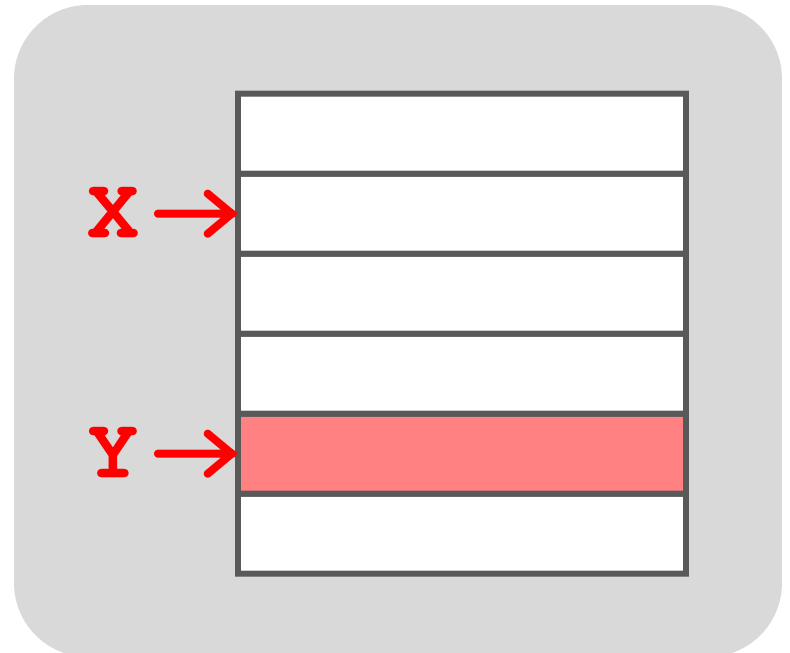
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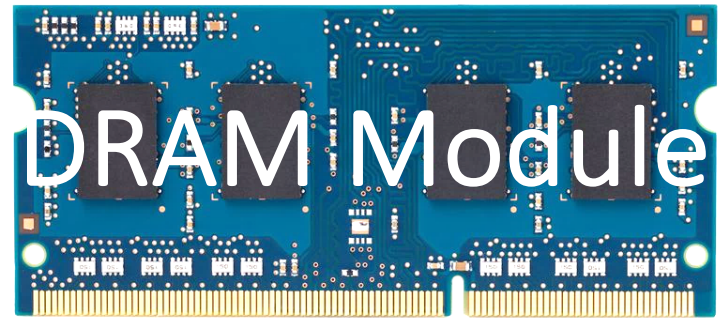
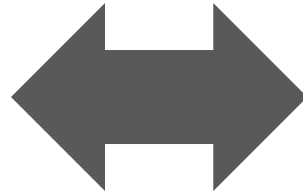
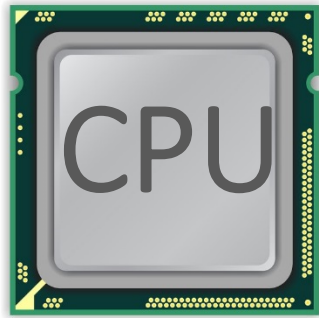
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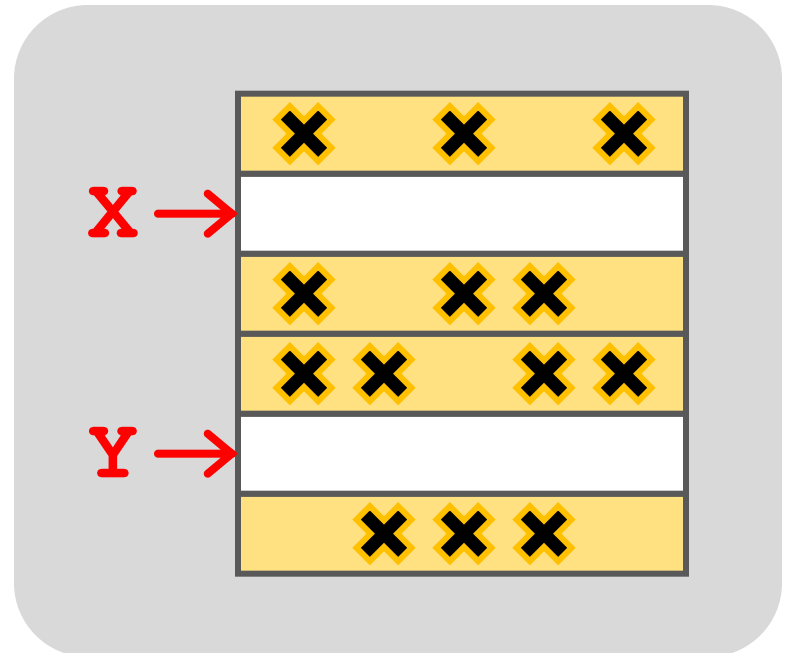
```
loop:  
  mov  (X), %eax  
  mov  (Y), %ebx  
  clflush (X)  
  clflush (Y)  
  mfence  
  jmp  loop
```



A Simple Program Can Induce Many Errors



```
loop:  
  mov  (X),  %eax  
  mov  (Y),  %ebx  
  clflush (X)  
  clflush (Y)  
  mfence  
  jmp  loop
```



Observed Errors in Real Systems

CPU Architecture	Errors	Access-Rate
Intel Haswell (2013)	22.9K	12.3M/sec
Intel Ivy Bridge (2012)	20.7K	11.7M/sec
Intel Sandy Bridge (2011)	16.1K	11.6M/sec
AMD Piledriver (2012)	59	6.1M/sec

A real reliability & security issue

One Can Take Over an Otherwise-Secure System

Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors

Abstract. Memory isolation is a key property of a reliable and secure computing system — an access to one memory address should not have unintended side effects on data stored in other addresses. However, as DRAM process technology

Project Zero

Flipping Bits in Memory Without Accessing Them:
An Experimental Study of DRAM Disturbance Errors
(Kim et al., ISCA 2014)

News and updates from the Project Zero team at Google

Exploiting the DRAM rowhammer bug to
gain kernel privileges (Seaborn, 2015)

Monday, March 9, 2015

Exploiting the DRAM rowhammer bug to gain kernel privileges

RowHammer Security Attack Example

- “Rowhammer” is a problem with some recent DRAM devices in which repeatedly accessing a row of memory can cause bit flips in adjacent rows (Kim et al., ISCA 2014).
 - Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors (Kim et al., ISCA 2014)
- We tested a selection of laptops and found that a subset of them exhibited the problem.
- We built two working privilege escalation exploits that use this effect.
 - Exploiting the DRAM rowhammer bug to gain kernel privileges (Seaborn+, 2015)
- One exploit uses rowhammer-induced bit flips to gain kernel privileges on x86-64 Linux when run as an unprivileged userland process.
- When run on a machine vulnerable to the rowhammer problem, the process was able to induce bit flips in page table entries (PTEs).
- It was able to use this to gain write access to its own page table, and hence gain read-write access to all of physical memory.

Security Implications



Security Implications



Rowhammer

It's like breaking into an apartment by repeatedly slamming a neighbor's door until the vibrations open the door you were after

More Security Implications (I)

“We can gain unrestricted access to systems of website visitors.”

www.iaik.tugraz.at ■

Not there yet, but ...



ROOT privileges for web apps!

29

Daniel Gruss (@lavados), Clémentine Maurice (@BloodyTangerine),
December 28, 2015 — 32c3, Hamburg, Germany



GATED
COMMUNITIES

Rowhammer.js: A Remote Software-Induced Fault Attack in JavaScript (DIMVA'16)

More Security Implications (II)

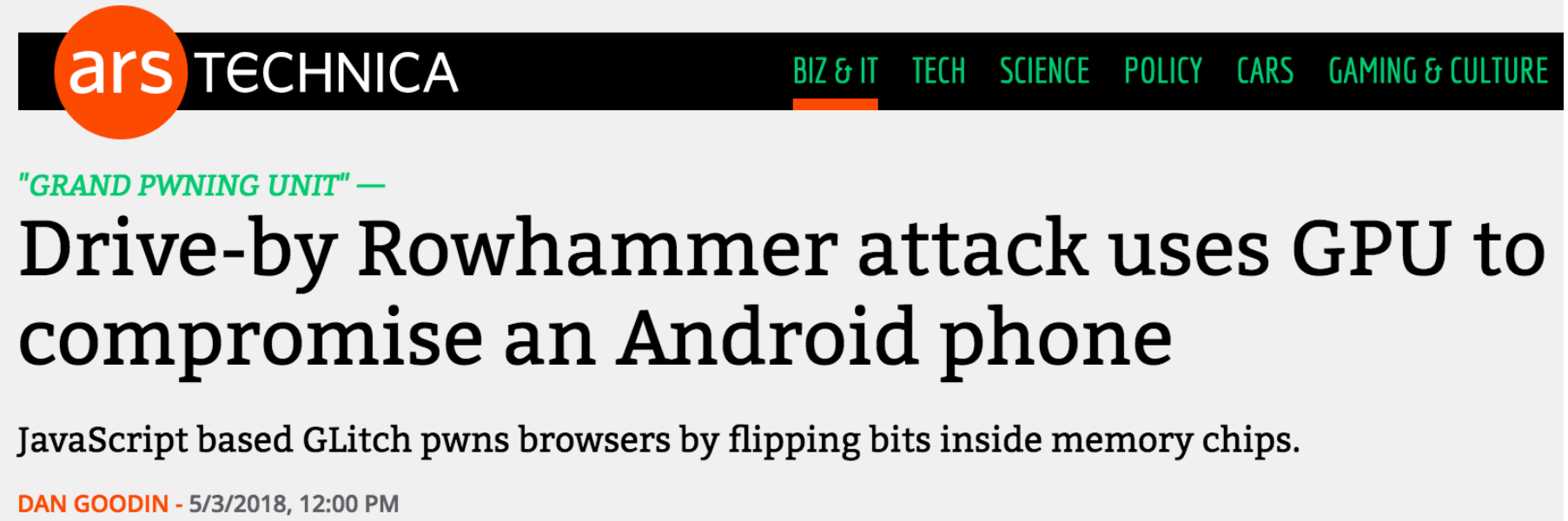
"Can gain control of a smart phone deterministically"



Drammer: Deterministic Rowhammer
Attacks on Mobile Platforms, CCS'16 ⁵⁹

More Security Implications (III)

- Using an integrated GPU in a mobile system to remotely escalate privilege via the WebGL interface



The screenshot shows the top of an Ars Technica article. The header includes the 'ars TECHNICA' logo and a navigation bar with categories: BIZ & IT, TECH, SCIENCE, POLICY, CARS, and GAMING & CULTURE. The article title is 'Drive-by Rowhammer attack uses GPU to compromise an Android phone', preceded by the sub-header '"GRAND PWINING UNIT" —'. A summary line reads: 'JavaScript based GLitch pwns browsers by flipping bits inside memory chips.' The author and date are listed as 'DAN GOODIN - 5/3/2018, 12:00 PM'.

Grand Pwning Unit: Accelerating Microarchitectural Attacks with the GPU

Pietro Frigo
Vrije Universiteit
Amsterdam
p.frigo@vu.nl

Cristiano Giuffrida
Vrije Universiteit
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Herbert Bos
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herbertb@cs.vu.nl

Kaveh Razavi
Vrije Universiteit
Amsterdam
kaveh@cs.vu.nl

More Security Implications (IV)

■ Rowhammer over RDMA (I)

ars TECHNICA

BIZ & IT TECH SCIENCE POLICY CARS GAMING & CULTURE

THROWHAMMER —

Packets over a LAN are all it takes to trigger serious Rowhammer bit flips

The bar for exploiting potentially serious DDR weakness keeps getting lower.

DAN GOODIN - 5/10/2018, 5:26 PM

Throwhammer: Rowhammer Attacks over the Network and Defenses

Andrei Tatar
VU Amsterdam

Radhesh Krishnan
VU Amsterdam

Elias Athanasopoulos
University of Cyprus

Cristiano Giuffrida
VU Amsterdam

Herbert Bos
VU Amsterdam

Kaveh Razavi
VU Amsterdam

More Security Implications (V)

■ Rowhammer over RDMA (II)



Nethammer—Exploiting DRAM Rowhammer Bug Through Network Requests



Nethammer: Inducing Rowhammer Faults through Network Requests

Moritz Lipp
Graz University of Technology

Daniel Gruss
Graz University of Technology

Misiker Tadesse Aga
University of Michigan

Clémentine Maurice
Univ Rennes, CNRS, IRISA

Michael Schwarz
Graz University of Technology

Lukas Raab
Graz University of Technology

Lukas Lamster
Graz University of Technology

More Security Implications?



Apple's Patch for RowHammer

- <https://support.apple.com/en-gb/HT204934>

Available for: OS X Mountain Lion v10.8.5, OS X Mavericks v10.9.5

Impact: A malicious application may induce memory corruption to escalate privileges

Description: A disturbance error, also known as Rowhammer, exists with some DDR3 RAM that could have led to memory corruption. This issue was mitigated by increasing memory refresh rates.

CVE-ID

CVE-2015-3693 : Mark Seaborn and Thomas Dullien of Google, working from original research by Yoongu Kim et al (2014)

HP, Lenovo, and other vendors released similar patches

Our Solution to RowHammer

- PARA: *Probabilistic Adjacent Row Activation*
- Key Idea
 - After closing a row, we activate (i.e., refresh) one of its neighbors with a low probability: $p = 0.005$
- Reliability Guarantee
 - When $p=0.005$, errors in one year: 9.4×10^{-14}
 - By adjusting the value of p , we can vary the strength of protection against errors

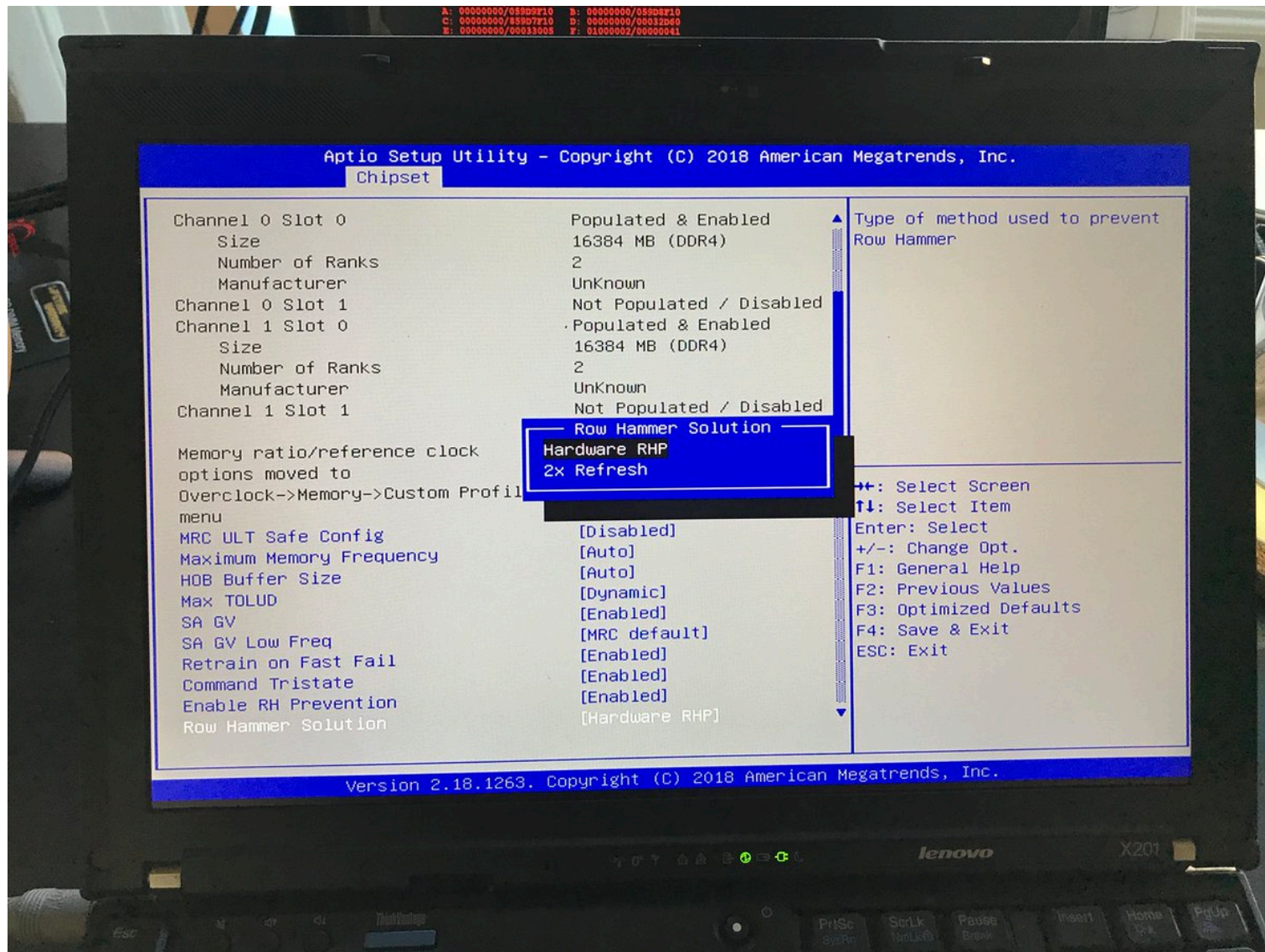
Advantages of PARA

- *PARA refreshes rows infrequently*
 - Low power
 - Low performance-overhead
 - Average slowdown: **0.20%** (for 29 benchmarks)
 - Maximum slowdown: **0.75%**
- *PARA is stateless*
 - Low cost
 - Low complexity
- *PARA is an effective and low-overhead solution to prevent disturbance errors*

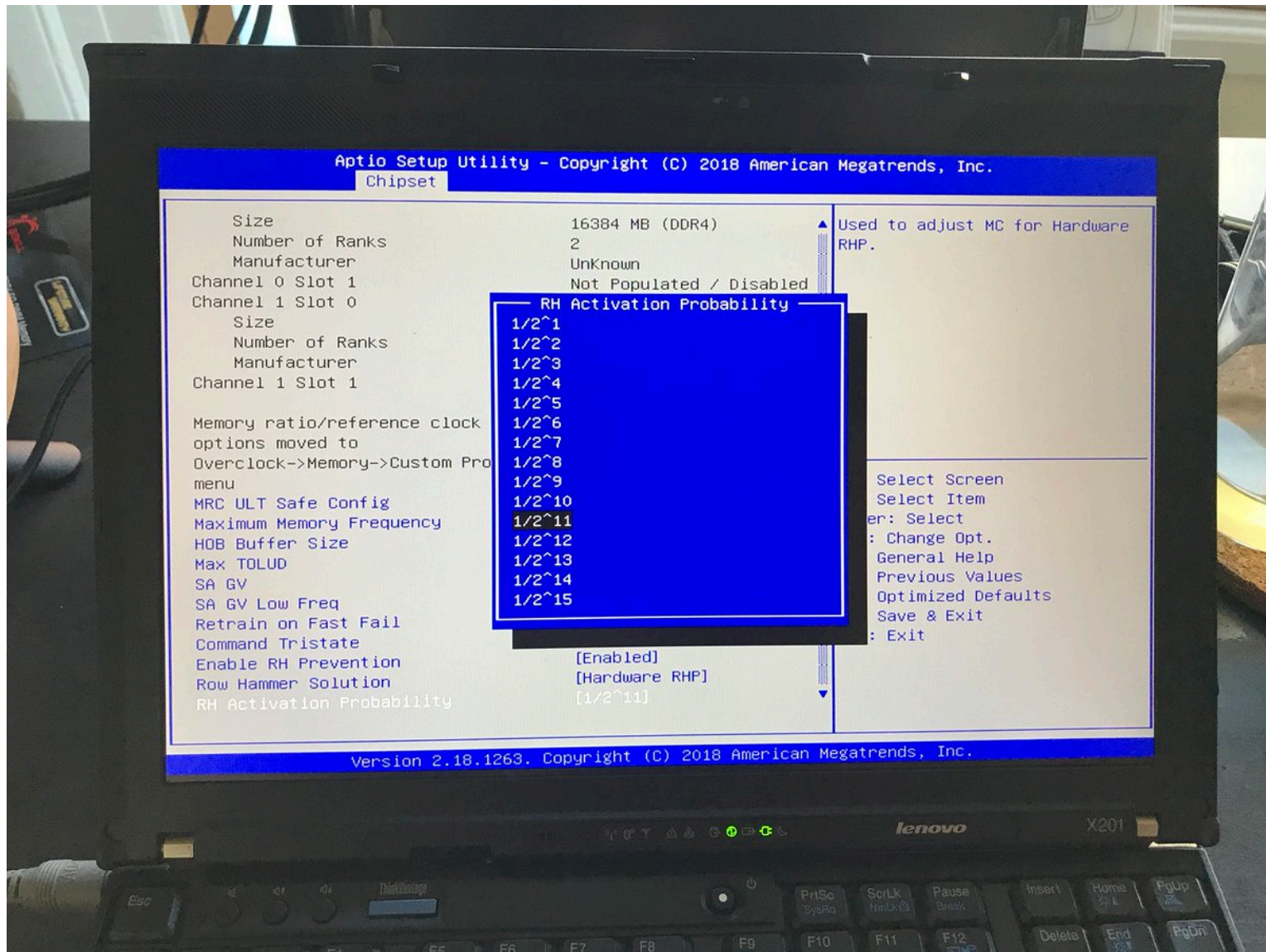
Requirements for PARA

- If implemented in **DRAM chip** (done today)
 - Enough slack in timing and refresh parameters
 - Plenty of slack today:
 - Lee et al., “**Adaptive-Latency DRAM: Optimizing DRAM Timing for the Common Case**,” HPCA 2015.
 - Chang et al., “**Understanding Latency Variation in Modern DRAM Chips**,” SIGMETRICS 2016.
 - Lee et al., “**Design-Induced Latency Variation in Modern DRAM Chips**,” SIGMETRICS 2017.
 - Chang et al., “**Understanding Reduced-Voltage Operation in Modern DRAM Devices**,” SIGMETRICS 2017.
 - Ghose et al., “**What Your DRAM Power Models Are Not Telling You: Lessons from a Detailed Experimental Study**,” SIGMETRICS 2018.
- If implemented in **memory controller**
 - Better coordination between memory controller and DRAM
 - Memory controller should know which rows are physically adjacent

Probabilistic Activation in Real Life (I)



Probabilistic Activation in Real Life (II)



More on RowHammer Analysis

- Yoongu Kim, Ross Daly, Jeremie Kim, Chris Fallin, Ji Hye Lee, Donghyuk Lee, Chris Wilkerson, Konrad Lai, and Onur Mutlu,
"Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors"
Proceedings of the 41st International Symposium on Computer Architecture (ISCA), Minneapolis, MN, June 2014.
[\[Slides \(pptx\) \(pdf\)\]](#) [\[Lightning Session Slides \(pptx\) \(pdf\)\]](#) [\[Source Code and Data\]](#)

Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors

Yoongu Kim¹ Ross Daly* Jeremie Kim¹ Chris Fallin* Ji Hye Lee¹
Donghyuk Lee¹ Chris Wilkerson² Konrad Lai Onur Mutlu¹

¹Carnegie Mellon University ²Intel Labs

Future of Memory Reliability

- Onur Mutlu,
"The RowHammer Problem and Other Issues We May Face as Memory Becomes Denser"
Invited Paper in Proceedings of the Design, Automation, and Test in Europe Conference (DATE), Lausanne, Switzerland, March 2017.
[[Slides \(pptx\)](#) ([pdf](#))]

The RowHammer Problem and Other Issues We May Face as Memory Becomes Denser

Onur Mutlu
ETH Zürich
onur.mutlu@inf.ethz.ch
<https://people.inf.ethz.ch/omutlu>

A RowHammer Retrospective

- Onur Mutlu and Jeremie Kim,
"RowHammer: A Retrospective"
IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems (TCAD) Special Issue on Top Picks in Hardware and Embedded Security, 2019.
[[Preliminary arXiv version](#)]

RowHammer: A Retrospective

Onur Mutlu^{§‡} Jeremie S. Kim^{‡§}
§ETH Zürich ‡Carnegie Mellon University

Main Memory Needs Intelligent Controllers

Industry Is Writing Papers About It, Too

DRAM Process Scaling Challenges

❖ Refresh

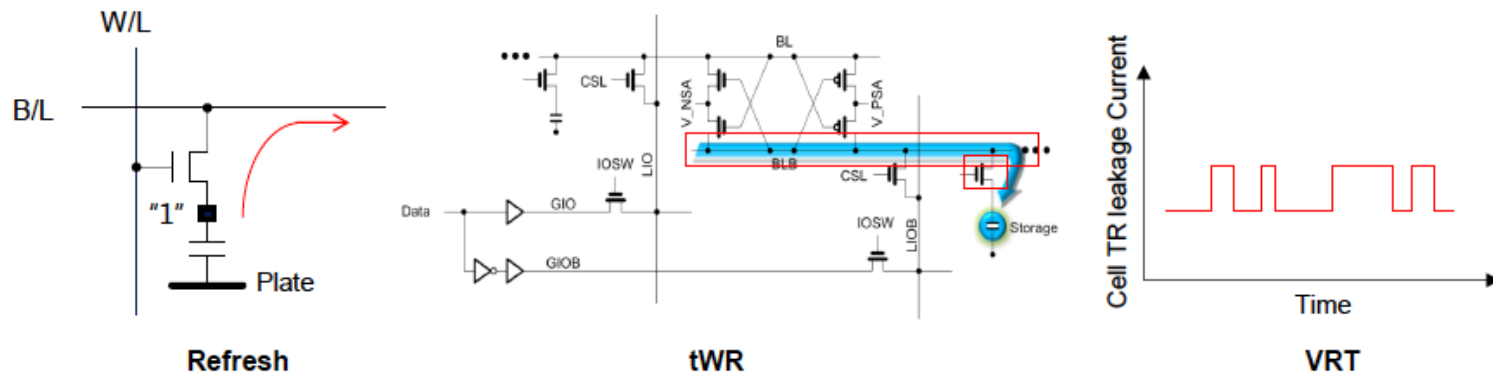
- Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance
- Leakage current of cell access transistors increasing

❖ tWR

- Contact resistance between the cell capacitor and access transistor increasing
- On-current of the cell access transistor decreasing
- Bit-line resistance increasing

❖ VRT

- Occurring more frequently with cell capacitance decreasing



Call for Intelligent Memory Controllers

DRAM Process Scaling Challenges

❖ Refresh

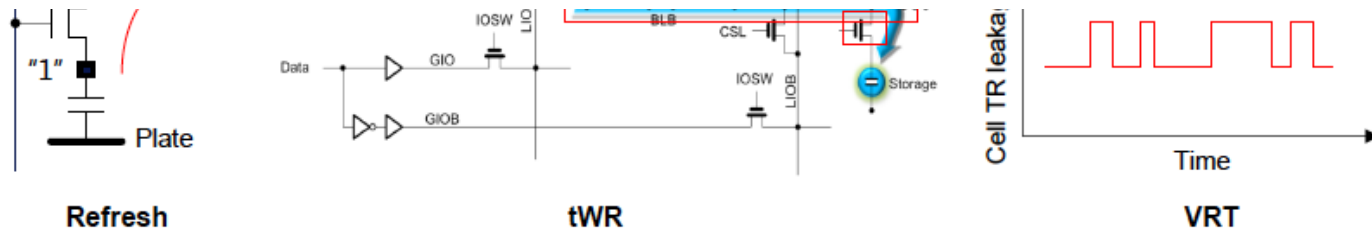
- Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance

THE MEMORY FORUM 2014

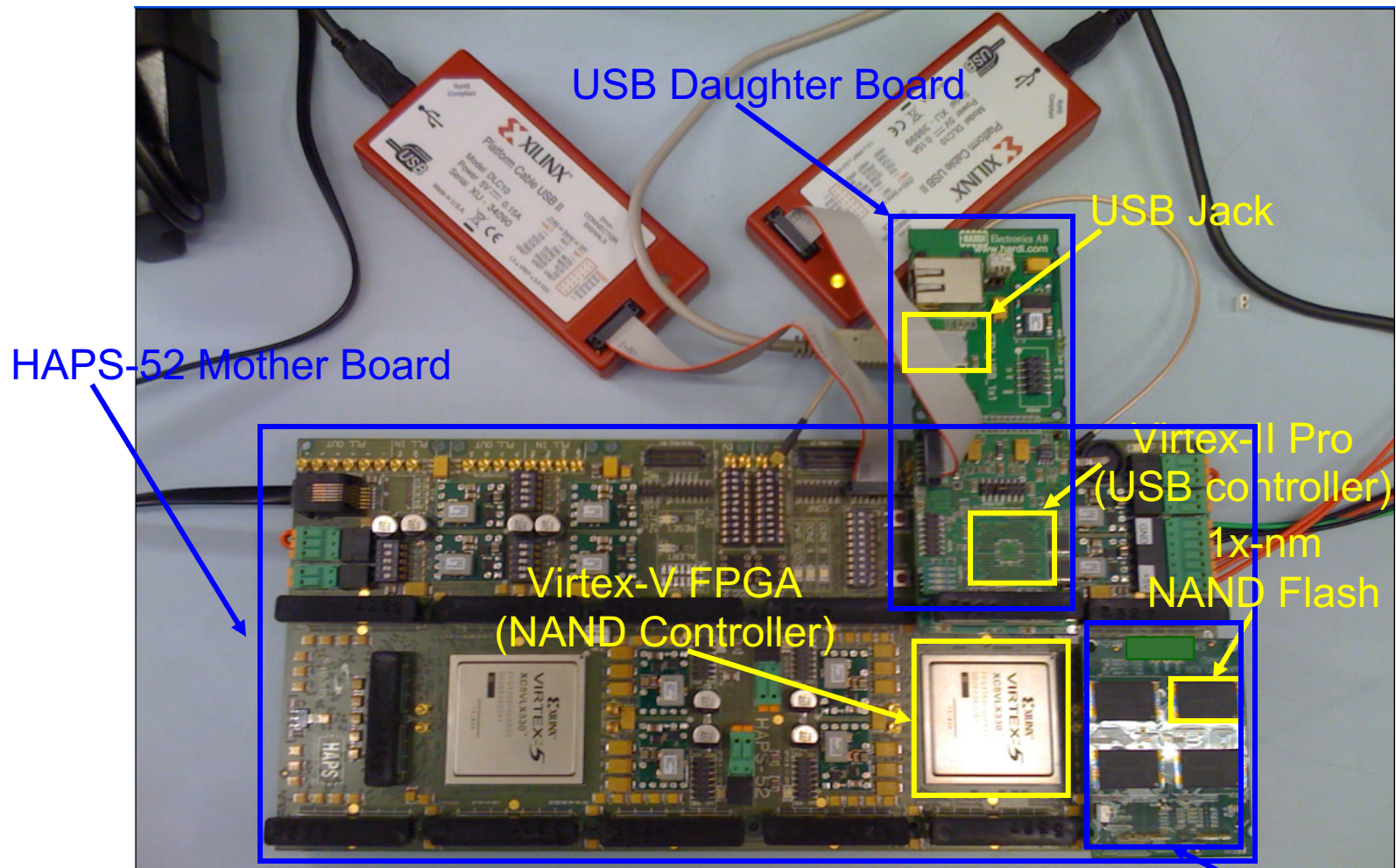
Co-Architecting Controllers and DRAM to Enhance DRAM Process Scaling

Uksong Kang, Hak-soo Yu, Churoo Park, *Hongzhong Zheng,
**John Halbert, **Kuljit Bains, SeongJin Jang, and Joo Sun Choi

*Samsung Electronics, Hwasung, Korea / *Samsung Electronics, San Jose / **Intel*



Aside: Intelligent Controller for NAND Flash



[DATE 2012, ICCD 2012, DATE 2013, ITJ 2013, ICCD 2013, SIGMETRICS 2014, HPCA 2015, DSN 2015, MSST 2015, JSAC 2016, HPCA 2017, DFRWS 2017, PIEEE 2017, HPCA 2018, SIGMETRICS 2018]

Cai+, "Error Characterization, Mitigation, and Recovery in Flash Memory Based Solid State Drives," Proc. IEEE 2017.



Proceedings of the IEEE, Sept. 2017



Error Characterization, Mitigation, and Recovery in Flash-Memory-Based Solid-State Drives

This paper reviews the most recent advances in solid-state drive (SSD) error characterization, mitigation, and data recovery techniques to improve both SSD's reliability and lifetime.

