Securing the Memory System The Story of RowHammer

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2 October 2023

Monte Verita SW Abstractions for HW Security Workshop

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RowHammer, RowPress & Beyond Can We Be Free of Bitflips (Soon)?

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2 October 2023

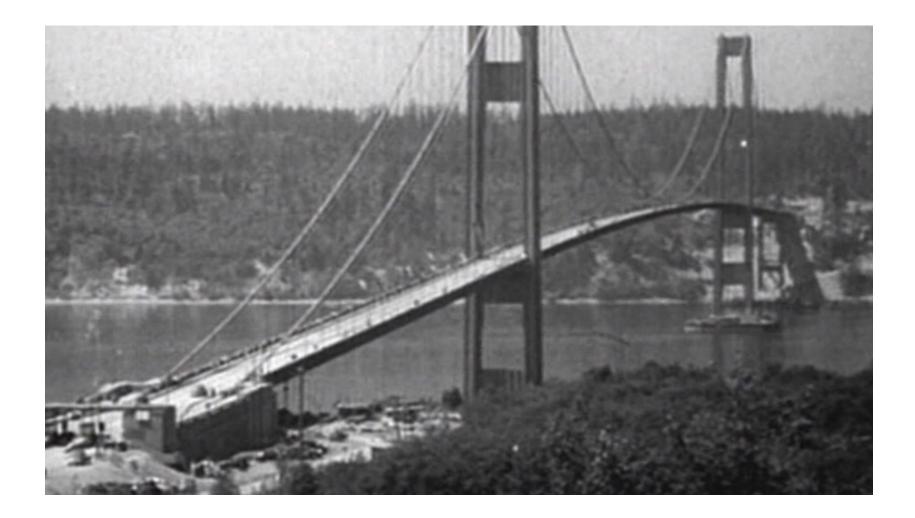
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How Reliable/Secure/Safe is This Bridge?



Collapse of the "Galloping Gertie"





How Secure Are These People?



Security is about preventing unforeseen consequences

Source: https://s-media-cache-ak0.pinimg.com/originals/48/09/54/4809543a9c7700246a0cf8acdae27abf.jpg

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How Safe & Secure Are Our Platforms?



Security is about preventing unforeseen consequences

SAFARI Source: https://taxistartup.com/wp-content/uploads/2015/03/UK-Self-Driving-Cars.jpg

What Is RowHammer?

- One can predictably induce bit flips in commodity DRAM chips
 >80% of the tested DRAM chips are vulnerable
- First example of how a simple hardware failure mechanism can create a widespread system security vulnerability



An "Early" Position Paper [IMW'13]

 Onur Mutlu,
 <u>"Memory Scaling: A Systems Architecture Perspective"</u> *Proceedings of the <u>5th International Memory</u> <i>Workshop (IMW)*, Monterey, CA, May 2013. <u>Slides</u> (pptx) (pdf)
 <u>EETimes Reprint</u>

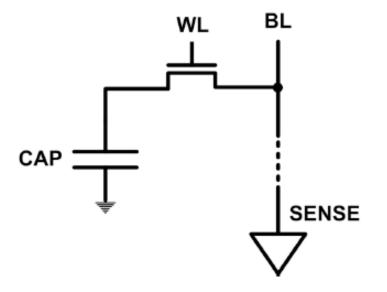
Memory Scaling: A Systems Architecture Perspective

Onur Mutlu Carnegie Mellon University onur@cmu.edu http://users.ece.cmu.edu/~omutlu/

https://people.inf.ethz.ch/omutlu/pub/memory-scaling_memcon13.pdf

The DRAM Scaling Problem

- DRAM stores charge in a capacitor (charge-based memory)
 - Capacitor must be large enough for reliable sensing
 - Access transistor should be large enough for low leakage and high retention time
 - Scaling beyond 40-35nm (2013) is challenging [ITRS, 2009]

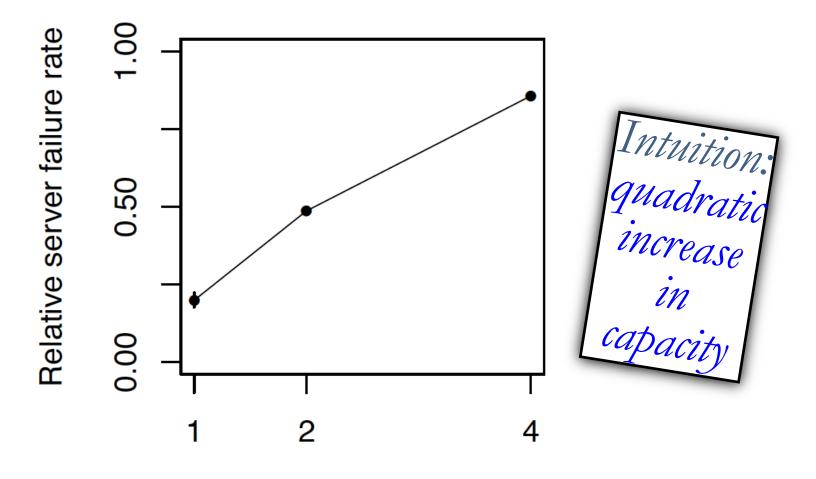


DRAM capacity, cost, and energy/power hard to scale

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As Memory Scales, It Becomes Unreliable

- Data from all of Facebook's servers worldwide
- Meza+, "Revisiting Memory Errors in Large-Scale Production Data Centers," DSN'15.



Chip density (Gb)

Large-Scale Failure Analysis of DRAM Chips

 Analysis and modeling of memory errors found in all of Facebook's server fleet

 Justin Meza, Qiang Wu, Sanjeev Kumar, and Onur Mutlu, "Revisiting Memory Errors in Large-Scale Production Data Centers: Analysis and Modeling of New Trends from the Field" Proceedings of the <u>45th Annual IEEE/IFIP International Conference on</u> Dependable Systems and Networks (DSN), Rio de Janeiro, Brazil, June 2015. [Slides (pptx) (pdf)] [DRAM Error Model]

Revisiting Memory Errors in Large-Scale Production Data Centers: Analysis and Modeling of New Trends from the Field

Justin Meza Qiang Wu* Sanjeev Kumar* Onur Mutlu

Carnegie Mellon University * Facebook, Inc.

Infrastructures to Understand Such Issues

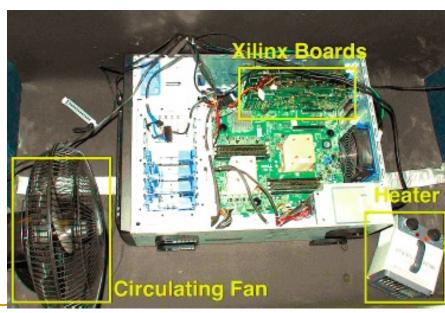


Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors (Kim et al., ISCA 2014)

Adaptive-Latency DRAM: Optimizing DRAM Timing for the Common-Case (Lee et al., HPCA 2015)

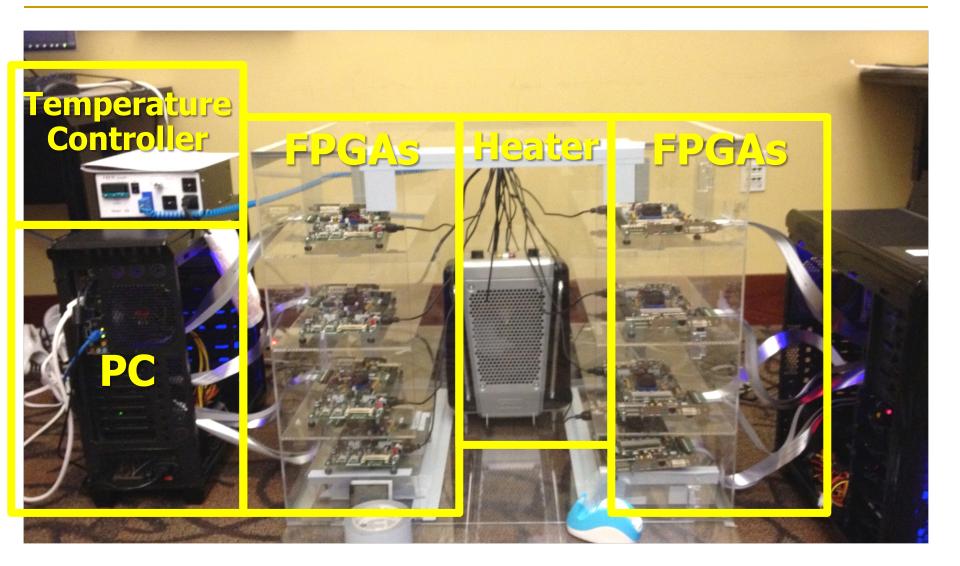
AVATAR: A Variable-Retention-Time (VRT) Aware Refresh for DRAM Systems (Qureshi et al., DSN 2015) An Experimental Study of Data Retention Behavior in Modern DRAM Devices: Implications for Retention Time Profiling Mechanisms (Liu et al., ISCA 2013)

The Efficacy of Error Mitigation Techniques for DRAM Retention Failures: A Comparative Experimental Study (Khan et al., SIGMETRICS 2014)



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Infrastructures to Understand Such Issues



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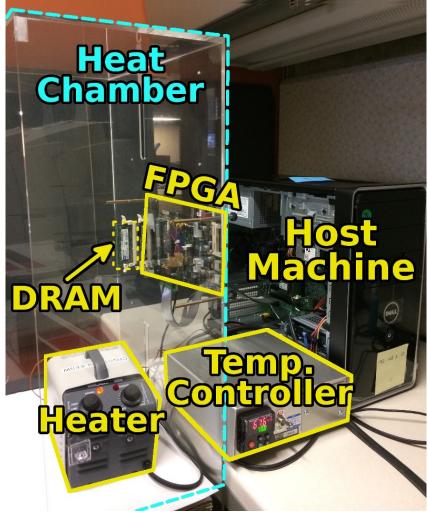
Kim+, "Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors," ISCA 2014.

SoftMC: Open Source DRAM Infrastructure

Hasan Hassan et al., "<u>SoftMC: A</u> <u>Flexible and Practical Open-</u> <u>Source Infrastructure for</u> <u>Enabling Experimental DRAM</u> <u>Studies</u>," HPCA 2017.

- Flexible
- Easy to Use (C++ API)
- Open-source

github.com/CMU-SAFARI/SoftMC



SoftMC: Open Source DRAM Infrastructure

 Hasan Hassan, Nandita Vijaykumar, Samira Khan, Saugata Ghose, Kevin Chang, Gennady Pekhimenko, Donghyuk Lee, Oguz Ergin, and Onur Mutlu, "SoftMC: A Flexible and Practical Open-Source Infrastructure for Enabling Experimental DRAM Studies" Proceedings of the 23rd International Symposium on High-Performance Computer Architecture (HPCA), Austin, TX, USA, February 2017.
 [Slides (pptx) (pdf)] [Lightning Session Slides (pptx) (pdf)] [Full Talk Lecture (39 minutes)]

SoftMC: A Flexible and Practical Open-Source Infrastructure for Enabling Experimental DRAM Studies

Hasan Hassan^{1,2,3} Nandita Vijaykumar³ Samira Khan^{4,3} Saugata Ghose³ Kevin Chang³ Gennady Pekhimenko^{5,3} Donghyuk Lee^{6,3} Oguz Ergin² Onur Mutlu^{1,3}

¹ETH Zürich ²TOBB University of Economics & Technology ³Carnegie Mellon University ⁴University of Virginia ⁵Microsoft Research ⁶NVIDIA Research

FARI https://github.com/CMU-SAFARI/SoftMC

DRAM Bender

 Ataberk Olgun, Hasan Hassan, A Giray Yağlıkçı, Yahya Can Tuğrul, Lois Orosa, Haocong Luo, Minesh Patel, Oğuz Ergin, and Onur Mutlu, "DRAM Bender: An Extensible and Versatile FPGA-based Infrastructure to Easily Test State-of-the-art DRAM Chips" *IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems (TCAD)*, 2023.
 [Extended arXiv version]
 [DRAM Bender Source Code]
 [DRAM Bender Tutorial Video (43 minutes)]

DRAM Bender: An Extensible and Versatile FPGA-based Infrastructure to Easily Test State-of-the-art DRAM Chips

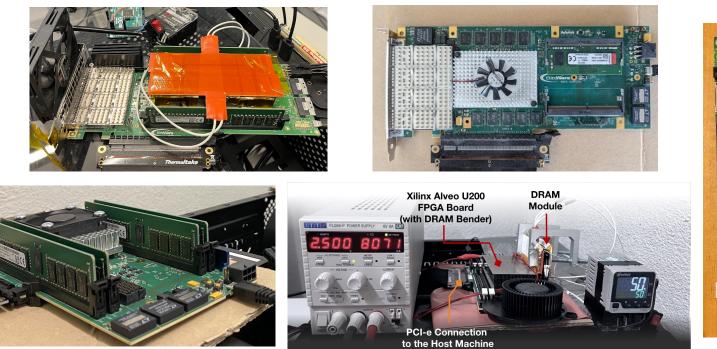
Ataberk Olgun§Hasan Hassan§A. Giray Yağlıkçı§Yahya Can Tuğrul§†Lois Orosa§⊙Haocong Luo§Minesh Patel§Oğuz Ergin†Onur Mutlu§§ETH Zürich†TOBB ETÜ⊙Galician Supercomputing Center

SAFARI <u>https://github.com/CMU-SAFARI/DRAM-Bender</u>

DRAM Bender: Prototypes

| Testing Infrastructure | Protocol Support | FPGA Support |
|-----------------------------------|------------------|------------------------|
| SoftMC [134] | DDR3 | One Prototype |
| LiteX RowHammer Tester (LRT) [17] | DDR3/4, LPDDR4 | Two Prototypes |
| DRAM Bender (this work) | DDR3/DDR4 | Five Prototypes |

Five out of the box FPGA-based prototypes



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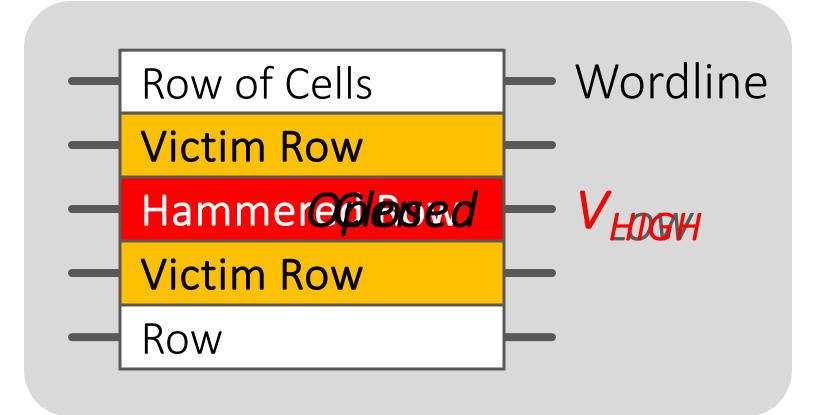


17

https://github.com/CMU-SAFARI/DRAM-Bender

A Curious Phenomenon

Modern DRAM is Prone to Disturbance Errors

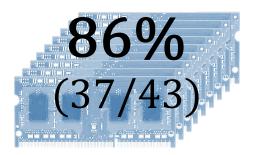


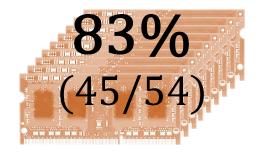
Repeatedly reading a row enough times (before memory gets refreshed) induces **disturbance errors** in **adjacent rows** in **most real DRAM chips you can buy today**

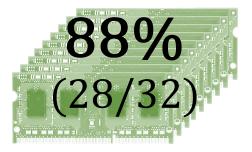
Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors, (Kim et al., ISCA 2014)

Most DRAM Modules Are Vulnerable

A company B company C company



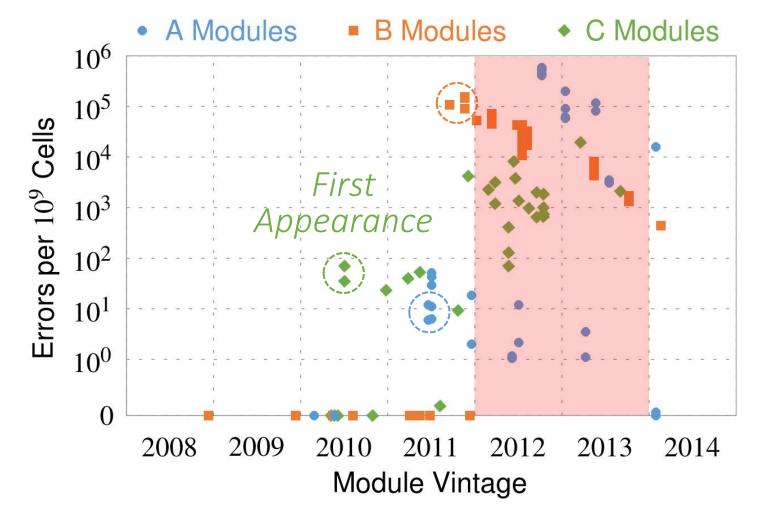




| Up to | Up to | Up to |
|---------------------|---------------------|---------------------|
| 1.0×10 ⁷ | 2.7×10 ⁶ | 3.3×10 ⁵ |
| errors | errors | errors |

<u>Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM</u> <u>Disturbance Errors</u>, (Kim et al., ISCA 2014)

Recent DRAM Is More Vulnerable



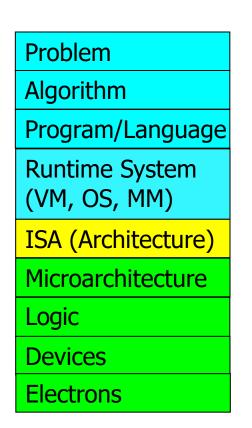
All modules from 2012–2013 are vulnerable

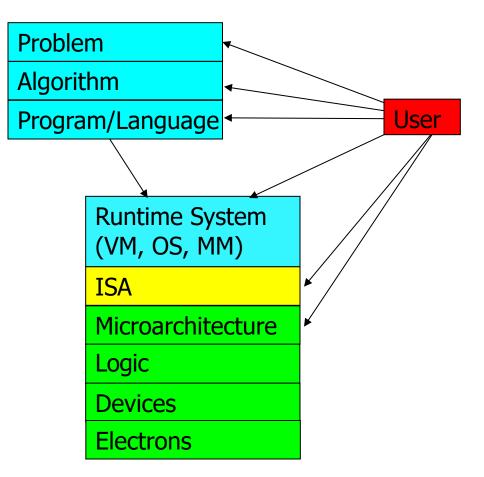
Why Is This Happening?