By YU CAI, SAUGATA GHOSE, ERICH F. HARATSCH, YIXIN LUO, AND ONUR MUTLU

<https://arxiv.org/pdf/1706.08642>

Agenda

- Major Trends Affecting Main Memory
- The Need for Intelligent Memory Controllers
 - Bottom Up: Push from Circuits and Devices
 - Top Down: Pull from Systems and Applications
- Processing in Memory: Two Directions
 - Minimally Changing Memory Chips
 - Exploiting 3D-Stacked Memory
- How to Enable Adoption of Processing in Memory
- Conclusion

Three Key Systems Trends

1. Data access is a major bottleneck

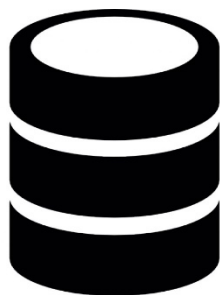
- ▣ Applications are increasingly data hungry

2. Energy consumption is a key limiter

3. Data movement energy dominates compute

- ▣ Especially true for off-chip to on-chip movement

The Need for More Memory Performance



In-memory Databases

[Mao+, EuroSys'12;
Clapp+ (Intel), IISWC'15]



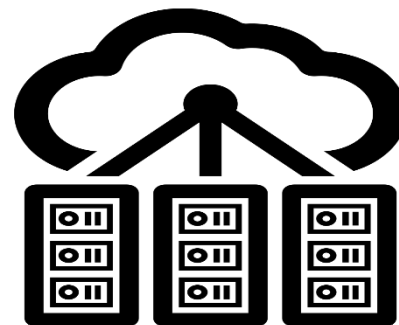
In-Memory Data Analytics

[Clapp+ (Intel), IISWC'15;
Awan+, BDCloud'15]



Graph/Tree Processing

[Xu+, IISWC'12; Umuroglu+, FPL'15]



Datacenter Workloads

[Kanev+ (Google), ISCA'15]

The Need for More Memory Performance



Chrome

Google's web browser



TensorFlow Mobile

Google's machine learning
framework

VP9



Video Playback

Google's **video codec**

VP9



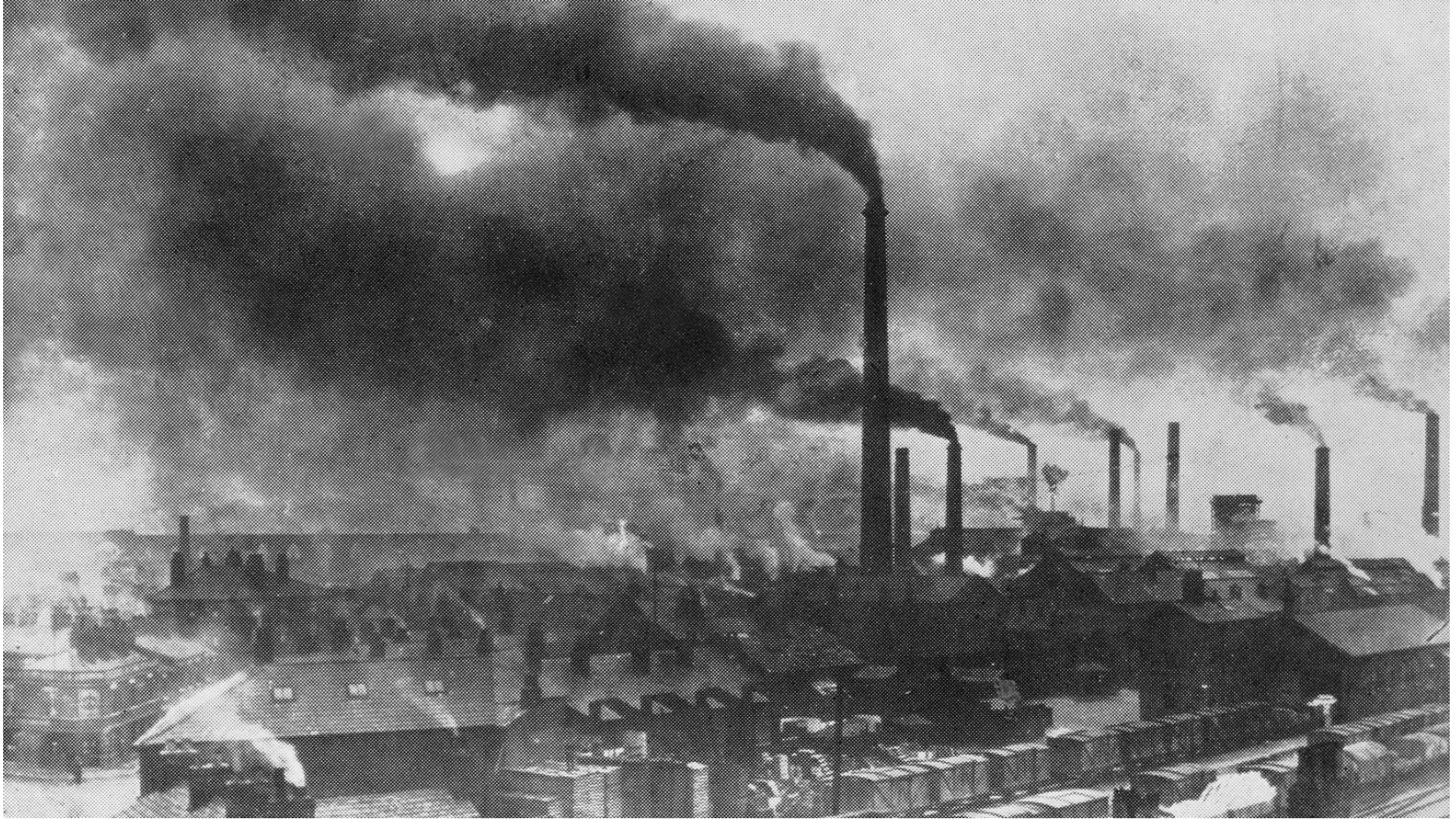
Video Capture

Google's **video codec**

Do We Want This?



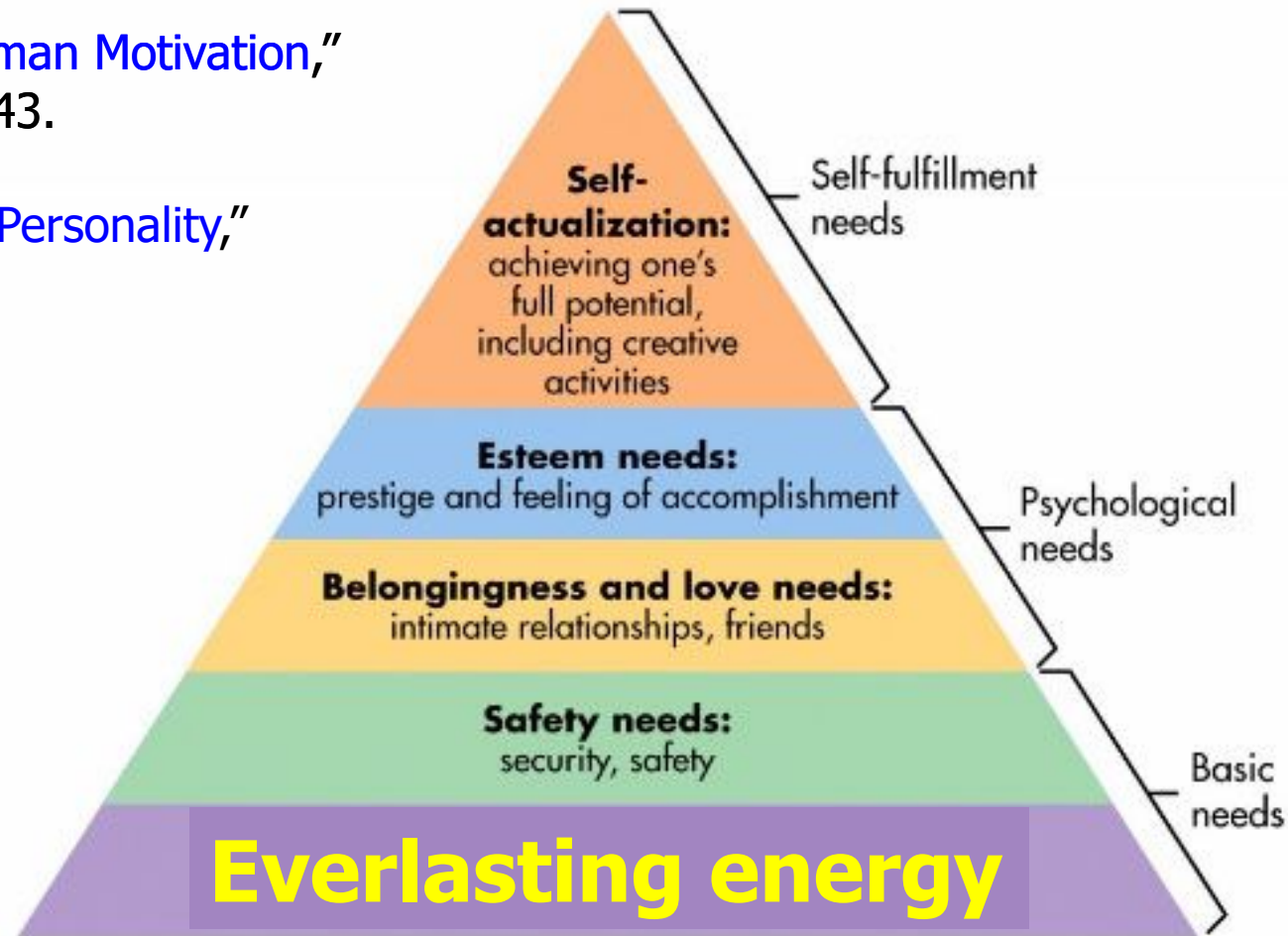
Or This?



Maslow's (Human) Hierarchy of Needs, Revisited

Maslow, "A Theory of Human Motivation,"
Psychological Review, 1943.

Maslow, "Motivation and Personality,"
Book, 1954-1970.



High Performance,
Energy Efficient,
Sustainable

The Problem

Data access is the major performance and energy bottleneck

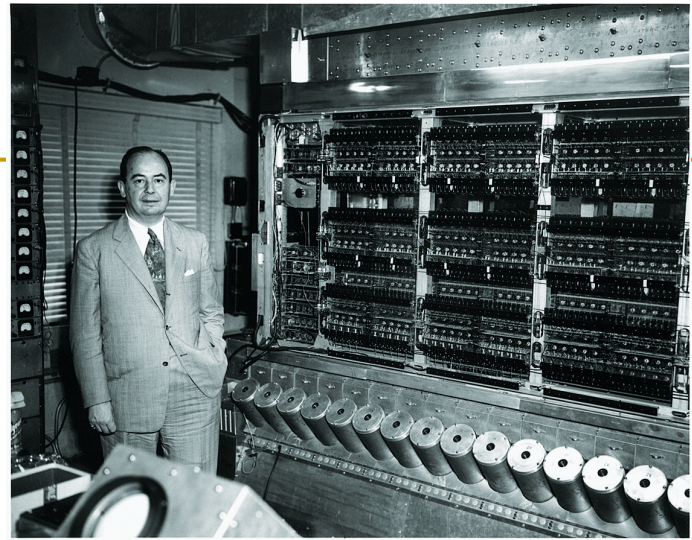
Our current
design principles
cause great energy waste
(and great performance loss)

The Problem

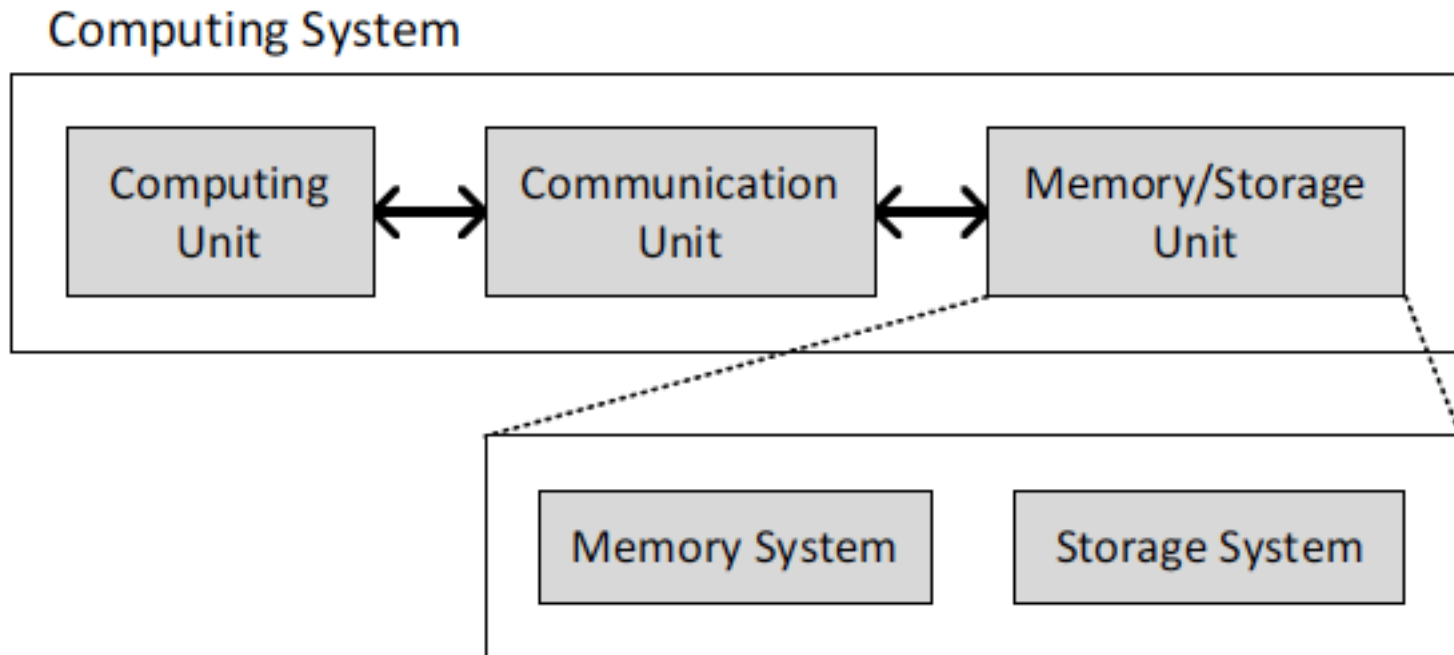
Processing of data
is performed
far away from the data

A Computing System

- Three key components
- Computation
- Communication
- Storage/memory

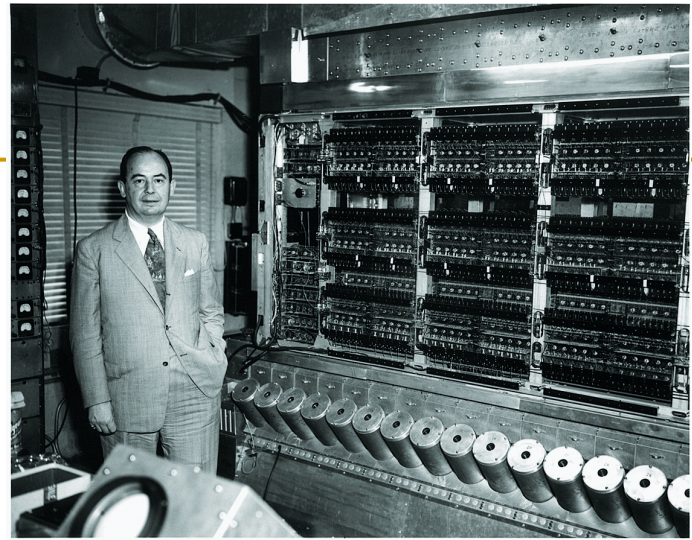


Burks, Goldstein, von Neumann, "Preliminary discussion of the logical design of an electronic computing instrument," 1946.



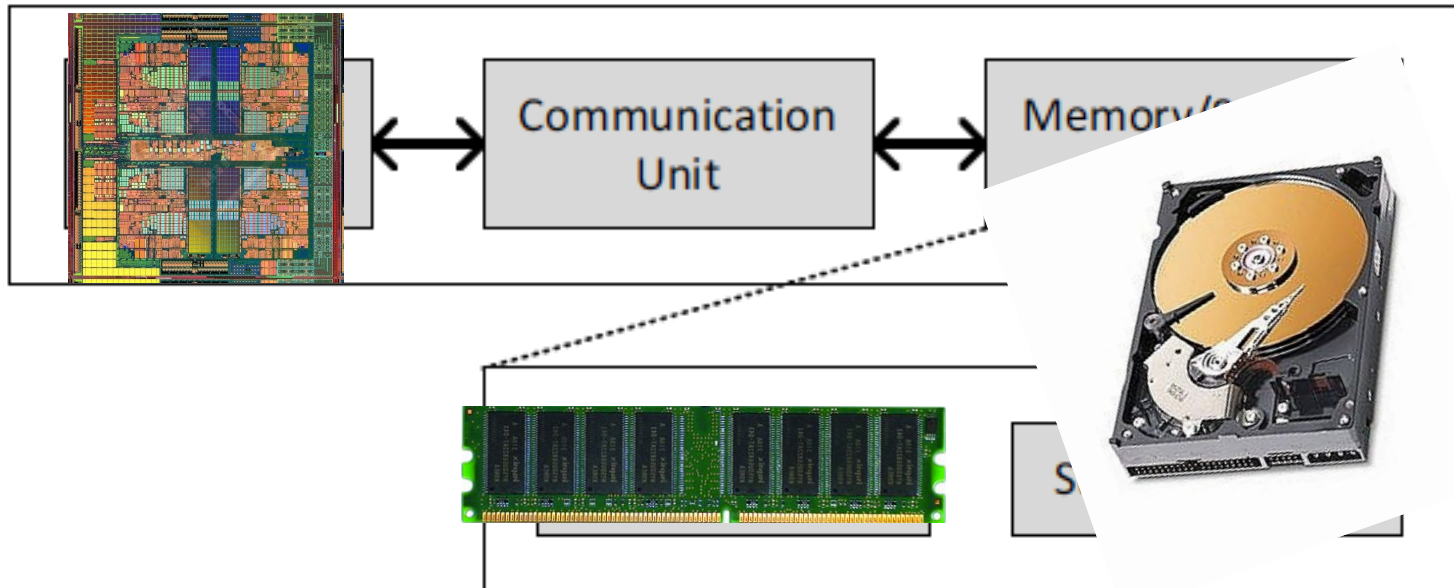
A Computing System

- Three key components
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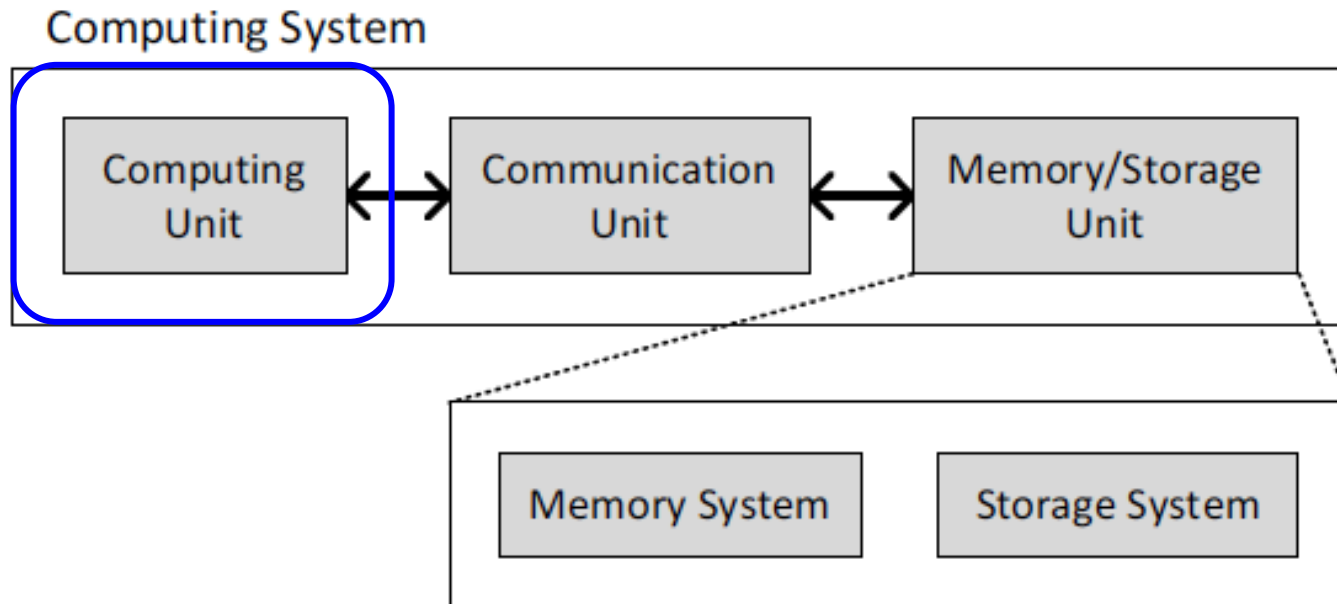
Burks, Goldstein, von Neumann, "Preliminary discussion of the logical design of an electronic computing instrument," 1946.

Computing System



Today's Computing Systems

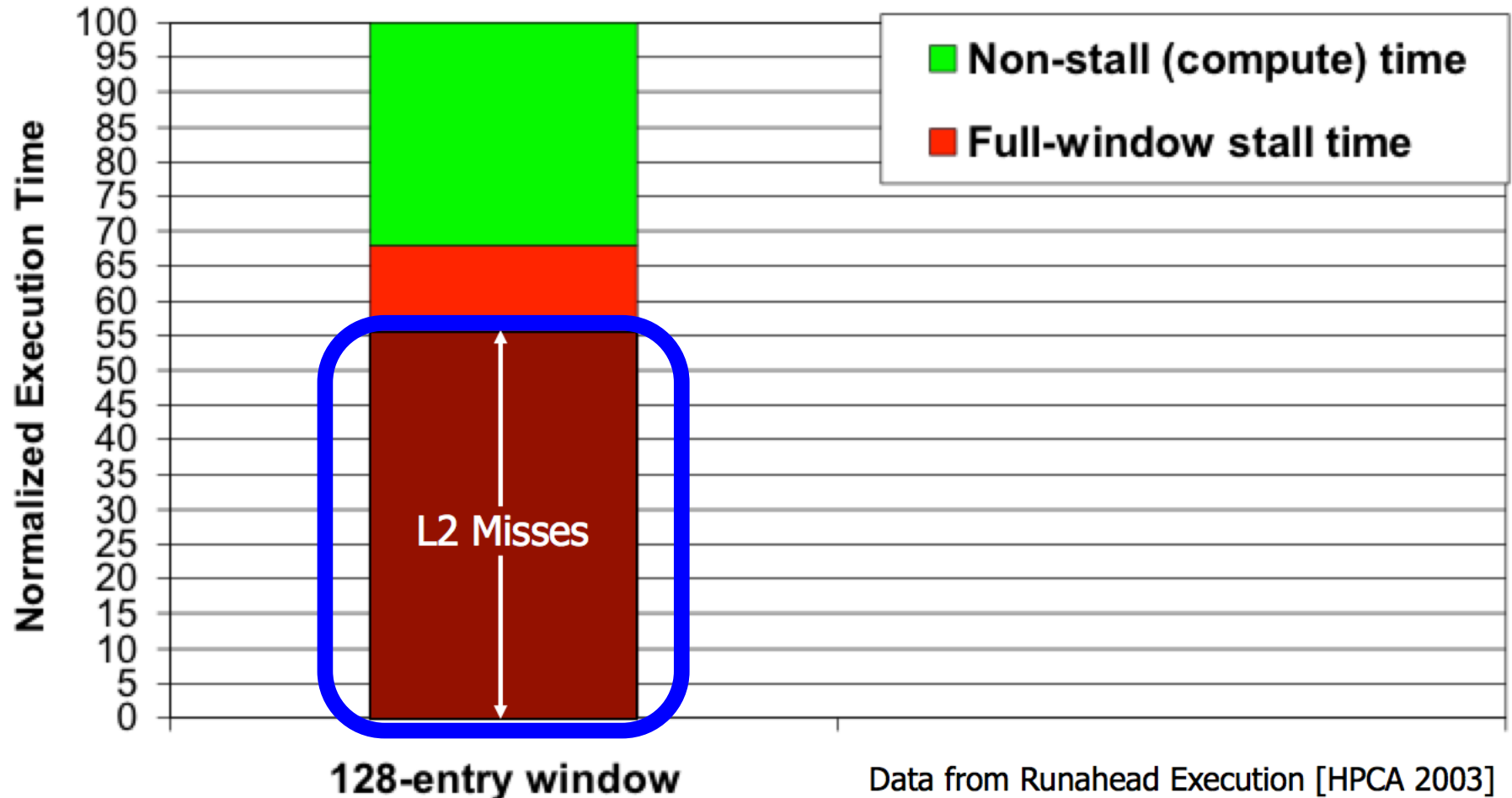
- Are overwhelmingly processor centric
- All data processed in the processor → at great system cost
- Processor is heavily optimized and is considered the master
- Data storage units are dumb and are largely unoptimized (except for some that are on the processor die)



Yet ...

I expect that over the coming decade memory subsystem design will be the *only* important design issue for microprocessors.

- **“It’s the Memory, Stupid!”** (Richard Sites, MPR, 1996)



The Performance Perspective

- Onur Mutlu, Jared Stark, Chris Wilkerson, and Yale N. Patt,
"Runahead Execution: An Alternative to Very Large Instruction Windows for Out-of-order Processors"
Proceedings of the 9th International Symposium on High-Performance Computer Architecture (HPCA), pages 129-140, Anaheim, CA, February 2003. [Slides \(pdf\)](#)

Runahead Execution: An Alternative to Very Large Instruction Windows for Out-of-order Processors

Onur Mutlu § Jared Stark † Chris Wilkerson ‡ Yale N. Patt §

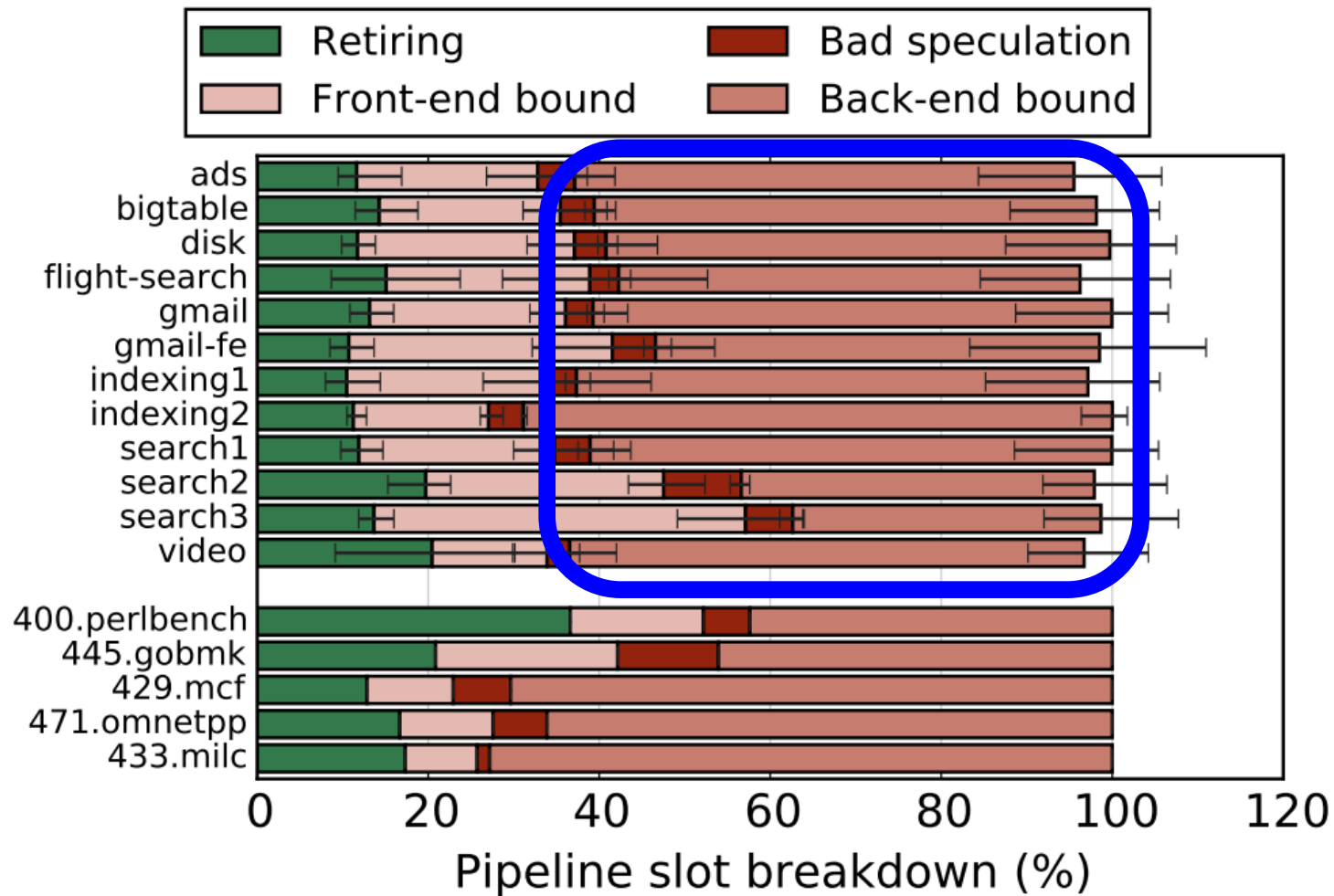
§ECE Department
The University of Texas at Austin
{onur,patt}@ece.utexas.edu

†Microprocessor Research
Intel Labs
jared.w.stark@intel.com

‡Desktop Platforms Group
Intel Corporation
chris.wilkerson@intel.com

The Performance Perspective (Today)

- All of Google's Data Center Workloads (2015):



The Performance Perspective (Today)

- All of Google's Data Center Workloads (2015):

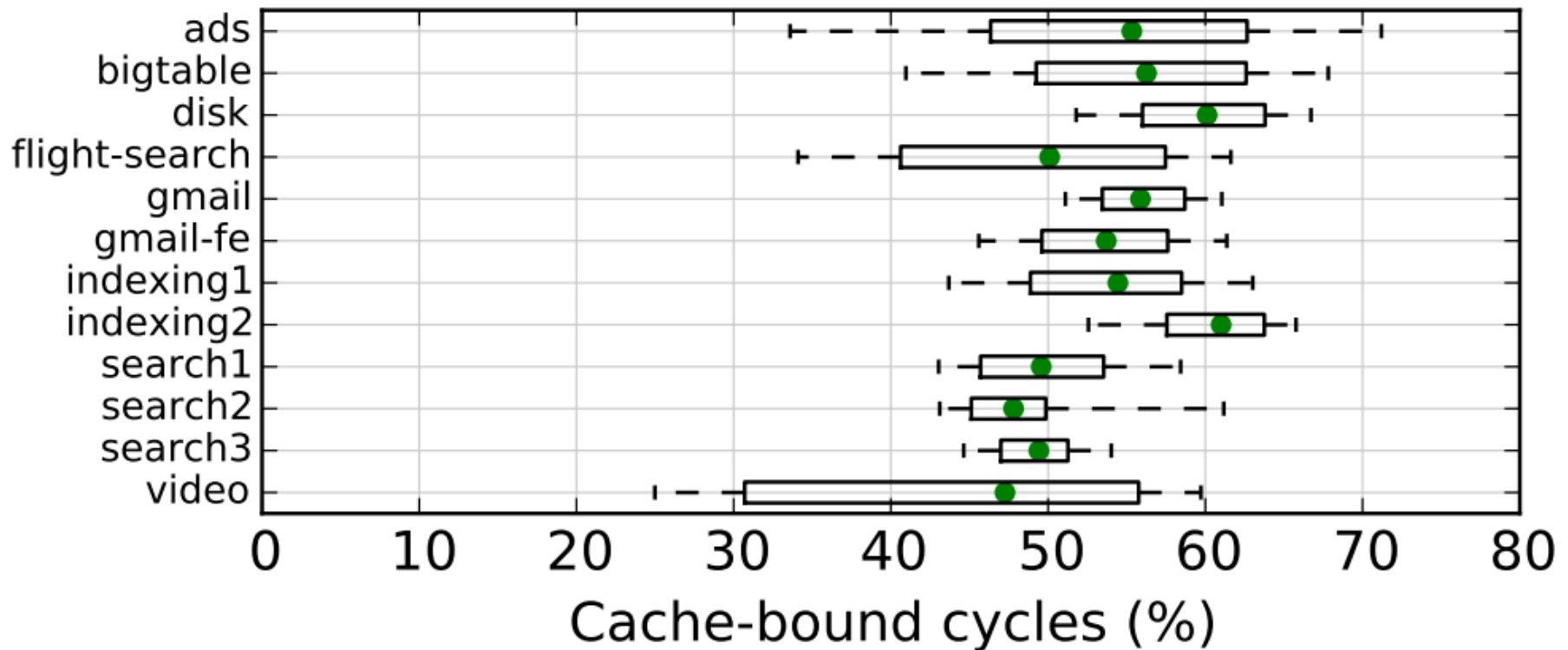


Figure 11: Half of cycles are spent stalled on caches.

Perils of Processor-Centric Design

■ Grossly-imbalanced systems

- ❑ Processing done only in **one place**
- ❑ Everything else just stores and moves data: **data moves a lot**
 - Energy inefficient
 - Low performance
 - Complex

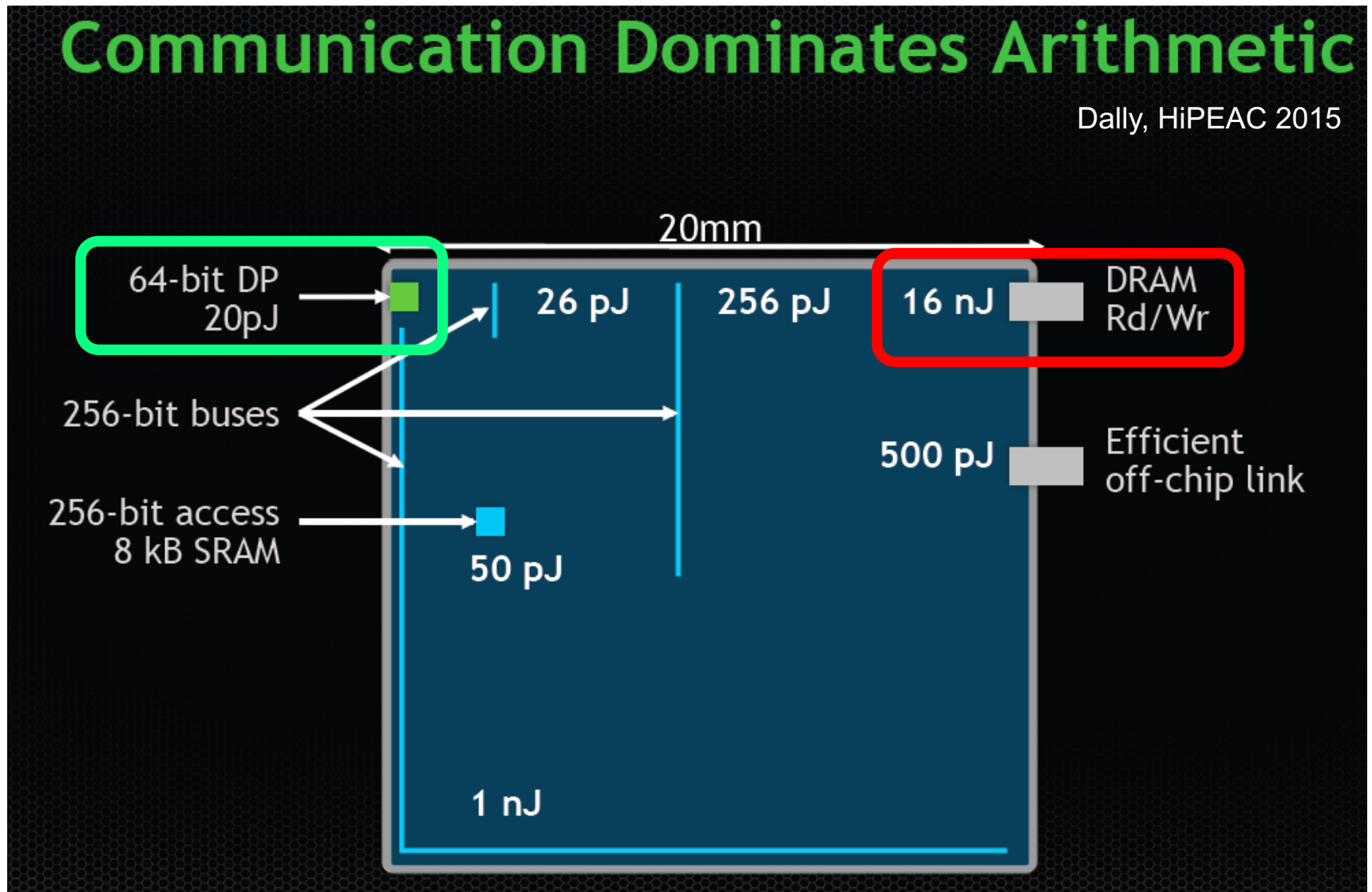
■ Overly complex and bloated processor (and accelerators)

- ❑ To tolerate data access from memory
- ❑ Complex hierarchies and mechanisms
 - Energy inefficient
 - Low performance
 - Complex

The Energy Perspective

Communication Dominates Arithmetic

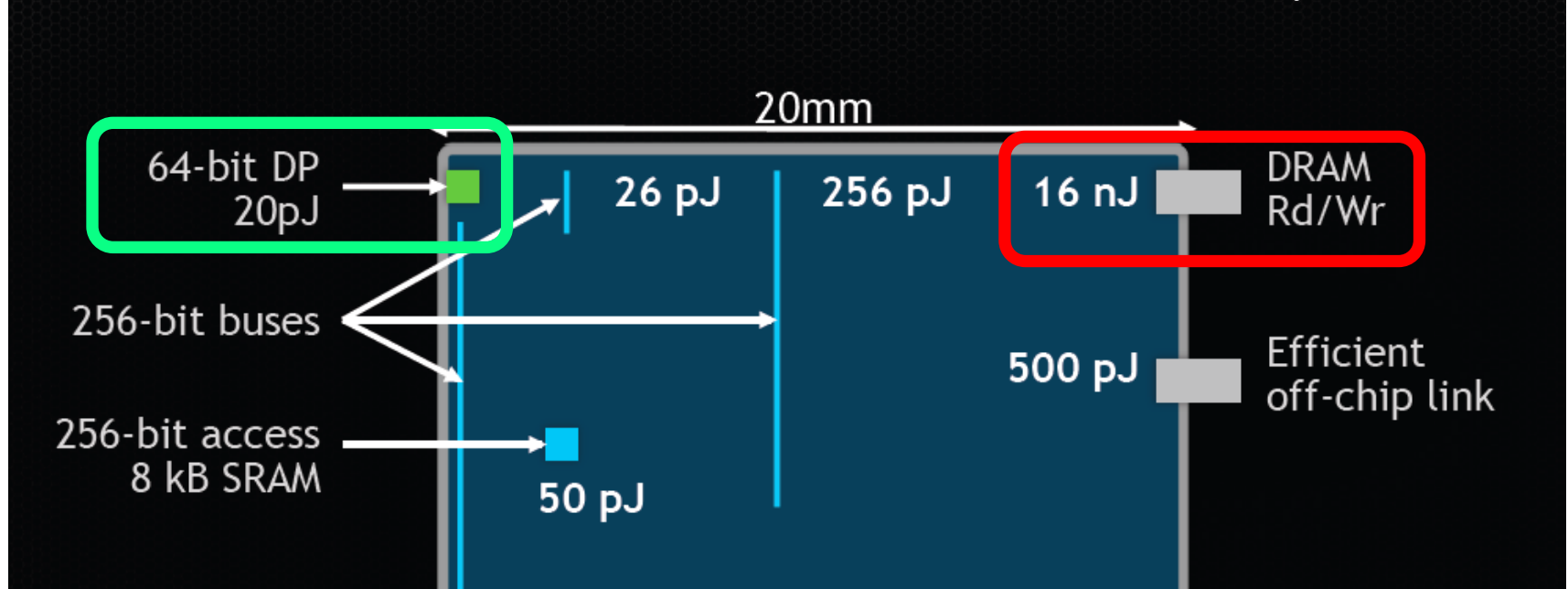
Dally, HiPEAC 2015



Data Movement vs. Computation Energy

Communication Dominates Arithmetic

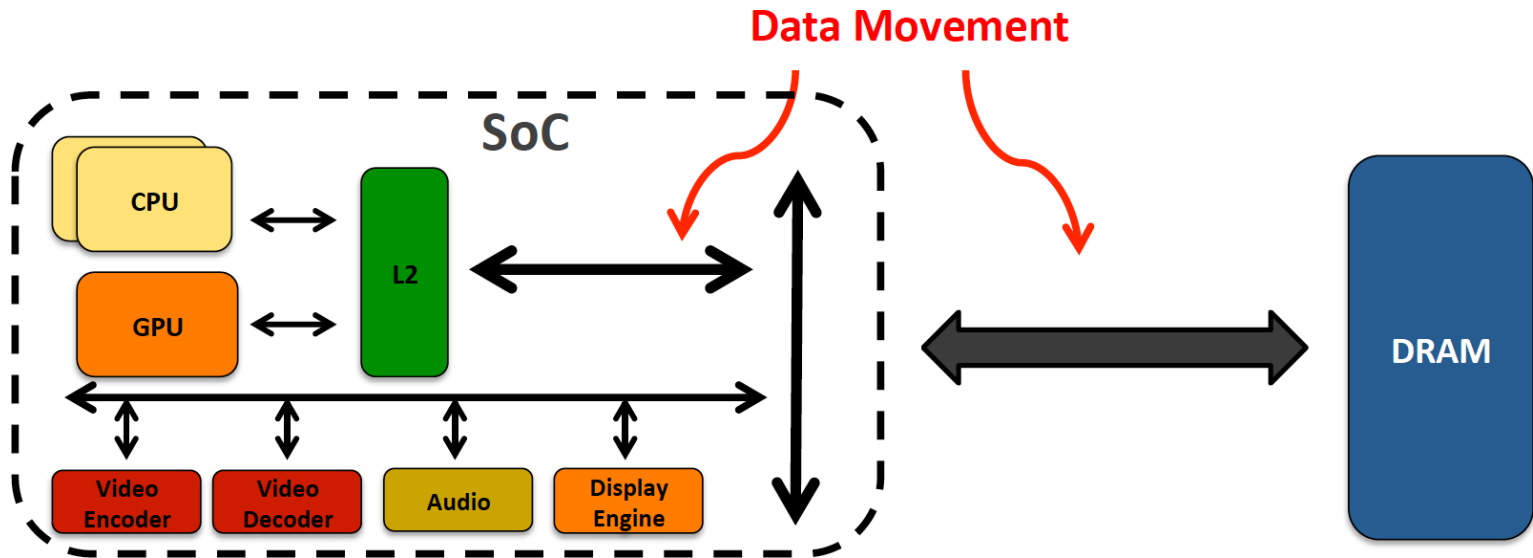
Dally, HiPEAC 2015



A memory access consumes $\sim 1000\times$ the energy of a complex addition

Data Movement vs. Computation Energy

- **Data movement** is a major system energy bottleneck
 - ❑ Comprises 41% of mobile system energy during web browsing [2]
 - ❑ Costs ~ 115 times as much energy as an ADD operation [1, 2]



[1]: Reducing data Movement Energy via Online Data Clustering and Encoding (MICRO'16)

[2]: Quantifying the energy cost of data movement for emerging smart phone workloads on mobile platforms (IISWC'14)

Energy Waste in Mobile Devices

- Amirali Boroumand, Saugata Ghose, Youngsok Kim, Rachata Ausavarungnirun, Eric Shiu, Rahul Thakur, Daehyun Kim, Aki Kuusela, Allan Knies, Parthasarathy Ranganathan, and Onur Mutlu, **"Google Workloads for Consumer Devices: Mitigating Data Movement Bottlenecks"** *Proceedings of the 23rd International Conference on Architectural Support for Programming Languages and Operating Systems (ASPLOS)*, Williamsburg, VA, USA, March 2018.

**62.7% of the total system energy
is spent on data movement**

Google Workloads for Consumer Devices: Mitigating Data Movement Bottlenecks

Amirali Boroumand¹

Saugata Ghose¹

Youngsok Kim²

Rachata Ausavarungnirun¹

Eric Shiu³

Rahul Thakur³

Daehyun Kim^{4,3}

Aki Kuusela³

Allan Knies³

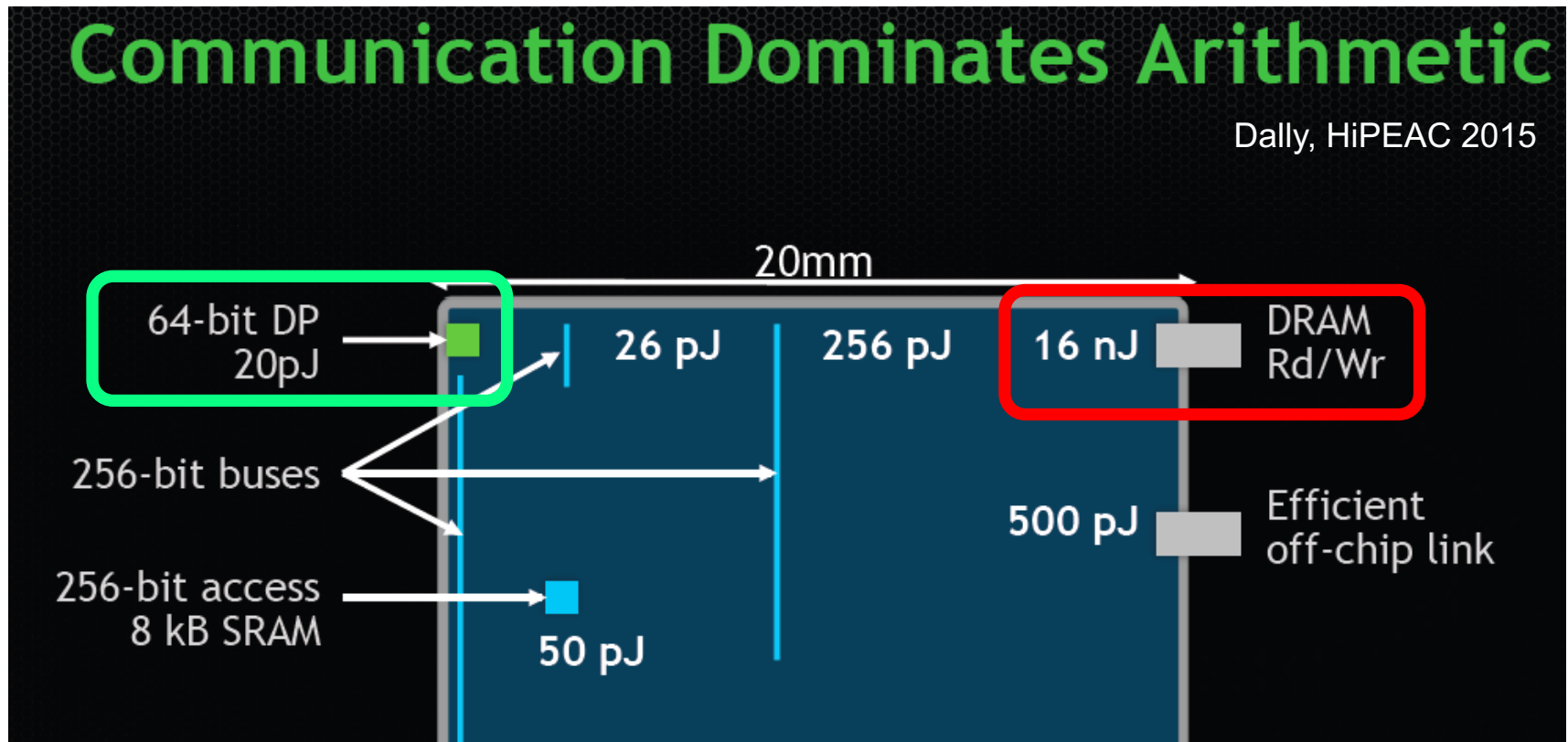
Parthasarathy Ranganathan³

Onur Mutlu^{5,1}

We Do Not Want to Move Data!

Communication Dominates Arithmetic

Dally, HiPEAC 2015

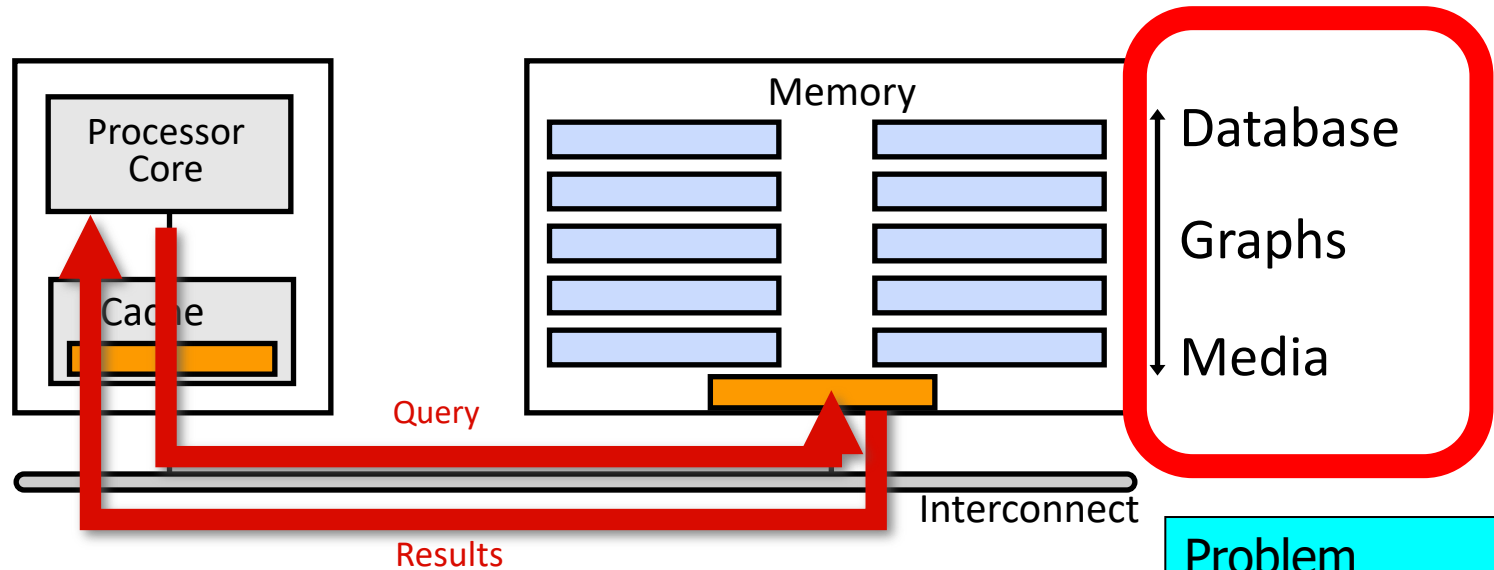


A memory access consumes $\sim 1000\times$
the energy of a complex addition

We Need A Paradigm Shift To ...

- Enable computation with minimal data movement
- Compute where it makes sense (where data resides)
- Make computing architectures more data-centric

Goal: Processing Inside Memory



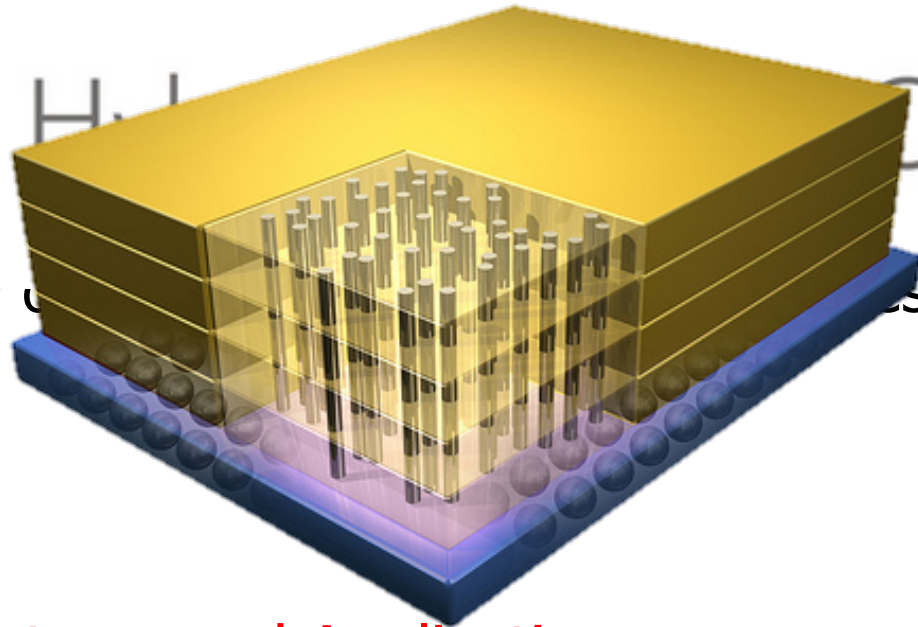
- Many questions ... How do we design the:
 - ❑ compute-capable memory & controllers?
 - ❑ processor chip and in-memory units?
 - ❑ software and hardware interfaces?
 - ❑ system software, compilers, languages?
 - ❑ algorithms and theoretical foundations?

Problem
Algorithm
Program/Language
System Software
SW/HW Interface
Micro-architecture
Logic
Devices
Electrons

Why In-Memory Computation Today?



→ Industry C



■ Pull from Systems and Applications

- ❑ Data access is a major system and application bottleneck
- ❑ Systems are energy limited
- ❑ Data movement much more energy-hungry than computation

Agenda

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- How to Enable Adoption of Processing in Memory
- Conclusion

Processing in Memory: Two Approaches

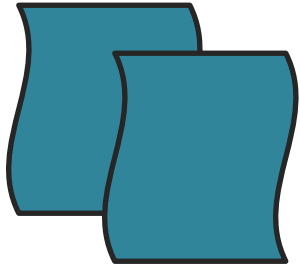
1. Minimally changing memory chips
2. Exploiting 3D-stacked memory

Approach 1: Minimally Changing DRAM

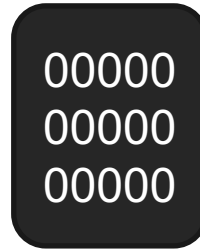
- DRAM has great capability to perform **bulk data movement and computation** internally with small changes
 - Can **exploit internal connectivity** to move data
 - Can **exploit analog computation capability**
 - ...
- Examples: RowClone, In-DRAM AND/OR, Gather/Scatter DRAM
 - RowClone: Fast and Efficient In-DRAM Copy and Initialization of Bulk Data (Seshadri et al., MICRO 2013)
 - Fast Bulk Bitwise AND and OR in DRAM (Seshadri et al., IEEE CAL 2015)
 - Gather-Scatter DRAM: In-DRAM Address Translation to Improve the Spatial Locality of Non-unit Strided Accesses (Seshadri et al., MICRO 2015)
 - "Ambit: In-Memory Accelerator for Bulk Bitwise Operations Using Commodity DRAM Technology" (Seshadri et al., MICRO 2017)

Starting Simple: Data Copy and Initialization

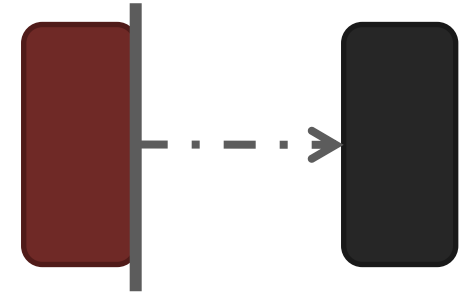
memmove & memcpy: 5% cycles in Google's datacenter [Kanev+ ISCA'15]



Forking



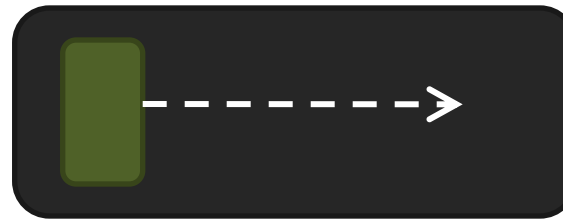
**Zero initialization
(e.g., security)**



Checkpointing



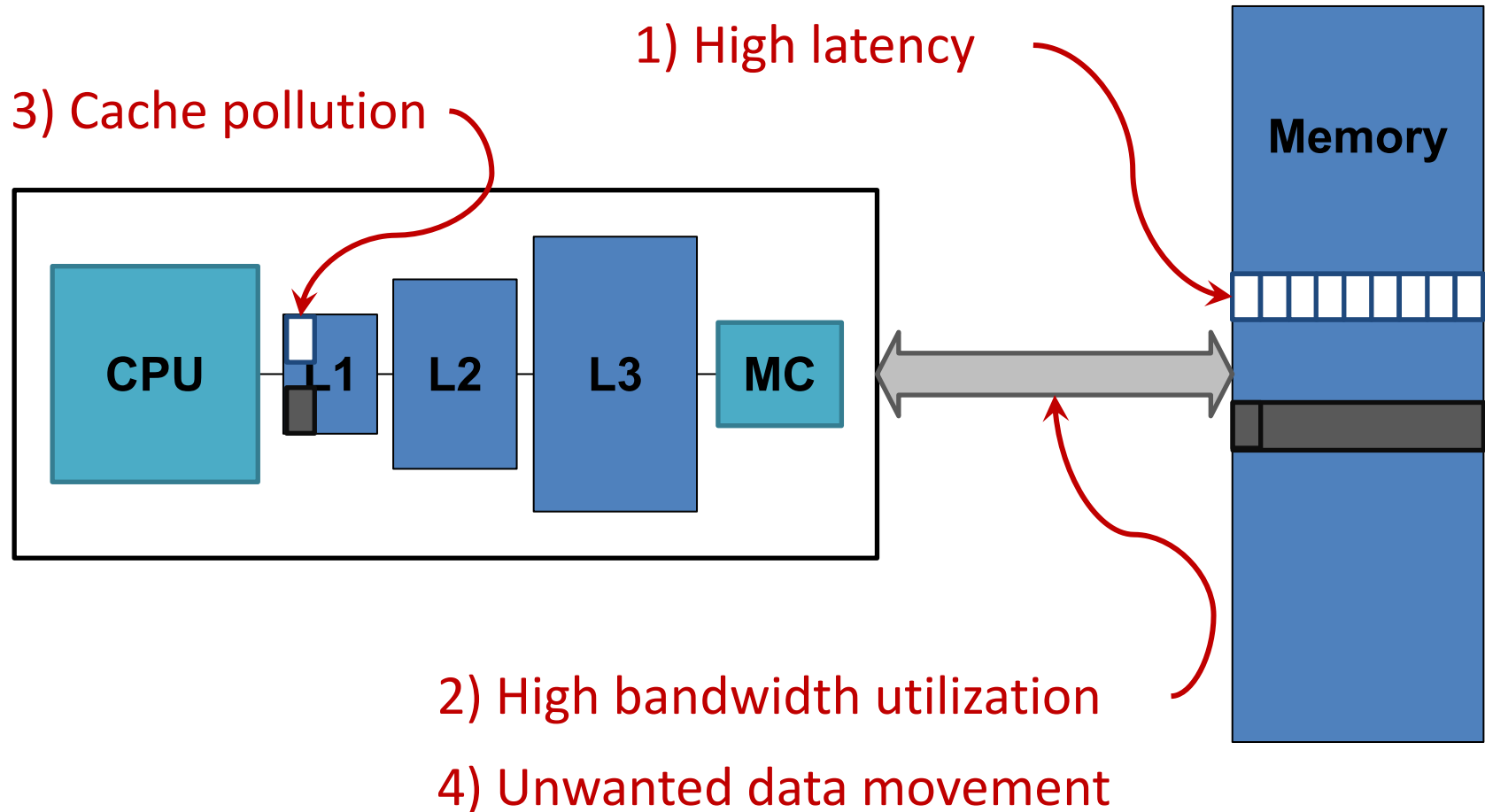
**VM Cloning
Deduplication**



Page Migration

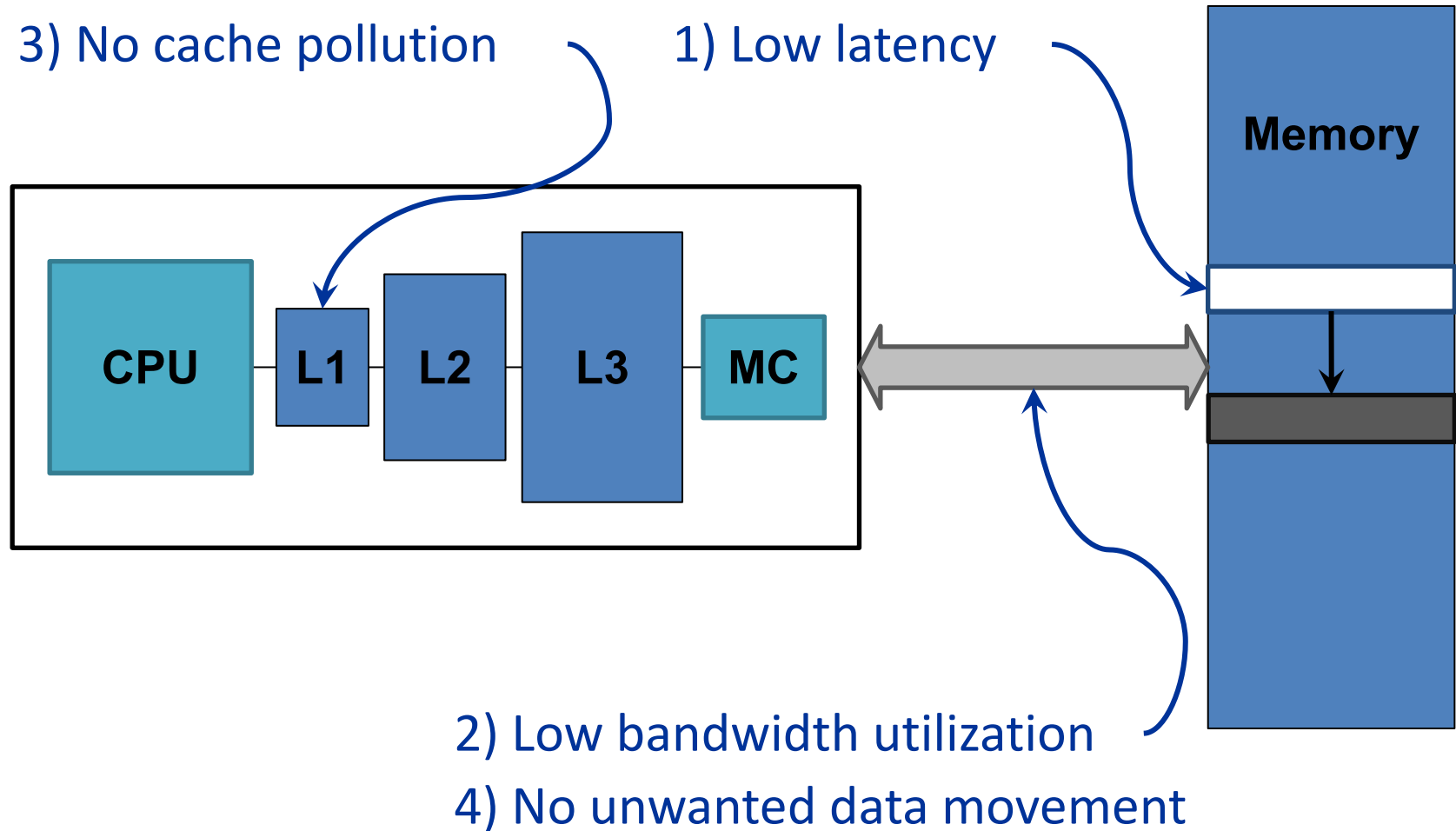
...
Many more

Today's Systems: Bulk Data Copy



1046ns, 3.6uJ (for 4KB page copy via DMA)

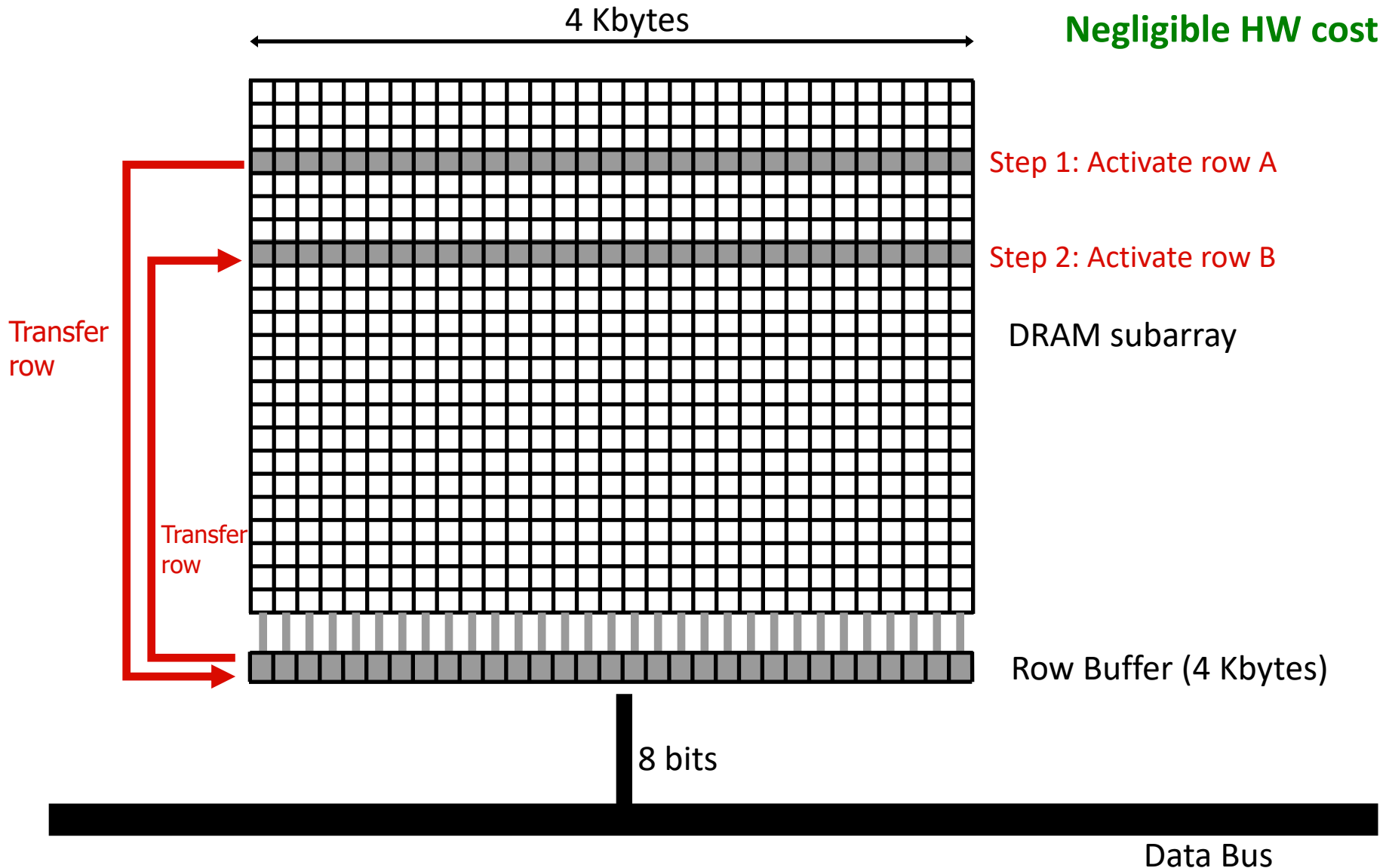
Future Systems: In-Memory Copy



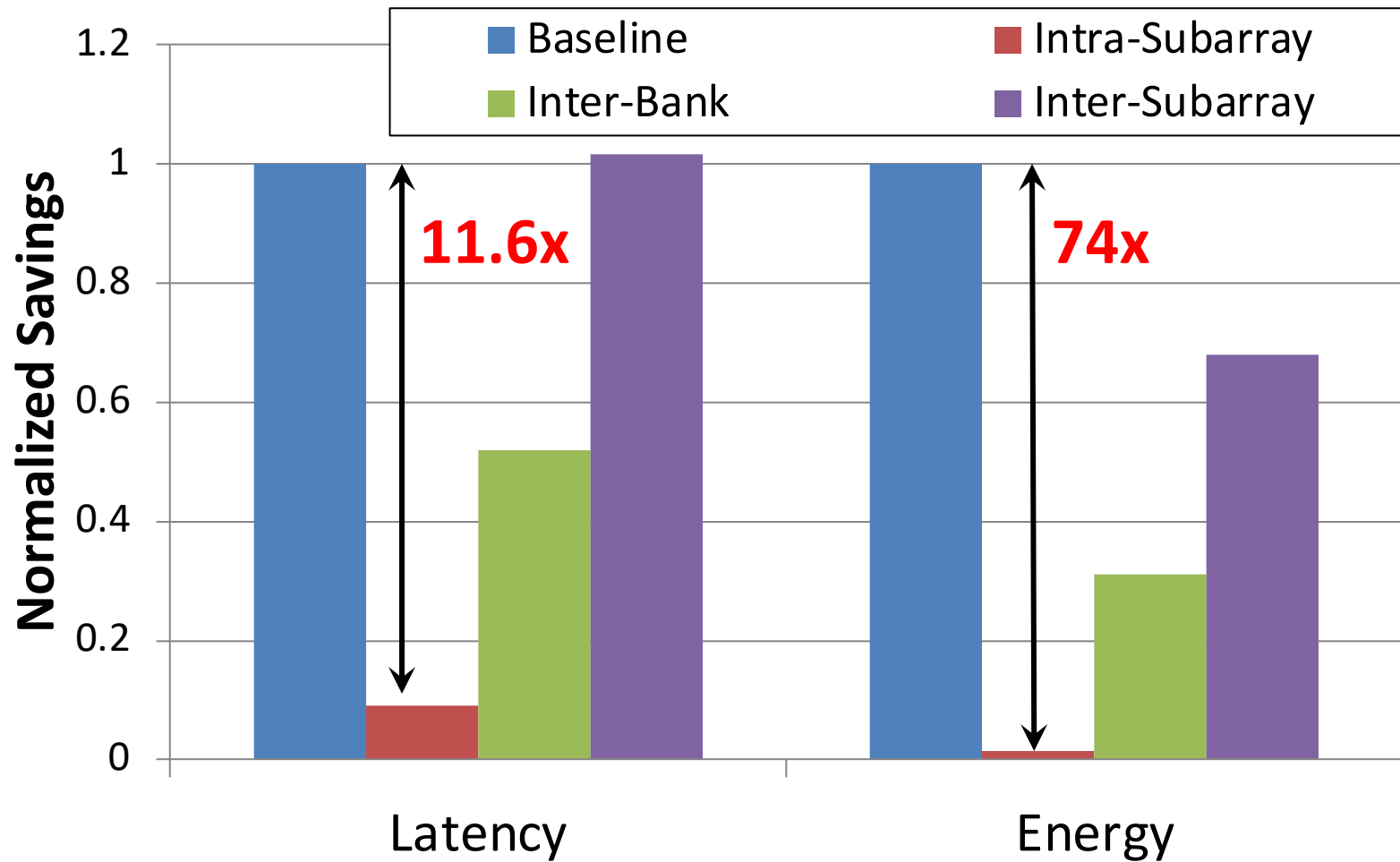
1046ns, 3.6uJ → 90ns, 0.04uJ

RowClone: In-DRAM Row Copy

Idea: Two consecutive ACTivates
Negligible HW cost



RowClone: Latency and Energy Savings



Seshadri et al., "RowClone: Fast and Efficient In-DRAM Copy and Initialization of Bulk Data," MICRO 2013.