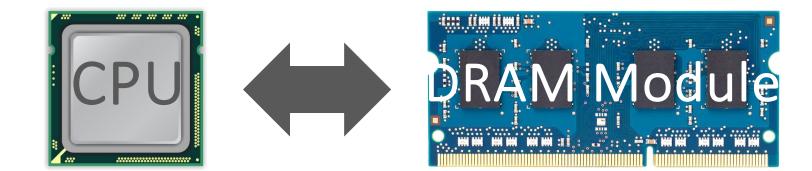
- DRAM cells are too close to each other!
 - They are not electrically isolated from each other
- Access to one cell affects the value in nearby cells
 - due to electrical interference between
 - the cells
 - wires used for accessing the cells
 - Also called cell-to-cell coupling/interference
- Example: When we activate (apply high voltage) to a row, an adjacent row gets slightly activated as well
 - Vulnerable cells in that slightly-activated row lose a little bit of charge
 - □ If RowHammer happens enough times, charge in such cells gets drained

Higher-Level Implications

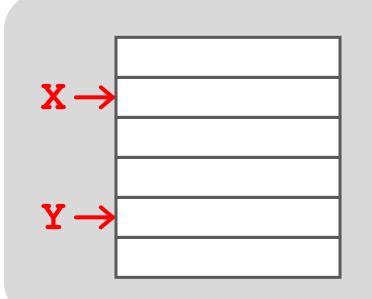
This simple circuit level failure mechanism has enormous implications on upper layers of the transformation hierarchy

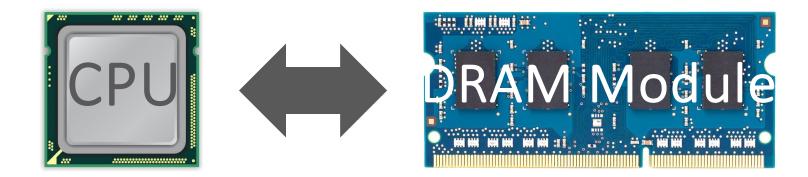




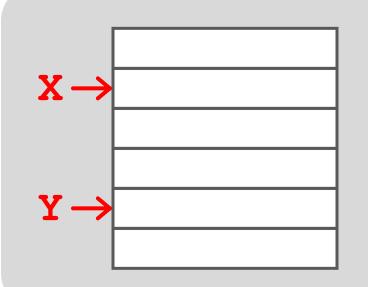


loop: mov (X), %eax mov (Y), %ebx clflush (X) clflush (Y) mfence jmp loop

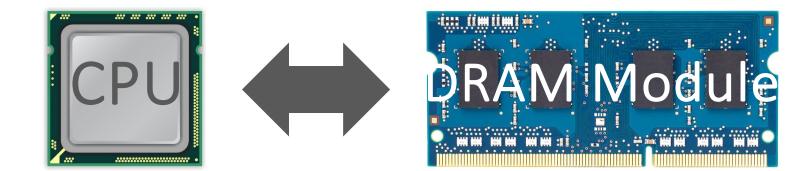




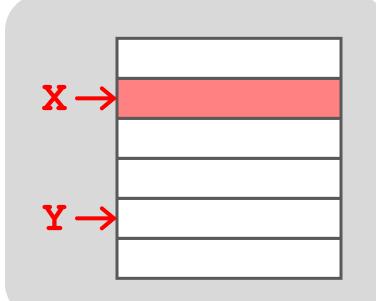
- Avoid *cache hits* Flush X from cache
- Avoid *row hits* to X
 Read Y in another row

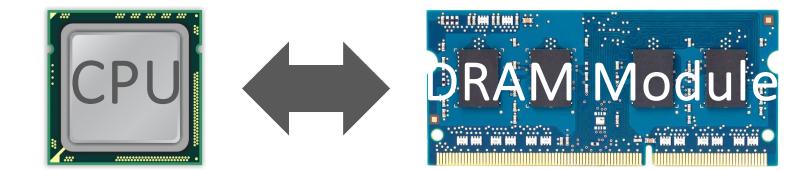


Download from: https://github.com/CMU-SAFARI/rowhammer

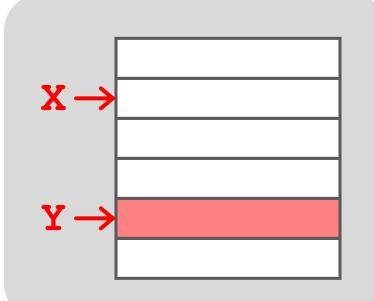


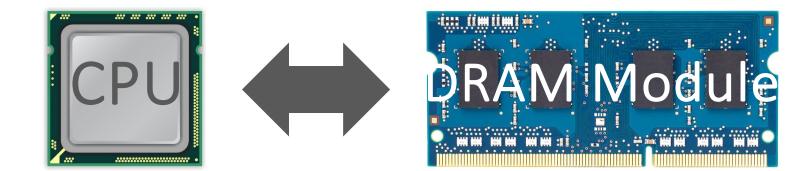
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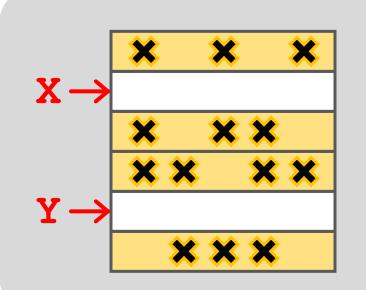


loop: mov (X), %eax mov (Y), %ebx clflush (X) clflush (Y) mfence jmp loop





loop: mov (X), %eax mov (Y), %ebx clflush (X) clflush (Y) mfence jmp loop



Observed Errors in Real Systems

| CPU Architecture | Errors | Access-Rate |
|---------------------------|----------------|-------------|
| Intel Haswell (2013) | 22.9K | 12.3M/sec |
| Intel Ivy Bridge (2012) | 20.7K | 11.7M/sec |
| Intel Sandy Bridge (2011) | 16 . 1K | 11.6M/sec |
| AMD Piledriver (2012) | 59 | 6.1M/sec |

A real reliability & security issue

Kim+, "Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors," ISCA 2014.

One Can Take Over an Otherwise-Secure System

Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors

Abstract. Memory isolation is a key property of a reliable and secure computing system — an access to one memory address should not have unintended side effects on data stored in other addresses. However, as DRAM process technology

Project Zero

<u>Flipping Bits in Memory Without Accessing Them:</u> <u>An Experimental Study of DRAM Disturbance Errors</u> (Kim et al., ISCA 2014)

News and updates from the Project Zero team at Google

Exploiting the DRAM rowhammer bug to gain kernel privileges (Seaborn, 2015)

Monday, March 9, 2015

Exploiting the DRAM rowhammer bug to gain kernel privileges

RowHammer Security Attack Example

- "Rowhammer" is a problem with some recent DRAM devices in which repeatedly accessing a row of memory can cause bit flips in adjacent rows (Kim et al., ISCA 2014).
 - Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors (Kim et al., ISCA 2014)
- We tested a selection of laptops and found that a subset of them exhibited the problem.
- We built two working privilege escalation exploits that use this effect.
 - Exploiting the DRAM rowhammer bug to gain kernel privileges (Seaborn+, 2015)
- One exploit uses rowhammer-induced bit flips to gain kernel privileges on x86-64 Linux when run as an unprivileged userland process.
- When run on a machine vulnerable to the rowhammer problem, the process was able to induce bit flips in page table entries (PTEs).
- It was able to use this to gain write access to its own page table, and hence gain read-write access to all of physical memory.

Exploiting the DRAM rowhammer bug to gain kernel privileges (Seaborn & Dullien, 2015)

Security Implications



Security Implications



It's like breaking into an apartment by repeatedly slamming a neighbor's door until the vibrations open the door you were after

More Security Implications (I)

"We can gain unrestricted access to systems of website visitors."

Not there yet, but ...



ROOT privileges for web apps!

Daniel Gruss (@lavados), Clémentine Maurice (@BloodyTangerine), December 28, 2015 - 32c3, Hamburg, Germany





Rowhammer.js: A Remote Software-Induced Fault Attack in JavaScript (DIMVA'16)

29

More Security Implications (II)

"Can gain control of a smart phone deterministically"

Hammer And Root

androids Millions of Androids

Drammer: Deterministic Rowhammer

Attacks on Mobile Platforms, CCS'16 35

Source: https://fossbytes.com/drammer-rowhammer-attack-android-root-devices/

More Security Implications (III)

Using an integrated GPU in a mobile system to remotely escalate privilege via the WebGL interface. IEEE S&P 2018

ars technica

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"GRAND PWNING UNIT" --

Drive-by Rowhammer attack uses GPU to compromise an Android phone

JavaScript based GLitch pwns browsers by flipping bits inside memory chips.

DAN GOODIN - 5/3/2018, 12:00 PM

Grand Pwning Unit: Accelerating Microarchitectural Attacks with the GPU

Pietro Frigo Vrije Universiteit Amsterdam p.frigo@vu.nl Cristiano Giuffrida Vrije Universiteit Amsterdam giuffrida@cs.vu.nl Herbert Bos Vrije Universiteit Amsterdam herbertb@cs.vu.nl Kaveh Razavi Vrije Universiteit Amsterdam kaveh@cs.vu.nl

More Security Implications (IV)

Rowhammer over RDMA (I) USENIX ATC 2018

ars TECHNICA

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THROWHAMMER —

Packets over a LAN are all it takes to trigger serious Rowhammer bit flips

The bar for exploiting potentially serious DDR weakness keeps getting lower.

DAN GOODIN - 5/10/2018, 5:26 PM

Throwhammer: Rowhammer Attacks over the Network and Defenses

Andrei Tatar VU Amsterdam Radhesh Krishnan VU Amsterdam Elias Athanasopoulos University of Cyprus

Herbert Bos VU Amsterdam Kaveh Razavi VU Amsterdam Cristiano Giuffrida VU Amsterdam

More Security Implications (V)

Rowhammer over RDMA (II)

Security in a serious way

Nethammer—Exploiting DRAM Rowhammer Bug Through Network Requests



Nethammer: Inducing Rowhammer Faults through Network Requests

Moritz Lipp Graz University of Technology

Daniel Gruss Graz University of Technology Misiker Tadesse Aga University of Michigan

Clémentine Maurice Univ Rennes, CNRS, IRISA

Lukas Lamster Graz University of Technology Michael Schwarz Graz University of Technology

Lukas Raab Graz University of Technology

More Security Implications (VI)

IEEE S&P 2020

RAMBleed: Reading Bits in Memory Without Accessing Them

Andrew Kwong University of Michigan ankwong@umich.edu Daniel Genkin University of Michigan genkin@umich.edu Daniel Gruss Graz University of Technology daniel.gruss@iaik.tugraz.at Yuval Yarom University of Adelaide and Data61 yval@cs.adelaide.edu.au

More Security Implications (VII)

USENIX Security 2019

Terminal Brain Damage: Exposing the Graceless Degradation in Deep Neural Networks Under Hardware Fault Attacks

Sanghyun Hong, Pietro Frigo[†], Yiğitcan Kaya, Cristiano Giuffrida[†], Tudor Dumitraş

University of Maryland, College Park [†]Vrije Universiteit Amsterdam



A Single Bit-flip Can Cause Terminal Brain Damage to DNNs One specific bit-flip in a DNN's representation leads to accuracy drop over 90%

Our research found that a specific bit-flip in a DNN's bitwise representation can cause the accuracy loss up to 90%, and the DNN has 40-50% parameters, on average, that can lead to the accuracy drop over 10% when individually subjected to such single bitwise corruptions...

Read More

More Security Implications (VIII)

USENIX Security 2020

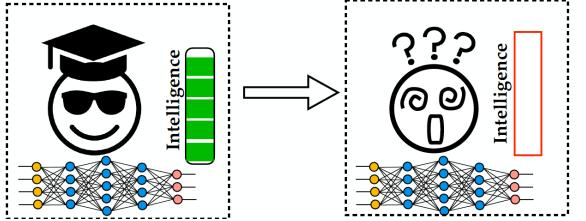
DeepHammer: Depleting the Intelligence of Deep Neural Networks through Targeted Chain of Bit Flips

Adnan Siraj RakinDeliang FanArizona State Universityasrakin@asu.edudfan@asu.edu

Degrade the **inference accuracy** to the level of **Random Guess**

Example: ResNet-20 for CIFAR-10, **10** output classes

Before attack, Accuracy: 90.2% After attack, Accuracy: ~10% (1/10)



More Security Implications?



A RowHammer Survey Across the Stack

Onur Mutlu and Jeremie Kim,
 "RowHammer: A Retrospective"
 IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems (TCAD) Special Issue on Top Picks in Hardware and Embedded Security, 2019.
 [Preliminary arXiv version]
 [Slides from COSADE 2019 (pptx)]
 [Slides from VLSI-SOC 2020 (pptx) (pdf)]
 [Talk Video (1 hr 15 minutes, with Q&A)]

RowHammer: A Retrospective

Onur Mutlu§‡Jeremie S. Kim‡§§ETH Zürich‡Carnegie Mellon University

A RowHammer Survey: Recent Update

 Onur Mutlu, Ataberk Olgun, and A. Giray Yaglikci, "Fundamentally Understanding and Solving RowHammer" Invited Special Session Paper at the <u>28th Asia and South Pacific Design</u> Automation Conference (ASP-DAC), Tokyo, Japan, January 2023. [arXiv version] [Slides (pptx) (pdf)] [Talk Video (26 minutes)]

Fundamentally Understanding and Solving RowHammer

Onur Mutlu onur.mutlu@safari.ethz.ch ETH Zürich Zürich, Switzerland Ataberk Olgun ataberk.olgun@safari.ethz.ch ETH Zürich Zürich, Switzerland A. Giray Yağlıkcı giray.yaglikci@safari.ethz.ch ETH Zürich Zürich, Switzerland

https://arxiv.org/pdf/2211.07613.pdf

Understanding RowHammer

First RowHammer Analysis [ISCA 2014]

Yoongu Kim, Ross Daly, Jeremie Kim, Chris Fallin, Ji Hye Lee, Donghyuk Lee, Chris Wilkerson, Konrad Lai, and Onur Mutlu,
 "Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors"
 Proceedings of the 41st International Symposium on Computer Architecture (ISCA), Minneapolis, MN, June 2014.
 [Slides (pptx) (pdf)] [Lightning Session Slides (pptx) (pdf)] [Source Code and Data] [Lecture Video (1 hr 49 mins), 25 September 2020]
 One of the 7 papers of 2012-2017 selected as Top Picks in Hardware and Embedded Security for IEEE TCAD (link).
 Selected to the ISCA-50 25-Year Retrospective Issue covering 1996-2020 in 2023 (Retrospective (pdf) Full Issue).

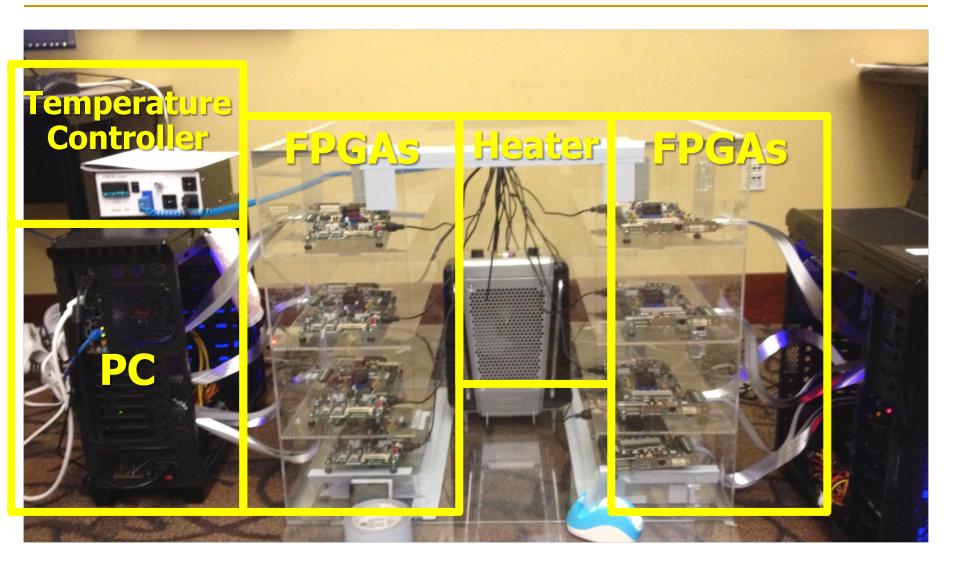
Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors

Yoongu Kim¹ Ross Daly^{*} Jeremie Kim¹ Chris Fallin^{*} Ji Hye Lee¹ Donghyuk Lee¹ Chris Wilkerson² Konrad Lai Onur Mutlu¹

¹Carnegie Mellon University ²Intel Labs

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RowHammer Infrastructure (2012-2014)



SAFARI

Kim+, "Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors," ISCA 2014.

| Manufacturer | Madala | Date* | Timing | g† | Organiz | ation | | Chip | | Vict | ims-per-Mo | dule | RI _{th} (ms) |
|------------------------|--|----------------|--------------|----------------------|-----------|-------|------------------------|------------|-------------------------|---|---|--|-----------------------|
| manujaciurer | Module | (yy-ww) | Freq (MT/s) | t _{RC} (ns) | Size (GB) | Chips | Size (Gb) [‡] | Pins | DieVersion [§] | Average | Minimum | Maximum | Min |
| | A ₁ | 10-08 | 1066 | 50.625 | 0.5 | 4 | 1 | ×16 | В | 0 | 0 | 0 | - |
| | A ₂ | 10-20 | 1066 | 50.625 | 1 | 8 | 1 | $\times 8$ | \mathcal{F} | 0 | 0 | 0 | - |
| | A ₃₋₅ | 10-20 | 1066 | 50.625 | 0.5 | 4 | 1 | ×16 | В | 0 | 0 | 0 | - |
| | A ₆₋₇ | 11-24 | 1066 | 49.125 | 1 | 4 | 2 | ×16 | \mathcal{D} | 7.8×10^{1} | 5.2×10^{1} | 1.0×10^{2} | 21.3 |
| | A ₈₋₁₂ | 11-26 | 1066 | 49.125 | 1 | 4 | 2 | ×16 | \mathcal{D} | 2.4×10^{2} | 5.4×10^{1} | 4.4×10^{2} | 16.4 |
| • | A ₁₃₋₁₄ | 11-50 | 1066 | 49.125 | 1 | 4 | 2 | ×16 | \mathcal{D} | 8.8×10^{1} | 1.7×10^{1} | 1.6×10^{2} | 26.2 |
| Α | A ₁₅₋₁₆ | 12-22 | 1600 | 50.625 | 1 | 4 | 2 | ×16 | \mathcal{D} | 9.5 | 9 | 1.0×10^{1} | 34.4 |
| Total of | A ₁₇₋₁₈ | 12-26 | 1600 | 49.125 | 2 | 8 | 2 | ×8 | M | 1.2×10^{2} | 3.7×10^{1} | 2.0×10^{2} | 21.3 |
| 43 Modules | A ₁₉₋₃₀ | 12-40 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | ĸ | 8.6×10^{6} | | 1.0×10^{7} | 8.2 |
| | A ₃₁₋₃₄ | 13-02 | 1600 | 48.125 | 2 2 | 8 | 2 2 | ×8 ×8 | - | 1.8×10^{6} 4.0×10^{1} | 1.0×10^{6} 1.9×10^{1} | 3.5×10^{6} | 11.5 21.3 |
| | A ₃₅₋₃₆ | 13-14 13-20 | 1600 1600 | 48.125 48.125 | 2 | 8 | 2 | ×8 | - <i>K</i> | 4.0×10^{-1} 1.7×10^{6} | 1.9×10^{-1} 1.4×10^{6} | 6.1×10^{1} 2.0×10^{6} | 9.8 |
| | A ₃₇₋₃₈ A ₃₉₋₄₀ | 13-20 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | ĸ | 5.7×10^4 | | 6.0×10^4 | 16.4 |
| | A 39-40 | 13-28 | 1600 | 49.125 | 2 | 8 | 2 | ×8 | - | 2.7×10^{5} | | 2.7×10^{5} | 18.0 |
| | A ₄₁ A ₄₂₋₄₃ | 14-04 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | κ | 0.5 | 0 | 1 | 62.3 |
| | B ₁ | 08-49 | 1066 | 50.625 | 1 | 8 | 1 | ×8 | D | 0 | 0 | 0 | _ |
| | B ₂ | 08-49 | 1066 | 50.625 | 1 | 8 | 1 | ×8 | ε | 0 | 0 | 0 | - |
| | B ₃ | 10-19 | 1066 | 50.625 | 1 | 8 | i | ×8 | F | 0 | 0 | 0 | - |
| | B ₄ | 10-31 | 1333 | 49.125 | 2 | 8 | 2 | $\times 8$ | С | 0 | 0 | 0 | - |
| | B ₅ | 11-13 | 1333 | 49.125 | 2 | 8 | 2 | $\times 8$ | С | 0 | 0 | 0 | - |
| | B ₆ | 11-16 | 1066 | 50.625 | 1 | 8 | 1 | $\times 8$ | \mathcal{F} | 0 | 0 | 0 | - |
| | B ₇ | 11-19 | 1066 | 50.625 | 1 | 8 | 1 | $\times 8$ | F | 0 | 0 | 0 | - |
| D | B ₈ | 11-25 | 1333 | 49.125 | 2 | 8 | 2 | $\times 8$ | С | 0 | 0 | 0 | - |
| В | B ₉ | 11-37 | 1333 | 49.125 | 2 | 8 | 2 | ×8 | \mathcal{D} | 1.9×10^{6} | 1.9×10^{6} | 1.9×10^{6} | 11.5 |
| Total of | B ₁₀₋₁₂ | 11-46 | 1333 | 49.125 | 2 | 8 | 2 | ×8 | D | | 1.5×10^{6} | | 11.5 |
| | B ₁₃ | 11-49 | 1333 1866 | 49.125 | 2 2 | 8 | 2 2 | ×8 ×8 | C D | $0 \\ 9.1 \times 10^5$ | $0 \\ 9.1 \times 10^5$ | $0 \\ 9.1 \times 10^5$ | 9.8 |
| | B ₁₄ | 12-01 12-10 | 1866 | 47.125 47.125 | 2 | 8 | 2 | ×8 | D | 9.1×10^{5} 9.8×10^{5} | 9.1×10^{5} 7.8×10^{5} | 9.1×10^{6} 1.2×10^{6} | 11.5 |
| | B ₁₅₋₃₁ | 12-10 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | E | 7.4×10^{5} | | 7.4×10^{5} | 11.5 |
| | B ₃₂ B ₃₃₋₄₂ | 12-28 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | ε | 5.2×10^{5} | 1.9×10^{5} | 7.3×10^{5} | 11.5 |
| | B ₄₃₋₄₇ | 12-31 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | ε | 4.0×10^{5} | | 5.5×10^{5} | 13.1 |
| | B ₄₈₋₅₁ | 13-19 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | Ē | 1.1×10^{5} | 7.4×10^{4} | 1.4×10^{5} | 14.7 |
| | B ₅₂₋₅₃ | 13-40 | 1333 | 49.125 | 2 | 8 | 2 | $\times 8$ | \mathcal{D} | 2.6×10^{4} | | 2.9×10^{4} | 21.3 |
| | B ₅₄ | 14-07 | 1333 | 49.125 | 2 | 8 | 2 | ×8 | \mathcal{D} | $7.5 	imes 10^3$ | $7.5 	imes 10^3$ | $7.5 	imes 10^3$ | 26.2 |
| | C ₁ | 10-18 | 1333 | 49.125 | 2 | 8 | 2 | ×8 | \mathcal{A} | 0 | 0 | 0 | - |
| | C ₂ | 10-20 | 1066 | 50.625 | 2 | 8 | 2 | $\times 8$ | \mathcal{A} | 0 | 0 | 0 | - |
| | C ₂ | 10-22 | 1066 | 50.625 | 2 | 8 | 2 | $\times 8$ | \mathcal{A} | 0 | 0 | 0 | - |
| | C | 10-26 | 1333 | 49.125 | 2 | 8 | 2 | $\times 8$ | B | 8.9×10^{2} | 6.0×10^{2} | 1.2×10^{3} | 29.5 |
| | U ₆ | 10-43 | 1333 | 49.125 | 1 | 8 | 1 | ×8 | τ | 0 | 0 | 0 | - |
| | U7 | 10-51 | 1333 | 49.125 | 2 | 8 | 2 | ×8 | В | 4.0×10^{2} | 4.0×10^{2} | 4.0×10^{2} | 29.5 |
| | C ₈ | 11-12 | 1333 | 46.25 | 2 | 8 | 2 | ×8 | B | 6.9×10^2 | 6.9×10^2 | 6.9×10^2 | 21.3 |
| | C, | 11-19 | 1333 | 46.25 | 2 2 | 8 | 2 | ×8 ×8 | B B | 9.2 × 10 ² 3 | 9.2 × 10 ² 3 | 9.2×10^2 | 27.9 39.3 |
| | C ₁₀ | 11-31 11-42 | 1333 1333 | 49.125 49.125 | 2 | 8 | 2 | ×8 | B | 1.6×10^2 | 1.6×10^2 | 1.6×10^2 | 39.3 |
| С | C ₁₁ C ₁₂ | 11-42 | 1555 | 49.125 | 2 | 8 | 2 | ×8 | C | | | 7.1×10^4 | 19.7 |
| The task of | C ₁₃ | 12-08 | 1333 | 49.125 | 2 | 8 | 2 | ×8 | C | 3.9×10^4 | 3.9×10^4 | 3.9×10^4 | 21.3 |
| Total of 32 Modules | C ₁₄₋₁₅ | 12-12 | 1333 | 49.125 | 2 | 8 | 2 | ×8 | Ĉ | | 2.1×10^4 | 5.4×10^{4} | 21.3 |
| 32 Modules | C ₁₆₋₁₈ | 12-20 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | C | 3.5×10^{3} | 1.2×10^{3} | 7.0×10^{3} | 27.9 |
| | C ₁₉ | 12-23 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | ε | 1.4×10^{5} | 1.4×10^{5} | 1.4×10^{5} | 18.0 |
| | C20 | 12-24 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | С | 6.5×10^{4} | 6.5×10^{4} | 6.5×10^{4} | 21.3 |
| | C ₂₁ | 12-26 | 1600 | 48.125 | 2 | 8 | 2 | $\times 8$ | С | $2.3 	imes 10^4$ | $2.3 	imes 10^4$ | $2.3 	imes 10^4$ | 24.6 |
| | C ₂₂ | 12-32 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | С | 1.7×10^{4} | 1.7×10^{4} | $1.7 	imes 10^4$ | 22.9 |
| | C ₂₃₋₂₄ | 12-37 | 1600 | 48.125 | 2 | 8 | 2 | $\times 8$ | С | | | | 18.0 |
| | G25.20 | 12-41 | 1600 | 48.125 | 2 | 8 | 2 | $\times 8$ | С | 2.0×10^{4} | 1.1×10^{4} | 3.2×10^{4} | 19.7 |
| | C ₃₁ C ₃₂ | 13-11 | 1600 | 48.125 | 2 | 8 | 2 | ×8 | С | 3.3×10^{5} | | 3.3×10^{5} | 14.7 |
| | | 13-35 | 1600 | 48.125 | 2 | 8 | 2 | $\times 8$ | С | 3.7×10^{4} | 3.7×10^{4} | 3.7×10^{4} | 21.3 |

* We report the manufacture date marked on the chip packages, which is more accurate than other dates that can be gleaned from a module. † We report timing constraints stored in the module's on-board ROM [33], which is read by the system BIOS to calibrate the memory controller. ‡ The maximum DRAM chip size supported by our testing platform is 2Gb.

§ We report DRAM die versions marked on the chip packages, which typically progress in the following manner: $\mathcal{M} \to \mathcal{A} \to \mathcal{B} \to \mathcal{C} \to \cdots$.

Table 3. Sample population of 129 DDR3 DRAM modules, categorized by manufacturer and sorted by manufacture date

Tested DRAM Modules from 2008-2014

(129 total)

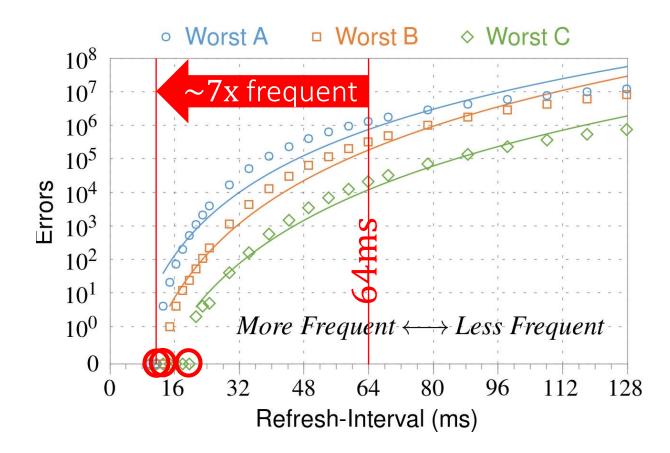
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RowHammer Characterization Results

- 1. Most Modules Are at Risk
- 2. Errors vs. Vintage
- 3. Error = Charge Loss
- 4. Adjacency: Aggressor & Victim
- 5. Sensitivity Studies
- 6. Other Results in Paper
- 7. Solution Space

Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors, (Kim et al., ISCA 2014)

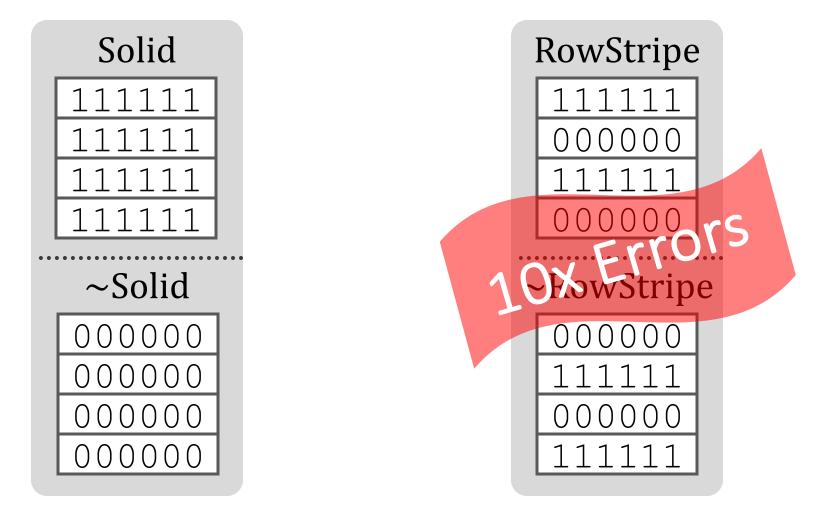
2 Refresh Interval



Note: Using three modules with the most errors (only first bank)

More frequent refreshes \rightarrow Fewer errors





Errors affected by data stored in other cells

6. Other Key Observations [ISCA'14]

- Victim Cells ≠ Retention-Weak Cells
 - Almost no overlap between them
- Errors are repeatable
 - Across ten iterations of testing, >70% of victim cells had errors in every iteration
- As many as 4 errors per cache-line

 Simple ECC (e.g., SECDED) cannot prevent all errors
- Cells affected by two aggressors on either side
 Double sided hammering

Major RowHammer Characteristics (2014)

Yoongu Kim, Ross Daly, Jeremie Kim, Chris Fallin, Ji Hye Lee, Donghyuk Lee, Chris Wilkerson, Konrad Lai, and Onur Mutlu,
 "Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors"
 Proceedings of the 41st International Symposium on Computer Architecture (ISCA), Minneapolis, MN, June 2014.
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Yoongu Kim¹ Ross Daly^{*} Jeremie Kim¹ Chris Fallin^{*} Ji Hye Lee¹ Donghyuk Lee¹ Chris Wilkerson² Konrad Lai Onur Mutlu¹

¹Carnegie Mellon University ²Intel Labs

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RowHammer is Getting Much Worse (2020)

 Jeremie S. Kim, Minesh Patel, A. Giray Yaglikci, Hasan Hassan, Roknoddin Azizi, Lois Orosa, and Onur Mutlu, "Revisiting RowHammer: An Experimental Analysis of Modern Devices and Mitigation Techniques" Proceedings of the <u>47th International Symposium on Computer</u> <u>Architecture</u> (ISCA), Valencia, Spain, June 2020.
 [Slides (pptx) (pdf)]
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Revisiting RowHammer: An Experimental Analysis of Modern DRAM Devices and Mitigation Techniques

Jeremie S. Kim^{§†} Minesh Patel[§] A. Giray Yağlıkçı[§] Hasan Hassan[§] Roknoddin Azizi[§] Lois Orosa[§] Onur Mutlu^{§†} [§]ETH Zürich [†]Carnegie Mellon University

RowHammer Has Many Dimensions (2021)

 Lois Orosa, Abdullah Giray Yaglikci, Haocong Luo, Ataberk Olgun, Jisung Park, Hasan Hassan, Minesh Patel, Jeremie S. Kim, and Onur Mutlu,
 "A Deeper Look into RowHammer's Sensitivities: Experimental Analysis of Real DRAM Chips and Implications on Future Attacks and Defenses" Proceedings of the <u>54th International Symposium on Microarchitecture</u> (MICRO), Virtual, October 2021.
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A Deeper Look into RowHammer's Sensitivities: Experimental Analysis of Real DRAM Chips and Implications on Future Attacks and Defenses

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ETH Zürich

ETH Zürich

ETH Zürich

ETH Zürich

RowHammer vs. Wordline Voltage (2022)

 A. Giray Yağlıkçı, Haocong Luo, Geraldo F. de Oliviera, Ataberk Olgun, Minesh Patel, Jisung Park, Hasan Hassan, Jeremie S. Kim, Lois Orosa, and Onur Mutlu, "Understanding RowHammer Under Reduced Wordline Voltage: An Experimental Study Using Real DRAM Devices" Proceedings of the 52nd Annual IEEE/IFIP International Conference on Dependable Systems and Networks (DSN), Baltimore, MD, USA, June 2022.
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 [Talk Video (34 minutes, including Q&A)]
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Understanding RowHammer Under Reduced Wordline Voltage: An Experimental Study Using Real DRAM Devices

A. Giray Yağlıkçı¹ Haocong Luo¹ Geraldo F. de Oliviera¹ Ataberk Olgun¹ Minesh Patel¹ Jisung Park¹ Hasan Hassan¹ Jeremie S. Kim¹ Lois Orosa^{1,2} Onur Mutlu¹ ¹ETH Zürich ²Galicia Supercomputing Center (CESGA)

RowHammer in HBM Chips (2023)

 Ataberk Olgun, Majd Osserian, A. Giray Yağlıkçı, Yahya Can Tugrul, Haocong Luo, Steve Rhyner, Behzad Salami, Juan Gomez-Luna, and Onur Mutlu,
 <u>"An Experimental Analysis of RowHammer in HBM2 DRAM Chips"</u> *Proceedings of the <u>53nd Annual IEEE/IFIP International Conference on</u> <i>Dependable Systems and Networks Disrupt Track (DSN Disrupt)*, Porto, Portugal, June 2023.
 [arXiv version]
 [Slides (pptx) (pdf)]
 [Talk Video (24 minutes, including Q&A)]

An Experimental Analysis of RowHammer in HBM2 DRAM Chips

Ataberk Olgun¹ Majd Osseiran^{1,2} A. Giray Yağlıkçı¹ Yahya Can Tuğrul¹ Haocong Luo¹ Steve Rhyner¹ Behzad Salami¹ Juan Gomez Luna¹ Onur Mutlu¹ ¹SAFARI Research Group, ETH Zürich ²American University of Beirut

https://arxiv.org/pdf/2305.17918.pdf

RowHammer Solutions

Two Types of RowHammer Solutions

Immediate

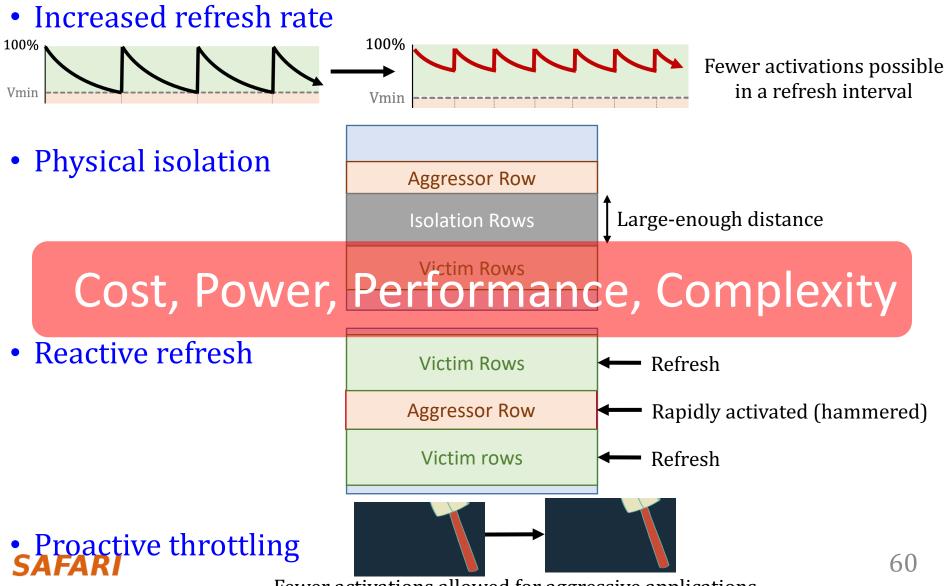
- To protect the vulnerable DRAM chips in the field
- Limited possibilities

- Longer-term
 - To protect future DRAM chips
 - Wider range of protection mechanisms

- Our ISCA 2014 paper proposes both types of solutions
 - Seven solutions in total
 - PARA proposed as best solution \rightarrow already employed in the field

RowHammer Solution Approaches

• More robust DRAM chips **and/or** error-correcting codes



Fewer activations allowed for aggressive applications

Apple's Security Patch for RowHammer

https://support.apple.com/en-gb/HT204934

Available for: OS X Mountain Lion v10.8.5, OS X Mavericks v10.9.5

Impact: A malicious application may induce memory corruption to escalate privileges

Description: A disturbance error, also known as Rowhammer, exists with some DDR3 RAM that could have led to memory corruption. This issue was mitigated by increasing memory refresh rates.