More on RowClone

- Vivek Seshadri, Yoongu Kim, Chris Fallin, Donghyuk Lee, Rachata Ausavarungnirun, Gennady Pekhimenko, Yixin Luo, Onur Mutlu, Michael A. Kozuch, Phillip B. Gibbons, and Todd C. Mowry,
"RowClone: Fast and Energy-Efficient In-DRAM Bulk Data Copy and Initialization"
Proceedings of the 46th International Symposium on Microarchitecture (MICRO), Davis, CA, December 2013. [[Slides \(pptx\)](#)] [[pdf](#)] [[Lightning Session Slides \(pptx\)](#)] [[pdf](#)] [[Poster \(pptx\)](#)] [[pdf](#)]

RowClone: Fast and Energy-Efficient In-DRAM Bulk Data Copy and Initialization

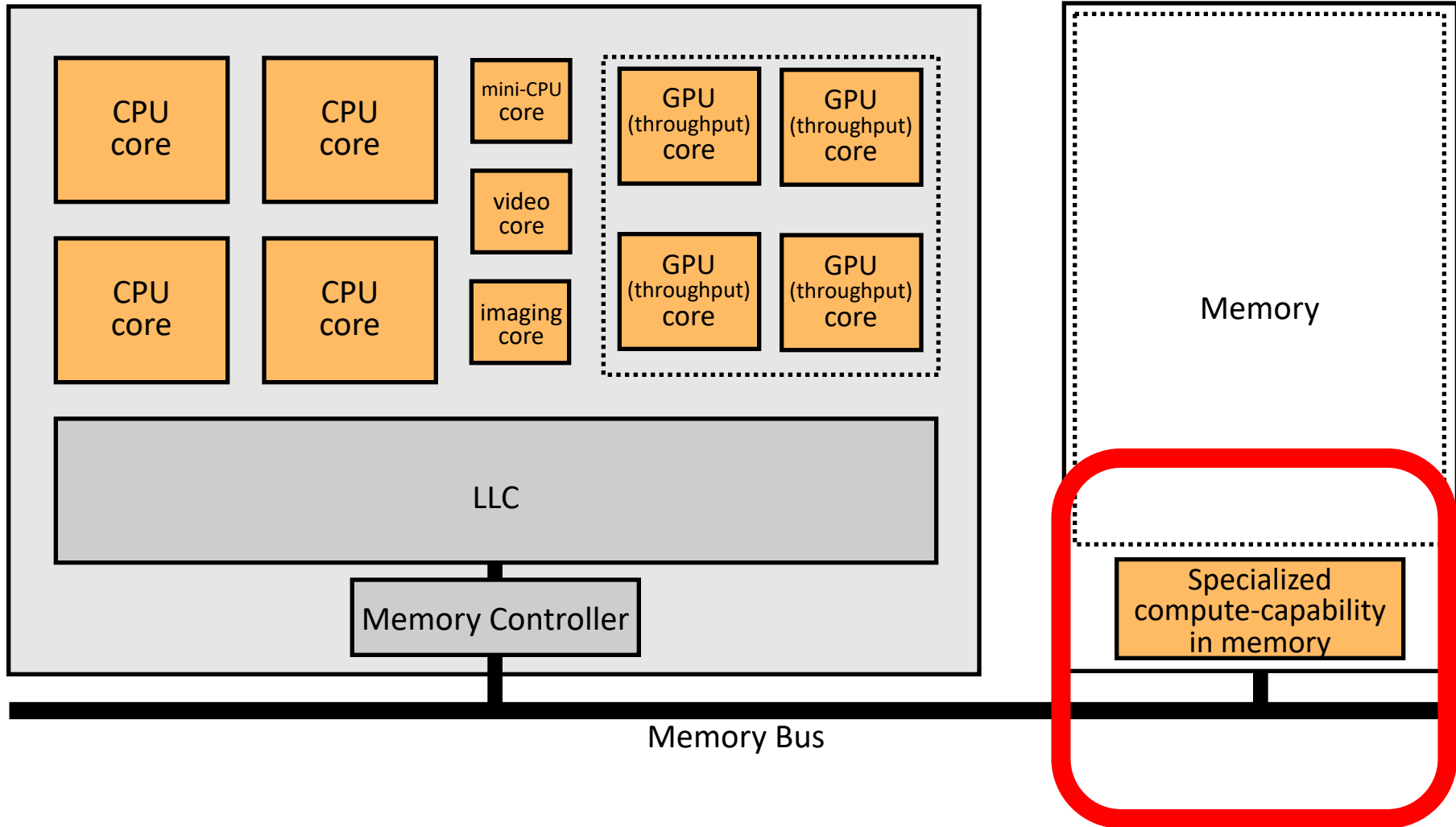
Vivek Seshadri vseshadr@cs.cmu.edu	Yoongu Kim yoongukim@cmu.edu	Chris Fallin* cfallin@c1f.net	Donghyuk Lee donghyuk1@cmu.edu
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Rachata Ausavarungnirun rachata@cmu.edu	Gennady Pekhimenko gpekhime@cs.cmu.edu	Yixin Luo yixinluo@andrew.cmu.edu
--	---	--------------------------------------

Onur Mutlu onur@cmu.edu	Phillip B. Gibbons† phillip.b.gibbons@intel.com	Michael A. Kozuch† michael.a.kozuch@intel.com	Todd C. Mowry tcm@cs.cmu.edu
----------------------------	--	--	---------------------------------

Carnegie Mellon University †Intel Pittsburgh

Memory as an Accelerator

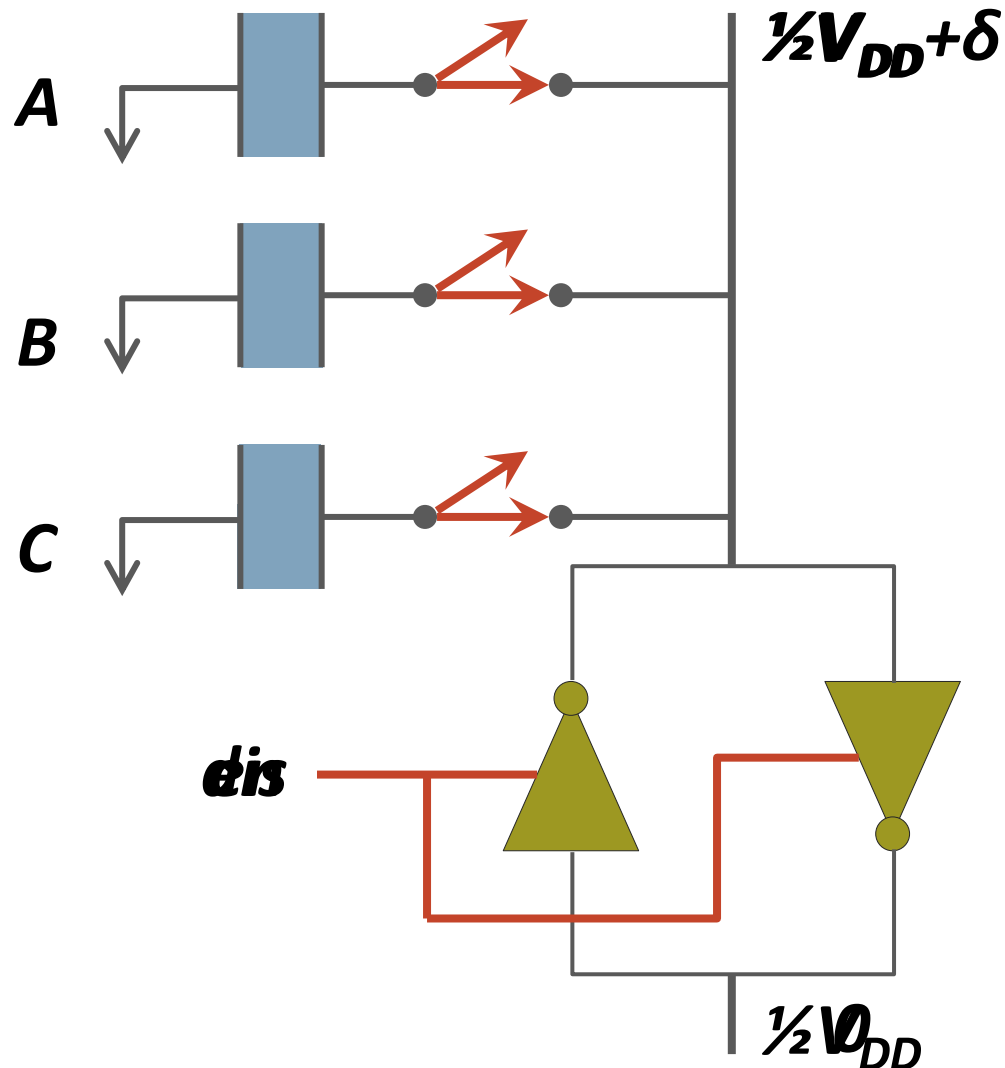


Memory similar to a "conventional" accelerator

In-Memory Bulk Bitwise Operations

- We can support in-DRAM COPY, ZERO, AND, OR, NOT, MAJ
- At low cost
- Using analog computation capability of DRAM
 - Idea: activating multiple rows performs computation
- 30-60X performance and energy improvement
 - Seshadri+, “Ambit: In-Memory Accelerator for Bulk Bitwise Operations Using Commodity DRAM Technology,” MICRO 2017.
- New memory technologies enable even more opportunities
 - Memristors, resistive RAM, phase change mem, STT-MRAM, ...
 - Can operate on data with minimal movement

In-DRAM AND/OR: Triple Row Activation



Final State
 $AB + BC + AC$

**$C(A + B) +$
 $\sim C(AB)$**

In-DRAM NOT: Dual Contact Cell

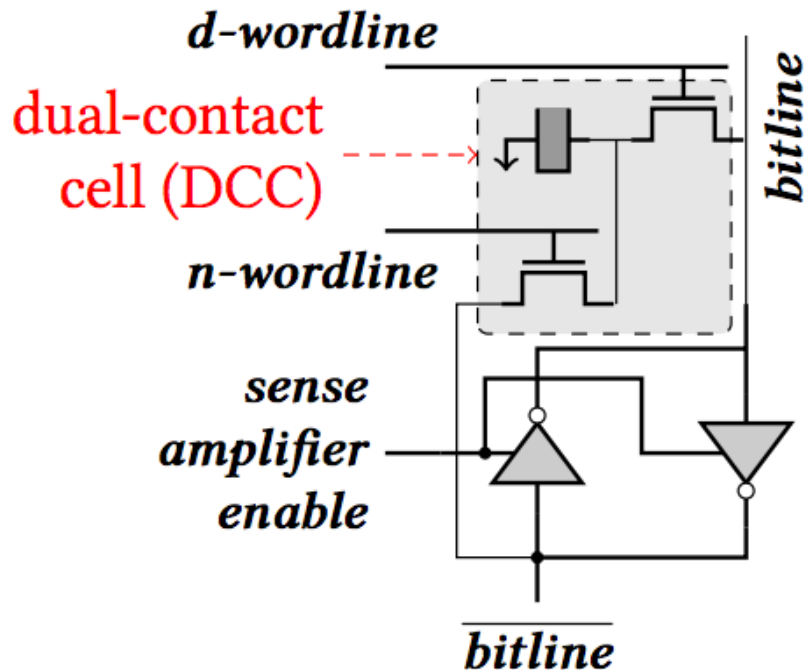


Figure 5: A dual-contact cell connected to both ends of a sense amplifier

Idea:
Feed the
negated value
in the sense amplifier
into a special row

Performance: In-DRAM Bitwise Operations

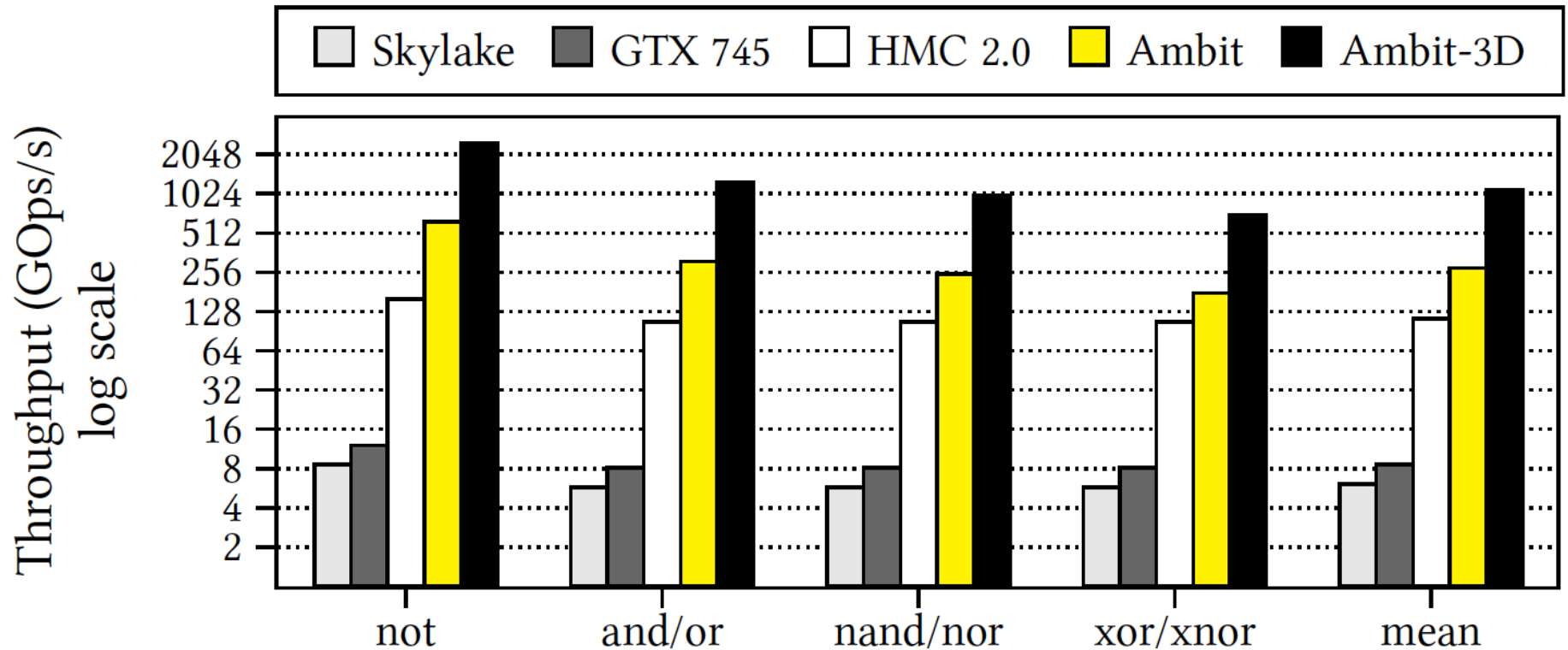


Figure 9: Throughput of bitwise operations on various systems.

Seshadri+, "Ambit: In-Memory Accelerator for Bulk Bitwise Operations using Commodity DRAM Technology," MICRO 2017.

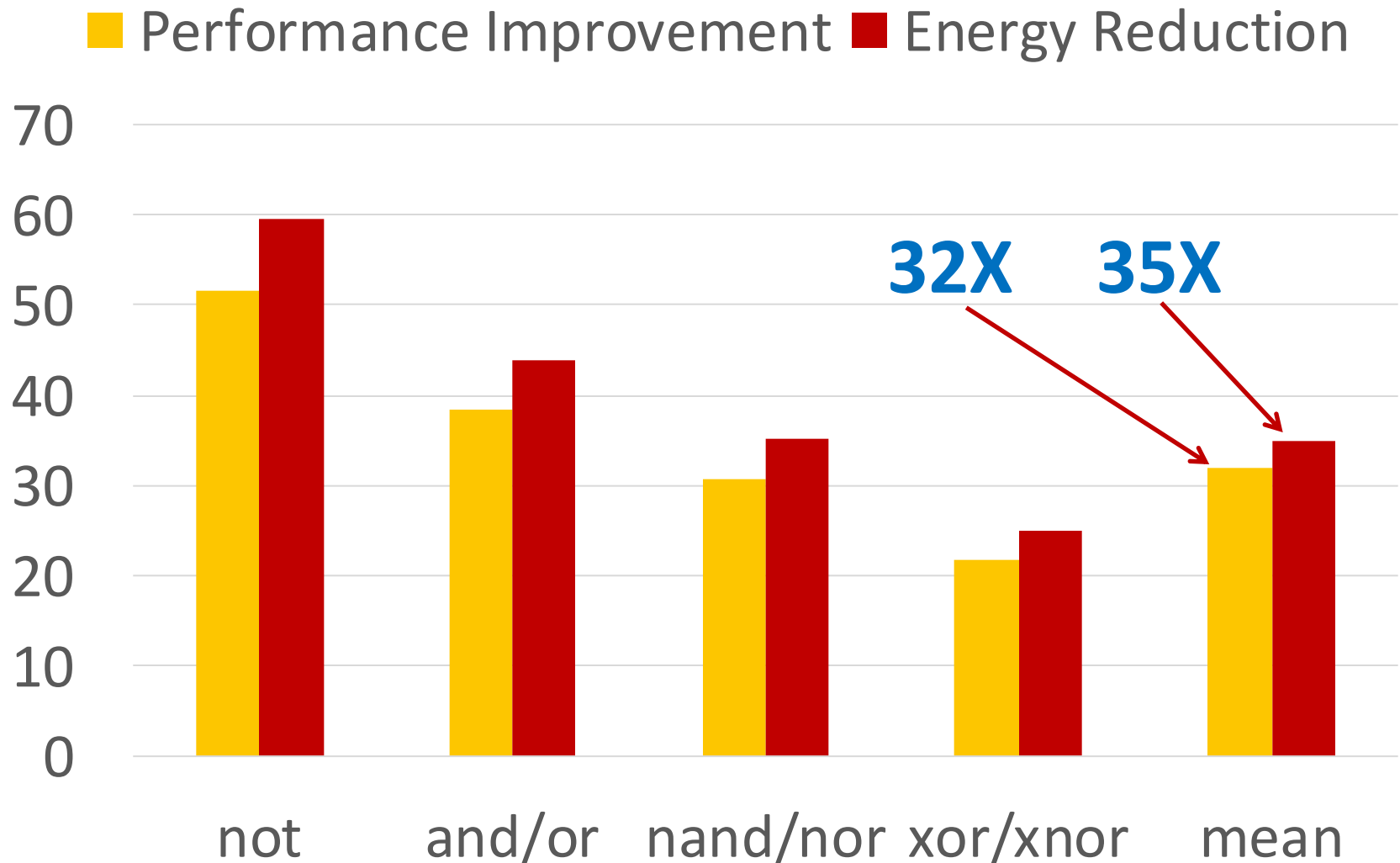
Energy of In-DRAM Bitwise Operations

	Design	not	and/or	nand/nor	xor/xnor
DRAM & Channel Energy (nJ/KB)	DDR3	93.7	137.9	137.9	137.9
	Ambit	1.6	3.2	4.0	5.5
	(↓)	59.5X	43.9X	35.1X	25.1X

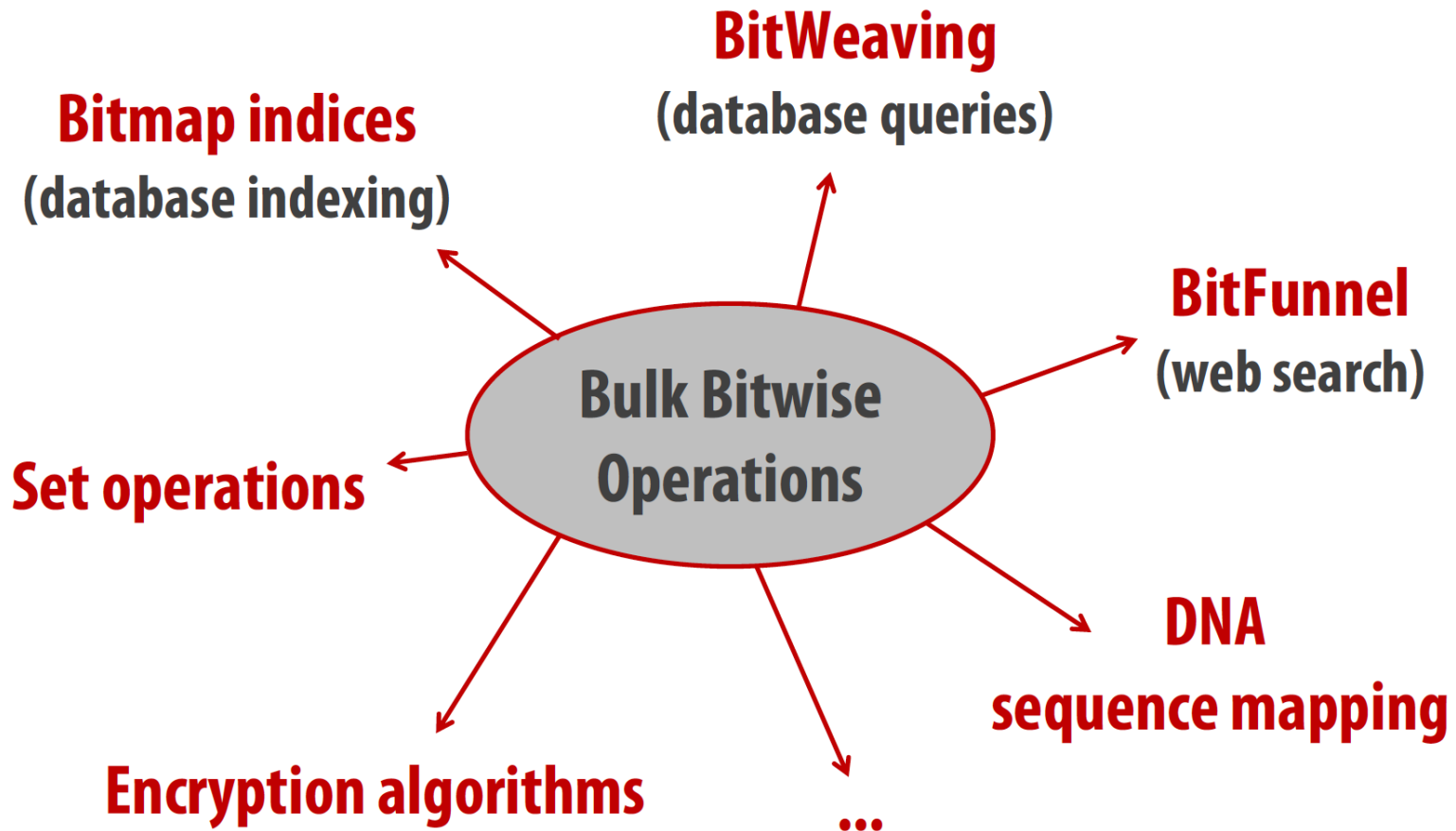
Table 3: Energy of bitwise operations. (↓) indicates energy reduction of Ambit over the traditional DDR3-based design.

Seshadri+, "Ambit: In-Memory Accelerator for Bulk Bitwise Operations using Commodity DRAM Technology," MICRO 2017.

Ambit vs. DDR3: Performance and Energy



Bulk Bitwise Operations in Workloads



Example Data Structure: Bitmap Index

- Alternative to B-tree and its variants
- Efficient for performing *range queries* and *joins*
- **Many bitwise operations to perform a query**

age < 18 18 < age < 25 25 < age < 60 age > 60

Bitmap 1

Bitmap 2

Bitmap 3

Bitmap 4

Performance: Bitmap Index on Ambit

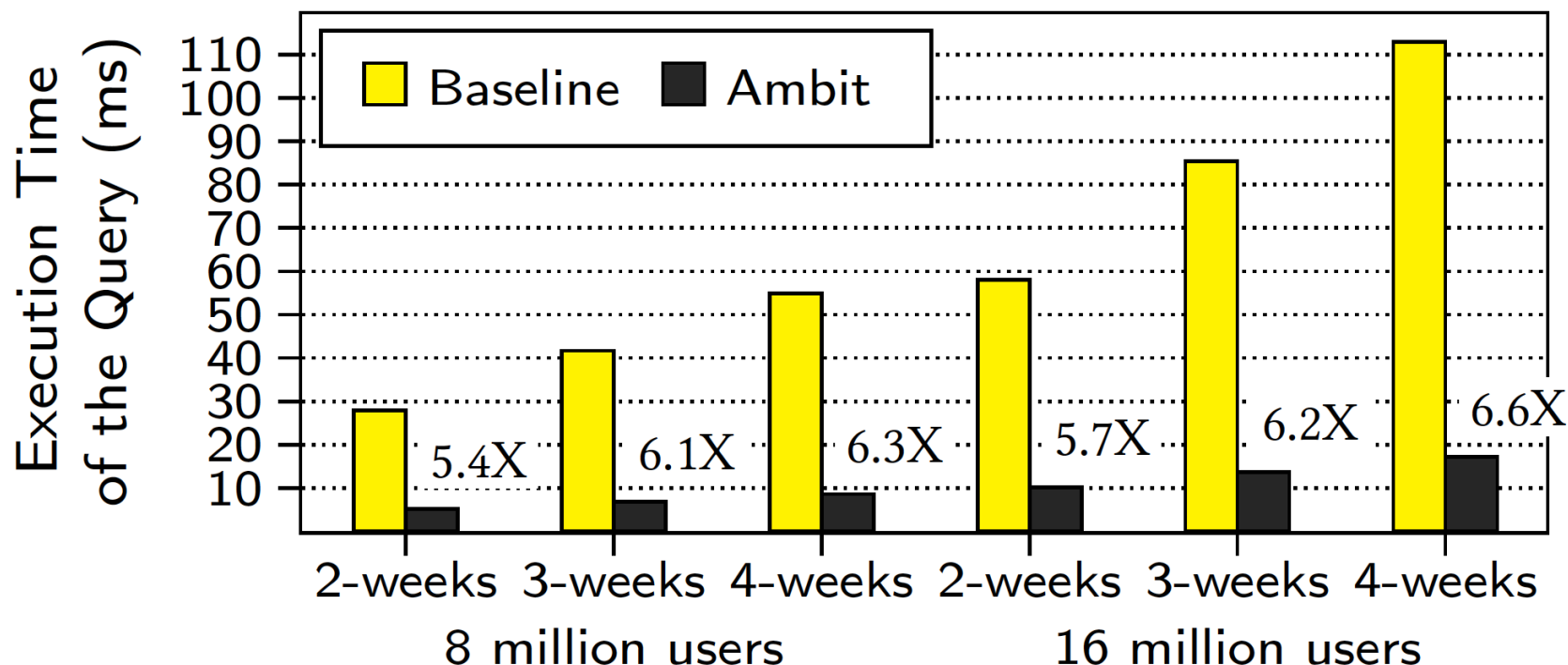


Figure 10: Bitmap index performance. The value above each bar indicates the reduction in execution time due to Ambit.

>5.4-6.6X Performance Improvement

Performance: BitWeaving on Ambit

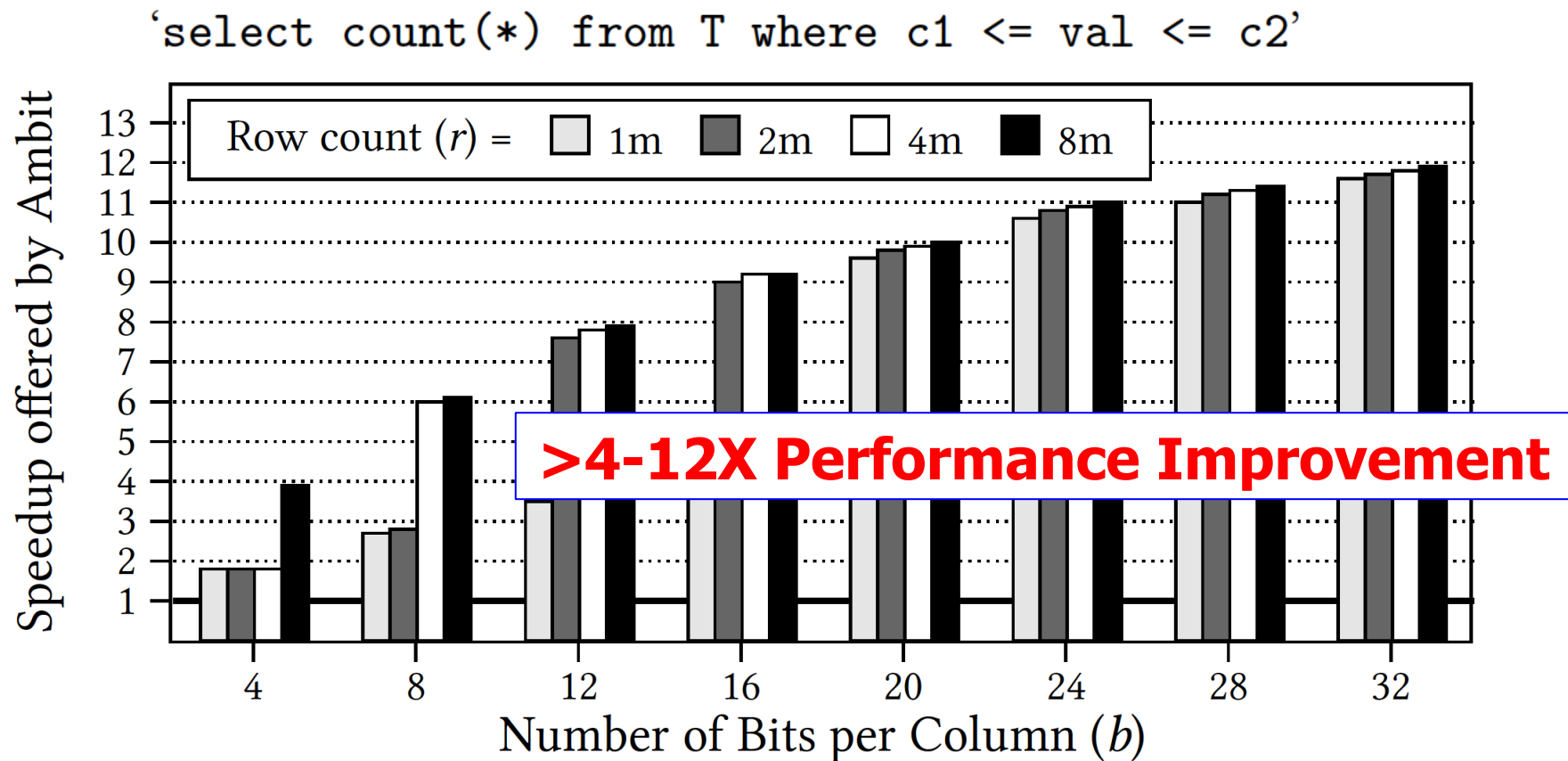


Figure 11: Speedup offered by Ambit over baseline CPU with SIMD for BitWeaving

Seshadri+, "Ambit: In-Memory Accelerator for Bulk Bitwise Operations using Commodity DRAM Technology," MICRO 2017.

More on In-DRAM Bulk AND/OR

- Vivek Seshadri, Kevin Hsieh, Amirali Boroumand, Donghyuk Lee, Michael A. Kozuch, Onur Mutlu, Phillip B. Gibbons, and Todd C. Mowry,
"Fast Bulk Bitwise AND and OR in DRAM"
IEEE Computer Architecture Letters (***CAL***), April 2015.

Fast Bulk Bitwise AND and OR in DRAM

Vivek Seshadri*, Kevin Hsieh*, Amirali Boroumand*, Donghyuk Lee*,
Michael A. Kozuch†, Onur Mutlu*, Phillip B. Gibbons†, Todd C. Mowry*

*Carnegie Mellon University

†Intel Pittsburgh

More on Ambit

- Vivek Seshadri et al., “**Ambit: In-Memory Accelerator for Bulk Bitwise Operations Using Commodity DRAM Technology**,” MICRO 2017.

Ambit: In-Memory Accelerator for Bulk Bitwise Operations Using Commodity DRAM Technology

Vivek Seshadri^{1,5} Donghyuk Lee^{2,5} Thomas Mullins^{3,5} Hasan Hassan⁴ Amirali Boroumand⁵
Jeremie Kim^{4,5} Michael A. Kozuch³ Onur Mutlu^{4,5} Phillip B. Gibbons⁵ Todd C. Mowry⁵

¹Microsoft Research India ²NVIDIA Research ³Intel ⁴ETH Zürich ⁵Carnegie Mellon University

More on In-DRAM Bulk Bitwise Execution

- Vivek Seshadri and Onur Mutlu,
"In-DRAM Bulk Bitwise Execution Engine"
Invited Book Chapter in Advances in Computers, to appear
in 2020.
[[Preliminary arXiv version](#)]

In-DRAM Bulk Bitwise Execution Engine

Vivek Seshadri
Microsoft Research India
`visesha@microsoft.com`

Onur Mutlu
ETH Zürich
`onur.mutlu@inf.ethz.ch`

Computing Architectures with Minimal Data Movement

Challenge: Intelligent Memory Device

Does **memory**
have to be
dumb?

Agenda

- Major Trends Affecting Main Memory
- The Need for Intelligent Memory Controllers
 - Bottom Up: Push from Circuits and Devices
 - Top Down: Pull from Systems and Applications
- Processing in Memory: Two Directions
 - Minimally Changing Memory Chips
 - Exploiting 3D-Stacked Memory
- How to Enable Adoption of Processing in Memory
- Conclusion

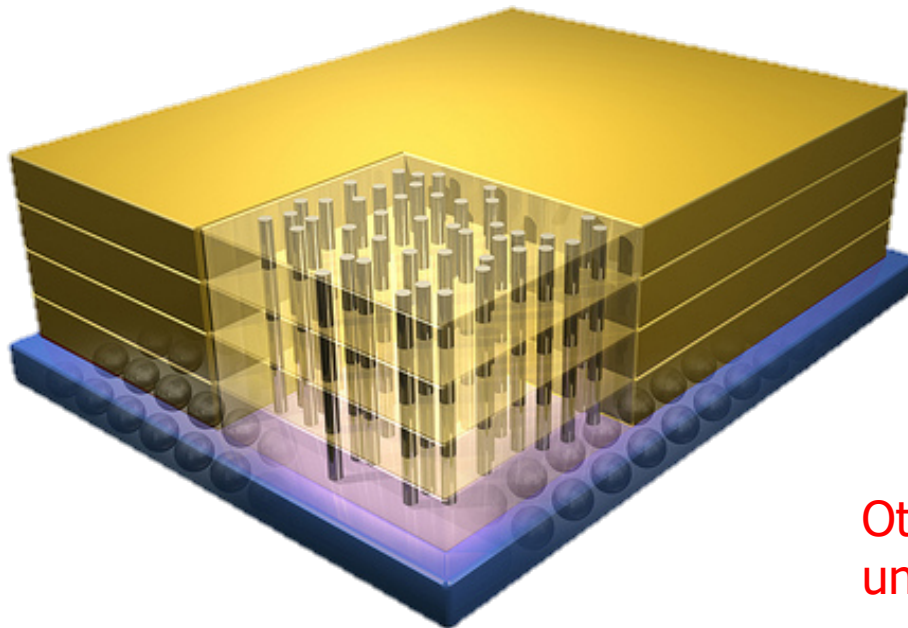
Processing in Memory: Two Approaches

1. Minimally changing memory chips
2. Exploiting 3D-stacked memory

Opportunity: 3D-Stacked Logic+Memory



Hybrid Memory Cube
C O N S O R T I U M



Memory

Logic

Other "True 3D" technologies
under development

DRAM Landscape (circa 2015)

<i>Segment</i>	<i>DRAM Standards & Architectures</i>
Commodity	DDR3 (2007) [14]; DDR4 (2012) [18]
Low-Power	LPDDR3 (2012) [17]; LPDDR4 (2014) [20]
Graphics	GDDR5 (2009) [15]
Performance	eDRAM [28], [32]; RLDram3 (2011) [29]
3D-Stacked	WIO (2011) [16]; WIO2 (2014) [21]; MCDRAM (2015) [13]; HBM (2013) [19]; HMC1.0 (2013) [10]; HMC1.1 (2014) [11]
Academic	SBA/SSA (2010) [38]; Staged Reads (2012) [8]; RAIDR (2012) [27]; SALP (2012) [24]; TL-DRAM (2013) [26]; RowClone (2013) [37]; Half-DRAM (2014) [39]; Row-Buffer Decoupling (2014) [33]; SARP (2014) [6]; AL-DRAM (2015) [25]

Table 1. Landscape of DRAM-based memory

Kim+, “[Ramulator: A Flexible and Extensible DRAM Simulator](#)”, IEEE CAL 2015.

Two Key Questions in 3D-Stacked PIM

- What are the performance and energy benefits of using 3D-stacked memory as a coarse-grained accelerator?
 - By changing the entire system
 - By performing simple function offloading
- What is the minimal processing-in-memory support we can provide?
 - With minimal changes to system and programming

Graph Processing

- Large graphs are everywhere (circa 2015)



36 Million
Wikipedia Pages



1.4 Billion
Facebook Users

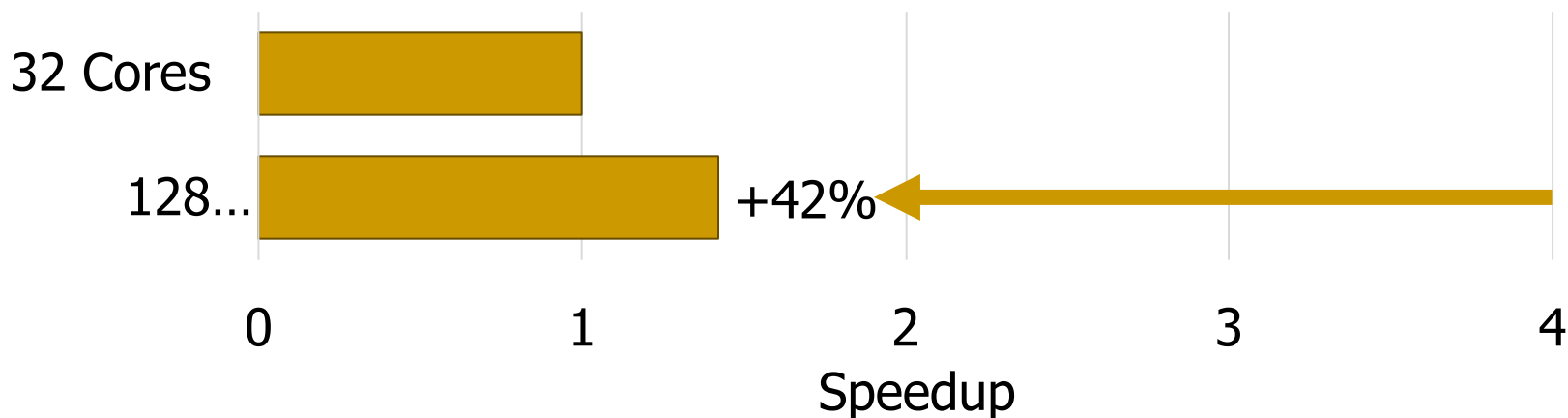


300 Million
Twitter Users



30 Billion
Instagram Photos

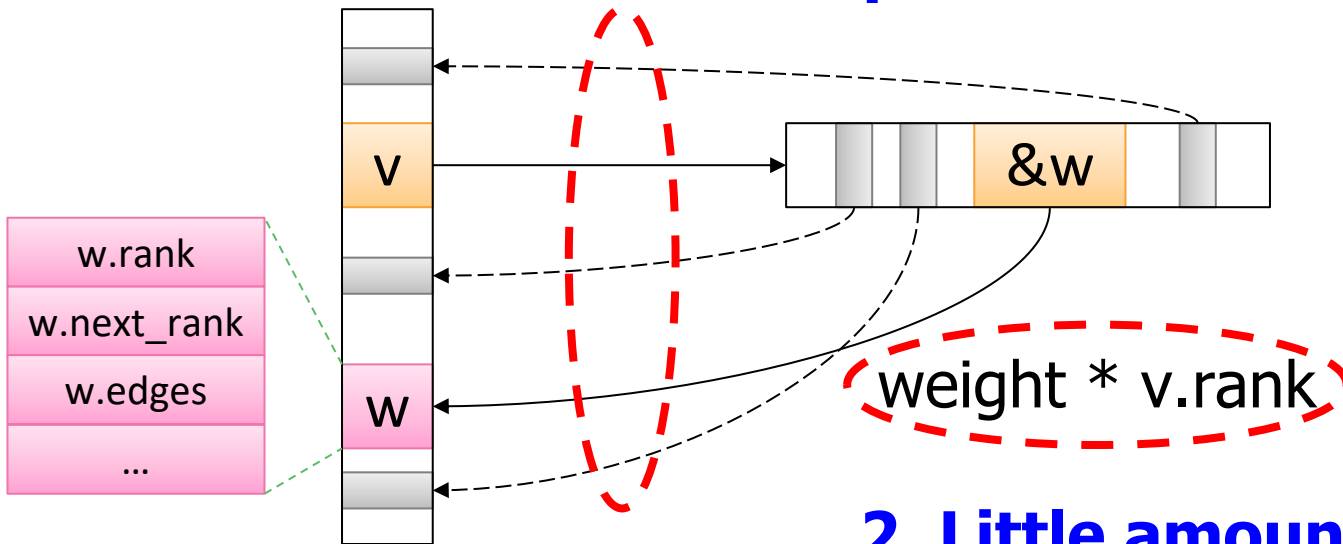
- Scalable large-scale graph processing is challenging



Key Bottlenecks in Graph Processing

```
for (v: graph.vertices) {  
  for (w: v.successors) {  
    w.next_rank += weight * v.rank;  
  }  
}
```

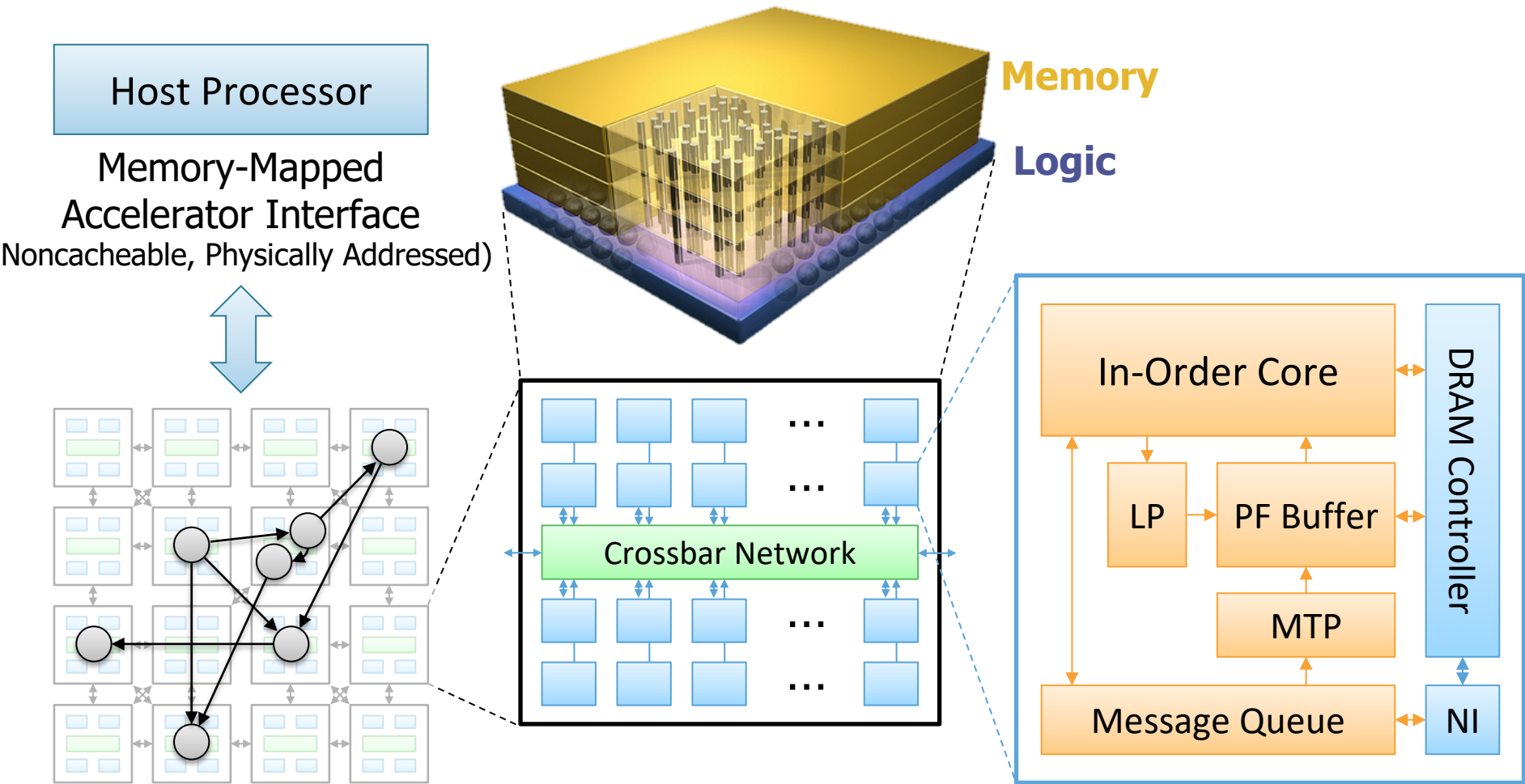
1. Frequent random memory accesses



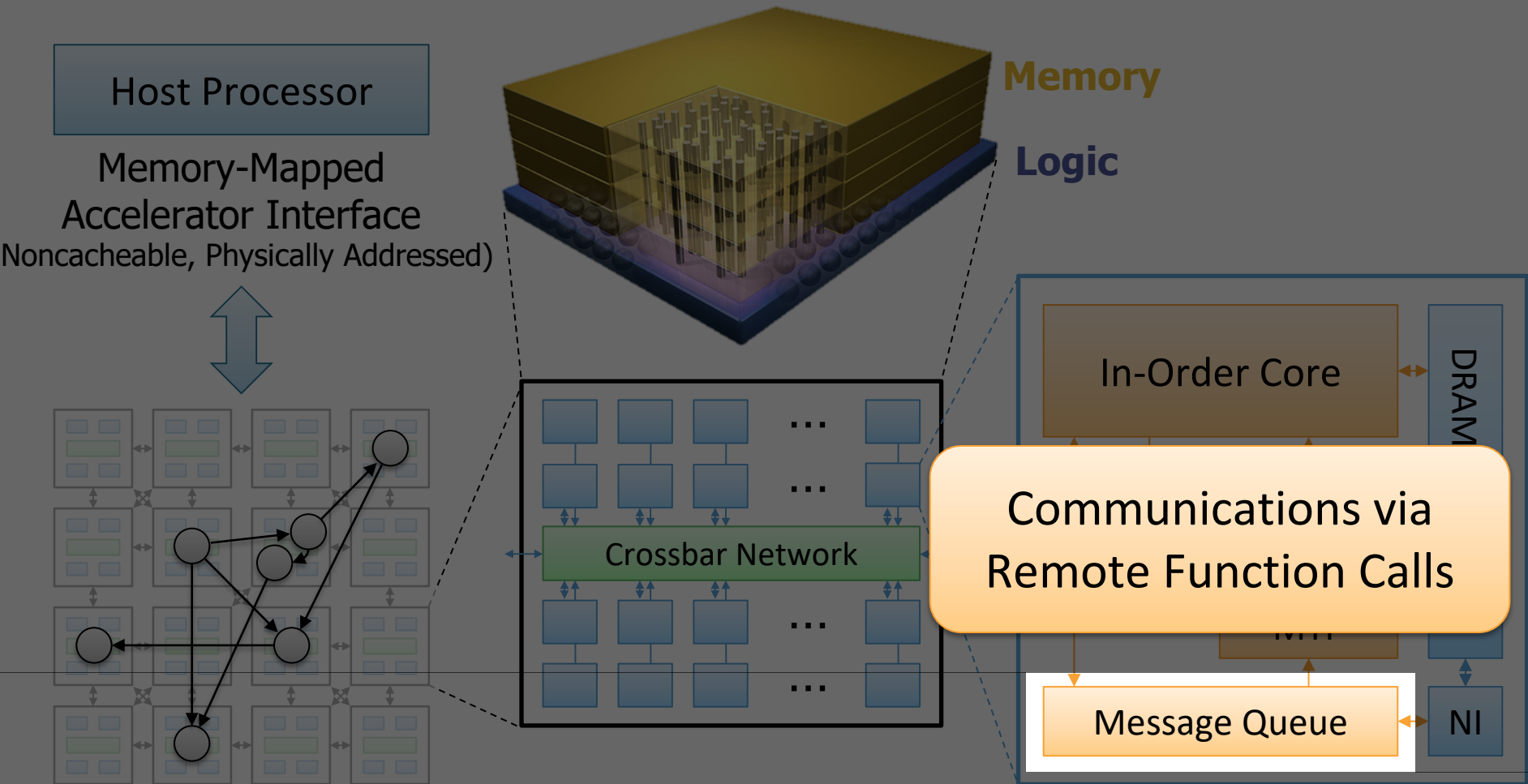
2. Little amount of computation

Tesseract System for Graph Processing

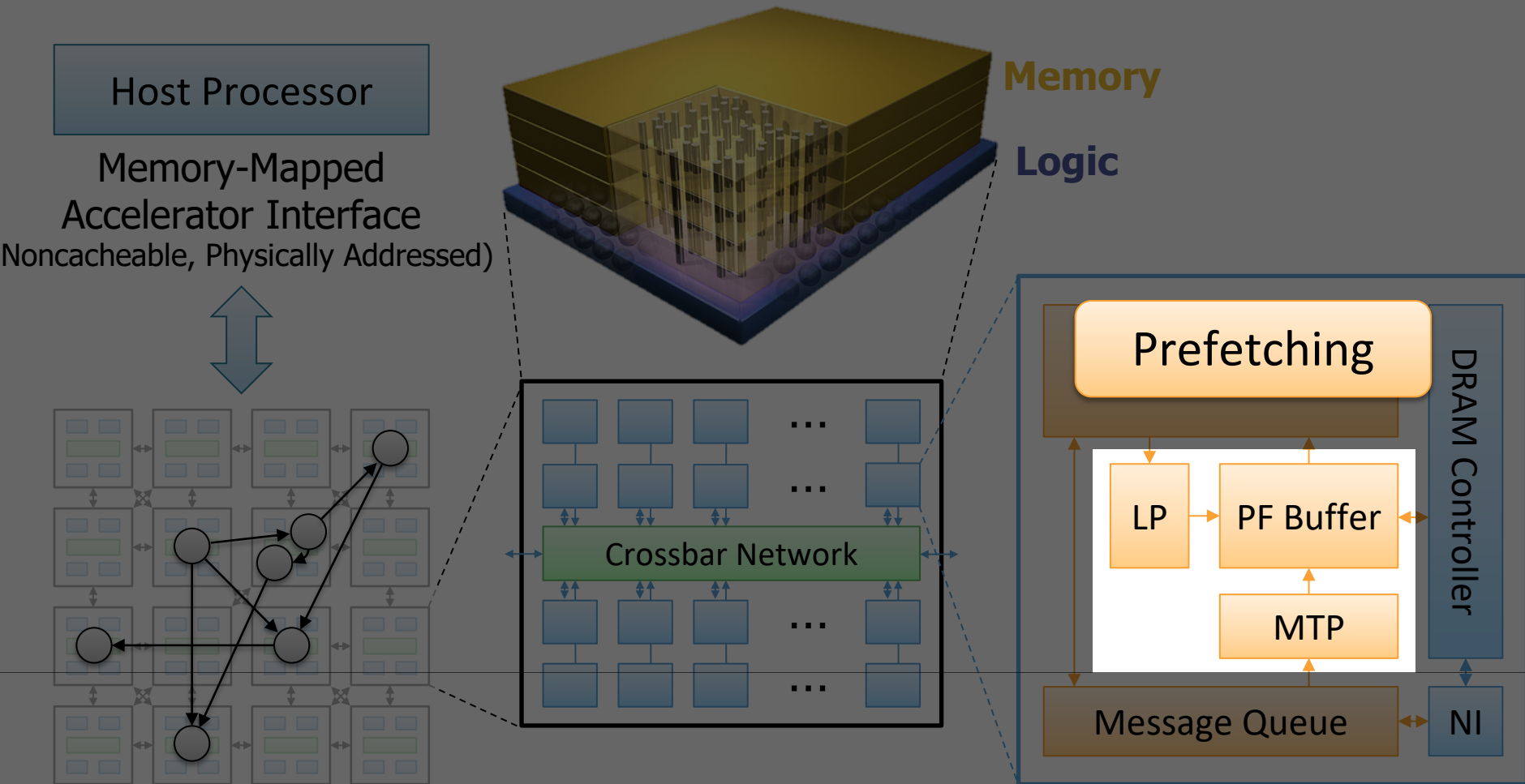
Interconnected set of 3D-stacked memory+logic chips with simple cores



Tesseract System for Graph Processing

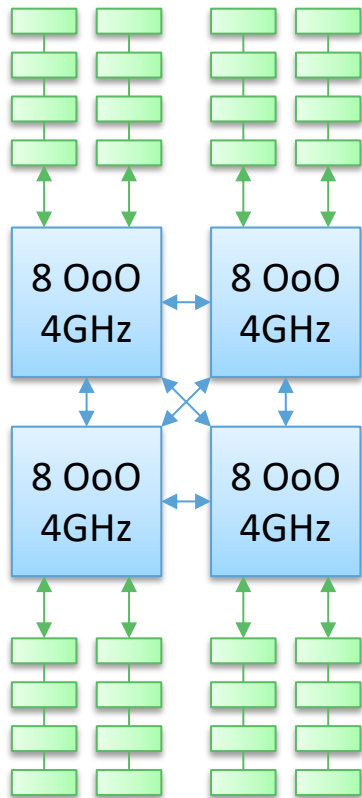


Tesseract System for Graph Processing



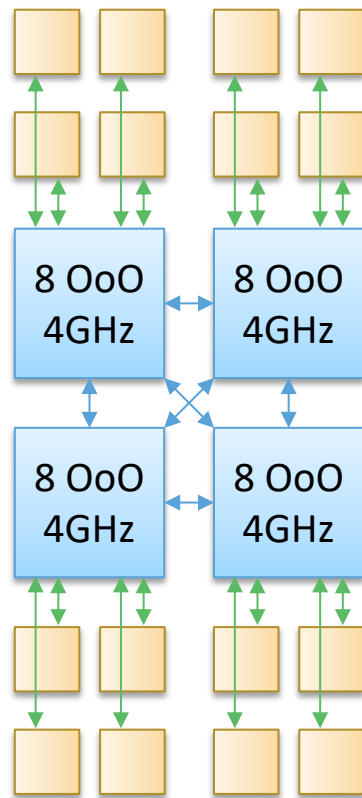
Evaluated Systems

DDR3-OoO



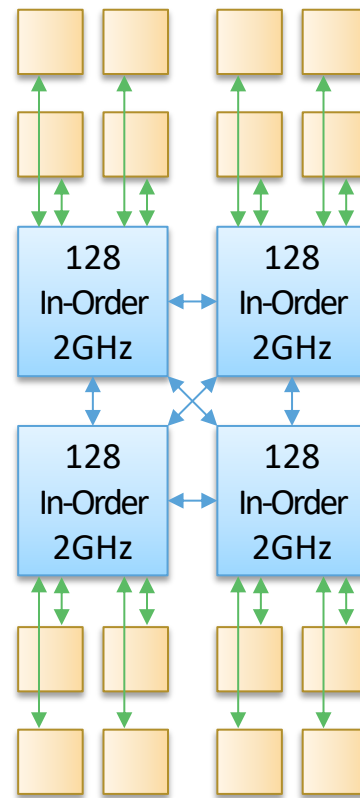
102.4GB/s

HMC-OoO



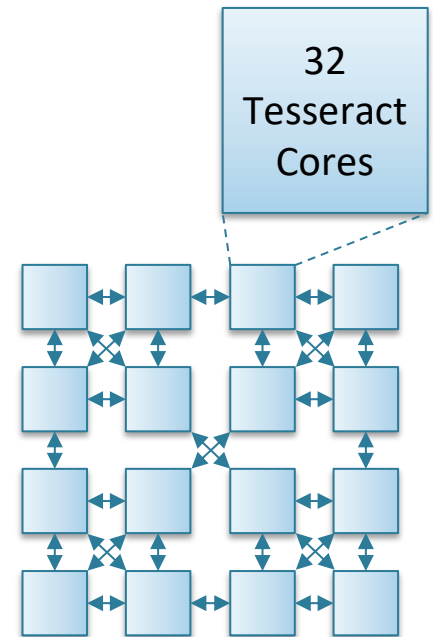
640GB/s

HMC-MC



640GB/s

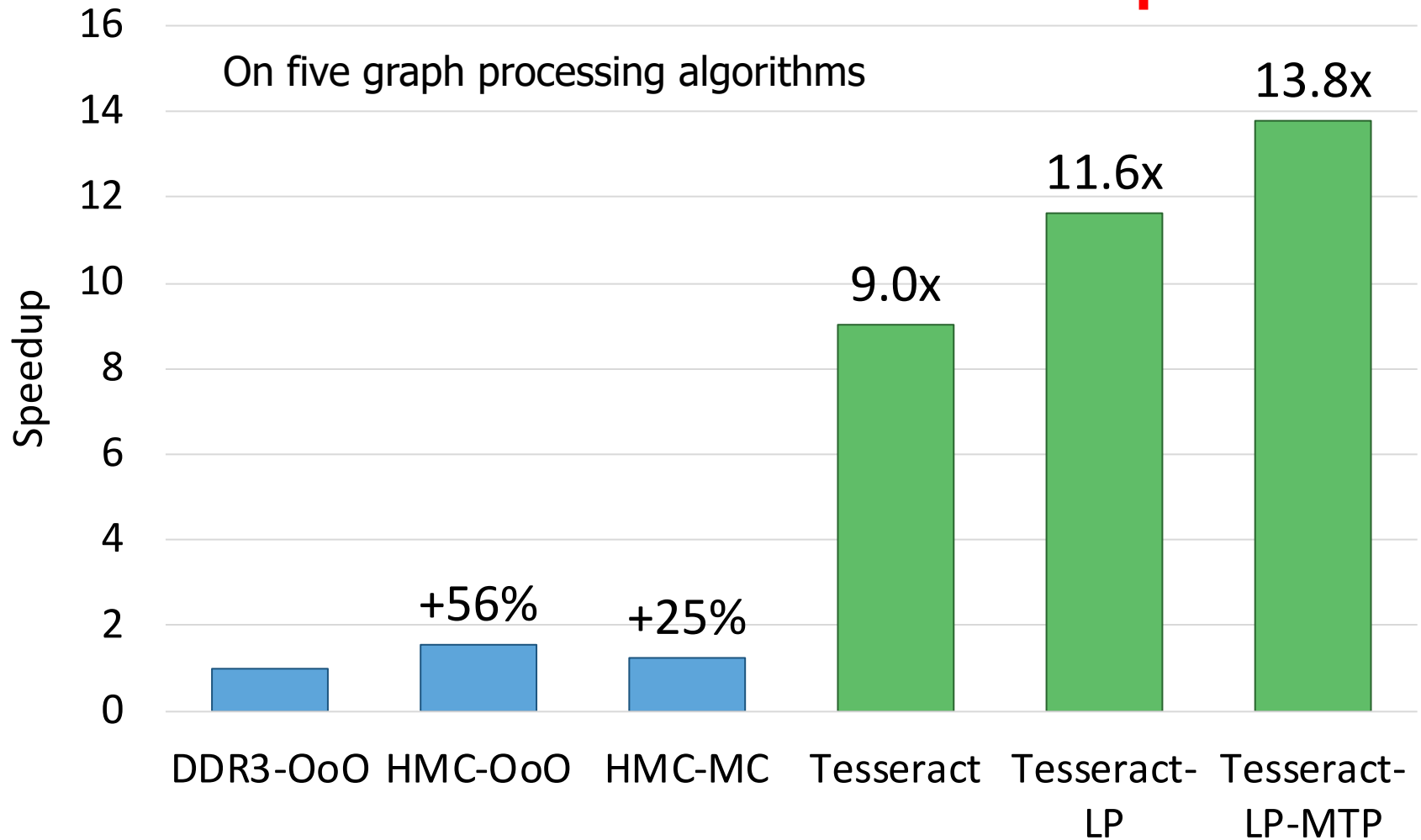
Tesseract



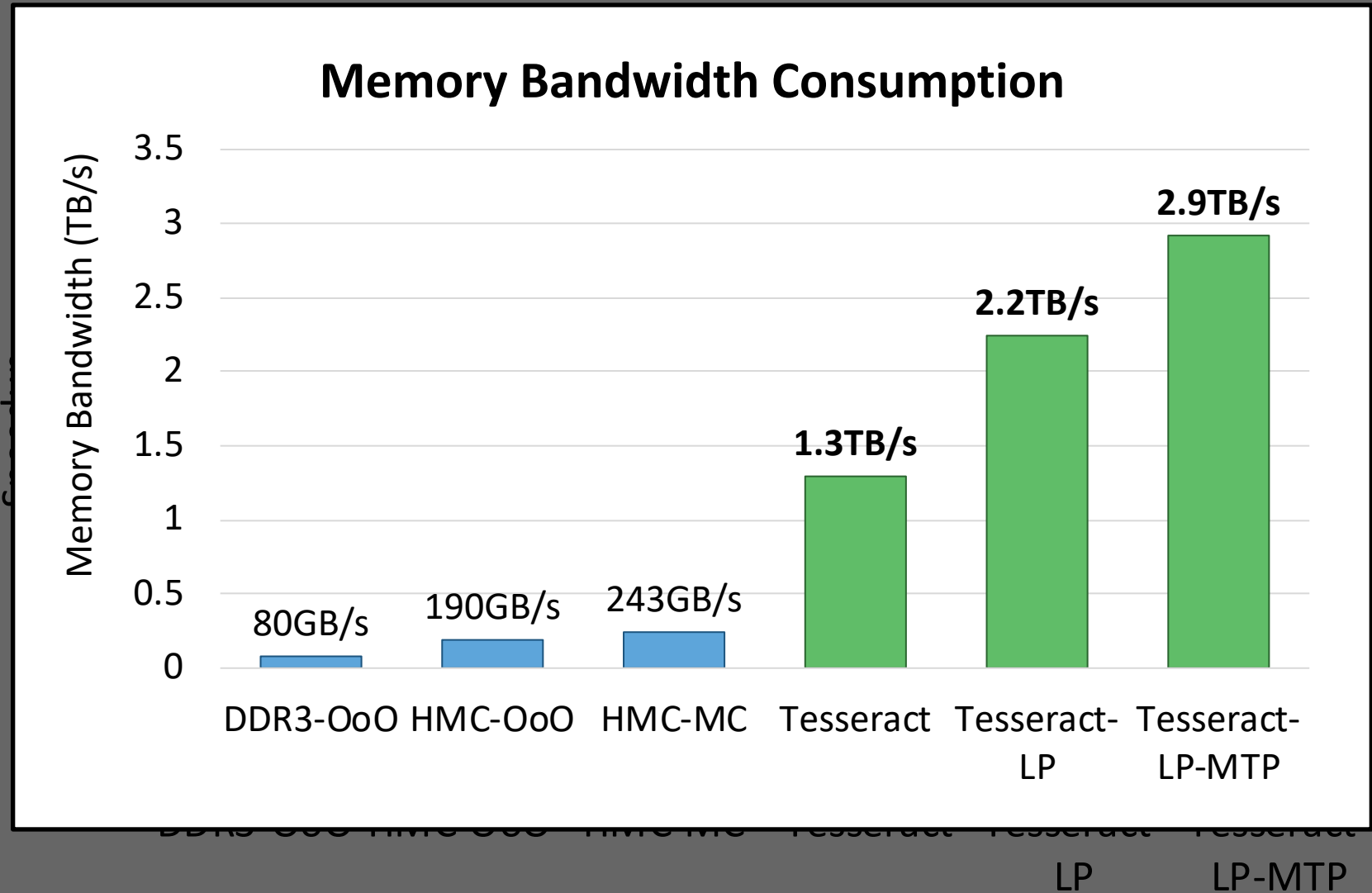
8TB/s

Tesseract Graph Processing Performance

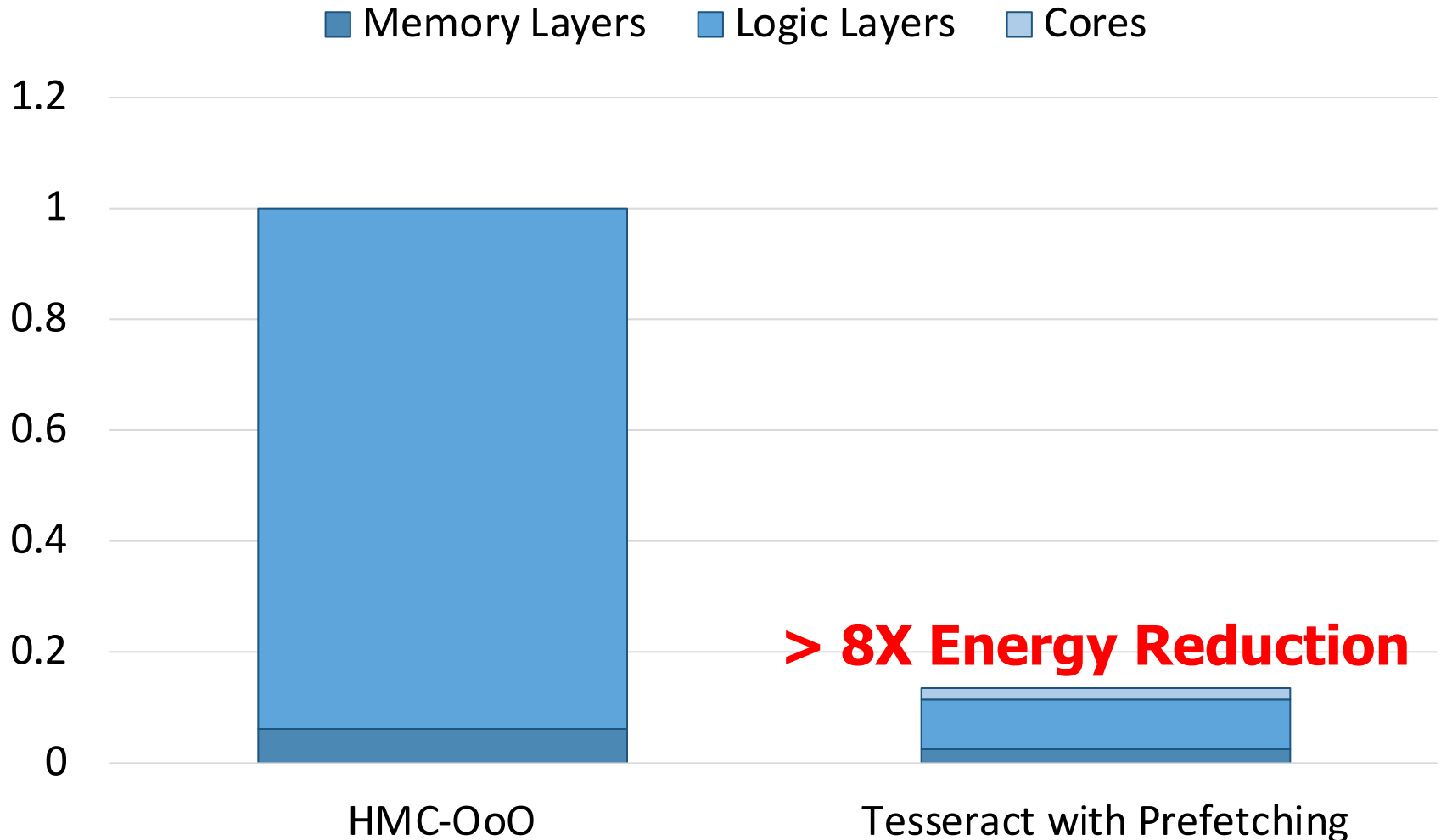
>13X Performance Improvement



Tesseract Graph Processing Performance



Tesseract Graph Processing System Energy



More on Tesseract

- Junwhan Ahn, Sungpack Hong, Sungjoo Yoo, Onur Mutlu, and Kiyoun Choi,
"A Scalable Processing-in-Memory Accelerator for Parallel Graph Processing"
Proceedings of the 42nd International Symposium on Computer Architecture (ISCA), Portland, OR, June 2015.
[\[Slides \(pdf\)\]](#) [\[Lightning Session Slides \(pdf\)\]](#)

A Scalable Processing-in-Memory Accelerator for Parallel Graph Processing

Junwhan Ahn Sungpack Hong[§] Sungjoo Yoo Onur Mutlu[†] Kiyoun Choi
junwhan@snu.ac.kr, sungpack.hong@oracle.com, sungjoo.yoo@gmail.com, onur@cmu.edu, kchoi@snu.ac.kr

Seoul National University

[§]Oracle Labs

[†]Carnegie Mellon University

Two Key Questions in 3D-Stacked PIM

- What are the performance and energy benefits of using 3D-stacked memory as a coarse-grained accelerator?
 - By changing the entire system
 - By performing simple function offloading

- What is the minimal processing-in-memory support we can provide?
 - With minimal changes to system and programming

PIM on Mobile Devices

- Amirali Boroumand, Saugata Ghose, Youngsok Kim, Rachata Ausavarungnirun, Eric Shiu, Rahul Thakur, Daehyun Kim, Aki Kuusela, Allan Knies, Parthasarathy Ranganathan, and Onur Mutlu, **"Google Workloads for Consumer Devices: Mitigating Data Movement Bottlenecks"**

Proceedings of the 23rd International Conference on Architectural Support for Programming Languages and Operating Systems (ASPLOS), Williamsburg, VA, USA, March 2018.