CVE-ID

CVE-2015-3693 : Mark Seaborn and Thomas Dullien of Google, working from original research by Yoongu Kim et al (2014)

HP, Lenovo, and many other vendors released similar patches

Our First Solution to RowHammer

- PARA: <u>Probabilistic Adjacent Row Activation</u>
- Key Idea
 - After closing a row, we activate (i.e., refresh) one of its neighbors with a low probability: p = 0.005
- Reliability Guarantee
 - When p=0.005, errors in one year: 9.4×10^{-14}
 - By adjusting the value of p, we can vary the strength of protection against errors

Advantages of PARA

- PARA refreshes rows infrequently
 - Low power
 - Low performance-overhead
 - Average slowdown: 0.20% (for 29 benchmarks)
 - Maximum slowdown: 0.75%
- PARA is stateless
 - Low cost
 - Low complexity
- PARA is an effective and low-overhead solution to prevent disturbance errors

Requirements for PARA

- If implemented in DRAM chip (done today)
 - Enough slack in timing and refresh parameters
 - Plenty of slack today:
 - Lee et al., "Adaptive-Latency DRAM: Optimizing DRAM Timing for the Common Case," HPCA 2015.
 - Chang et al., "Understanding Latency Variation in Modern DRAM Chips," SIGMETRICS 2016.
 - Lee et al., "Design-Induced Latency Variation in Modern DRAM Chips," SIGMETRICS 2017.
 - Chang et al., "Understanding Reduced-Voltage Operation in Modern DRAM Devices," SIGMETRICS 2017.
 - Ghose et al., "What Your DRAM Power Models Are Not Telling You: Lessons from a Detailed Experimental Study," SIGMETRICS 2018.
 - Kim et al., "Solar-DRAM: Reducing DRAM Access Latency by Exploiting the Variation in Local Bitlines," ICCD 2018.
- If implemented in memory controller
 - Need coordination between controller and DRAM
 - Memory controller should know which rows are physically adjacent

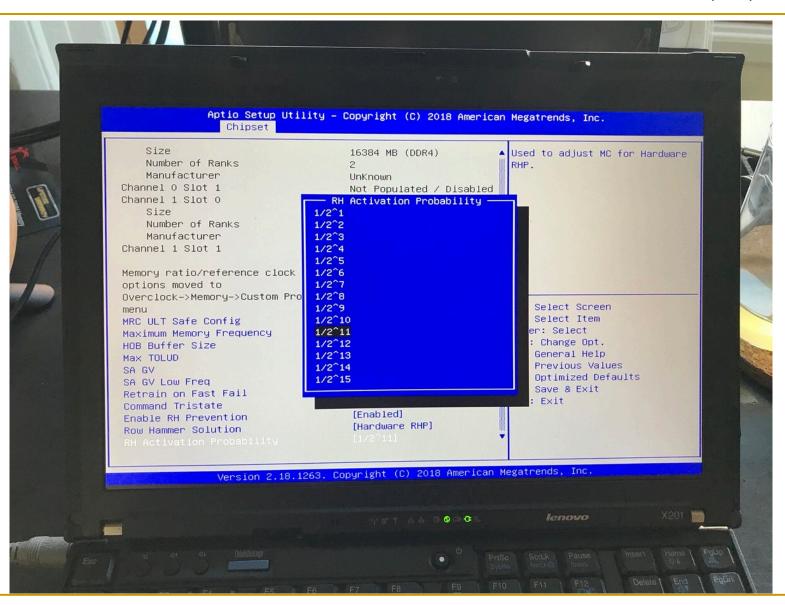
Probabilistic Activation in Real Life (I)

| Aptio Setup Utili Chipset | ty – Copyright (C) 2018 Americ | can Megatrends, Inc. | |
|--|--|---|--|
| Channel 0 Slot 0 Size Number of Ranks Manufacturer Channel 0 Slot 1 Channel 1 Slot 0 Size Number of Ranks Manufacturer Channel 1 Slot 1 Memory ratio/reference clock options moved to Overclock->Memory->Custom Prof: menu MRC ULT Safe Config Maximum Memory Frequency HOB Buffer Size Max TOLUD SA GV SA GV Low Freq Retrain on Fast Fail Command Tristate Enable RH Prevention Row Hammer Solution | [Disabled] [Auto] [Auto] [Dynamic] [Enabled] [MRC default] [Enabled] [Enabled] [Enabled] [Hardware RHP] | <pre>++: Select Screen fl: Select Item Enter: Select +/-: Change Opt. F1: General Help F2: Previous Values F3: Optimized Defaults F4: Save & Exit ESC: Exit</pre> | |
| Version 2.18.12 | 63. Copyright (C) 2018 America | n Megatrends, Inc. | |

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https://twitter.com/isislovecruft/status/1021939922754723841

Probabilistic Activation in Real Life (II)



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https://twitter.com/isislovecruft/status/1021939922754723841

Seven RowHammer Solutions Proposed

Yoongu Kim, Ross Daly, Jeremie Kim, Chris Fallin, Ji Hye Lee, Donghyuk Lee, Chris Wilkerson, Konrad Lai, and Onur Mutlu,
 "Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors"
 Proceedings of the <u>41st International Symposium on Computer Architecture</u> (ISCA), Minneapolis, MN, June 2014.
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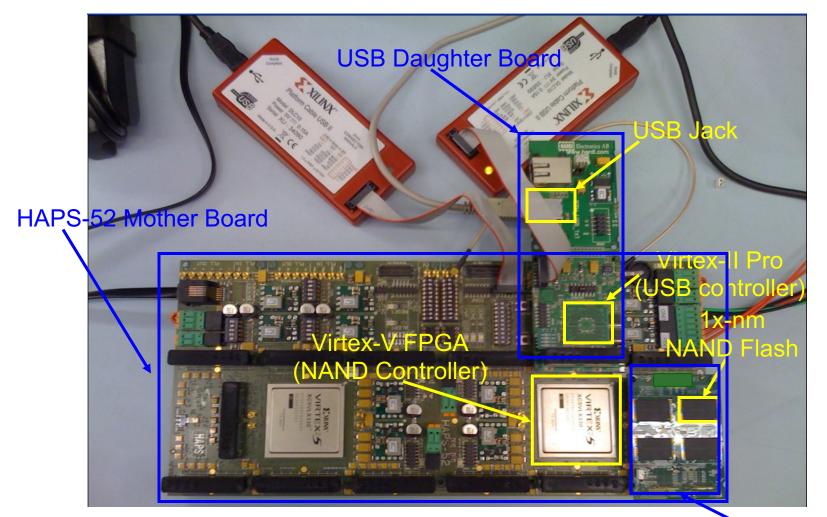
¹Carnegie Mellon University ²Intel Labs

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Main Memory Needs **Intelligent Controllers** for Security, Safety, Reliability, Scaling

Aside: Intelligent Controller for NAND Flash



[DATE 2012, ICCD 2012, DATE 2013, ITJ 2013, ICCD 2013, SIGMETRICS 2014, HPCA 2015, DSN 2015, MSST 2015, JSAC 2016, HPCA 2017, DFRWS 2017, PIEEE 2017, HPCA 2018, SIGMETRICS 2018]

NAND Daughter Board

Cai+, "Error Characterization, Mitigation, and Recovery in Flash Memory Based Solid State Drives," Proc. IEEE 2017.

Intelligent Flash Controllers [PIEEE'17]



Proceedings of the IEEE, Sept. 2017

Error Characterization, Mitigation, and Recovery in Flash-Memory-Based Solid-State Drives



This paper reviews the most recent advances in solid-state drive (SSD) error characterization, mitigation, and data recovery techniques to improve both SSD's reliability and lifetime.

By Yu Cai, Saugata Ghose, Erich F. Haratsch, Yixin Luo, and Onur Mutlu

https://arxiv.org/pdf/1706.08642

Two Major RowHammer Directions

Understanding RowHammer

- Many effects still need to be rigorously examined
 - Aging of DRAM Chips
 - Environmental Conditions (e.g., Process, Voltage, Temperature)
 - Memory Access Patterns
 - Memory Controller & System Design Decisions

...

Solving RowHammer

- Flexible and efficient RowHammer solutions are necessary
 - In-field patchable / reconfigurable / programmable solutions
- Co-architecting System and Memory is important
 - To avoid performance and denial-of-service problems

RowHammer in 2020-2023

Revisiting RowHammer

RowHammer is Getting Much Worse

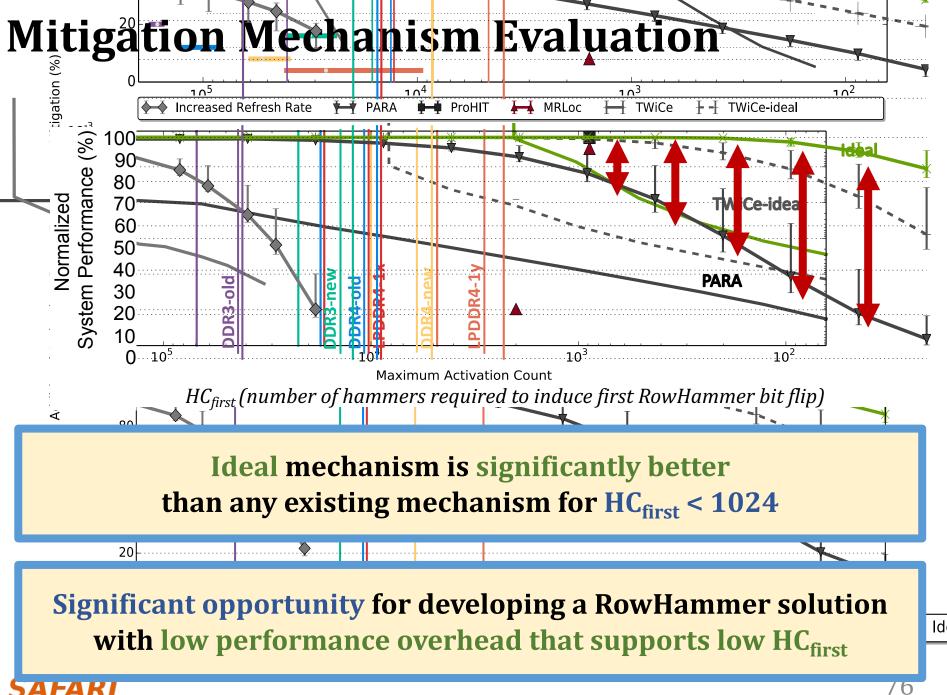
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Key Takeaways from 1580 Chips

- Newer DRAM chips are much more vulnerable to **RowHammer (more bit flips, happening earlier)**
- There are new chips whose weakest cells fail after **only** 4800 hammers
- Chips of newer DRAM technology nodes can exhibit RowHammer bit flips 1) in more rows and 2) farther away from the victim row.
- Existing mitigation mechanisms are NOT effective at future technology nodes SAFARI



New RowHammer Characteristics

RowHammer Has Many Dimensions

 Lois Orosa, Abdullah Giray Yaglikci, Haocong Luo, Ataberk Olgun, Jisung Park, Hasan Hassan, Minesh Patel, Jeremie S. Kim, and Onur Mutlu,
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Our Goal

Provide insights into three fundamental properties



To find **effective and efficient** attacks and defenses



Summary of The Study & Key Results

- 272 DRAM chips from four major manufacturers
- 6 major takeaways from 16 novel observations
- A RowHammer bit flip is more likely to occur
 1) in a bounded range of temperature
 2) if the aggressor row is active for longer time
 3) in certain physical regions of the DRAM module under attack
- Our novel observations can inspire and aid future work
 - Craft more effective attacks
 - Design more effective and efficient defenses

Example Attack Improvement 3: Bypassing Defenses with Aggressor Row Active Time

Activating aggressor rows as frequently as possible:

Row A is
activeRow B is
activeRow A is
activeRow B is
activeTime

Keeping aggressor rows active for a longer time:

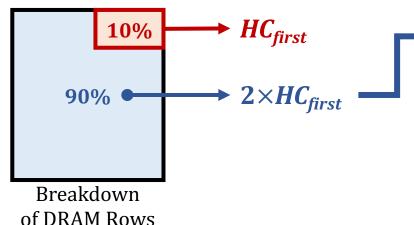


Reduces the minimum activation count to induce a bit flip by 36%

Bypasses defenses that do not account for this reduction

Example Defense Improvements

• Example 1: Leveraging variation across DRAM rows



Aggressiveness can be reduced:
 33% area reduction
 for BlockHammer [Yağlıkçı+, HPCA'21]

 80% area reduction
 for Graphene [Park+, MICRO'20]

• Example 2: Leveraging variation with temperature

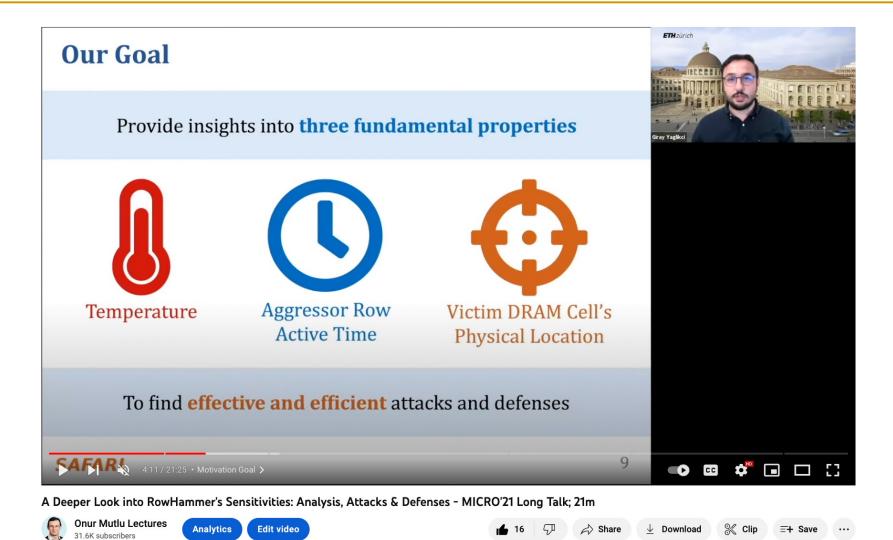
• A DRAM cell experiences **bit flips** within **a bounded temperature range**



• A row can be **disabled** within the row's **vulnerable temperature range**



Deeper Look into RowHammer: Talk Video



https://youtube.com/watch?v=fkM32oA0u6U&si=EnSIkaIECMiOmarE

More RowHammer Analysis

RowHammer vs. Wordline Voltage (2022)

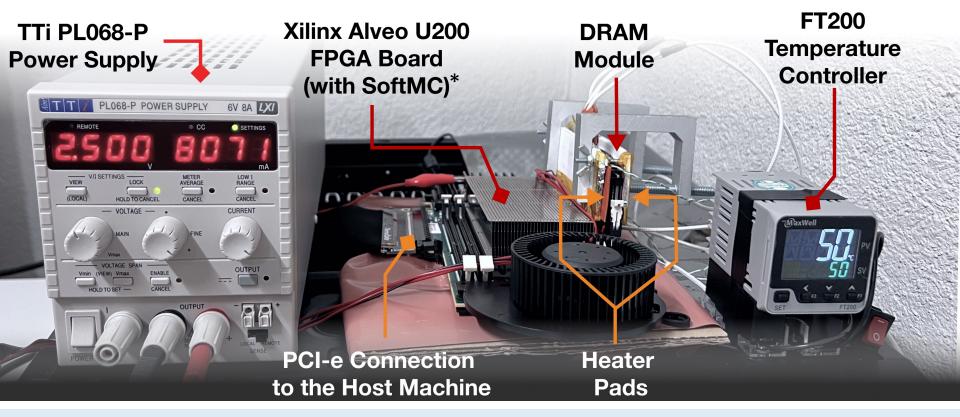
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Updated DRAM Testing Infrastructure

FPGA-based SoftMC (Xilinx Virtex UltraScale+ XCU200)



Fine-grained control over DRAM commands, timing parameters (±1.5ns), temperature (±0.1°C), and wordline voltage (±1mV)

SAFARI *Hassan et al., "<u>SoftMC: A Flexible and Practical Open-Source Infrastructure for Enabling Experimental</u> 86 <u>DRAM Studies</u>," in HPCA, 2017. [Available on GitHub: <u>https://github.com/CMU-SAFARI/SoftMC]</u>

Summary

We provide *the first* RowHammer characterization **under reduced wordline voltage**

Experimental results with 272 real DRAM chips show that reducing wordline voltage:

1. Reduces RowHammer vulnerability

- Bit error rate caused by a RowHammer attack reduces by 15.2% (66.9% max)
- A row needs to be activated 7.4% more times (85.8% max) to induce the first bit flip

2. Increases row activation latency

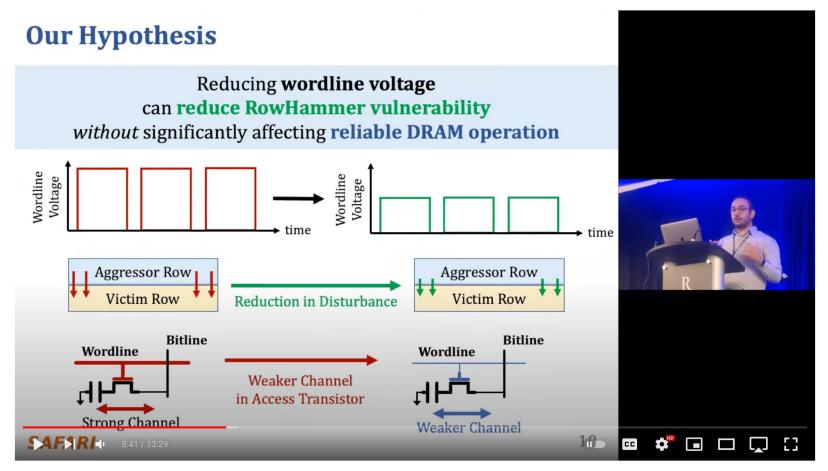
- More than **76%** of the tested DRAM chips **reliably operate** using **nominal** timing parameters
- Remaining **24% reliably operate** with **increased** (up to 24ns) row activation latency

3. Reduces data retention time

- 80% of the tested DRAM chips reliably operate using nominal refresh rate
- Remaining 20% reliably operate by
 - Using single error correcting codes
 - **Doubling the refresh rate** for a small fraction (16.4%) of DRAM rows

Reducing wordline voltage can **reduce RowHammer vulnerability** *without* significantly affecting **reliable DRAM operation**

RowHammer vs. Wordline Voltage: Talk Video



Understanding RowHammer Under Reduced Wordline Voltage - Live Talk in DSN'22 by Giray Yaglikci



https://youtube.com/watch?v=fkM32oA0u6U&si=EnSIkaIECMiOmarE

•••

RowHammer in HBM Chips (2023)

 Ataberk Olgun, Majd Osserian, A. Giray Yağlıkçı, Yahya Can Tugrul, Haocong Luo, Steve Rhyner, Behzad Salami, Juan Gomez-Luna, and Onur Mutlu, "An Experimental Analysis of RowHammer in HBM2 DRAM Chips" Proceedings of the <u>53nd Annual IEEE/IFIP International Conference on</u> <u>Dependable Systems and Networks</u> Disrupt Track (DSN Disrupt), Porto, Portugal, June 2023.
 [arXiv version]
 [Slides (pptx) (pdf)]
 [Talk Video (24 minutes, including Q&A)]

An Experimental Analysis of RowHammer in HBM2 DRAM Chips

Ataberk Olgun¹ Majd Osseiran^{1,2} A. Giray Yağlıkçı¹ Yahya Can Tuğrul¹ Haocong Luo¹ Steve Rhyner¹ Behzad Salami¹ Juan Gomez Luna¹ Onur Mutlu¹ ¹SAFARI Research Group, ETH Zürich ²American University of Beirut

https://arxiv.org/pdf/2305.17918.pdf

New RowHammer Solutions

Industry-Adopted Solutions Do Not Work

 Pietro Frigo, Emanuele Vannacci, Hasan Hassan, Victor van der Veen, Onur Mutlu, Cristiano Giuffrida, Herbert Bos, and Kaveh Razavi,
 "TRRespass: Exploiting the Many Sides of Target Row Refresh"
 Proceedings of the <u>41st IEEE Symposium on Security and Privacy</u> (S&P), San Francisco, CA, USA, May 2020.
 [Slides (pptx) (pdf)]
 [Lecture Slides (pptx) (pdf)]
 [Lecture Slides (pptx) (pdf)]
 [Lecture Video (17 minutes)]
 [Lecture Video (59 minutes)]
 [Source Code]
 [Web Article]
 Best paper award.
 Pwnie Award 2020 for Most Innovative Research. Pwnie Awards 2020

TRRespass: Exploiting the Many Sides of Target Row Refresh

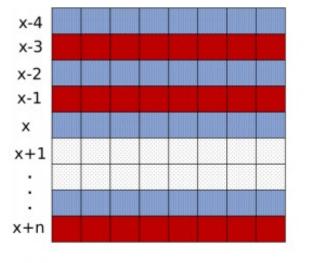
Pietro Frigo^{*†} Emanuele Vannacci^{*†} Hasan Hassan[§] Victor van der Veen[¶] Onur Mutlu[§] Cristiano Giuffrida^{*} Herbert Bos^{*} Kaveh Razavi^{*}

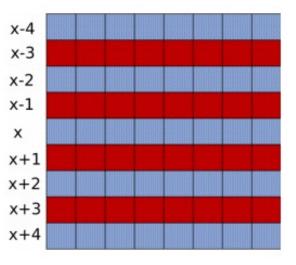
*Vrije Universiteit Amsterdam

[§]ETH Zürich

[¶]Qualcomm Technologies Inc.