Google Workloads for Consumer Devices: Mitigating Data Movement Bottlenecks

Amirali Boroumand¹

Saugata Ghose¹

Youngsok Kim²

Rachata Ausavarungnirun¹

Eric Shiu³

Rahul Thakur³

Daehyun Kim^{4,3}

Aki Kuusela³

Allan Knies³

Parthasarathy Ranganathan³

Onur Mutlu^{5,1}

Consumer Devices



Consumer devices are everywhere!

**Energy consumption is
a first-class concern in consumer devices**



Four Important Workloads



Chrome

Google's web browser



TensorFlow Mobile

Google's machine learning framework

VP9



Video Playback

Google's **video codec**

VP9

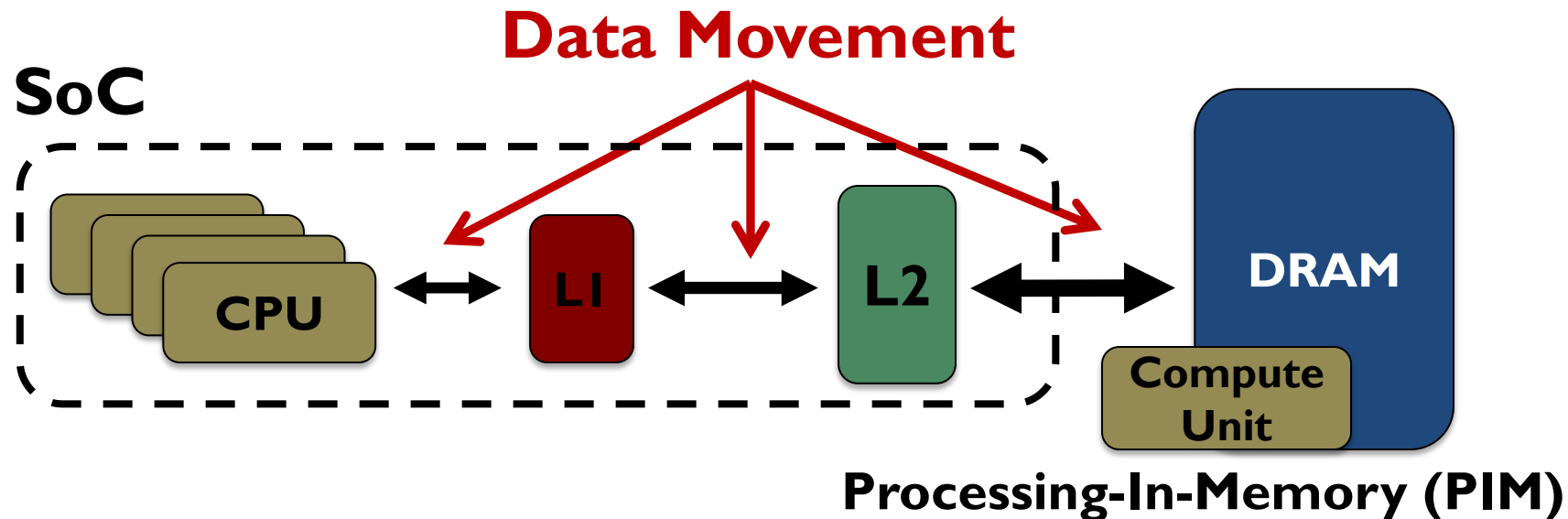


Video Capture

Google's **video codec**

Energy Cost of Data Movement

1st key observation: **62.7%** of the total system energy is spent on **data movement**



Potential solution: move computation **close to data**

Challenge: limited area and energy budget

Using PIM to Reduce Data Movement

2nd key observation: a significant fraction of the **data movement** often comes from **simple functions**

We can design lightweight logic to implement these simple functions in **memory**

Small embedded
low-power core



Small fixed-function
accelerators



Offloading to PIM logic reduces energy and improves performance, on average, by 55.4% and 54.2%

Workload Analysis



Chrome

Google's web browser



TensorFlow Mobile

Google's machine learning
framework

VP9



Video Playback

Google's **video codec**

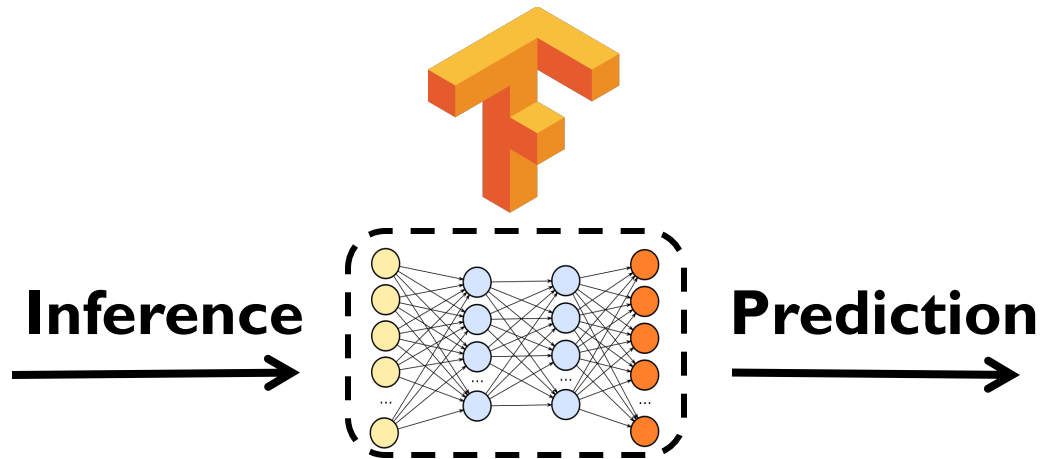
VP9



Video Capture

Google's **video codec**

TensorFlow Mobile



57.3% of the inference energy is spent on data movement



54.4% of the **data movement** energy comes from packing/unpacking and quantization

Packing



Reorders elements of matrices to minimize **cache misses** during **matrix multiplication**



Up to **40%** of the inference **energy** and **31%** of inference **execution time**



Packing's data movement accounts for up to **35.3%** of the inference **energy**

A simple **data reorganization** process that requires **simple arithmetic**

Quantization



Converts 32-bit floating point to 8-bit integers to improve inference execution time and energy consumption



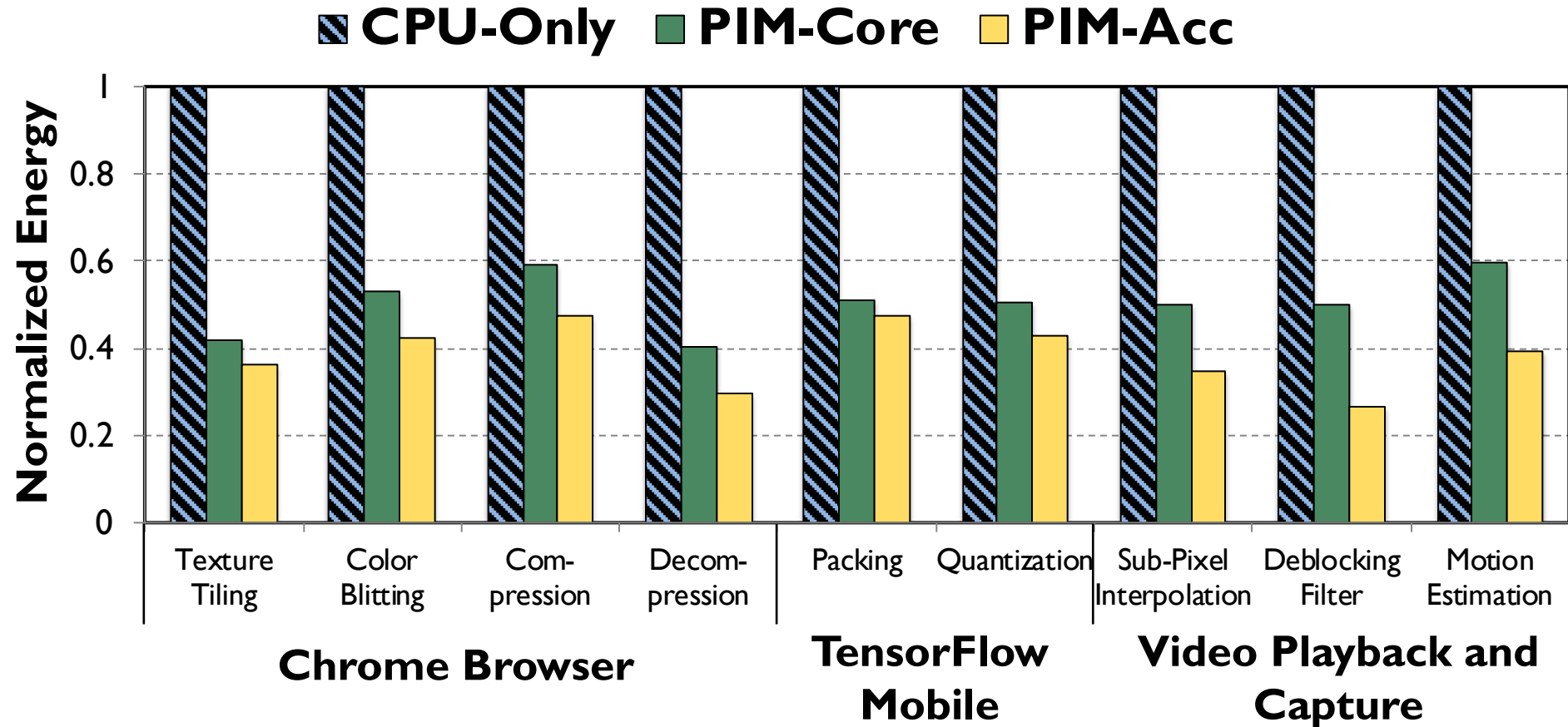
Up to **16.8%** of the inference **energy** and **16.1%** of inference **execution time**



Majority of **quantization** energy comes from **data movement**

A simple **data conversion** operation that requires **shift, addition, and multiplication** operations

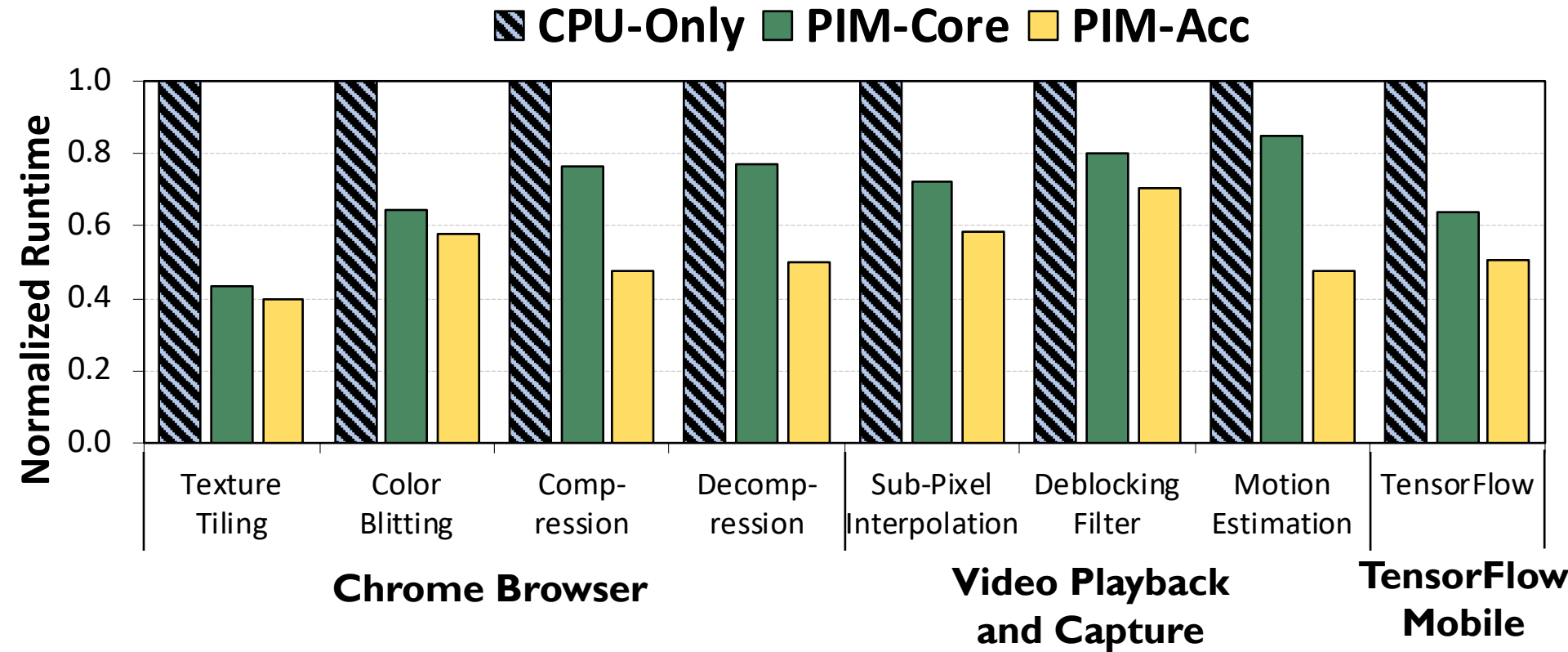
Normalized Energy



PIM core and PIM accelerator reduce
energy consumption on average by 49.1% and 55.4%

SAFARI

Normalized Runtime



Offloading these kernels to **PIM core** and **PIM accelerator** improves **performance** on average by **44.6%** and **54.2%**

More on PIM for Mobile Devices

- Amirali Boroumand, Saugata Ghose, Youngsok Kim, Rachata Ausavarungnirun, Eric Shiu, Rahul Thakur, Daehyun Kim, Aki Kuusela, Allan Knies, Parthasarathy Ranganathan, and Onur Mutlu, **"Google Workloads for Consumer Devices: Mitigating Data Movement Bottlenecks"** *Proceedings of the 23rd International Conference on Architectural Support for Programming Languages and Operating Systems (ASPLOS)*, Williamsburg, VA, USA, March 2018.

**62.7% of the total system energy
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Google Workloads for Consumer Devices: Mitigating Data Movement Bottlenecks

Amirali Boroumand¹

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Daehyun Kim^{4,3}

Aki Kuusela³

Allan Knies³

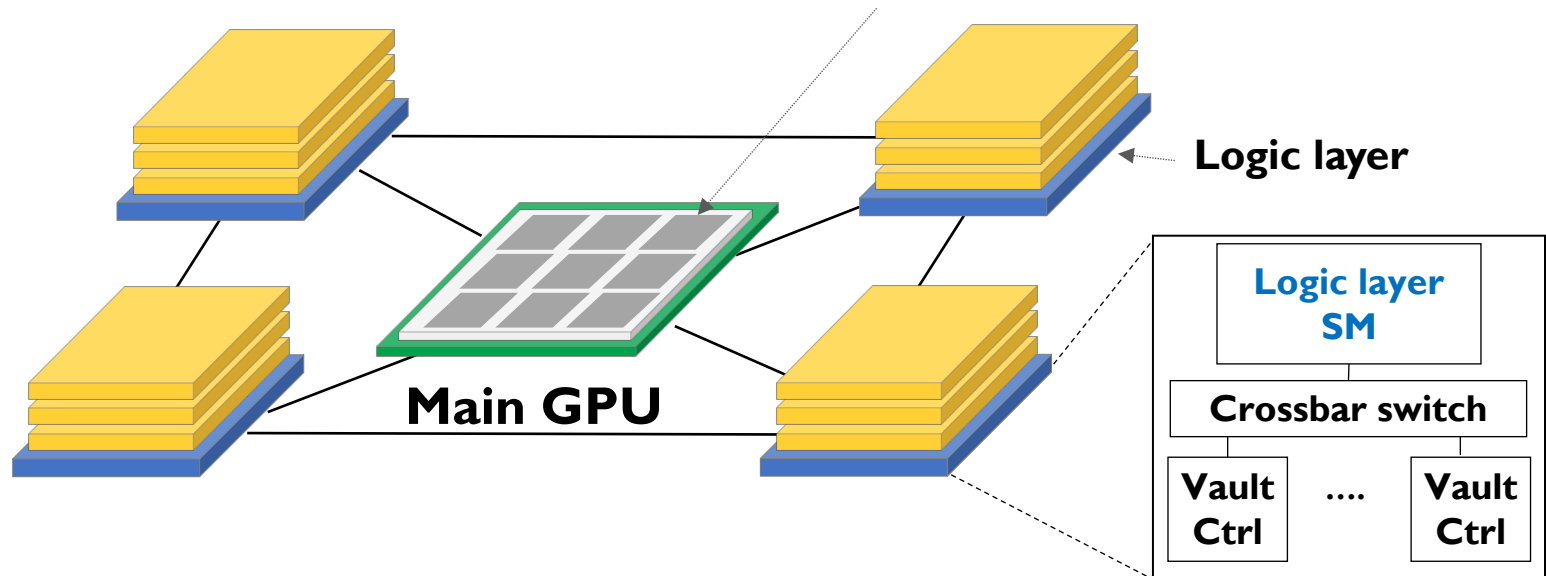
Parthasarathy Ranganathan³

Onur Mutlu^{5,1}

Truly Distributed GPU Processing with PIM?

**3D-stacked memory
(memory stack)**

SM (Streaming Multiprocessor)



```
__global__  
void applyScaleFactorsKernel( uint8_T * const out,  
    uint8_T const * const in, const double *factor,  
    size_t const numRows, size_t const numCols )  
{  
    // Work out which pixel we are working on.  
    const int rowIdx = blockIdx.x * blockDim.x + threadIdx.x;  
    const int colIdx = blockIdx.y;  
    const int sliceIdx = threadIdx.z;  
  
    // Check this thread isn't off the image  
    if( rowIdx >= numRows ) return;  
  
    // Compute the index of my element  
    size_t linearIdx = rowIdx + colIdx*numRows +  
        sliceIdx*numRows*numCols;
```


Accelerating GPU Execution with PIM (I)

- Kevin Hsieh, Eiman Ebrahimi, Gwangsun Kim, Niladrish Chatterjee, Mike O'Connor, Nandita Vijaykumar, Onur Mutlu, and Stephen W. Keckler, **"Transparent Offloading and Mapping (TOM): Enabling Programmer-Transparent Near-Data Processing in GPU Systems"**
Proceedings of the 43rd International Symposium on Computer Architecture (ISCA), Seoul, South Korea, June 2016.
[[Slides \(pptx\)](#)] [[pdf](#)]
[[Lightning Session Slides \(pptx\)](#)] [[pdf](#)]

Transparent Offloading and Mapping (TOM): Enabling Programmer-Transparent Near-Data Processing in GPU Systems

Kevin Hsieh[‡] Eiman Ebrahimi[†] Gwangsun Kim^{*} Niladrish Chatterjee[†] Mike O'Connor[†]
Nandita Vijaykumar[‡] Onur Mutlu^{§‡} Stephen W. Keckler[†]

[‡]Carnegie Mellon University [†]NVIDIA ^{*}KAIST [§]ETH Zürich

Accelerating GPU Execution with PIM (II)

- Ashutosh Pattnaik, Xulong Tang, Adwait Jog, Onur Kayiran, Asit K. Mishra, Mahmut T. Kandemir, Onur Mutlu, and Chita R. Das,
"Scheduling Techniques for GPU Architectures with Processing-In-Memory Capabilities"
Proceedings of the 25th International Conference on Parallel Architectures and Compilation Techniques (PACT), Haifa, Israel, September 2016.

Scheduling Techniques for GPU Architectures with Processing-In-Memory Capabilities

Ashutosh Pattnaik¹ Xulong Tang¹ Adwait Jog² Onur Kayiran³
Asit K. Mishra⁴ Mahmut T. Kandemir¹ Onur Mutlu^{5,6} Chita R. Das¹
¹Pennsylvania State University ²College of William and Mary
³Advanced Micro Devices, Inc. ⁴Intel Labs ⁵ETH Zürich ⁶Carnegie Mellon University

Accelerating Linked Data Structures

- Kevin Hsieh, Samira Khan, Nandita Vijaykumar, Kevin K. Chang, Amirali Boroumand, Saugata Ghose, and Onur Mutlu,
"Accelerating Pointer Chasing in 3D-Stacked Memory: Challenges, Mechanisms, Evaluation"
Proceedings of the 34th IEEE International Conference on Computer Design (ICCD), Phoenix, AZ, USA, October 2016.

Accelerating Pointer Chasing in 3D-Stacked Memory: Challenges, Mechanisms, Evaluation

Kevin Hsieh[†] Samira Khan[‡] Nandita Vijaykumar[†]
Kevin K. Chang[†] Amirali Boroumand[†] Saugata Ghose[†] Onur Mutlu^{§†}
[†]*Carnegie Mellon University* [‡]*University of Virginia* [§]*ETH Zürich*

Accelerating Dependent Cache Misses

- Milad Hashemi, Khubaib, Eiman Ebrahimi, Onur Mutlu, and Yale N. Patt, **"Accelerating Dependent Cache Misses with an Enhanced Memory Controller"**

Proceedings of the 43rd International Symposium on Computer Architecture (ISCA), Seoul, South Korea, June 2016.

[[Slides \(pptx\)](#) ([pdf](#))]

[[Lightning Session Slides \(pptx\)](#) ([pdf](#))]

Accelerating Dependent Cache Misses with an Enhanced Memory Controller

Milad Hashemi*, Khubaib[†], Eiman Ebrahimi[‡], Onur Mutlu[§], Yale N. Patt*

*The University of Texas at Austin [†]Apple [‡]NVIDIA [§]ETH Zürich & Carnegie Mellon University

Two Key Questions in 3D-Stacked PIM

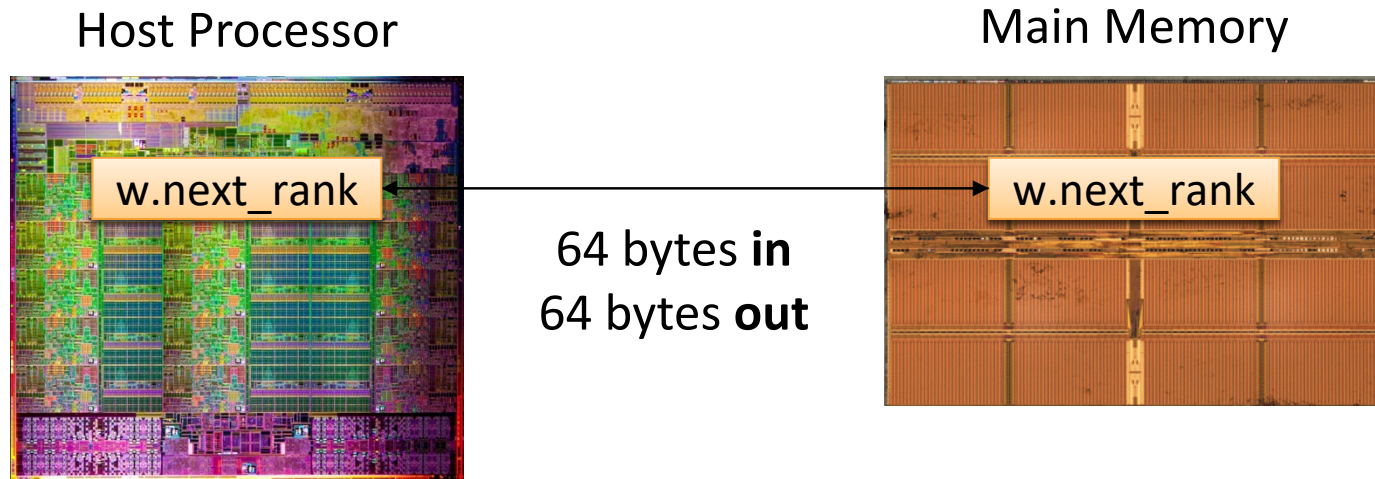
- What are the performance and energy benefits of using 3D-stacked memory as a coarse-grained accelerator?
 - By changing the entire system
 - By performing simple function offloading
- What is the minimal processing-in-memory support we can provide?
 - With minimal changes to system and programming

PEI: PIM-Enabled Instructions (Ideas)

- **Goal:** Develop mechanisms to get the most out of near-data processing with **minimal cost, minimal changes to the system, no changes to the programming model**
- **Key Idea 1:** Expose each PIM operation as a **cache-coherent, virtually-addressed host processor instruction** (called PEI) that operates on **only a single cache block**
 - e.g., `__pim_add(&w.next_rank, value) → pim.add r1, (r2)`
 - No changes sequential execution/programming model
 - No changes to virtual memory
 - Minimal changes to cache coherence
 - No need for data mapping: Each PEI restricted to a single memory module
- **Key Idea 2:** **Dynamically decide where to execute a PEI** (i.e., the host processor or PIM accelerator) based on simple locality characteristics and simple hardware predictors
 - Execute each operation at the location that provides the best performance

Simple PIM Operations as ISA Extensions (II)

```
for (v: graph.vertices) {  
    value = weight * v.rank;  
    for (w: v.successors) {  
        w.next_rank += value;  
    }  
}
```



Conventional Architecture

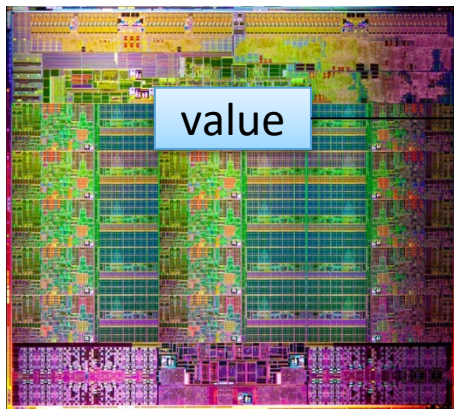
Simple PIM Operations as ISA Extensions (III)

```
for (v: graph.vertices) {  
    value = weight * v.rank;  
    for (w: v.successors) {  
        __pim_add(&w.next_rank, value);  
    }  
}
```

pim.add r1, (r2)

__pim_add(&w.next_rank, value);

Host Processor



Main Memory



8 bytes in
0 bytes out

In-Memory Addition

PEI: PIM-Enabled Instructions (Example)

```
for (v: graph.vertices) {  
    value = weight * v.rank;  
    for (w: v.successors) {
```

pim.add r1, (r2)

```
        __pim_add(&w.next_rank, value);  
    }
```

pfence

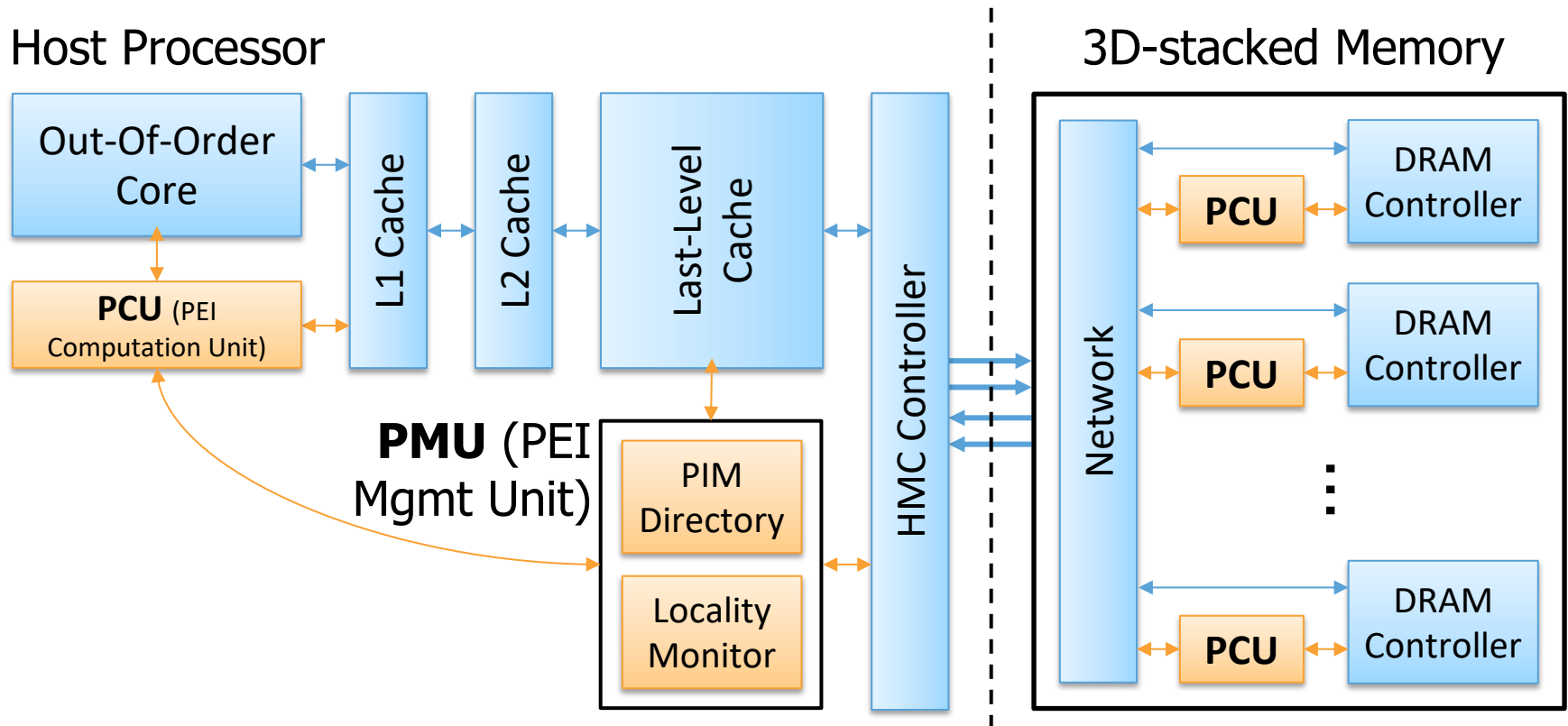
```
pfence();  
}
```

Table 1: Summary of Supported PIM Operations

Operation	R	W	Input	Output	Applications
8-byte integer increment	O	O	0 bytes	0 bytes	AT
8-byte integer min	O	O	8 bytes	0 bytes	BFS, SP, WCC
Floating-point add	O	O	8 bytes	0 bytes	PR
Hash table probing	O	X	8 bytes	9 bytes	HJ
Histogram bin index	O	X	1 byte	16 bytes	HG, RP
Euclidean distance	O	X	64 bytes	4 bytes	SC
Dot product	O	X	32 bytes	8 bytes	SVM

- Executed either in memory or in the processor: dynamic decision
 - ❑ Low-cost locality monitoring for a single instruction
- Cache-coherent, virtually-addressed, single cache block only
- Atomic between different PEIs
- *Not* atomic with normal instructions (use *pfence* for ordering)

Example (Abstract) PEI uArchitecture



Example PEI uArchitecture

PEI: Initial Evaluation Results

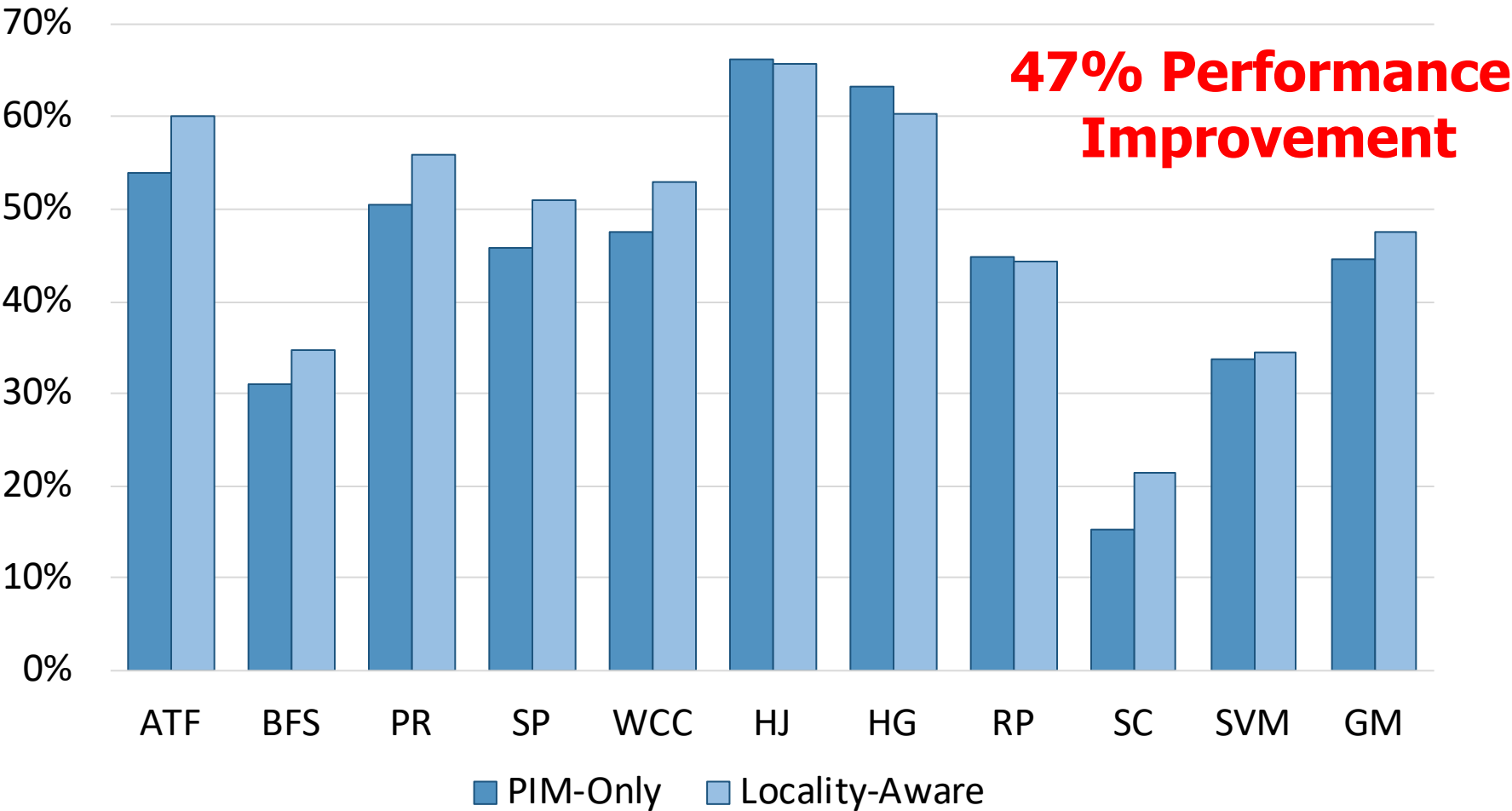
- Initial evaluations with **10 emerging data-intensive workloads**
 - ❑ Large-scale graph processing
 - ❑ In-memory data analytics
 - ❑ Machine learning and data mining
 - ❑ Three input sets (small, medium, large) for each workload to analyze the impact of data locality
- Pin-based cycle-level x86-64 simulation
- **Performance Improvement and Energy Reduction:**
 - 47% average speedup with large input data sets
 - 32% speedup with small input data sets
 - 25% avg. energy reduction in a single node with large input data sets

Table 2: Baseline Simulation Configuration

Component	Configuration
Core	16 out-of-order cores, 4 GHz, 4-issue
L1 I/D-Cache	Private, 32 KB, 4/8-way, 64 B blocks, 16 MSHRs
L2 Cache	Private, 256 KB, 8-way, 64 B blocks, 16 MSHRs
L3 Cache	Shared, 16 MB, 16-way, 64 B blocks, 64 MSHRs
On-Chip Network	Crossbar, 2 GHz, 144-bit links
Main Memory	32 GB, 8 HMCs, daisy-chain (80 GB/s full-duplex)
HMC	4 GB, 16 vaults, 256 DRAM banks [20]
– DRAM	FR-FCFS, tCL = tRCD = tRP = 13.75 ns [27]
– Vertical Links	64 TSVs per vault with 2 Gb/s signaling rate [23]

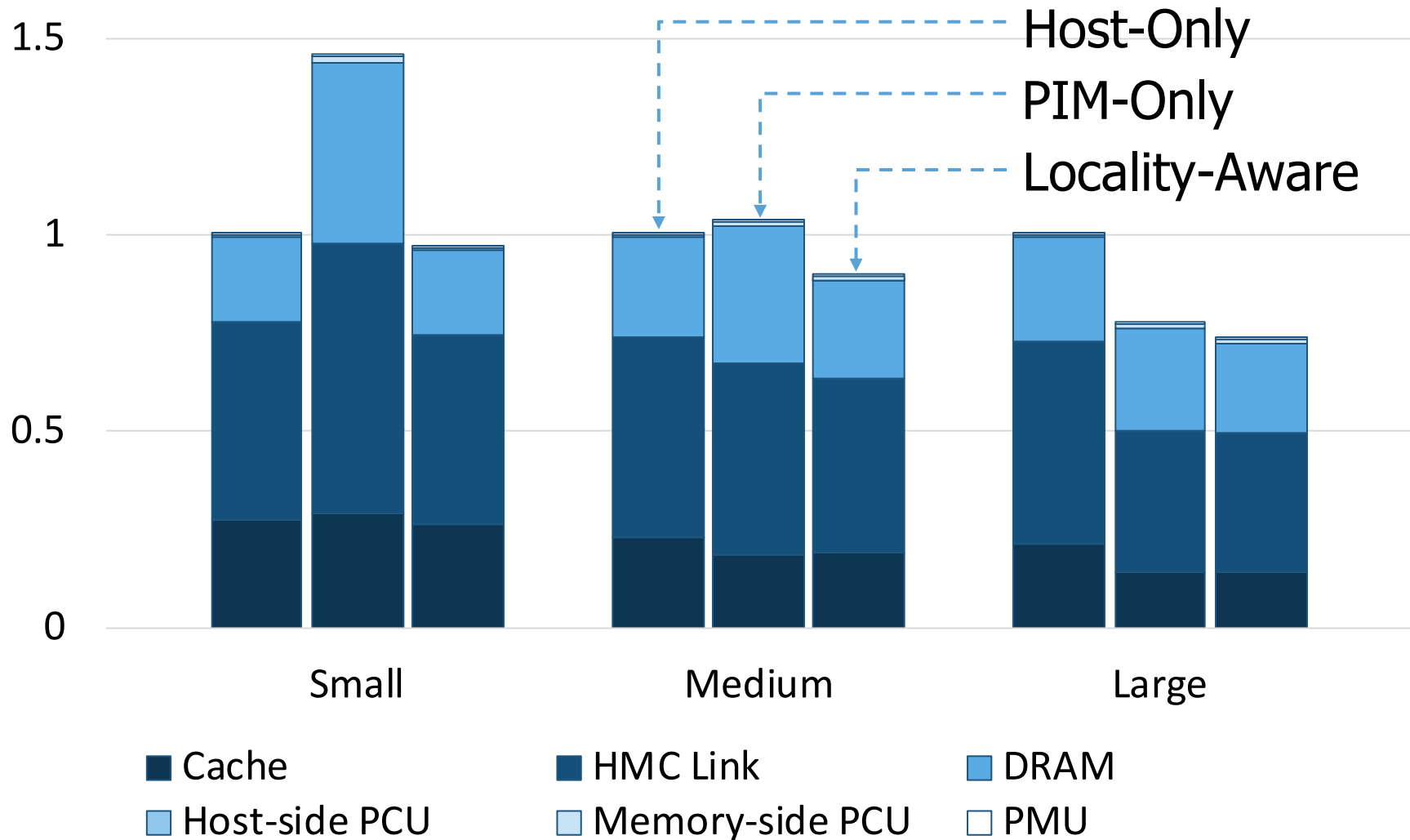
PEI Performance Delta: Large Data Sets

(Large Inputs, Baseline: Host-Only)



PEI Energy Consumption

25% Energy Reduction



Simpler PIM: PIM-Enabled Instructions

- Junwhan Ahn, Sungjoo Yoo, Onur Mutlu, and Kiyoungh Choi, **"PIM-Enabled Instructions: A Low-Overhead, Locality-Aware Processing-in-Memory Architecture"** *Proceedings of the 42nd International Symposium on Computer Architecture (ISCA)*, Portland, OR, June 2015.
[[Slides \(pdf\)](#)] [[Lightning Session Slides \(pdf\)](#)]

PIM-Enabled Instructions: A Low-Overhead, Locality-Aware Processing-in-Memory Architecture

Junwhan Ahn Sungjoo Yoo Onur Mutlu[†] Kiyoungh Choi

junwhan@snu.ac.kr, sungjoo.yoo@gmail.com, onur@cmu.edu, kchoi@snu.ac.kr

Seoul National University

[†]Carnegie Mellon University

Automatic Code and Data Mapping

- Kevin Hsieh, Eiman Ebrahimi, Gwangsun Kim, Niladrish Chatterjee, Mike O'Connor, Nandita Vijaykumar, Onur Mutlu, and Stephen W. Keckler, **"Transparent Offloading and Mapping (TOM): Enabling Programmer-Transparent Near-Data Processing in GPU Systems"**
Proceedings of the 43rd International Symposium on Computer Architecture (ISCA), Seoul, South Korea, June 2016.
[[Slides \(pptx\)](#)] [[pdf](#)]
[[Lightning Session Slides \(pptx\)](#)] [[pdf](#)]

Transparent Offloading and Mapping (TOM): Enabling Programmer-Transparent Near-Data Processing in GPU Systems

Kevin Hsieh[‡] Eiman Ebrahimi[†] Gwangsun Kim^{*} Niladrish Chatterjee[†] Mike O'Connor[†]
Nandita Vijaykumar[‡] Onur Mutlu^{§‡} Stephen W. Keckler[†]

[‡]Carnegie Mellon University [†]NVIDIA ^{*}KAIST [§]ETH Zürich

Automatic Offloading of Critical Code

- Milad Hashemi, Khubaib, Eiman Ebrahimi, Onur Mutlu, and Yale N. Patt, **"Accelerating Dependent Cache Misses with an Enhanced Memory Controller"**

Proceedings of the 43rd International Symposium on Computer Architecture (ISCA), Seoul, South Korea, June 2016.

[[Slides \(pptx\)](#) ([pdf](#))]

[[Lightning Session Slides \(pptx\)](#) ([pdf](#))]

Accelerating Dependent Cache Misses with an Enhanced Memory Controller

Milad Hashemi*, Khubaib[†], Eiman Ebrahimi[‡], Onur Mutlu[§], Yale N. Patt*

*The University of Texas at Austin [†]Apple [‡]NVIDIA [§]ETH Zürich & Carnegie Mellon University

Automatic Offloading of Prefetch Mechanisms

- Milad Hashemi, Onur Mutlu, and Yale N. Patt,
"Continuous Runahead: Transparent Hardware Acceleration for Memory Intensive Workloads"
Proceedings of the 49th International Symposium on Microarchitecture (MICRO), Taipei, Taiwan, October 2016.
[[Slides \(pptx\)](#)] [[pdf](#)] [[Lightning Session Slides \(pdf\)](#)] [[Poster \(pptx\)](#)] [[pdf](#)]

Continuous Runahead: Transparent Hardware Acceleration for Memory Intensive Workloads

Milad Hashemi*, Onur Mutlu[§], Yale N. Patt*

**The University of Texas at Austin* [§]*ETH Zürich*

Efficient Automatic Data Coherence Support

- Amirali Boroumand, Saugata Ghose, Minesh Patel, Hasan Hassan, Brandon Lucia, Kevin Hsieh, Krishna T. Malladi, Hongzhong Zheng, and Onur Mutlu,
"LazyPIM: An Efficient Cache Coherence Mechanism for Processing-in-Memory"
IEEE Computer Architecture Letters (CAL), June 2016.

LazyPIM: An Efficient Cache Coherence Mechanism for Processing-in-Memory

Amirali Boroumand[†], Saugata Ghose[†], Minesh Patel[†], Hasan Hassan^{†§}, Brandon Lucia[†],
Kevin Hsieh[†], Krishna T. Malladi^{*}, Hongzhong Zheng^{*}, and Onur Mutlu^{††}

[†]Carnegie Mellon University ^{*}Samsung Semiconductor, Inc. [§]TOBB ETÜ [‡]ETH Zürich

Efficient Automatic Data Coherence Support

- Amirali Boroumand, Saugata Ghose, Minesh Patel, Hasan Hassan, Brandon Lucia, Kevin Hsieh, Krishna T. Malladi, Hongzhong Zheng, and Onur Mutlu,

"CoNDA: Efficient Cache Coherence Support for Near-Data Accelerators"

Proceedings of the 46th International Symposium on Computer Architecture (ISCA), Phoenix, AZ, USA, June 2019.

CoNDA: Efficient Cache Coherence Support for Near-Data Accelerators

Amirali Boroumand[†]

Saugata Ghose[†]

Minesh Patel[★]

Hasan Hassan[★]

Brandon Lucia[†]

Rachata Ausavarungnirun^{†‡}

Kevin Hsieh[†]

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[‡]KMUTNB

[◇]Simon Fraser University

[§]Samsung Semiconductor, Inc.

Computing Architectures with Minimal Data Movement

Agenda

- Major Trends Affecting Main Memory
- The Need for Intelligent Memory Controllers
 - Bottom Up: Push from Circuits and Devices
 - Top Down: Pull from Systems and Applications
- Processing in Memory: Two Directions
 - Minimally Changing Memory Chips
 - Exploiting 3D-Stacked Memory
- How to Enable Adoption of Processing in Memory
- Conclusion

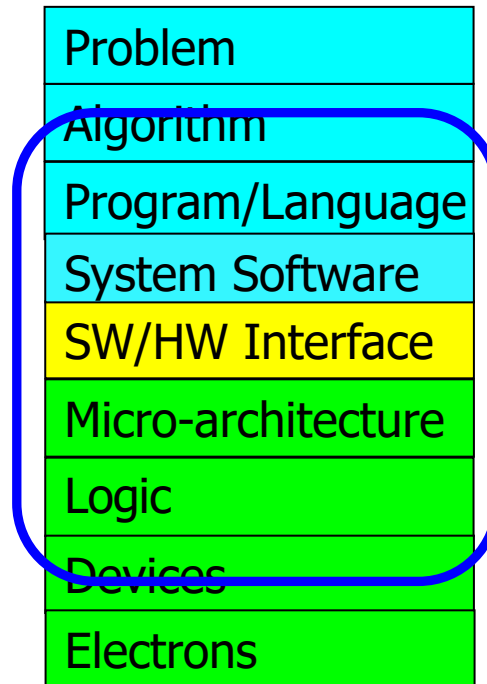
How to Enable Adoption of Processing in Memory

Barriers to Adoption of PIM

1. Functionality of and applications & software for PIM
2. Ease of programming (interfaces and compiler/HW support)
3. System support: coherence & virtual memory
4. Runtime and compilation systems for adaptive scheduling, data mapping, access/sharing control
5. Infrastructures and models to assess benefits and feasibility

All can be solved with change of mindset

We Need to Revisit the Entire Stack



We can get there step by step

PIM Review and Open Problems

Processing Data Where It Makes Sense: Enabling In-Memory Computation

Onur Mutlu^{a,b}, Saugata Ghose^b, Juan Gómez-Luna^a, Rachata Ausavarungnirun^{b,c}

^a*ETH Zürich*

^b*Carnegie Mellon University*

^c*King Mongkut's University of Technology North Bangkok*

Onur Mutlu, Saugata Ghose, Juan Gomez-Luna, and Rachata Ausavarungnirun,
**"Processing Data Where It Makes Sense: Enabling In-Memory
Computation"**

*Invited paper in Microprocessors and Microsystems (**MICPRO**), June 2019.
[arXiv version]*

PIM Review and Open Problems (II)

A Workload and Programming Ease Driven Perspective of Processing-in-Memory

Saugata Ghose[†] Amirali Boroumand[†] Jeremie S. Kim^{†§} Juan Gómez-Luna[§] Onur Mutlu^{§†}

[†]*Carnegie Mellon University*

[§]*ETH Zürich*

Saugata Ghose, Amirali Boroumand, Jeremie S. Kim, Juan Gomez-Luna, and Onur Mutlu,

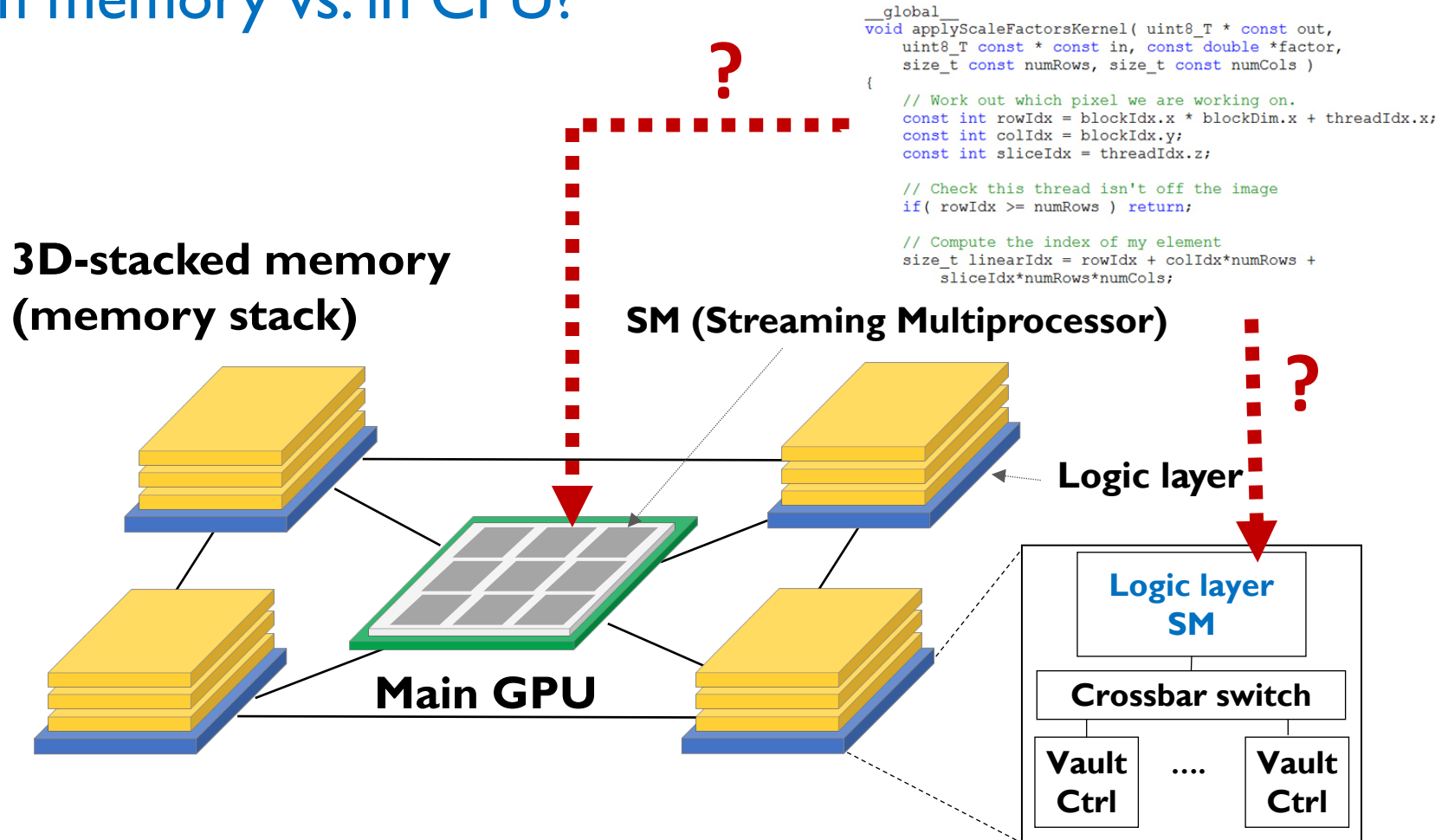
"Processing-in-Memory: A Workload-Driven Perspective"

Invited Article in IBM Journal of Research & Development, Special Issue on Hardware for Artificial Intelligence, to appear in November 2019.

[Preliminary arXiv version]

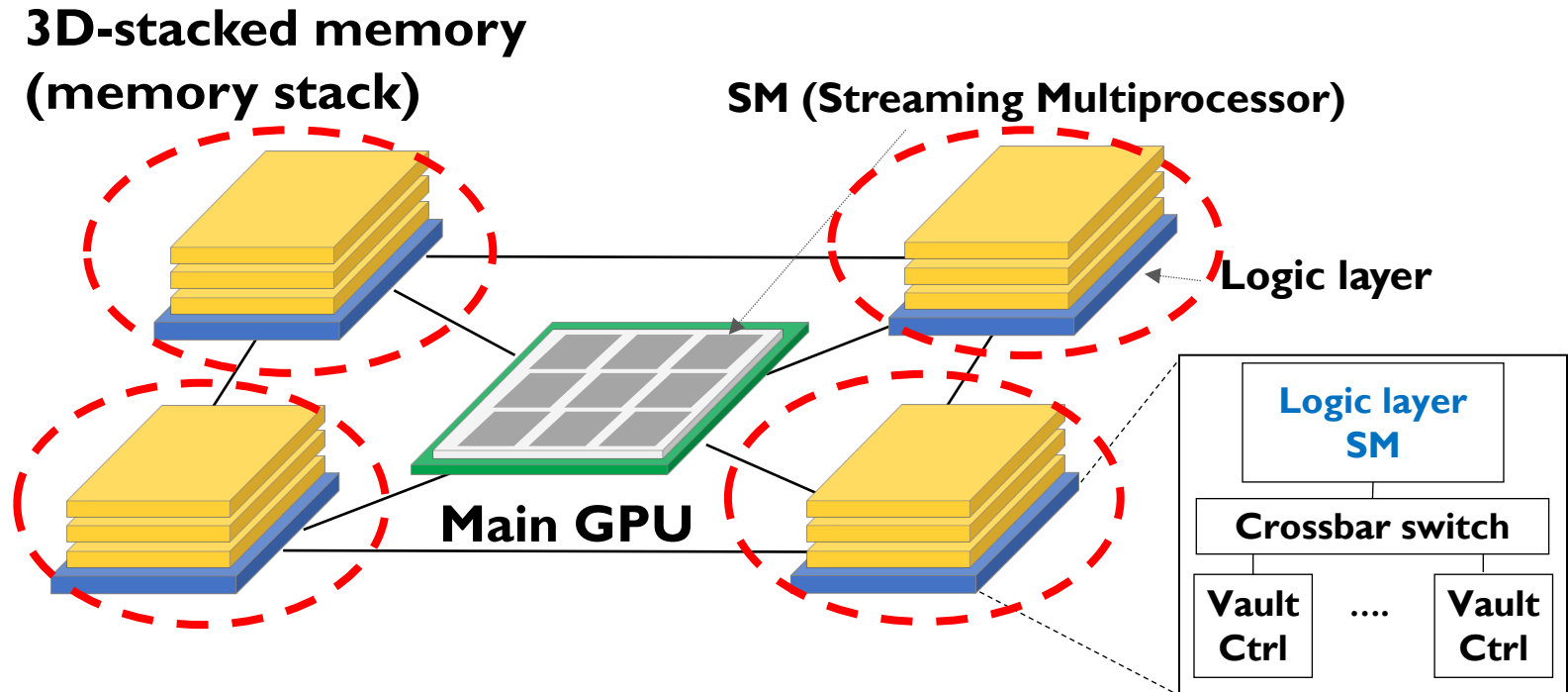
Key Challenge 1: Code Mapping

- **Challenge 1:** Which operations should be executed in memory vs. in CPU?



Key Challenge 2: Data Mapping

- **Challenge 2:** How should data be mapped to different 3D memory stacks?



How to Do the Code and Data Mapping?

- Kevin Hsieh, Eiman Ebrahimi, Gwangsun Kim, Niladrish Chatterjee, Mike O'Connor, Nandita Vijaykumar, Onur Mutlu, and Stephen W. Keckler, **"Transparent Offloading and Mapping (TOM): Enabling Programmer-Transparent Near-Data Processing in GPU Systems"**
Proceedings of the 43rd International Symposium on Computer Architecture (ISCA), Seoul, South Korea, June 2016.
[[Slides \(pptx\)](#)] [[pdf](#)]
[[Lightning Session Slides \(pptx\)](#)] [[pdf](#)]

Transparent Offloading and Mapping (TOM): Enabling Programmer-Transparent Near-Data Processing in GPU Systems

Kevin Hsieh[‡] Eiman Ebrahimi[†] Gwangsun Kim* Niladrish Chatterjee[†] Mike O'Connor[†]
Nandita Vijaykumar[‡] Onur Mutlu^{§‡} Stephen W. Keckler[†]

[‡]Carnegie Mellon University [†]NVIDIA *KAIST [§]ETH Zürich

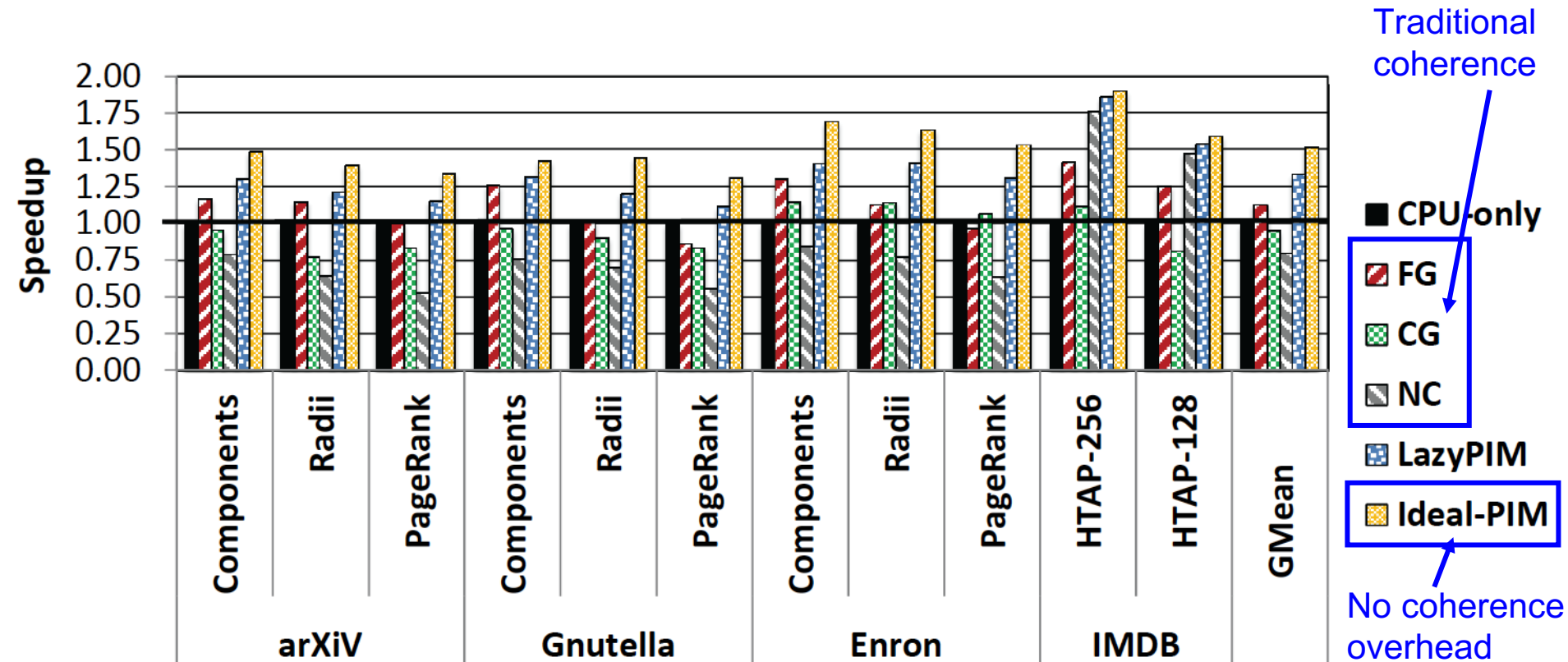
How to Schedule Code?

- Ashutosh Pattnaik, Xulong Tang, Adwait Jog, Onur Kayiran, Asit K. Mishra, Mahmut T. Kandemir, Onur Mutlu, and Chita R. Das, **"Scheduling Techniques for GPU Architectures with Processing-In-Memory Capabilities"**
Proceedings of the 25th International Conference on Parallel Architectures and Compilation Techniques (PACT), Haifa, Israel, September 2016.

Scheduling Techniques for GPU Architectures with Processing-In-Memory Capabilities

Ashutosh Pattnaik¹ Xulong Tang¹ Adwait Jog² Onur Kayiran³
Asit K. Mishra⁴ Mahmut T. Kandemir¹ Onur Mutlu^{5,6} Chita R. Das¹
¹Pennsylvania State University ²College of William and Mary
³Advanced Micro Devices, Inc. ⁴Intel Labs ⁵ETH Zürich ⁶Carnegie Mellon University

Challenge: Coherence for Hybrid CPU-PIM Apps



How to Maintain Coherence?

- Amirali Boroumand, Saugata Ghose, Minesh Patel, Hasan Hassan, Brandon Lucia, Kevin Hsieh, Krishna T. Malladi, Hongzhong Zheng, and Onur Mutlu,
"LazyPIM: An Efficient Cache Coherence Mechanism for Processing-in-Memory"
***IEEE Computer Architecture Letters* (**CAL**), June 2016.**

LazyPIM: An Efficient Cache Coherence Mechanism for Processing-in-Memory

Amirali Boroumand[†], Saugata Ghose[†], Minesh Patel[†], Hasan Hassan^{†§}, Brandon Lucia[†],
Kevin Hsieh[†], Krishna T. Malladi^{*}, Hongzhong Zheng^{*}, and Onur Mutlu^{‡†}

[†]Carnegie Mellon University ^{*}Samsung Semiconductor, Inc. [§]TOBB ETÜ [‡]ETH Zürich

How to Support Virtual Memory?

- Kevin Hsieh, Samira Khan, Nandita Vijaykumar, Kevin K. Chang, Amirali Boroumand, Saugata Ghose, and Onur Mutlu,
"Accelerating Pointer Chasing in 3D-Stacked Memory: Challenges, Mechanisms, Evaluation"
Proceedings of the 34th IEEE International Conference on Computer Design (ICCD), Phoenix, AZ, USA, October 2016.

Accelerating Pointer Chasing in 3D-Stacked Memory: Challenges, Mechanisms, Evaluation

Kevin Hsieh[†] Samira Khan[‡] Nandita Vijaykumar[†]
Kevin K. Chang[†] Amirali Boroumand[†] Saugata Ghose[†] Onur Mutlu^{§†}
[†]*Carnegie Mellon University* [‡]*University of Virginia* [§]*ETH Zürich*

How to Design Data Structures for PIM?

- Zhiyu Liu, Irina Calciu, Maurice Herlihy, and Onur Mutlu,
"Concurrent Data Structures for Near-Memory Computing"
*Proceedings of the 29th ACM Symposium on Parallelism in Algorithms
and Architectures (SPAA)*, Washington, DC, USA, July 2017.
[[Slides \(pptx\)](#)] [[pdf](#)]

Concurrent Data Structures for Near-Memory Computing

Zhiyu Liu

Computer Science Department
Brown University
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Irina Calciu

VMware Research Group
icalciu@vmware.com

Maurice Herlihy

Computer Science Department
Brown University
mph@cs.brown.edu

Onur Mutlu

Computer Science Department
ETH Zürich
onur.mutlu@inf.ethz.ch

Simulation Infrastructures for PIM

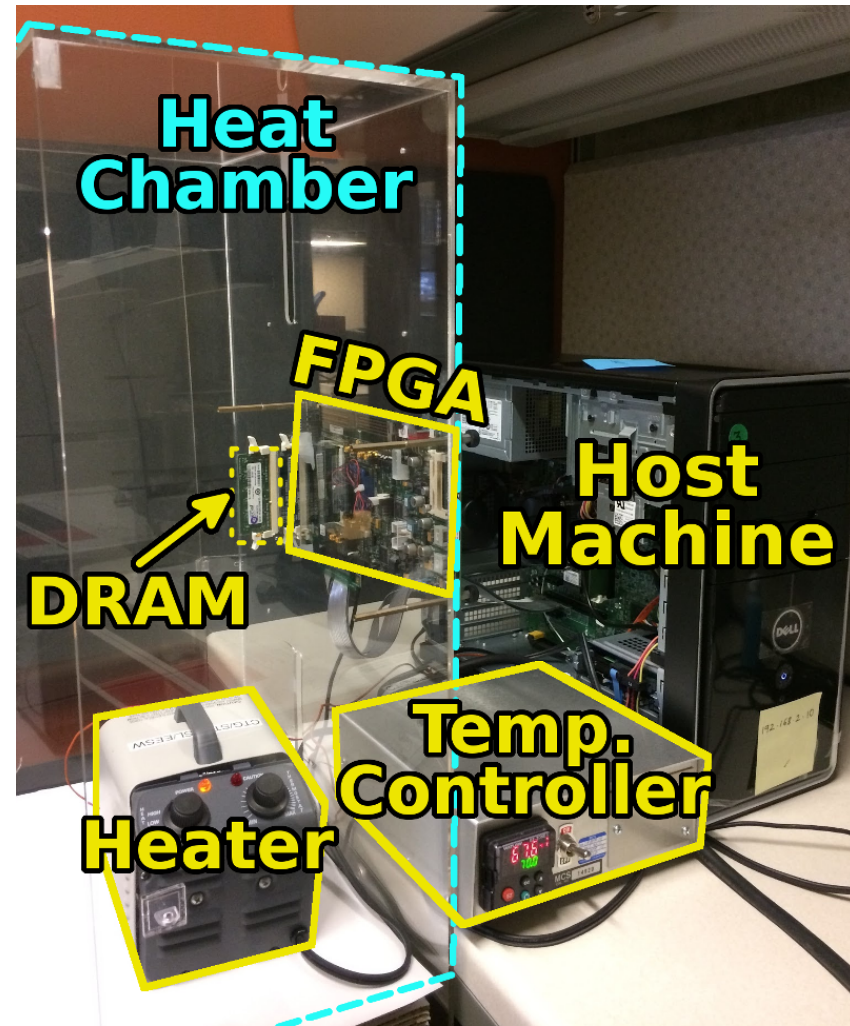
- **Ramulator** extended for PIM
 - Flexible and extensible DRAM simulator
 - Can model many different memory standards and proposals
 - Kim+, “**Ramulator: A Flexible and Extensible DRAM Simulator**”, IEEE CAL 2015.
 - <https://github.com/CMU-SAFARI/ramulator-pim>
 - <https://github.com/CMU-SAFARI/ramulator>
 - [[Source Code for Ramulator-PIM](#)]

Ramulator: A Fast and Extensible DRAM Simulator

Yoongu Kim¹ Weikun Yang^{1,2} Onur Mutlu¹
¹Carnegie Mellon University ²Peking University

An FPGA-based Test-bed for PIM?