Example Many-Sided Hammering Patterns





(a) Assisted double-sided

(b) 4-sided

Fig. 12: Hammering patterns discovered by *TRRespass*. Aggressor rows are in red () and victim rows are in blue ().

TRRespass Key Takeaways

RowHammer is still an open problem

Security by obscurity is likely not a good solution

Industry-Adopted Solutions Are Very Poor

 Hasan Hassan, Yahya Can Tugrul, Jeremie S. Kim, Victor van der Veen, Kaveh Razavi, and Onur Mutlu,
 "Uncovering In-DRAM RowHammer Protection Mechanisms: A New Methodology, Custom RowHammer Patterns, and Implications" Proceedings of the <u>54th International Symposium on Microarchitecture</u> (MICRO), Virtual, October 2021.
 [Slides (pptx) (pdf)]
 [Short Talk Slides (pptx) (pdf)]
 [Lightning Talk Slides (pptx) (pdf)]
 [Lightning Talk Video (100 seconds)]
 [arXiv version]

Uncovering In-DRAM RowHammer Protection Mechanisms: A New Methodology, Custom RowHammer Patterns, and Implications

| Hasan Hassan † | Yahya Can Tuğrul ^{†‡} | Jeremie S. Kin | Victor van der Veen $^{\sigma}$ | | |
|---------------------------|---|----------------|--|--|--|
| | Kaveh Razavi † | Onur Mutlu | Ť | | |
| †ETH Zürich | [‡] TOBB University of Economics | & Technology | $^{\sigma}$ Qualcomm Technologies Inc. | | |

Key Takeaways

All 45 modules we test are vulnerable

99.9% of rows in a DRAM bank experience **at least one RowHammer bit flip**

ECC is ineffective: up to 7 RowHammer bit flips in an 8-byte dataword

| Module | Date (yy-ww) | Chip Density (Gbit) | Organization | | | Our Key TRR Observations and Results | | | | | | | | |
|--------|-----------------|---------------------------|--------------|-------|------|--------------------------------------|------------|------------------------|-----------------------|-----------------|---------------------|------------------------|----------------------------|---------------------------------------|
| | | | Ranks | Banks | Pins | $ HC_{first}^{\dagger}$ | Version | Aggressor Detection | Aggressor Capacity | Per-Bank TRR | TRR-to-REF Ratio | Neighbors Refreshed | % Vulnerable DRAM Rows† | Max. Bit Flips per Row per Hammer‡ |
| A0 | 19-50 | 8 | 1 | 16 | 8 | 16K | A_{TRR1} | Counter-based | 16 | 1 | 1/9 | 4 | 73.3% | 1.16 |
| A1-5 | 19-36 | 8 | 1 | 8 | 16 | 13K - 15K | A_{TRR1} | Counter-based | 16 | 1 | 1/9 | 4 | 99.2% - 99.4% | 2.32 - 4.73 |
| A6-7 | 19-45 | 8 | 1 | 8 | 16 | 13K - 15K | A_{TRR1} | Counter-based | 16 | 1 | 1/9 | 4 | 99.3% - 99.4% | 2.12 - 3.86 |
| A8-9 | 20-07 | 8 | 1 | 16 | 8 | 12K - 14K | A_{TRR1} | Counter-based | 16 | 1 | 1/9 | 4 | 74.6% - 75.0% | 1.96 - 2.96 |
| A10-12 | 19-51 | 8 | 1 | 16 | 8 | 12K - 13K | A_{TRR1} | Counter-based | 16 | 1 | 1/9 | 4 | 74.6% - 75.0% | 1.48 - 2.86 |
| A13-14 | 20-31 | 8 | 1 | 8 | 16 | 11K-14K | A_{TRR2} | Counter-based | 16 | 1 | 1/9 | 2 | 94.3% - 98.6% | 1.53 - 2.78 |
| B0 | 18-22 | 4 | 1 | 16 | 8 | 44K | B_{TRR1} | Sampling-based | 1 | × | 1/4 | 2 | 99.9% | 2.13 |
| B1-4 | 20-17 | 4 | 1 | 16 | 8 | 159K - 192K | B_{TRR1} | Sampling-based | 1 | × | 1/4 | 2 | 23.3% - 51.2% | 0.06 - 0.11 |
| B5-6 | 16-48 | 4 | 1 | 16 | 8 | 44K-50K | B_{TRR1} | Sampling-based | 1 | × | 1/4 | 2 | 99.9% | 1.85 - 2.03 |
| B7 | 19-06 | 8 | 2 | 16 | 8 | 20K | B_{TRR1} | Sampling-based | 1 | × | 1/4 | 2 | 99.9% | 31.14 |
| B8 | 18-03 | 4 | 1 | 16 | 8 | 43K | B_{TRR1} | Sampling-based | 1 | × | 1/4 | 2 | 99.9% | 2.57 |
| B9-12 | 19-48 | 8 | 1 | 16 | 8 | 42K-65K | B_{TRR2} | Sampling-based | 1 | × | 1/9 | 2 | 36.3% - 38.9% | 16.83 - 24.26 |
| B13-14 | 20-08 | 4 | 1 | 16 | 8 | 11K-14K | B_{TRR3} | Sampling-based | 1 | 1 | 1/2 | 4 | 99.9% | 16.20 - 18.12 |
| C0-3 | 16-48 | 4 | 1 | 16 | x8 | 137K-194K | C_{TRR1} | Mix | Unknown | 1 | 1/17 | 2 | 1.0% - 23.2% | 0.05 - 0.15 |
| C4-6 | 17-12 | 8 | 1 | 16 | x8 | 130K - 150K | C_{TRR1} | Mix | Unknown | 1 | 1/17 | 2 | 7.8% - 12.0% | 0.06 - 0.08 |
| C7-8 | 20-31 | 8 | 1 | 8 | x16 | 40K-44K | C_{TRR1} | Mix | Unknown | 1 | 1/17 | 2 | 39.8% - 41.8% | 9.66 - 14.56 |
| C9-11 | 20-31 | 8 | 1 | 8 | x16 | 42K-53K | C_{TRR2} | Mix | Unknown | 1 | 1/9 | 2 | 99.7% | 9.30 - 32.04 |
| C12-14 | 20-46 | 16 | 1 | 8 | x16 | 6K-7K | C_{TRR3} | Mix | Unknown | 1 | 1/8 | 2 | 99.9% | 4.91 - 12.64 |
| | | | | | | | | | | | | | | |

Google's Half-Double RowHammer Attack (May 2021)

Google Security Blog

The latest news and insights from Google on security and safety on the Internet

Introducing Half-Double: New hammering technique for DRAM Rowhammer bug

May 25, 2021

Research Team: Salman Qazi, Yoongu Kim, Nicolas Boichat, Eric Shiu & Mattias Nissler

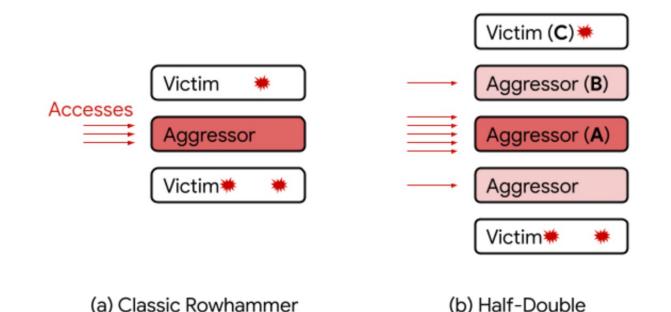
Today, we are sharing details around our discovery of Half-Double, a new Rowhammer technique that capitalizes on the worsening physics of some of the newer DRAM chips to alter the contents of memory.

Rowhammer is a DRAM vulnerability whereby repeated accesses to one address can tamper with the data stored at other addresses. Much like speculative execution vulnerabilities in CPUs, Rowhammer is a breach of the security guarantees made by the underlying hardware. As an electrical coupling phenomenon within the silicon itself, Rowhammer allows the potential bypass of hardware and software memory protection policies. This can allow untrusted code to break out of its sandbox and take full control of the system.

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https://security.googleblog.com/2021/05/introducing-half-double-new-hammering.html

Google's Half-Double RowHammer Attack (May 2021)



 Given three consecutive rows A, B, and C, we were able to attack C by directing a very large number of accesses to A, along with just a handful (~dozens) to B.

- Based on our experiments, accesses to B have a non-linear gating effect, in which they appear to "transport" the Rowhammer effect of A onto C.
- This is likely an indication that the electrical coupling responsible for Rowhammer is a property of distance, effectively becoming stronger and longer-ranged as cell geometries shrink down.

Google's Half-Double RowHammer Attack

Appears at USENIX Security 2022

Half-Double: Hammering From the Next Row Over

 Andreas Kogler¹ Jonas Juffinger^{1,2} Salman Qazi³ Yoongu Kim³ Moritz Lipp^{4*} Nicolas Boichat³ Eric Shiu⁵ Mattias Nissler³ Daniel Gruss¹
 ¹Graz University of Technology ²Lamarr Security Research ³Google ⁴Amazon Web Services ⁵Rivos

BlockHammer Solution in 2021

A. Giray Yaglikci, Minesh Patel, Jeremie S. Kim, Roknoddin Azizi, Ataberk Olgun, Lois Orosa, Hasan Hassan, Jisung Park, Konstantinos Kanellopoulos, Taha Shahroodi, Saugata Ghose, and Onur Mutlu, "BlockHammer: Preventing RowHammer at Low Cost by Blacklisting Rapidly-Accessed DRAM Rows" Proceedings of the 27th International Symposium on High-Performance Computer Architecture (HPCA), Virtual, February-March 2021. [Slides (pptx) (pdf)] [Short Talk Slides (pptx) (pdf)] [Intel Hardware Security Academic Awards Short Talk Slides (pptx) (pdf)] [Talk Video (22 minutes)] [Short Talk Video (7 minutes)] [Intel Hardware Security Academic Awards Short Talk Video (2 minutes)] [BlockHammer Source Code] Intel Hardware Security Academic Award Finalist (one of 4 finalists out of 34 nominations)

BlockHammer: Preventing RowHammer at Low Cost by Blacklisting Rapidly-Accessed DRAM Rows

A. Giray Yağlıkçı¹ Minesh Patel¹ Jeremie S. Kim¹ Roknoddin Azizi¹ Ataberk Olgun¹ Lois Orosa¹ Hasan Hassan¹ Jisung Park¹ Konstantinos Kanellopoulos¹ Taha Shahroodi¹ Saugata Ghose² Onur Mutlu¹ ¹ETH Zürich ²University of Illinois at Urbana–Champaign

Two Key Challenges

1 Scalability with worsening RowHammer vulnerability

2 Compatibility with commodity DRAM chips



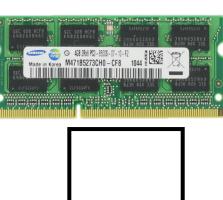
BlockHammer: Practical Throttling-based Mechanism

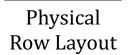
SLOW





- BlockHammer detects a RowHammer attack using area-efficient Bloom filters
- BlockHammer selectively throttles accesses from within the memory controller
- Bit flips **do not** occur





Row A

• BlockHammer can optionally inform the system software about the attack

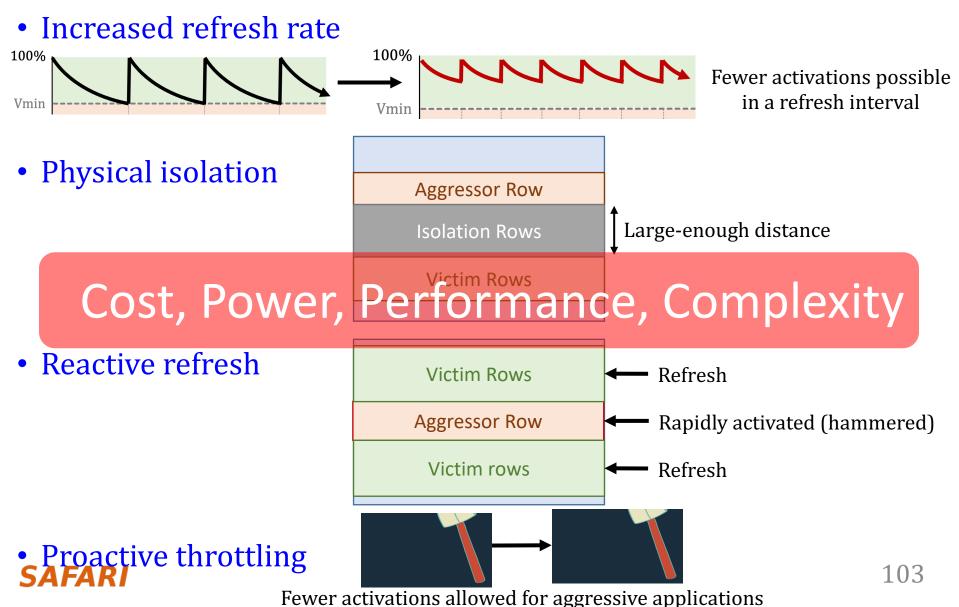
BlockHammer is compatible with commodity DRAM chips No need for proprietary info of or modifications to DRAM chips



Main Memory Needs **Intelligent Controllers** for Security, Safety, Reliability, Scaling

RowHammer Solution Approaches

• More robust DRAM chips **and/or** error-correcting codes



RowHammer in 2023: SK Hynix

ISSCC 2023 / SESSION 28 / HIGH-DENSITY MEMORIES

28.8 A 1.1V 16Gb DDR5 DRAM with Probabilistic-Aggressor Tracking, Refresh-Management Functionality, Per-Row Hammer Tracking, a Multi-Step Precharge, and Core-Bias Modulation for Security and Reliability Enhancement

Woongrae Kim, Chulmoon Jung, Seongnyuh Yoo, Duckhwa Hong, Jeongjin Hwang, Jungmin Yoon, Ohyong Jung, Joonwoo Choi, Sanga Hyun, Mankeun Kang, Sangho Lee, Dohong Kim, Sanghyun Ku, Donhyun Choi, Nogeun Joo, Sangwoo Yoon, Junseok Noh, Byeongyong Go, Cheolhoe Kim, Sunil Hwang, Mihyun Hwang, Seol-Min Yi, Hyungmin Kim, Sanghyuk Heo, Yeonsu Jang, Kyoungchul Jang, Shinho Chu, Yoonna Oh, Kwidong Kim, Junghyun Kim, Soohwan Kim, Jeongtae Hwang, Sangil Park, Junphyo Lee, Inchul Jeong, Joohwan Cho, Jonghwan Kim

SK hynix Semiconductor, Icheon, Korea

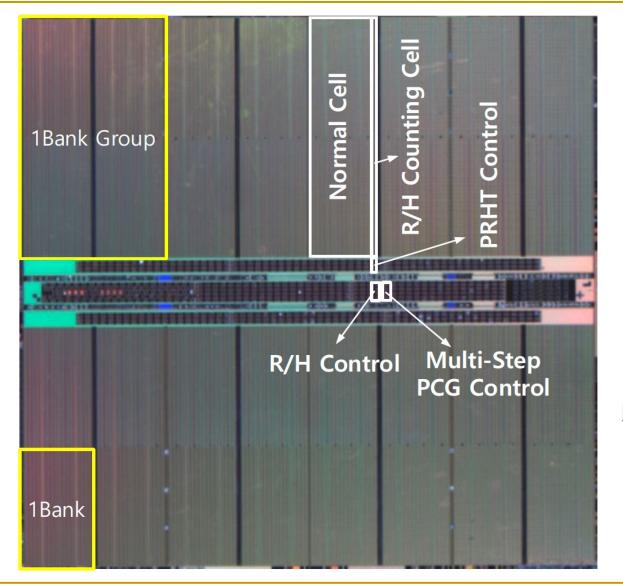


Industry's RowHammer Solutions (I)

SK hynix Semiconductor, Icheon, Korea

DRAM products have been recently adopted in a wide range of high-performance computing applications: such as in cloud computing, in big data systems, and IoT devices. This demand creates larger memory capacity requirements, thereby requiring aggressive DRAM technology node scaling to reduce the cost per bit [1,2]. However, DRAM manufacturers are facing technology scaling challenges due to row hammer and refresh retention time beyond 1a-nm [2]. Row hammer is a failure mechanism, where repeatedly activating a DRAM row disturbs data in adjacent rows. Scaling down severely threatens reliability since a reduction of DRAM cell size leads to a reduction in the intrinsic row hammer tolerance [2,3]. To improve row hammer tolerance, there is a need to probabilistically activate adjacent rows with carefully sampled active addresses and to improve intrinsic row hammer tolerance [2]. In this paper, row-hammer-protection and refresh-management schemes are presented to guarantee DRAM security and reliability despite the aggressive scaling from 1a-nm to sub 10-nm nodes. The probabilisticaggressor-tracking scheme with a refresh-management function (RFM) and per-row hammer tracking (PRHT) improve DRAM resilience. A multi-step precharge reinforces intrinsic row-hammer tolerance and a core-bias modulation improves retention time: even in the face of cell-transistor degradation due to technology scaling. This comprehensive scheme leads to a reduced probability of failure, due to row hammer attacks, by 93.1% and an improvement in retention time by 17%.

Industry's RowHammer Solutions (II)



ISSCC 2023 / SESSION 28 / HIGH-DENSITY MEMORIES

28.8 A 1.1V 16Gb DDR5 DRAM with Probabilistic-Aggressor Tracking, Refresh-Management Functionality, Per-Row Hammer Tracking, a Multi-Step Precharge, and Core-Bias Modulation for Security and Reliability Enhancement

Woongrae Kim, Chulmoon Jung, Seongnyuh Yoo, Duckhwa Hong, Jeongjin Hwang, Jungmin Yoon, Ohyong Jung, Joonwoo Choi, Sanga Hyun, Mankeun Kang, Sangho Lee, Dohong Kim, Sanghyun Ku, Donhyun Choi, Nogeun Joo, Sangwoo Yoon, Junseok Noh, Byeongyong Go, Cheolhoe Kim, Sunil Hwang, Mihyun Hwang, Seol-Min Yi, Hyungmin Kim, Sanghyuk Heo, Yeonsu Jang, Kyoungchul Jang, Shinho Chu, Yoonna Oh, Kwidong Kim, Junghyun Kim, Soohwan Kim, Jeongtae Hwang, Sangil Park, Junphyo Lee, Inchul Jeong, Joohwan Cho, Jonghwan Kim

SK hynix Semiconductor, Icheon, Korea

DSAC: Low-Cost Rowhammer Mitigation Using In-DRAM Stochastic and Approximate Counting Algorithm

Seungki Hong Dongha Kim Jaehyung Lee Reum Oh Changsik Yoo Sangjoon Hwang Jooyoung Lee

DRAM Design Team, Memory Division, Samsung Electronics

https://arxiv.org/pdf/2302.03591v1.pdf



Are we now RowHammer-free in 2023 and Beyond?

Are We Now RowHammer Free in 2023?

Appeared at ISCA in June 2023

RowPress: Amplifying Read-Disturbance in Modern DRAM Chips

Haocong Luo Ataberk Olgun A. Giray Yağlıkçı Yahya Can Tuğrul Steve Rhyner Meryem Banu Cavlak Joël Lindegger Mohammad Sadrosadati Onur Mutlu *ETH Zürich*

https://arxiv.org/pdf/2306.17061.pdf

SAFARI https://github.com/CMU-SAFARI/RowPress

RowPress

RowPress [ISCA 2023]



 Haocong Luo, Ataberk Olgun, Giray Yaglikci, Yahya Can Tugrul, Steve Rhyner, M. Banu Cavlak, Joel Lindegger, Mohammad Sadrosadati, and Onur Mutlu, "RowPress: Amplifying Read Disturbance in Modern DRAM Chips" Proceedings of the 50th International Symposium on Computer Architecture (ISCA), Orlando, FL, USA, June 2023.
 [Slides (pptx) (pdf)]
 [Lightning Talk Slides (pptx) (pdf)]
 [Lightning Talk Video (3 minutes)]
 [RowPress Source Code and Datasets (Officially Artifact Evaluated with All Badges)]
 Officially artifact evaluated as available, reusable and reproducible. Best artifact award at ISCA 2023.

RowPress: Amplifying Read-Disturbance in Modern DRAM Chips

Haocong Luo Ataberk Olgun A. Giray Yağlıkçı Yahya Can Tuğrul Steve Rhyner Meryem Banu Cavlak Joël Lindegger Mohammad Sadrosadati Onur Mutlu *ETH Zürich*



RowPress Amplifying Read Disturbance in Modern DRAM Chips

ISCA 2023 Session 2B: Monday 19 June, 2:15 PM EDT

Haocong Luo

Ataberk Olgun

A. Giray Yağlıkçı Yahya Can Tuğrul Steve Rhyner Meryem Banu Cavlak Joël Lindegger Mohammad Sadrosadati Onur Mutlu





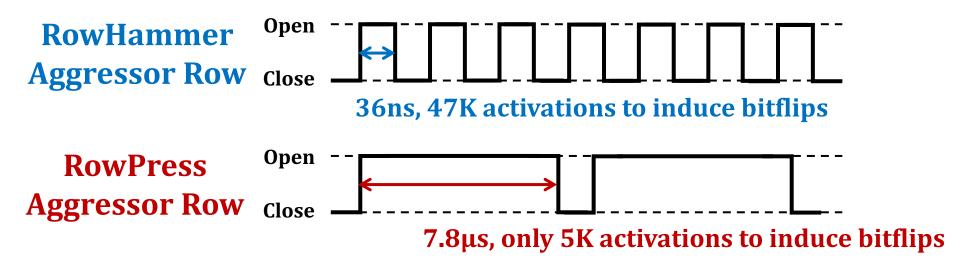
High-Level Summary

- We demonstrate and analyze RowPress, a new read disturbance phenomenon that causes bitflips in real DRAM chips
- We show that RowPress is **different from the RowHammer vulnerability**
- We demonstrate RowPress **using a user-level program** on a real Intel system with real DRAM chips
- We provide **effective solutions** to RowPress

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RowPress vs. RowHammer

Instead of using a high activation count, increase the time that the aggressor row stays open



We observe bitflips even with **ONLY ONE activation** in extreme cases where the row stays open for 30ms

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Major Takeaways from Real DRAM Chips

RowPress significantly **amplifies** DRAM's vulnerability to read disturbance

RowPress has a **different** underlying error **mechanism** from RowHammer



Real-System Demonstration (I)



Intel Core i5-10400

(Comet Lake)

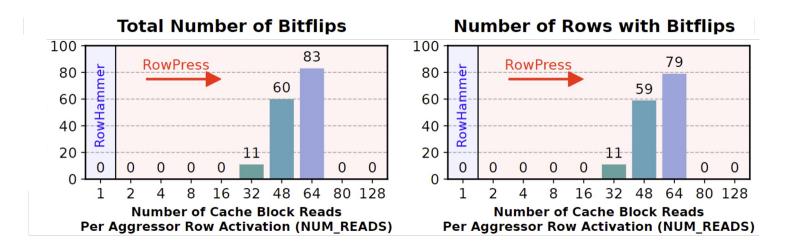


Samsung DDR4 Module M378A2K43CB1-CTD (Date Code: 20-10) w/ TRR RowHammer Mitigation

Key Idea: A proof-of-concept RowPress program keeps a DRAM row open for a longer period by **keeping on accessing different cache blocks in the row**



Real-System Demonstration (II)



Key Idea: A proof-of-concept RowPress program keeps a DRAM row open for a longer period by **keeping on accessing different cache blocks in the row**

Leveraging RowPress, a user-level program induces bitflips when RowHammer cannot

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More Results & Source Code

Many more results & analyses in the paper

- ➢ 6 major takeaways
- ▶ 19 major empirical observations
- ➢ 3 more potential mitigations



Fully open source and artifact evaluated

https://github.com/CMU-SAFARI/RowPress



SAFARI https://arxiv.org/pdf/2306.17061.pdf

RowPress [ISCA 2023]



 Haocong Luo, Ataberk Olgun, Giray Yaglikci, Yahya Can Tugrul, Steve Rhyner, M. Banu Cavlak, Joel Lindegger, Mohammad Sadrosadati, and Onur Mutlu, "RowPress: Amplifying Read Disturbance in Modern DRAM Chips" Proceedings of the 50th International Symposium on Computer Architecture (ISCA), Orlando, FL, USA, June 2023.
 [Slides (pptx) (pdf)]
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 [RowPress Source Code and Datasets (Officially Artifact Evaluated with All Badges)]
 Officially artifact evaluated as available, reusable and reproducible. Best artifact award at ISCA 2023.

RowPress: Amplifying Read-Disturbance in Modern DRAM Chips

Haocong Luo Ataberk Olgun A. Giray Yağlıkçı Yahya Can Tuğrul Steve Rhyner Meryem Banu Cavlak Joël Lindegger Mohammad Sadrosadati Onur Mutlu *ETH Zürich*

More to Come...

Two Major Directions

Understanding Bitflips (Hardware errors in general)

- Many effects on bitflips still need to be rigorously examined
 - Aging of DRAM Chips
 - Environmental Conditions (e.g., Process, Voltage, Temperature)
 - Memory Access Patterns
 - Memory Controller & System Design Decisions
 - **...**

Solving Bitflips (Hardware errors in general)

- Flexible and efficient solutions are necessary
 - In-field patchable / reconfigurable / programmable solutions
- Co-architecting across the system stack/components is important
 - To avoid performance and denial-of-service problems

A RowHammer Survey: Recent Update

 Onur Mutlu, Ataberk Olgun, and A. Giray Yaglikci, "Fundamentally Understanding and Solving RowHammer" Invited Special Session Paper at the <u>28th Asia and South Pacific Design</u> Automation Conference (ASP-DAC), Tokyo, Japan, January 2023. [arXiv version] [Slides (pptx) (pdf)] [Talk Video (26 minutes)]

Fundamentally Understanding and Solving RowHammer

Onur Mutlu onur.mutlu@safari.ethz.ch ETH Zürich Zürich, Switzerland Ataberk Olgun ataberk.olgun@safari.ethz.ch ETH Zürich Zürich, Switzerland A. Giray Yağlıkcı giray.yaglikci@safari.ethz.ch ETH Zürich Zürich, Switzerland

https://arxiv.org/pdf/2211.07613.pdf

A Case for Transparent Reliability in DRAM Systems

Minesh Patel[†] Taha Shahroodi^{‡†} Aditya Manglik[†] A. Giray Yağlıkçı[†] Ataberk Olgun[†] Haocong Luo[†] Onur Mutlu[†] [†]ETH Zürich [‡]TU Delft

https://arxiv.org/pdf/2204.10378.pdf

A Case for Self-Managing DRAM Chips: Improving Performance, Efficiency, Reliability, and Security via Autonomous in-DRAM Maintenance Operations

Hasan Hassan

Ataberk Olgun

A. Giray Yağlıkçı

Haocong Luo Onur Mutlu

ETH Zürich

https://arxiv.org/pdf/2207.13358.pdf

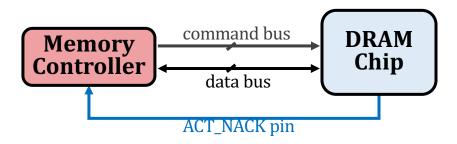
Self-Managing DRAM: Overview

Self-Managing DRAM (SMD)

enables autonomous in-DRAM maintenance operations

Key Idea:

Prevent the memory controller from accessing DRAM regions that are under maintenance by rejecting row activation (ACT) commands



Leveraging the ability to *reject an ACT*, a maintenance operation can be implemented *completely* within a DRAM chip

SAFAR

SMD-Based Maintenance Mechanisms

DRAM Refresh

Fixed Rate (SMD-FR)

uniformly refreshes *all* DRAM rows with a *fixed* refresh period

Variable Rate (SMD-VR)

skips refreshing rows that can **retain their data for longer** than the default refresh period

Probabilistic (SMD-PRP)

RowHammer Protection

Performs **neighbor row refresh** with **a small probability** on every row activation

Deterministic (SMD-DRP)

keeps track of most *frequently activated* rows and performs *neighbor* row refresh when activation count threshold is exceeded



Periodic Scrubbing (SMD-MS) periodically scans the entire DRAM for errors and corrects them

Self-Managing DRAM: Summary

The three major DRAM maintenance operations:

- *Refresh
- RowHammer Protection
- Memory Scrubbing

Implementing new **maintenance mechanisms** often requires difficult-to-realize changes

Our Goal

Ease the process of enabling new DRAM maintenance operations

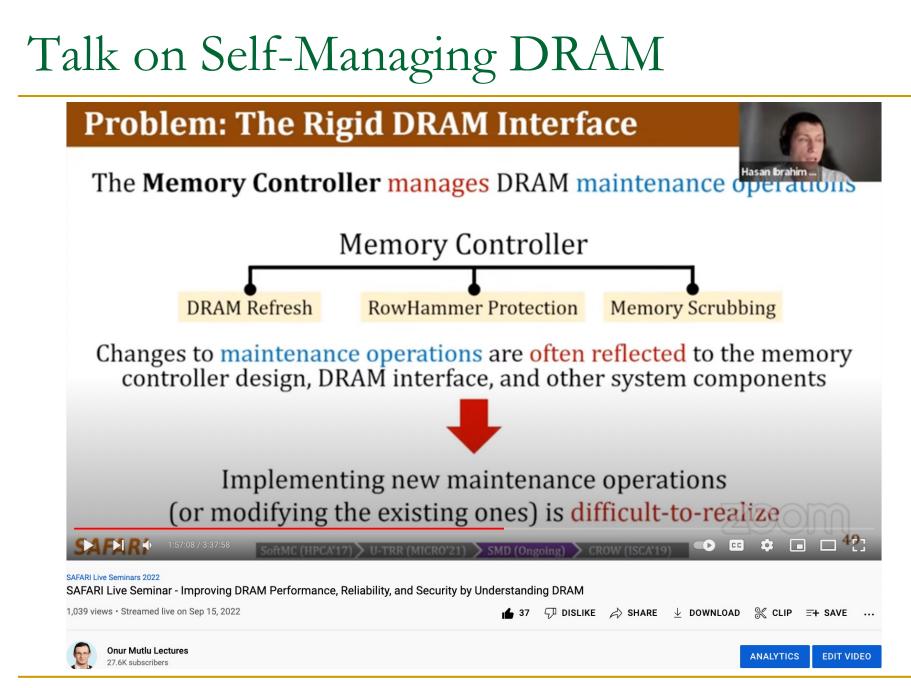
Enable more efficient in-DRAM maintenance operations

Self-Managing DRAM (SMD)

Enables implementing new **in-DRAM** maintenance mechanisms with **no further changes** in the *DRAM interface* and *memory controller*

SMD-based *refresh*, *RowHammer protection*, and *scrubbing* achieve 9.2% speedup and 6.2% lower DRAM energy vs. conventional DRAM





https://www.youtube.com/watch?v=mGa6-vpExbE

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Future Memory Reliability/Security Challenges

Collapse of the "Galloping Gertie" (1940)





Another Example (1994)



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Yet Another Example (2007)



Source: Morry Gash/AP, https://www.npr.org/2017/08/01/540669701/10-years-after-bridge-collapse-america-is-still-crumbling?t=1535427165809

A More Recent Example (2018)









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https://usa.streetsblog.org/2022/01/28/pittsburgh-bridge-collapse-underscores-urgent-need-for-fix-it-first-policy/



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https://www.npr.org/2022/01/28/1076343656/pittsburgh-bridge-collapse-biden-visit

Intelligent Memory Controllers Can Avoid Such Failures



Main Memory Needs **Intelligent Controllers** for Security, Safety, Reliability, Scaling

Using Memory Errors to Attack a Virtual Machine

Sudhakar Govindavajhala * Andrew W. Appel Princeton University {sudhakar,appel}@cs.princeton.edu

We present an experimental study showing that soft memory errors can lead to serious security vulnerabilities in Java and .NET virtual machines, or in any system that relies on type-checking of untrusted programs as a protection mechanism. Our attack works by sending to the JVM a Java program that is designed so that almost any memory error in its address space will allow it to take control of the JVM. All conventional Java and .NET virtual machines are vulnerable to this attack. The technique of the attack is broadly applicable against other language-based security schemes such as proof-carrying code.

We measured the attack on two commercial Java Virtual Machines: Sun's and IBM's. We show that a singlebit error in the Java program's data space can be exploited to execute arbitrary code with a probability of about 70%, and multiple-bit errors with a lower probability.

Our attack is particularly relevant against smart cards or tamper-resistant computers, where the user has physical access (to the outside of the computer) and can use various means to induce faults; we have successfully used heat. Fortunately, there are some straightforward defenses against this attack.

7 Physical fault injection

If the attacker has physical access to the outside of the machine, as in the case of a smart card or other tamperresistant computer, the attacker can induce memory errors. We considered attacks on boxes in form factors ranging from a credit card to a palmtop to a desktop PC.

We considered several ways in which the attacker could induce errors.⁴

IEEE S&P 2003

Before RowHammer (II)

Using Memory Errors to Attack a Virtual Machine

Sudhakar Govindavajhala * Andrew W. Appel Princeton University {sudhakar,appel}@cs.princeton.edu



Figure 3. Experimental setup to induce memory errors, showing a PC built from surplus components, clip-on gooseneck lamp, 50-watt spotlight bulb, and digital thermometer. Not shown is the variable AC power supply for the lamp.

IEEE S&P 2003

https://www.cs.princeton.edu/~appel/papers/memerr.pdf

A simple, exploitable memory error can be induced by software



Acknowledgments

SAFARI Research Group safari.ethz.ch



SAFARI Research Group

Computer architecture, HW/SW, systems, bioinformatics, security, memory

https://safari.ethz.ch/safari-newsletter-january-2021/



SAFARI Research Group: December 2021

<u>https://safari.ethz.ch/safari-newsletter-december-2021/</u>



Think Big, Aim High



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View in your browser December 2021



SAFARI Newsletter June 2023 Edition

<u>https://safari.ethz.ch/safari-newsletter-june-2023/</u>



Think Big, Aim High



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SAFARI Introduction & Research

Computer architecture, HW/SW, systems, bioinformatics, security, memory



Seminar in Computer Architecture - Lecture 5: Potpourri of Research Topics (Spring 2023)



Open Source Tools: SAFARI GitHub SAFARI Research Group at ETH Zurich and Carnegie Mellon University. Site for source code and tools distribution from SAFARI Research Group at ETH Zurich and Carnegie Mellon University.

Projects

Overview

Repositories 80

📯 241 followers 📀 ETH Zurich and Carnegie Mellon U... 🧬 https://safari.ethz.ch/ 🏹 omutlu@gmail.com

8 People 13

Packages

Pinned Customize pins aramulator Public ... prim-benchmarks Public A Fast and Extensible DRAM Simulator, with built-in support for PrIM (Processing-In-Memory benchmarks) is the first benchmark suite modeling many different DRAM technologies including DDRx, LPDDRx, for a real-world processing-in-memory (PIM) architecture. PrIM is GDDRx, WIOx, HBMx, and various academic proposals. Described in developed to evaluate, analyze, and characterize the first publ... the... ● C++ ☆ 440 ¥ 194 ¥ 38 MQSim Public **rowhammer** Public MQSim is a fast and accurate simulator modeling the performance of Source code for testing the Row Hammer error mechanism in DRAM modern multi-queue (MQ) SSDs as well as traditional SATA based devices. Described in the ISCA 2014 paper by Kim et al. at SSDs. MQSim faithfully models new high-bandwidth protocol http://users.ece.cmu.edu/~omutlu/pub/dram-row-hammer_isca14.pdf. implement... ● C++ ☆ 213 ♀ 121 · 208 ¥ 41 C C ... SoftMC Public ... SparseP (Public SoftMC is an experimental FPGA-based memory controller design that SparseP is the first open-source Sparse Matrix Vector Multiplication (SpMV) software package for real-world Processing-In-Memory (PIM) can be used to develop tests for DDR3 SODIMMs using a C++ based architectures. SparseP is developed to evaluate and characteri... API. The design, the interface, and its capabilities and limitatio... ●C ☆63 ¥11 Veriloa ☆ 104 % 26

https://github.com/CMU-SAFARI/

RowHammer, RowPress & Beyond Can We Be Free of Bitflips (Soon)?