- Hasan Hassan et al., **SoftMC: A Flexible and Practical Open-Source Infrastructure for Enabling Experimental DRAM Studies** HPCA 2017.
- Flexible
- Easy to Use (C++ API)
- Open-source
github.com/CMU-SAFARI/SoftMC



Simulation Infrastructures for PIM (in SSDs)

- Arash Tavakkol, Juan Gomez-Luna, Mohammad Sadrosadati, Saugata Ghose, and Onur Mutlu,
"MQSim: A Framework for Enabling Realistic Studies of Modern Multi-Queue SSD Devices"
Proceedings of the 16th USENIX Conference on File and Storage Technologies (FAST), Oakland, CA, USA, February 2018.
[[Slides \(pptx\)](#)] [[pdf](#)]
[[Source Code](#)]

MQSim: A Framework for Enabling Realistic Studies of Modern Multi-Queue SSD Devices

Arash Tavakkol[†], Juan Gómez-Luna[†], Mohammad Sadrosadati[†], Saugata Ghose[‡], Onur Mutlu^{†‡}
[†]*ETH Zürich* [‡]*Carnegie Mellon University*

Performance & Energy Models for PIM

- Gagandeep Singh, Juan Gomez-Luna, Giovanni Mariani, Geraldo F. Oliveira, Stefano Corda, Sander Stujik, Onur Mutlu, and Henk Corporaal, **"NAPEL: Near-Memory Computing Application Performance Prediction via Ensemble Learning"**
Proceedings of the 56th Design Automation Conference (DAC), Las Vegas, NV, USA, June 2019.
[[Slides \(pptx\)](#)] [[pdf](#)]
[[Poster \(pptx\)](#)] [[pdf](#)]
[[Source Code for Ramulator-PIM](#)]

NAPEL: Near-Memory Computing Application Performance Prediction via Ensemble Learning

Gagandeep Singh ^{a,c}	Juan Gómez-Luna ^b	Giovanni Mariani ^c	Geraldo F. Oliveira ^b
Stefano Corda ^{a,c}	Sander Stuijk ^a	Onur Mutlu ^b	Henk Corporaal ^a
^a Eindhoven University of Technology		^b ETH Zürich	^c IBM Research - Zurich

New Applications and Use Cases for PIM

- Jeremie S. Kim, Damla Senol Cali, Hongyi Xin, Donghyuk Lee, Saugata Ghose, Mohammed Alser, Hasan Hassan, Oguz Ergin, Can Alkan, and Onur Mutlu, **"GRIM-Filter: Fast Seed Location Filtering in DNA Read Mapping Using Processing-in-Memory Technologies"** ***BMC Genomics***, 2018.
Proceedings of the 16th Asia Pacific Bioinformatics Conference (APBC), Yokohama, Japan, January 2018.
[arxiv.org Version \(pdf\)](#)

GRIM-Filter: Fast seed location filtering in DNA read mapping using processing-in-memory technologies

Jeremie S. Kim^{1,6*}, Damla Senol Cali¹, Hongyi Xin², Donghyuk Lee³, Saugata Ghose¹, Mohammed Alser⁴, Hasan Hassan⁶, Oguz Ergin⁵, Can Alkan^{4*} and Onur Mutlu^{6,1*}

From The Sixteenth Asia Pacific Bioinformatics Conference 2018
Yokohama, Japan. 15-17 January 2018

Genome Read In-Memory (GRIM) Filter:

Fast Seed Location Filtering in DNA Read Mapping
using Processing-in-Memory Technologies

Jeremie Kim,

Damla Senol, Hongyi Xin, Donghyuk Lee,
Saugata Ghose, Mohammed Alser, Hasan Hassan,
Oguz Ergin, Can Alkan, and Onur Mutlu

Carnegie Mellon



ETH zürich

Executive Summary

- **Genome Read Mapping** is a very important problem and is the first step in many types of genomic analysis
 - Could lead to improved health care, medicine, quality of life
- Read mapping is an **approximate string matching** problem
 - Find the best fit of 100 character strings into a 3 billion character dictionary
 - **Alignment** is currently the best method for determining the similarity between two strings, but is **very expensive**
- We propose an in-memory processing algorithm **GRIM-Filter** for accelerating read mapping, by reducing the number of required alignments
- We implement GRIM-Filter using **in-memory processing** within **3D-stacked memory** and show up to **3.7x speedup**.

Google Workloads for Consumer Devices: Mitigating Data Movement Bottlenecks

Amirali Boroumand

Saugata Ghose, Youngsok Kim, Rachata Ausavarungnirun,
Eric Shiu, Rahul Thakur, Daehyun Kim, Aki Kuusela,
Allan Knies, Parthasarathy Ranganathan, Onur Mutlu

SAFARI

Carnegie Mellon

Google



SEOUL
NATIONAL
UNIVERSITY

ETH zürich

Open Problems: PIM Adoption

Enabling the Adoption of Processing-in-Memory: Challenges, Mechanisms, Future Research Directions

SAUGATA GHOSE, KEVIN HSIEH, AMIRALI BOROUMAND,
RACHATA AUSAVARUNGNIRUN

Carnegie Mellon University

ONUR MUTLU

ETH Zürich and Carnegie Mellon University

Saugata Ghose, Kevin Hsieh, Amirali Boroumand, Rachata Ausavarungnirun, Onur Mutlu,
**"Enabling the Adoption of Processing-in-Memory: Challenges, Mechanisms,
Future Research Directions"**

Invited Book Chapter, to appear in 2018.

[[Preliminary arxiv.org version](https://arxiv.org/pdf/1802.00320.pdf)]

PIM Review and Open Problems

Processing Data Where It Makes Sense: Enabling In-Memory Computation

Onur Mutlu^{a,b}, Saugata Ghose^b, Juan Gómez-Luna^a, Rachata Ausavarungnirun^{b,c}

^a*ETH Zürich*

^b*Carnegie Mellon University*

^c*King Mongkut's University of Technology North Bangkok*

Onur Mutlu, Saugata Ghose, Juan Gomez-Luna, and Rachata Ausavarungnirun,
**"Processing Data Where It Makes Sense: Enabling In-Memory
Computation"**

*Invited paper in Microprocessors and Microsystems (**MICPRO**), June 2019.
[arXiv version]*

PIM Review and Open Problems (II)

A Workload and Programming Ease Driven Perspective of Processing-in-Memory

Saugata Ghose[†] Amirali Boroumand[†] Jeremie S. Kim^{†§} Juan Gómez-Luna[§] Onur Mutlu^{§†}

[†]*Carnegie Mellon University*

[§]*ETH Zürich*

Saugata Ghose, Amirali Boroumand, Jeremie S. Kim, Juan Gomez-Luna, and Onur Mutlu,

"Processing-in-Memory: A Workload-Driven Perspective"

Invited Article in IBM Journal of Research & Development, Special Issue on Hardware for Artificial Intelligence, to appear in November 2019.

[Preliminary arXiv version]

Agenda

- Major Trends Affecting Main Memory
- The Need for Intelligent Memory Controllers
 - Bottom Up: Push from Circuits and Devices
 - Top Down: Pull from Systems and Applications
- Processing in Memory: Two Directions
 - Minimally Changing Memory Chips
 - Exploiting 3D-Stacked Memory
- How to Enable Adoption of Processing in Memory
- Conclusion

Fundamentally Energy-Efficient (Data-Centric) Computing Architectures

Fundamentally Low-Latency (Data-Centric) Computing Architectures

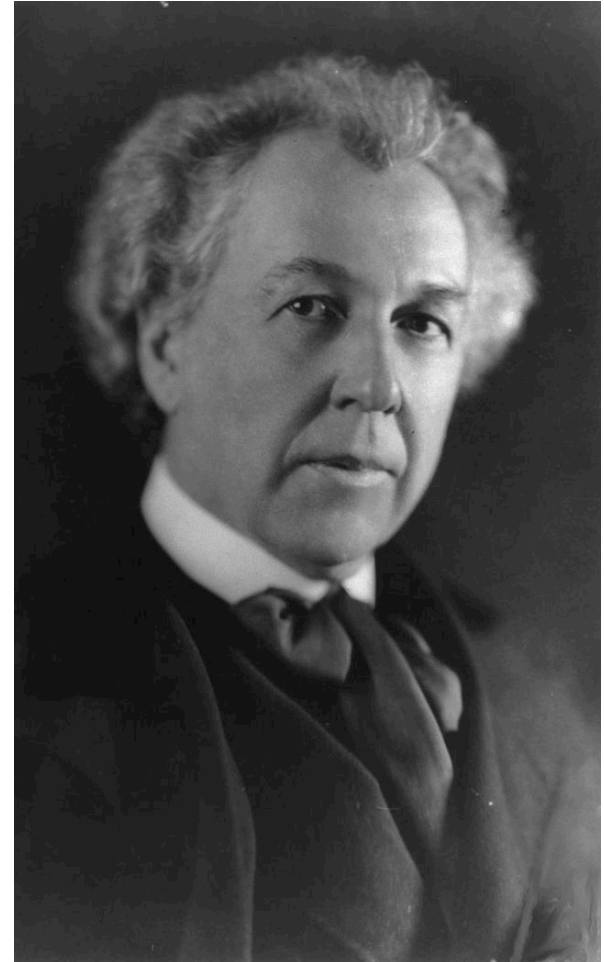
Computing Architectures with Minimal Data Movement

Main Memory Needs
Intelligent Controllers

Concluding Remarks

A Quote from A Famous Architect

- “architecture [...] based upon **principle**, and not upon **precedent**”



Precedent-Based Design?

- “architecture [...] based upon **principle**, and not upon **precedent**”



Principled Design

- “architecture [...] based upon **principle**, and not upon **precedent**”





The Overarching Principle

Organic architecture

From Wikipedia, the free encyclopedia

Organic architecture is a [philosophy](#) of [architecture](#) which promotes harmony between human habitation and the natural world through design approaches so sympathetic and well integrated with its site, that buildings, furnishings, and surroundings become part of a unified, interrelated composition.

A well-known example of organic architecture is [Fallingwater](#), the residence Frank Lloyd Wright designed for the Kaufmann family in rural Pennsylvania. Wright had many choices to locate a home on this large site, but chose to place the home directly over the waterfall and creek creating a close, yet noisy dialog with the rushing water and the steep site. The horizontal striations of stone masonry with daring [cantilevers](#) of colored beige concrete blend with native rock outcroppings and the wooded environment.

Another Example: Precedent-Based Design



Principled Design



Another Principled Design



Source: By Martín Gómez Tagle - Lisbon, Portugal, CC BY-SA 3.0, <https://commons.wikimedia.org/w/index.php?curid=13764903>

Source: <http://www.arcspace.com/exhibitions/unsorted/santiago-calatrava/>

Another Principled Design



Principle Applied to Another Structure



Source: By 準建築人手札網站 Forgemind ArchiMedia - Flickr: IMG_2489.JPG, CC BY 2.0

Source: <https://www.dezeen.com/2016/08/29/santiago-calatrava-oculus-world-trade-center-transportation-hub-new-york-photographs-hufton-crow/>
<https://commons.wikimedia.org/wiki/index.php?curid=91498396>, https://en.wikipedia.org/wiki/Santiago_Calatrava

The Overarching Principle

Zoomorphic architecture

From Wikipedia, the free encyclopedia

Zoomorphic architecture is the practice of using animal forms as the inspirational basis and blueprint for architectural design. "While animal forms have always played a role adding some of the deepest layers of meaning in architecture, it is now becoming evident that a new strand of **biomorphism** is emerging where the meaning derives not from any specific representation but from a more general allusion to biological processes."^[1]

Some well-known examples of Zoomorphic architecture can be found in the **TWA Flight Center** building in **New York City**, by **Eero Saarinen**, or the **Milwaukee Art Museum** by **Santiago Calatrava**, both inspired by the form of a bird's wings.^[3]

Overarching Principle for Computing?



Concluding Remarks

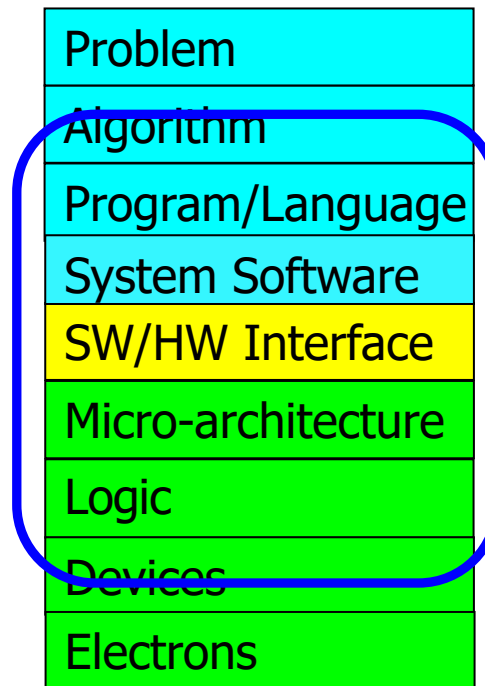
- It is time to design **principled system architectures** to solve the **memory problem**
- Design complete systems to be balanced, high-performance, and energy-efficient, i.e., **data-centric (or memory-centric)**
- Enable computation capability inside and close to memory
- **This** can
 - ❑ Lead to **orders-of-magnitude** improvements
 - ❑ **Enable new applications & computing platforms**
 - ❑ **Enable better understanding of nature**
 - ❑ ...

The Future of Processing in Memory is Bright

- Regardless of challenges
 - in underlying technology and overlying problems/requirements

Can enable:

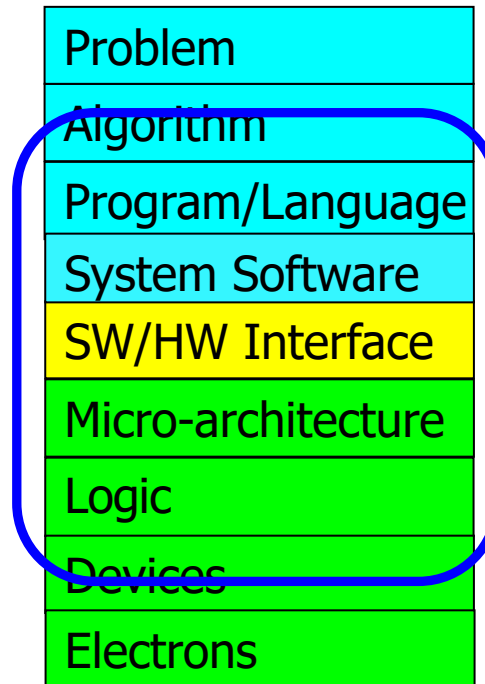
- Orders of magnitude improvements
- New applications and computing systems



Yet, we have to

- Think across the stack
- Design enabling systems

We Need to Revisit the Entire Stack



We can get there step by step

If In Doubt, See Other Doubtful Technologies

- A very “doubtful” emerging technology
 - for at least two decades



Proceedings of the IEEE, Sept. 2017

Error Characterization, Mitigation, and Recovery in Flash-Memory-Based Solid-State Drives

This paper reviews the most recent advances in solid-state drive (SSD) error characterization, mitigation, and data recovery techniques to improve both SSD's reliability and lifetime.

By YU CAI, SAUGATA GHOSE, ERICH F. HARATSCH, YIXIN LUO, AND ONUR MUTLU

PIM Review and Open Problems

Processing Data Where It Makes Sense: Enabling In-Memory Computation

Onur Mutlu^{a,b}, Saugata Ghose^b, Juan Gómez-Luna^a, Rachata Ausavarungnirun^{b,c}

^a*ETH Zürich*

^b*Carnegie Mellon University*

^c*King Mongkut's University of Technology North Bangkok*

Onur Mutlu, Saugata Ghose, Juan Gomez-Luna, and Rachata Ausavarungnirun,
**"Processing Data Where It Makes Sense: Enabling In-Memory
Computation"**

*Invited paper in Microprocessors and Microsystems (**MICPRO**), June 2019.
[arXiv version]*

PIM Review and Open Problems (II)

A Workload and Programming Ease Driven Perspective of Processing-in-Memory

Saugata Ghose[†] Amirali Boroumand[†] Jeremie S. Kim^{†§} Juan Gómez-Luna[§] Onur Mutlu^{§†}

[†]*Carnegie Mellon University*

[§]*ETH Zürich*

Saugata Ghose, Amirali Boroumand, Jeremie S. Kim, Juan Gomez-Luna, and Onur Mutlu,

"Processing-in-Memory: A Workload-Driven Perspective"

Invited Article in IBM Journal of Research & Development, Special Issue on Hardware for Artificial Intelligence, to appear in November 2019.

[Preliminary arXiv version]

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■ My collaborators

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- NIH
- GSRC
- SRC
- CyLab

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SAFARI Research Group

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Think BIG, Aim HIGH!

<https://safari.ethz.ch>

Processing Data Where It Makes Sense in Modern Computing Systems: Enabling In-Memory Computation

Onur Mutlu

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19 November 2019

ICCD Keynote

SAFARI

ETH zürich

Carnegie Mellon

Slides Not Covered
But Could Be Useful

Readings, Videos, Reference Materials

Accelerated Memory Course (~6.5 hours)

■ ACACES 2018

- ❑ Memory Systems and Memory-Centric Computing Systems
- ❑ Taught by Onur Mutlu July 9-13, 2018
- ❑ ~6.5 hours of lectures

■ Website for the Course including Videos, Slides, Papers

- ❑ https://safari.ethz.ch/memory_systems/ACACES2018/
- ❑ <https://www.youtube.com/playlist?list=PL5Q2soXY2Zi-HXxomthrpDpMJm05P6J9x>

■ All Papers are at:

- ❑ <https://people.inf.ethz.ch/omutlu/projects.htm>
- ❑ Final lecture notes and readings (for all topics)

Longer Memory Course (~18 hours)

■ Tu Wien 2019

- ❑ Memory Systems and Memory-Centric Computing Systems
- ❑ Taught by Onur Mutlu June 12-19, 2019
- ❑ ~18 hours of lectures

■ Website for the Course including Videos, Slides, Papers

- ❑ https://safari.ethz.ch/memory_systems/TUWien2019
- ❑ https://www.youtube.com/playlist?list=PL5Q2soXY2Zi_gntM55VoMIKlw7YrXOhbl

■ All Papers are at:

- ❑ <https://people.inf.ethz.ch/omutlu/projects.htm>
- ❑ Final lecture notes and readings (for all topics)

Some Overview Talks

https://www.youtube.com/watch?v=kgiZISOcGFM&list=PL5Q2soXY2Zi8D_5MGV6EnXEJHnV2YFBJI

■ Future Computing Architectures

- https://www.youtube.com/watch?v=kgiZISOcGFM&list=PL5Q2soXY2Zi8D_5MGV6EnXEJHnV2YFBJI&index=1

■ Enabling In-Memory Computation

- https://www.youtube.com/watch?v=oHqsNbxgdzM&list=PL5Q2soXY2Zi8D_5MGV6EnXEJHnV2YFBJI&index=7

■ Accelerating Genome Analysis

- https://www.youtube.com/watch?v=hPnSmfwu2-A&list=PL5Q2soXY2Zi8D_5MGV6EnXEJHnV2YFBJI&index=9

■ Rethinking Memory System Design

- https://www.youtube.com/watch?v=F7xZLNMIY1E&list=PL5Q2soXY2Zi8D_5MGV6EnXEJHnV2YFBJI&index=3

Reference Overview Paper I

Processing Data Where It Makes Sense: Enabling In-Memory Computation

Onur Mutlu^{a,b}, Saugata Ghose^b, Juan Gómez-Luna^a, Rachata Ausavarungnirun^{b,c}

^a*ETH Zürich*

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^c*King Mongkut's University of Technology North Bangkok*

Onur Mutlu, Saugata Ghose, Juan Gomez-Luna, and Rachata Ausavarungnirun,
**"Processing Data Where It Makes Sense: Enabling In-Memory
Computation"**

*Invited paper in Microprocessors and Microsystems (**MICPRO**), June 2019.
[arXiv version]*

Reference Overview Paper II

Enabling the Adoption of Processing-in-Memory: Challenges, Mechanisms, Future Research Directions

SAUGATA GHOSE, KEVIN HSIEH, AMIRALI BOROUMAND,
RACHATA AUSAVARUNGNIRUN

Carnegie Mellon University

ONUR MUTLU

ETH Zürich and Carnegie Mellon University

Saugata Ghose, Kevin Hsieh, Amirali Boroumand, Rachata Ausavarungnirun, Onur Mutlu,
**"Enabling the Adoption of Processing-in-Memory: Challenges, Mechanisms,
Future Research Directions"**

Invited Book Chapter, to appear in 2018.

[[Preliminary arxiv.org version](https://arxiv.org/pdf/1802.00320.pdf)]

Reference Overview Paper III

- Onur Mutlu and Lavanya Subramanian,
"Research Problems and Opportunities in Memory Systems"
Invited Article in Supercomputing Frontiers and Innovations (SUPERFRI), 2014/2015.

Research Problems and Opportunities in Memory Systems

Onur Mutlu¹, Lavanya Subramanian¹

Reference Overview Paper IV

- Onur Mutlu,
"The RowHammer Problem and Other Issues We May Face as Memory Becomes Denser"
*Invited Paper in Proceedings of the Design, Automation, and Test in Europe Conference (**DATE**), Lausanne, Switzerland, March 2017.*
[[Slides \(pptx\)](#) ([pdf](#))]

The RowHammer Problem and Other Issues We May Face as Memory Becomes Denser

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ETH Zürich
onur.mutlu@inf.ethz.ch
<https://people.inf.ethz.ch/omutlu>

Reference Overview Paper V

- Onur Mutlu,
"Memory Scaling: A Systems Architecture Perspective"

*Technical talk at MemCon 2013 (**MEMCON**), Santa Clara, CA, August 2013. [[Slides \(pptx\)](#)] [[pdf](#)]
[[Video](#)] [[Coverage on StorageSearch](#)]*

Memory Scaling: A Systems Architecture Perspective

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Proceedings of the IEEE, Sept. 2017

Error Characterization, Mitigation, and Recovery in Flash-Memory-Based Solid-State Drives

This paper reviews the most recent advances in solid-state drive (SSD) error characterization, mitigation, and data recovery techniques to improve both SSD's reliability and lifetime.

By YU CAI, SAUGATA GHOSE, ERICH F. HARATSCH, YIXIN LUO, AND ONUR MUTLU

Reference Overview Paper VII

- Onur Mutlu and Jeremie Kim,
"RowHammer: A Retrospective"
IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems (TCAD) Special Issue on Top Picks in Hardware and Embedded Security, 2019.
[[Preliminary arXiv version](#)]

RowHammer: A Retrospective

Onur Mutlu^{§‡} Jeremie S. Kim^{‡§}
§ETH Zürich ‡Carnegie Mellon University

Reference Overview Paper VIII

A Workload and Programming Ease Driven Perspective of Processing-in-Memory

Saugata Ghose[†] Amirali Boroumand[†] Jeremie S. Kim^{†§} Juan Gómez-Luna[§] Onur Mutlu^{§†}

[†]*Carnegie Mellon University*

[§]*ETH Zürich*

Saugata Ghose, Amirali Boroumand, Jeremie S. Kim, Juan Gomez-Luna, and Onur Mutlu,

"Processing-in-Memory: A Workload-Driven Perspective"

Invited Article in IBM Journal of Research & Development, Special Issue on Hardware for Artificial Intelligence, to appear in November 2019.

[Preliminary arXiv version]

Related Videos and Course Materials (I)

- **Undergraduate Computer Architecture Course Lecture Videos (2015, 2014, 2013)**
- **Undergraduate Computer Architecture Course Materials (2015, 2014, 2013)**

- **Graduate Computer Architecture Course Lecture Videos (2018, 2017, 2015, 2013)**
- **Graduate Computer Architecture Course Materials (2018, 2017, 2015, 2013)**

- **Parallel Computer Architecture Course Materials (Lecture Videos)**

Related Videos and Course Materials (II)

- **Freshman Digital Circuits and Computer Architecture Course Lecture Videos (2018, 2017)**
- **Freshman Digital Circuits and Computer Architecture Course Materials (2018)**
- **Memory Systems Short Course Materials (Lecture Video on Main Memory and DRAM Basics)**

Some Open Source Tools (I)

- Rowhammer – Program to Induce RowHammer Errors
 - <https://github.com/CMU-SAFARI/rowhammer>
- Ramulator – Fast and Extensible DRAM Simulator
 - <https://github.com/CMU-SAFARI/ramulator>
- MemSim – Simple Memory Simulator
 - <https://github.com/CMU-SAFARI/memsim>
- NOCulator – Flexible Network-on-Chip Simulator
 - <https://github.com/CMU-SAFARI/NOCulator>
- SoftMC – FPGA-Based DRAM Testing Infrastructure
 - <https://github.com/CMU-SAFARI/SoftMC>
- Other open-source software from my group
 - <https://github.com/CMU-SAFARI/>
 - <http://www.ece.cmu.edu/~safari/tools.html>

Some Open Source Tools (II)

- MQSim – A Fast Modern SSD Simulator
 - <https://github.com/CMU-SAFARI/MQSim>
- Mosaic – GPU Simulator Supporting Concurrent Applications
 - <https://github.com/CMU-SAFARI/Mosaic>
- IMPICA – Processing in 3D-Stacked Memory Simulator
 - <https://github.com/CMU-SAFARI/IMPICA>
- SMLA – Detailed 3D-Stacked Memory Simulator
 - <https://github.com/CMU-SAFARI/SMLA>
- HWASim – Simulator for Heterogeneous CPU-HWA Systems
 - <https://github.com/CMU-SAFARI/HWASim>
- Other open-source software from my group
 - <https://github.com/CMU-SAFARI/>
 - <http://www.ece.cmu.edu/~safari/tools.html>

More Open Source Tools (III)

- A lot more open-source software from my group
 - ❑ <https://github.com/CMU-SAFARI/>
 - ❑ <http://www.ece.cmu.edu/~safari/tools.html>



SAFARI Research Group at ETH Zurich and Carnegie Mellon University

Site for source code and tools distribution from SAFARI Research Group at ETH Zurich and Carnegie Mellon University.

📍 ETH Zurich and Carnegi... 🔗 <http://www.ece.cmu.ed...> ✉ omutlu@gmail.com

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MQSim

MQSim is a fast and accurate simulator modeling the performance of modern multi-queue (MQ) SSDs as well as traditional SATA based SSDs. MQSim faithfully models new high-bandwidth protocol implementations, steady-state SSD conditions, and the full end-to-end latency of requests in modern SSDs. It is described in detail in the FAST 2018 paper by A...

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Top languages

● C++ ● C ● C# ● AGS Script
● Verilog

Most used topics

Manage

dram reliability

Referenced Papers

- All are available at

<https://people.inf.ethz.ch/omutlu/projects.htm>

<http://scholar.google.com/citations?user=7XyGUGkAAAAJ&hl=en>

<https://people.inf.ethz.ch/omutlu/acaces2018.html>

Ramulator: A Fast and Extensible DRAM Simulator

[IEEE Comp Arch Letters'15]

Ramulator Motivation

- DRAM and Memory Controller landscape is changing
- Many new and upcoming standards
- Many new controller designs
- A fast and easy-to-extend simulator is very much needed

<i>Segment</i>	<i>DRAM Standards & Architectures</i>
Commodity	DDR3 (2007) [14]; DDR4 (2012) [18]
Low-Power	LPDDR3 (2012) [17]; LPDDR4 (2014) [20]
Graphics	GDDR5 (2009) [15]
Performance	eDRAM [28], [32]; RLDram3 (2011) [29]
3D-Stacked	WIO (2011) [16]; WIO2 (2014) [21]; MCDRAM (2015) [13]; HBM (2013) [19]; HMC1.0 (2013) [10]; HMC1.1 (2014) [11]
Academic	SBA/SSA (2010) [38]; Staged Reads (2012) [8]; RAIDR (2012) [27]; SALP (2012) [24]; TL-DRAM (2013) [26]; RowClone (2013) [37]; Half-DRAM (2014) [39]; Row-Buffer Decoupling (2014) [33]; SARP (2014) [6]; AL-DRAM (2015) [25]

Table 1. Landscape of DRAM-based memory

Ramulator

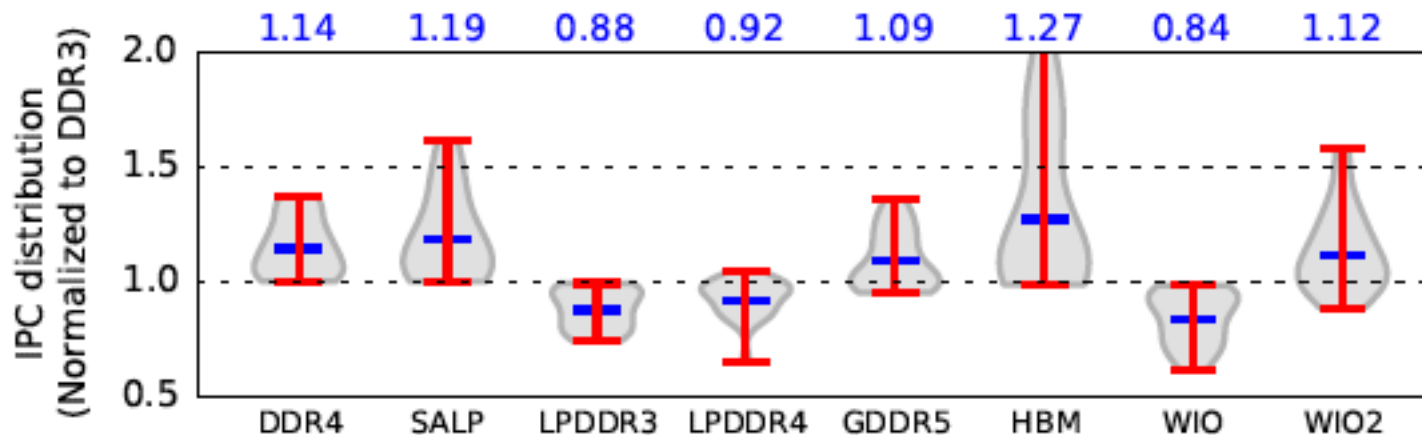
- Provides out-of-the box support for many DRAM standards:
 - ❑ DDR3/4, LPDDR3/4, GDDR5, WIO1/2, HBM, plus new proposals (SALP, AL-DRAM, TLDRAM, RowClone, and SARP)
- ~2.5X faster than fastest open-source simulator
- Modular and extensible to different standards

<i>Simulator</i> (clang -O3)	<i>Cycles (10⁶)</i>		<i>Runtime (sec.)</i>		<i>Req/sec (10³)</i>		<i>Memory</i> (MB)
	<i>Random</i>	<i>Stream</i>	<i>Random</i>	<i>Stream</i>	<i>Random</i>	<i>Stream</i>	
Ramulator	652	411	752	249	133	402	2.1
DRAMSim2	645	413	2,030	876	49	114	1.2
USIMM	661	409	1,880	750	53	133	4.5
DrSim	647	406	18,109	12,984	6	8	1.6
NVMain	666	413	6,881	5,023	15	20	4,230.0

Table 3. Comparison of five simulators using two traces

Case Study: Comparison of DRAM Standards

<i>Standard</i>	<i>Rate (MT/s)</i>	<i>Timing (CL-RCD-RP)</i>	<i>Data-Bus (Width×Chan.)</i>	<i>Rank-per-Chan</i>	<i>BW (GB/s)</i>
DDR3	1,600	11-11-11	64-bit × 1	1	11.9
DDR4	2,400	16-16-16	64-bit × 1	1	17.9
SALP [†]	1,600	11-11-11	64-bit × 1	1	11.9
LPDDR3	1,600	12-15-15	64-bit × 1	1	11.9
LPDDR4	2,400	22-22-22	32-bit × 2*	1	17.9
GDDR5 [12]	6,000	18-18-18	64-bit × 1	1	44.7
HBM	1,000	7-7-7	128-bit × 8*	1	119.2
WIO	266	7-7-7	128-bit × 4*	1	15.9
WIO2	1,066	9-10-10	128-bit × 8*	1	127.2



Across 22 workloads, simple CPU model

Figure 2. Performance comparison of DRAM standards

Ramulator Paper and Source Code

- Yoongu Kim, Weikun Yang, and Onur Mutlu,
"Ramulator: A Fast and Extensible DRAM Simulator"
IEEE Computer Architecture Letters (CAL), March 2015.
[[Source Code](#)]
- Source code is released under the liberal MIT License
 - <https://github.com/CMU-SAFARI/ramulator>

Ramulator: A Fast and Extensible DRAM Simulator

Yoongu Kim¹ Weikun Yang^{1,2} Onur Mutlu¹
¹Carnegie Mellon University ²Peking University