Onur Mutlu

omutlu@gmail.com

https://people.inf.ethz.ch/omutlu

2 October 2023

Monte Verita SW Abstractions for HW Security Workshop

SAFARI

ETH zürich

Carnegie Mellon

Securing the Memory System The Story of RowHammer

Onur Mutlu

omutlu@gmail.com

https://people.inf.ethz.ch/omutlu

2 October 2023

Monte Verita SW Abstractions for HW Security Workshop

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ETH zürich

Carnegie Mellon

DDCA (Spring 2022)

Spring 2022 Edition:

https://safari.ethz.ch/digitaltechnik/spring2022/do ku.php?id=schedule

Spring 2021 Edition:

https://safari.ethz.ch/digitaltechnik/spring2021/do ku.php?id=schedule

Youtube Livestream (Spring 2022):

https://www.youtube.com/watch?v=cpXdE3HwvK 0&list=PL5Q2soXY2Zi97Ya5DEUpMpO2bbAoaG7c6

Youtube Livestream (Spring 2021):

- https://www.youtube.com/watch?v=LbC0EZY8yw 4&list=PL5O2soXY2Zi uei3aY39YB5pfW4SJ7LIN
- Bachelor's course
 - 2nd semester at ETH Zurich
 - Rigorous introduction into "How Computers Work"
 - Digital Design/Logic
 - **Computer Architecture**
 - 10 FPGA Lab Assignments

https://www.youtube.com/onurmutlulectures



Digital Design and Computer Architecture -Spring 2021

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sched

Trace: - schedule

Lectures/Schedule Lecture Buzzwords

Announcements

Materials

Lecture Video Plavlist on YouTube

Subject to Livestream Lecture Playlist





Spring 2021 Lectures/Schedule

| Week | Date | Livestream | Lecture | Readings | Lab | HW |
|-----------------|----------------|---|--|------------------------------------|-----|----|
| W1 | 25.02 Thu. | You Tube Live | L1: Introduction and Basics | Required Suggested Mentioned | | |
| 26.02 Y Fri. | 2 YouTube Live | L2a: Tradeoffs, Metrics, Mindset | Required | | | |
| | | L2b: Mysteries in Computer Architecture | Required Mentioned | | | |
| W2 | 04.03 Thu. | You Tube Live | L3a: Mysteries in Computer Architecture II | Required Suggested Mentioned | | |

 Readings Ontional HWs Lahs Extra Assignments Exams Technical Docs Resources Computer Architecture (CMU) SS15: Lecture Videos Somputer Architecture (CMU) SS15: Course Website S Digitaltechnik SS18: Lecture Videos

- Digitaltechnik SS18: Course Website S Digitaltechnik SS19: Lecture Videos
- Sigitaltechnik SS19: Course Website
- Digitaltechnik SS20: Lecture Videos
- Source State Course Signature Signat Website
 - S Moodle
- Secorded Lecture Plavlist

Comp Arch (Fall 2022)

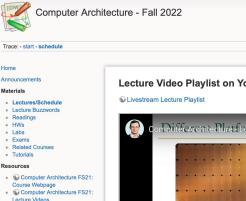
Fall 2022 Edition:

- https://safari.ethz.ch/architecture/fall2022/doku. php?id=schedule
- Fall 2021 Edition:
 - https://safari.ethz.ch/architecture/fall2021/doku. php?id=schedule

Youtube Livestream (2022):

- https://www.youtube.com/watch?v=4yfkM 5EFq o&list=PL5Q2soXY2Zi-Mnk1PxjEIG32HAGILkTOF
- Youtube Livestream (2021):
 - https://www.youtube.com/watch?v=4yfkM 5EFq o&list=PL5O2soXY2Zi-Mnk1PxjEIG32HAGILkTOF
- Master's level course
 - Taken by Bachelor's/Masters/PhD students
 - Cutting-edge research topics + fundamentals in **Computer Architecture**
 - 5 Simulator-based Lab Assignments
 - Potential research exploration
 - Many research readings

https://www.youtube.com/onurmutlulectures



Materials

= HWs

Labs

Resources

Webpage

Videos Moodle

- SHotCRP

(HDLBits)





Secture Playlist from Fall 2021



Fall 2022 Lectures & Schedule

| Week | Date | Livestream | Lecture | Readings | Lab | нw |
|---------------|---------------|---|--|------------------------|-----------------|----------------|
| W1 | 29.09 Thu. | You Tube Live | L1: Introduction and Basics | Required Mentioned | Lab 1 Out | HW 0 Out |
| 30.09 Fri. | You Tube Live | L2a: Memory Systems: Challenges and Opportunities a (PDF) a (PPT) | Described Suggested | | | |
| | | | L2b: Course Info & Logistics (PDF) | | | |
| W2 | 06.10 Thu. | You Tube Live | L3: Processing using Memory (PDF) (PPT) | Described Suggested | | HW 1 Out |

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RowHammer & DRAM Exploration (Fall 2022)

Fall 2022 Edition:

https://safari.ethz.ch/projects and seminars/fall2 022/doku.php?id=softmc

Spring 2022 Edition:

https://safari.ethz.ch/projects and seminars/sprin q2022/doku.php?id=softmc

Youtube Livestream (Spring 2022):

https://www.youtube.com/watch?v=r5QxuoJWttg &list=PL5O2soXY2Zi 1trfCckr6PTN8WR72icUO

Bachelor's course

- Elective at ETH Zurich
- Introduction to DRAM organization & operation
- Tutorial on using FPGA-based infrastructure
- Verilog & C++
- Potential research exploration

Lecture Video Playlist on YouTube Lecture Playlist C SoftMC Course: Meeting 1: Logistics & Intro ... P&S SoftMC Understanding and Improving Modern DRAM Performance, Reliability, and Security with Hands-On Experiments Hasan Hassan

Prof. Onur Mutlu

llŕ

ETH Zürich

Watch on 🕞 YouTube

2022 Meetings/Schedule (Tentative)

| Week | Date | Livestream | Meeting | Learning Materials | Assignments |
|------|---------------|----------------|--|--|-------------|
| W0 | 23.02 Wed. | You Tube Video | P&S SoftMC Tutorial | SoftMC Tutorial Slides (PDF) (PPT) | |
| W1 | 08.03 Tue. | You Tube Video | M1: Logistics & Intro to DRAM and SoftMC (PDF) = (PPT) | Required Materials Recommended Materials | HW0 |
| W2 | 15.03 Tue. | You Tube Video | M2: Revisiting RowHammer | left (Paper PDF) | |
| W3 | 22.03 Tue. | You Tube Video | M3: Uncovering in-DRAM TRR & TRRespass @ (PDF) # (PPT) | | |
| W4 | 29.03 Tue. | You Tube Video | M4: Deeper Look Into RowHammer's Sensitivities @ (PDF) # (PPT) | | |
| W5 | 05.04 Tue. | You Tube Video | M5: QUAC-TRNG | | |
| W6 | 12.04 Tue. | You Tube Video | M6: PiDRAM | | |

https://www.youtube.com/onurmutlulectures

Exploration of Emerging Memory Systems (Fall 2022)

Fall 2022 Edition:

<u>https://safari.ethz.ch/projects_and_seminars/fall2</u> 022/doku.php?id=ramulator

Spring 2022 Edition:

https://safari.ethz.ch/projects and seminars/sprin g2022/doku.php?id=ramulator

Youtube Livestream (Spring 2022):

- https://www.youtube.com/watch?v=aMllXRQd3s&list=PL5Q2soXY2Zi_TlmLGw_Z8hBo292 5ZApqV
- Bachelor's course
 - Elective at ETH Zurich
 - Introduction to memory system simulation
 - Tutorial on using Ramulator
 - C++
 - Potential research exploration

Lecture Video Playlist on YouTube

Secture Playlist



2022 Meetings/Schedule (Tentative)

| Week | Date | Livestream | Meeting | Learning Materials | Assignments |
|------|---------------|----------------|--|-----------------------|-------------|
| W1 | 09.03 Wed. | You Tube Video | M1: Logistics & Intro to Simulating Memory Systems Using Ramulator (PDF) (PPT) | | HWO |
| W2 | 16.03 Fri. | You Tube Video | M2: Tutorial on Using Ramulator (PDF) (PPT) | | |
| W3 | 25.02 Fri. | You Tube Video | M3: BlockHammer | | |
| W4 | 01.04 Fri. | You Tube Video | M4: CLR-DRAM | | |
| W5 | 08.04 Fri. | You Tube Video | M5: SIMDRAM (PDF) and (PPT) | | |
| W6 | 29.04 Fri. | You Tube Video | M6: DAMOV (PDF) 200 (PPT) | | |
| W7 | 06.05 Fri. | You Tube Video | M7: Syncron (PDF) | | |

https://www.youtube.com/onurmutlulectures



Industry-Adopted Solutions Do Not Work

 Pietro Frigo, Emanuele Vannacci, Hasan Hassan, Victor van der Veen, Onur Mutlu, Cristiano Giuffrida, Herbert Bos, and Kaveh Razavi,
 "TRRespass: Exploiting the Many Sides of Target Row Refresh"
 Proceedings of the <u>41st IEEE Symposium on Security and Privacy</u> (S&P), San Francisco, CA, USA, May 2020.
 [Slides (pptx) (pdf)]
 [Lecture Slides (pptx) (pdf)]
 [Lecture Slides (pptx) (pdf)]
 [Lecture Video (17 minutes)]
 [Lecture Video (59 minutes)]
 [Source Code]
 [Web Article]
 Best paper award.
 Pwnie Award 2020 for Most Innovative Research. Pwnie Awards 2020

TRRespass: Exploiting the Many Sides of Target Row Refresh

Pietro Frigo^{*†} Emanuele Vannacci^{*†} Hasan Hassan[§] Victor van der Veen[¶] Onur Mutlu[§] Cristiano Giuffrida^{*} Herbert Bos^{*} Kaveh Razavi^{*}

*Vrije Universiteit Amsterdam

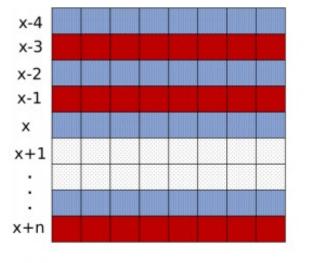
[§]ETH Zürich

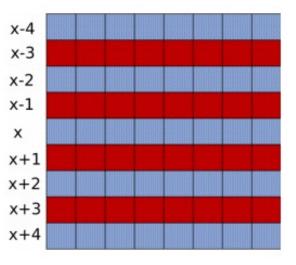
[¶]Qualcomm Technologies Inc.

TRRespass

- First work to show that TRR-protected DRAM chips are vulnerable to RowHammer in the field
 - Mitigations advertised as secure are not secure
- Introduces the Many-sided RowHammer attack
 - Idea: Hammer many rows to bypass TRR mitigations (e.g., by overflowing proprietary TRR tables that detect aggressor rows)
- (Partially) reverse-engineers the TRR and pTRR mitigation mechanisms implemented in DRAM chips and memory controllers
- Provides an automatic tool that can effectively create manysided RowHammer attacks in DDR4 and LPDDR4(X) chips

Example Many-Sided Hammering Patterns





(a) Assisted double-sided

(b) 4-sided

Fig. 12: Hammering patterns discovered by *TRRespass*. Aggressor rows are in red () and victim rows are in blue ().

TRRespass Vulnerable DRAM Modules

| 14-1-1- | Date | Freq. | Size | Or | ganizatio | n | MAG | Found | Dent Detterry | 0 | Corruption | ls | Double |
|-------------------------------|--------------------|-------|------|-------|-----------|------------|-----|-------------------|----------------|------------|----------------|-------------------|--------------|
| Module | (yy-ww) | (MHz) | (GB) | Ranks | Banks | Pins | MAC | Patterns | Best Pattern | Total | 1 ightarrow 0 | $0 \rightarrow 1$ | Refresh |
| $\mathcal{A}_{0,1,2,3}$ | 16-37 | 2132 | 4 | 1 | 16 | $\times 8$ | UL | 1 <u></u> | _ | | | | |
| \mathcal{A}_4 | 16-51 | 2132 | 4 | 1 | 16 | $\times 8$ | UL | 4 | 9-sided | 7956 | 4008 | 3948 | _ |
| \mathcal{A}_5 | 18-51 | 2400 | 4 | 1 | 8 | ×16 | UL | _ | | _ | | _ | |
| $A_{6,7}$ | 18-15 | 2666 | 4 | 1 | 8 | ×16 | UL | 10 - 10 L | | 81 10 | | 1.1 | 27 - 20 |
| \mathcal{A}_8 | 17-09 | 2400 | 8 | 1 | 16 | $\times 8$ | UL | 33 | 19-sided | 20808 | 10289 | 10519 | |
| \mathcal{A}_9 | 17-31 | 2400 | 8 | 1 | 16 | $\times 8$ | UL | 33 | 19-sided | 24854 | 12580 | 12274 | |
| \mathcal{A}_{10} | 19-02 | 2400 | 16 | 2 | 16 | $\times 8$ | UL | 488 | 10-sided | 11342 | 1809 | 11533 | \checkmark |
| \mathcal{A}_{11} | 19-02 | 2400 | 16 | 2 | 16 | $\times 8$ | UL | 523 | 10-sided | 12830 | 1682 | 11148 | ~ |
| $A_{12,13}$ | 18-50 | 2666 | 8 | 1 | 16 | $\times 8$ | UL | | | — | — | — | |
| \mathcal{A}_{14} | 19-08 | 3200 | 16 | 2 | 16 | $\times 8$ | UL | 120 | 14-sided | 32723 | 16490 | 16233 | |
| $\mathcal{A}_{15}^{\ddagger}$ | 17-08 | 2132 | 4 | 1 | 16 | $\times 8$ | UL | 2 | 9-sided | 22397 | 12351 | 10046 | _ |
| \mathcal{B}_0 | 18-11 | 2666 | 16 | 2 | 16 | $\times 8$ | UL | 2 | 3-sided | 17 | 10 | 7 | |
| \mathcal{B}_1 | 18-11 | 2666 | 16 | 2 | 16 | $\times 8$ | UL | 2 | 3-sided | 22 | 16 | 6 | — |
| \mathcal{B}_2 | 18-49 | 3000 | 16 | 2 | 16 | $\times 8$ | UL | 2 | 3-sided | 5 | 2 | 3 | |
| \mathcal{B}_3 | 19-08† | 3000 | 8 | 1 | 16 | $\times 8$ | UL | | 2 <u></u> | _ | | _ | _ |
| $\mathcal{B}_{4,5}$ | 19-08† | 2666 | 8 | 2 | 16 | $\times 8$ | UL | <u></u> | <u> </u> | <u> 19</u> | 0 | <u> </u> | <u> </u> |
| $\mathcal{B}_{6,7}$ | 19-08† | 2400 | 4 | 1 | 16 | $\times 8$ | UL | | <u></u> | | _ | _ | |
| \mathcal{B}_8^{\diamond} | 19-08† | 2400 | 8 | 1 | 16 | $\times 8$ | UL | | | - | | 10.00 | |
| \mathcal{B}_9^\diamond | 19-08† | 2400 | 8 | 1 | 16 | $\times 8$ | UL | 2 | 3-sided | 12 | | 12 | ~ |
| $\mathcal{B}_{10,11}$ | 16-13† | 2132 | 8 | 2 | 16 | $\times 8$ | UL | — | _ | _ | -2 | _ | _ |
| $\mathcal{C}_{0,1}$ | 18-46 | 2666 | 16 | 2 | 16 | $\times 8$ | UL | | | _ | 3 | _ | |
| $C_{2,3}$ | 19-08 [†] | 2800 | 4 | 1 | 16 | $\times 8$ | UL | <u> 10 - 10</u> - | <u> </u> | 2 <u></u> | | 1 <u>0 - 1</u> 1 | <u></u> |
| $C_{4,5}$ | 19-08† | 3000 | 8 | 1 | 16 | $\times 8$ | UL | _ | | | - | | _ |
| $C_{6,7}$ | 19-08 [†] | 3000 | 16 | 2 | 16 | $\times 8$ | UL | | | | | _ | _ |
| C_8 | 19-08† | 3200 | 16 | 2 | 16 | $\times 8$ | UL | | 5. | — | | — | _ |
| C_9 | 18-47 | 2666 | 16 | 2 | 16 | $\times 8$ | UL | _ | | _ | | _ | _ |
| $C_{10,11}$ | 19-04 | 2933 | 8 | 1 | 16 | $\times 8$ | UL | — | · · · · · | | | — | — |
| $\mathcal{C}_{12}^{\ddagger}$ | 15-01 [†] | 2132 | 4 | 1 | 16 | $\times 8$ | UT | 25 | 10-sided | 190037 | 63904 | 126133 | \checkmark |
| $\mathcal{C}_{13}^{\ddagger}$ | 18-49 | 2132 | 4 | 1 | 16 | $\times 8$ | UT | 3 | 9-sided | 694 | 239 | 455 | _ |

TABLE II: TRRespass results. We report the number of patterns found and bit flips detected for the 42 DRAM modules in our set.

The module does not report manufacturing date. Therefore, we report purchase date as an approximation. Analyzed using the FPGA-based SoftMC.

SAFA

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The system runs with double refresh frequency in standard conditions. We configured the refresh interval to be 64 ms in the BIOS settings.

TRRespass Vulnerable Mobile Phones

TABLE III: LPDDR4(X) results. Mobile phones tested against *TRRespass* on ARMv8 sorted by production date. We found bit flip inducing RowHammer patterns on 5 out of 13mobile phones.

| Mobile Phone | Year | SoC | Memory (GB) | Found Patterns |
|----------------------------|------|----------------|----------------|-------------------|
| Google Pixel | 2016 | MSM8996 | 4^{\dagger} | \checkmark |
| Google Pixel 2 | 2017 | MSM8998 | 4 | |
| Samsung G960F/DS | 2018 | Exynos 9810 | 4 | — |
| Huawei P20 DS | 2018 | Kirin 970 | 4 | |
| Sony XZ3 | 2018 | SDM845 | 4 | _ |
| HTC U12+ | 2018 | SDM845 | 6 | _ |
| LG G7 ThinQ | 2018 | SDM845 | 4† | \checkmark |
| Google Pixel 3 | 2018 | SDM845 | 4 | \checkmark |
| Google Pixel 4 | 2019 | SM8150 | 6 | _ |
| OnePlus 7 | 2019 | SM8150 | 8 | \checkmark |
| Samsung G970F/DS | 2019 | Exynos 9820 | 6 | \checkmark |
| Huawei P30 DS | 2019 | Kirin 980 | 6 | — |
| Xiaomi Redmi Note 8 Pro | 2019 | Helio G90T | 6 | _ |

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LPDDR4 (not LPDDR4X)

TRRespass Based RowHammer Attack

TABLE IV: Time to exploit. Time to find the first exploitable template on two sample modules from each DRAM vendor.

| Module | τ (ms) | <i>PTE</i> [81] | RSA-2048 [79] | sudo [27] |
|--------------------|-------------|-----------------|---------------|-----------|
| \mathcal{A}_{14} | 188.7 | 4.9s | 6m 27s | _ |
| \mathcal{A}_4 | 180.8 | 38.8s | 39m 28s | _ |
| \mathcal{B}_1 | 360.7 | _ | _ | _ |
| \mathcal{B}_2 | 331.2 | | — | — |
| \mathcal{C}_{12} | 300.0 | 2.3s | 74.6s | 54m16s |
| \mathcal{C}_{13} | 180.9 | 3h 15m | | _ |

 τ : Time to template a single row: time to fill the victim and aggressor rows + hammer time + time to scan the row.

TRRespass Key Results

- 13 out of 42 tested DDR4 DRAM modules are vulnerable
 - □ From all 3 major manufacturers
 - □ 3-, 9-, 10-, 14-, 19-sided hammer attacks needed
- 5 out of 13 mobile phones tested vulnerable
 - From 4 major manufacturers
 - With LPDDR4(X) DRAM chips
- These results are scratching the surface
 - TRRespass tool is not exhaustive
 - There is a lot of room for uncovering more vulnerable chips and phones

TRRespass Key Takeaways

RowHammer is still an open problem

Security by obscurity is likely not a good solution

Uncovering TRR Almost Completely

Industry-Adopted Solutions Are Very Poor

 Hasan Hassan, Yahya Can Tugrul, Jeremie S. Kim, Victor van der Veen, Kaveh Razavi, and Onur Mutlu,
 "Uncovering In-DRAM RowHammer Protection Mechanisms: A New Methodology, Custom RowHammer Patterns, and Implications" Proceedings of the <u>54th International Symposium on Microarchitecture</u> (MICRO), Virtual, October 2021.
 [Slides (pptx) (pdf)]
 [Short Talk Slides (pptx) (pdf)]
 [Lightning Talk Slides (pptx) (pdf)]
 [Lightning Talk Video (100 seconds)]
 [arXiv version]

Uncovering In-DRAM RowHammer Protection Mechanisms: A New Methodology, Custom RowHammer Patterns, and Implications

| Hasan Hassan † | Yahya Can Tuğrul ^{†‡} | Jeremie S. Kin | n^{\dagger} Victor van der Veen ^{σ} |
|---------------------------|---|----------------|--|
| | Kaveh Razavi † | Onur Mutlu | Ť |
| †ETH Zürich | [‡] TOBB University of Economics | & Technology | $^{\sigma}$ Qualcomm Technologies Inc. |

Key Takeaways

All 45 modules we test are vulnerable

99.9% of rows in a DRAM bank experience **at least one RowHammer bit flip**

ECC is ineffective: up to 7 RowHammer bit flips in an 8-byte dataword

| $ \begin{array}{c c c c c c c c c c c c c c c c c c c $ | |
|--|-------------------------------------|
| A1-519-368181613K-15K A_{TRR1} Counter-based16 \checkmark 1/9499.2% - 99.4%A6-719-458181613K-15K A_{TRR1} Counter-based16 \checkmark 1/9499.3% - 99.4%A8-920-078116812K-14K A_{TRR1} Counter-based16 \checkmark 1/9474.6% - 75.0%A10-1219-518116812K-13K A_{TRR1} Counter-based16 \checkmark 1/9474.6% - 75.0%A13-1420-318181611K-14K A_{TRR2} Counter-based16 \checkmark 1/9294.3% - 98.6%B018-224116844K B_{TRR1} Sampling-based1 X 1/4299.9% | Max. Bit Flips r Row per Hammer† |
| A1-519-368181613K-15K A_{TRR1} Counter-based16 \checkmark 1/9499.2% - 99.4%A6-719-458181613K-15K A_{TRR1} Counter-based16 \checkmark 1/9499.3% - 99.4%A8-920-078116812K-14K A_{TRR1} Counter-based16 \checkmark 1/9474.6% - 75.0%A10-1219-518116812K-13K A_{TRR1} Counter-based16 \checkmark 1/9474.6% - 75.0%A13-1420-318181611K-14K A_{TRR2} Counter-based16 \checkmark 1/9294.3% - 98.6%B018-224116844K B_{TRR1} Sampling-based1 X 1/4299.9% | 1.16 |
| A6-719-458181613K-15K A_{TRR1} Counter-based16 \checkmark 1/9499.3% - 99.4%A8-920-078116812K-14K A_{TRR1} Counter-based16 \checkmark 1/9474.6% - 75.0%A10-1219-518116812K-13K A_{TRR1} Counter-based16 \checkmark 1/9474.6% - 75.0%A13-1420-318181611K-14K A_{TRR2} Counter-based16 \checkmark 1/9294.3% - 98.6%B018-224116844K B_{TRR1} Sampling-based1 X 1/4299.9% | 2.32 - 4.73 |
| A10-12 19-51 8 1 16 8 12K-13K A_{TRR1} Counter-based 16 Image: V 1/9 4 74.6% - 75.0% A13-14 20-31 8 1 8 16 11K-14K A_{TRR2} Counter-based 16 Image: V 1/9 4 74.6% - 75.0% B0 18-22 4 1 16 8 44K B_{TRR1} Sampling-based 1 Image: V 1/4 2 99.9% | 2.12 - 3.86 |
| A13-14 20-31 8 1 8 16 11K-14K A_{TRR2} Counter-based 16 I/9 2 94.3% - 98.6% B0 18-22 4 1 16 8 44K B_{TRR1} Sampling-based 1 X 1/4 2 99.9% | 1.96 - 2.96 |
| B0 18-22 4 1 16 8 44K B_{TRR1} Sampling-based 1 X $1/4$ 2 99.9% | 1.48 - 2.86 |
| | 1.53 - 2.78 |
| | 2.13 |
| | 0.06 - 0.11 |
| B5-6 16-48 4 1 16 8 44K-50K B_{TRR1} Sampling-based 1 X 1/4 2 99.9% | 1.85 - 2.03 |
| B7 19-06 8 2 16 8 20K B_{TRR1} Sampling-based 1 X 1/4 2 99.9% | 31.14 |
| B8 18-03 4 1 16 8 43K B _{TRR1} Sampling-based 1 X 1/4 2 99.9% | 2.57 |
| B9-12 19-48 8 1 16 8 42K-65K B_{TRR2} Sampling-based 1 X 1/9 2 36.3% - 38.9% | 16.83 - 24.26 |
| B13-14 20-08 4 1 16 8 11K-14K B_{TRR3} Sampling-based 1 \checkmark 1/2 4 99.9% | 16.20 - 18.12 |
| C0-3 16-48 4 1 16 x8 137K-194K C_{TRR1} Mix Unknown \checkmark 1/17 2 1.0% - 23.2% | 0.05 - 0.15 |
| C4-6 17-12 8 1 16 x8 130K-150K C_{TRR1} Mix Unknown \checkmark 1/17 2 7.8% - 12.0% | 0.06 - 0.08 |
| C7-8 20-31 8 1 8 x16 40K-44K C _{TRR1} Mix Unknown ✓ 1/17 2 39.8% - 41.8% | 9.66 - 14.56 |
| C9-11 20-31 8 1 8 x16 42K-53K C _{TRR2} Mix Unknown ✓ 1/9 2 99.7% | 9.30 - 32.04 |
| C12-14 20-46 16 1 8 x16 6K-7K C _{TRR3} Mix Unknown 🖌 1/8 2 99.9% | 4.91 - 12.64 |

Google's Half-Double RowHammer Attack (May 2021)

Google Security Blog

The latest news and insights from Google on security and safety on the Internet

Introducing Half-Double: New hammering technique for DRAM Rowhammer bug

May 25, 2021

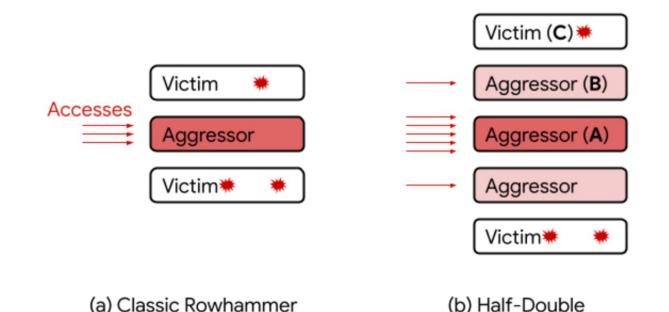
Research Team: Salman Qazi, Yoongu Kim, Nicolas Boichat, Eric Shiu & Mattias Nissler

Today, we are sharing details around our discovery of Half-Double, a new Rowhammer technique that capitalizes on the worsening physics of some of the newer DRAM chips to alter the contents of memory.

Rowhammer is a DRAM vulnerability whereby repeated accesses to one address can tamper with the data stored at other addresses. Much like speculative execution vulnerabilities in CPUs, Rowhammer is a breach of the security guarantees made by the underlying hardware. As an electrical coupling phenomenon within the silicon itself, Rowhammer allows the potential bypass of hardware and software memory protection policies. This can allow untrusted code to break out of its sandbox and take full control of the system.

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Google's Half-Double RowHammer Attack (May 2021)



- Given three consecutive rows A, B, and C, we were able to attack C by directing a very large number of accesses to A, along with just a handful (~dozens) to B.
- Based on our experiments, accesses to B have a non-linear gating effect, in which they appear to "transport" the Rowhammer effect of A onto C.
- This is likely an indication that the electrical coupling responsible for Rowhammer is a property of distance, effectively becoming stronger and longer-ranged as cell geometries shrink down.

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Google's Half-Double RowHammer Attack

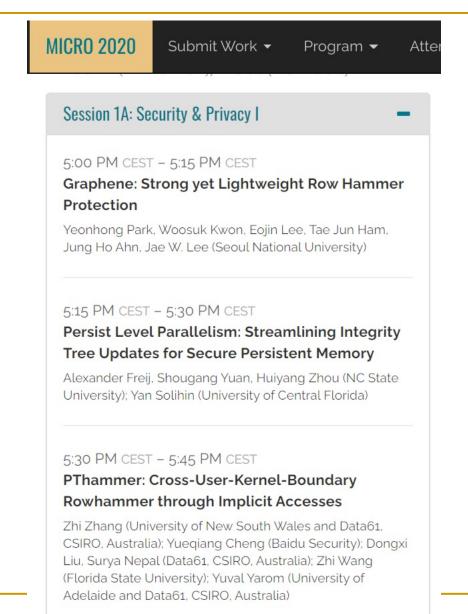
Appears at USENIX Security 2022

Half-Double: Hammering From the Next Row Over

 Andreas Kogler¹ Jonas Juffinger^{1,2} Salman Qazi³ Yoongu Kim³ Moritz Lipp^{4*} Nicolas Boichat³ Eric Shiu⁵ Mattias Nissler³ Daniel Gruss¹
 ¹Graz University of Technology ²Lamarr Security Research ³Google ⁴Amazon Web Services ⁵Rivos

More RowHammer in 2020-2023

RowHammer in 2020 (I)



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RowHammer in 2020 (II)

| S & P | 🖀 Home | Program - | Call For 🔻 | Attend - | Workshops 🔻 | | |
|-------|--|-----------------------------------|------------------------------------|------------------|--|--------------------|--|
| | Session #5 | : Rowhamn | ner | | | Room 2 | |
| | Session chair: Michael Franz (UC Irvine) | | | | | | |
| | RAMBleed: Reading Bits in Memory Without Accessing Them Andrew Kwong (University of Michigan), Daniel Genkin (University of Michigan), Daniel Gruss Data61) | | | | | | |
| | Are We Susceptible to Rowhammer? An End-to-End Methodology for Cloud Providers Lucian Cojocar (Microsoft Research), Jeremie Kim (ETH Zurich, CMU), Minesh Patel (ETH Zu (Microsoft Research), Onur Mutlu (ETH Zurich, CMU) | | | | | | |
| | Zhenkai Zh | ang (Texas Tec | h University), Z | ihao Zhan (Vande | owhammer Attacks erbilt University), Da s (Vanderbilt Univer | niel Balasubrar | |
| | Pietro Frigo Veen (Qual | o (Vrije Universi comm Technol | teit Amsterdam ogies, Inc.), On | ur Mutlu (ETH Zü | Refresh s), Emanuele Vanna rich), Cristiano Giuf dam, The Netherland | frida (Vrije Unive | |

RowHammer in 2020 (III)



DeepHammer: Depleting the Intelligence of Deep Neural Networks through Targeted Chain of Bit Flips Fan Yao, *University of Central Florida*; Adnan Siraj Rakin and Deliang Fan, *Arizona State University* AVAILABLE MEDIA 🗋 🗊 🕥 Show details 🕨

RowHammer in 2021 (I)

HotOS XVIII

The 18th Workshop on Hot Topics in Operating Systems

31 May 1 June–3 June 2021, Cyberspace, People's Couches, and Zoom

Stop! Hammer Time: Rethinking Our Approach to Rowhammer Mitigations

RowHammer in 2021 (II)



SMASH: Synchronized Many-sided Rowhammer Attacks from JavaScript

RowHammer in 2021 (III)



Session 10A: Security & Privacy III

Session Chair: Hoda Naghibijouybari (Binghamton)

9:00 PM CEST - 9:15 PM CEST

A Deeper Look into RowHammer's Sensitivities: Experimental Analysis of Real DRAM Chips and Implications on Future Attacks and Defenses

Lois Orosa, Abdullah Giray Yaglikci, Haocong Luo (ETH Zurich); Ataberk Olgun (TOBB University of Economics and Technology); Jisung Park, Hasan Hassan, Minesh Patel, Jeremie S. Kim, Onur Mutlu (ETH Zurich)

🖹 Paper

9:15 PM CEST - 9:30 PM CEST

Uncovering In-DRAM RowHammer Protection Mechanisms: A New Methodology, Custom RowHammer Patterns, and Implications

Hasan Hassan (ETH Zurich); Yahya Can Tugrul (TOBB University of Economics and Technology); Jeremie S. Kim (ETH Zurich); Victor van der Veen (Qualcomm); Kaveh Razavi, Onur Mutlu (ETH Zurich)

Paper

RowHammer in 2022 (I)

MAY 22-26, 2022 AT THE HYATT REGENCY, SAN FRANCISCO, CA 43rd IEEE Symposium on Security and Privacy

BLACKSMITH: Scalable Rowhammering in the Frequency Domain

SpecHammer: Combining Spectre and Rowhammer for New Speculative Attacks

PROTRR: Principled yet Optimal In-DRAM Target Row Refresh

DeepSteal: Advanced Model Extractions Leveraging Efficient Weight Stealing in Memories

RowHammer in 2022 (II)



Randomized Row-Swap: Mitigating Row Hammer by Breaking Spatial Correlation between Aggressor and Victim Rows

RowHammer in 2022 (III)

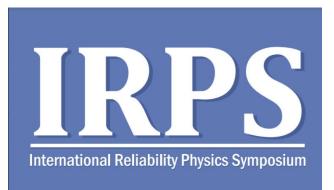
HPCA 2022

The 28th IEEE International Symposium on High-Performance Computer Architecture (HPCA-28), Seoul, South Korea

SafeGuard: Reducing the Security Risk from Row-Hammer via Low-Cost Integrity Protection

Mithril: Cooperative Row Hammer Protection on Commodity DRAM Leveraging Managed Refresh

RowHammer in 2022 (IV)



IRPS 2022

The Price of Secrecy: How Hiding Internal DRAM Topologies Hurts Rowhammer Defenses

Stefan Saroiu, Alec Wolman, Lucian Cojocar Microsoft

RowHammer in 2022 (V)

31ST USENIX SECURITY SYMPOSIUM AUGUST 10-12, 2022 BOSTON, MA, USA

Half-Double: Hammering From the Next Row Over

Andreas Kogler¹ Jonas Juffinger^{1,2} Salman Qazi³ Yoongu Kim³ Moritz Lipp^{4*} Nicolas Boichat³ Eric Shiu⁵ Mattias Nissler³ Daniel Gruss¹

¹Graz University of Technology ²Lamarr Security Research ³Google ⁴Amazon Web Services ⁵Rivos

RowHammer in 2022 (VI)



HAMMERSCOPE: Observing DRAM Power Consumption Using Rowhammer

When Frodo Flips: End-to-End Key Recovery on FrodoKEM via Rowhammer

RowHammer in 2022 (VII)



AQUA: Scalable Rowhammer Mitigation by Quarantining Aggressor Rows at Runtime

Anish Saxena, Gururaj Saileshwar (Georgia Institute of Technology); Prashant J. Nair (University of British Columbia); Moinuddin Qureshi (Georgia Institute of Technology)

HiRA: Hidden Row Activation for Reducing Refresh Latency of Off-the-Shelf DRAM Chips

Abdullah Giray Yaglikci (ETH Zürich); Ataberk Olgun (TOBB University of Economics and Technology); Lois Orosa, Minesh Patel, Haocong Luo, Hasan Hassan (ETH Zürich); Oguz Ergin (TOBB University of Economics and Technology); Onur Mutlu (ETH Zürich)

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RowHammer in 2022 (VII)

 A. Giray Yaglıkcı, Ataberk Olgun, Minesh Patel, Haocong Luo, Hasan Hassan, Lois Orosa, Oguz Ergin, and Onur Mutlu, "HiRA: Hidden Row Activation for Reducing Refresh Latency of Off-the-Shelf DRAM Chips" Proceedings of the <u>55th International Symposium on Microarchitecture</u> (MICRO), Chicago, IL, USA, October 2022.
 [Slides (pptx) (pdf)]
 [Longer Lecture Slides (pptx) (pdf)]
 [Lecture Video (36 minutes)]
 [arXiv version]

HiRA: Hidden Row Activation

for Reducing Refresh Latency of Off-the-Shelf DRAM Chips

A. Giray Yağlıkçı¹ Ataberk Olgun^{1,2} Minesh Patel¹ Haocong Luo¹ Hasan Hassan¹ Lois Orosa^{1,3} Oğuz Ergin² Onur Mutlu¹ ¹ETH Zürich ²TOBB University of Economics and Technology ³Galicia Supercomputing Center (CESGA)

https://arxiv.org/pdf/2209.10198.pdf

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A Case for Transparent Reliability in DRAM Systems

Minesh Patel[†] Taha Shahroodi^{‡†} Aditya Manglik[†] A. Giray Yağlıkçı[†] Ataberk Olgun[†] Haocong Luo[†] Onur Mutlu[†] [†]ETH Zürich [‡]TU Delft

https://arxiv.org/pdf/2204.10378.pdf

A Case for Self-Managing DRAM Chips: Improving Performance, Efficiency, Reliability, and Security via Autonomous in-DRAM Maintenance Operations

Hasan Hassan

Ataberk Olgun

A. Giray Yağlıkçı

Haocong Luo

ETH Zürich

Onur Mutlu

https://arxiv.org/pdf/2207.13358.pdf

RowHammer in 2023 (I)



Session 6C: Rowhammer and spectre

Bayview AB

11:00 AM - 12:15 PM

Session Chair: Eyal Ronen

REGA: Scalable Rowhammer Mitigation with Refresh-Generating Activations Michele Marazzi (ETH Zurich), Flavien Solt (ETH Zurich), Patrick Jattke (ETH Zurich), Kubo Takashi (Zentel Japan), Kaveh Razavi (ETH Zurich)

CSI:Rowhammer - Cryptographic Security and Integrity against Rowhammer

Jonas Juffinger (Lamarr Security Research, Graz University of Technology, Austria), Lukas Lamster (Graz University of Technology, Austria), Andreas Kogler (Graz University of Technology, Austria), Moritz Lipp (Amazon Web Services, Austria), Daniel Gruss (Graz University of Technology, Austria), Moritz Lipp (Amazon Web Services, Austria), Daniel Gruss (Graz University of Technology, Austria)

Jolt: Recovering TLS Signing Keys via Rowhammer Faults

Koksal Mus (Worcester Polytechnic Institute), Yarkın Doröz (Worcester Polytechnic Institute), M. Caner Tol (Worcester Polytechnic Institute), Kristi Rahman (Worcester Polytechnic Institute), Berk Sunar (Worcester Polytechnic Institute)

RowHammer in 2023 (II)

HPCA 2023

The 29th IEEE International Symposium on High-Performance Computer Architecture (HPCA-29)

Scalable and Secure Row-Swap: Efficient and Safe Row Hammer Mitigation in Memory Systems

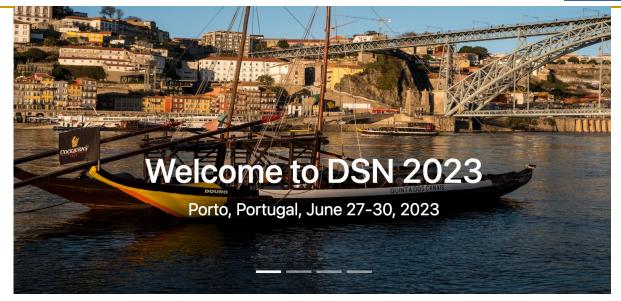
Jeonghyun Woo (University of British Columbia), Gururaj Saileshwar (Georgia Institute of Technology), Prashant J. Nair (University of British Columbia)

SHADOW: Preventing Row Hammer in DRAM with Intra-Subarray Row Shuffling Minbok Wi (Seoul National University), Jaehyun Park (Seoul National University), Seoyoung Ko (Seoul National University), Michael Jaemin Kim (Seoul National University), Nam Sung Kim (UIUC), Eojin Lee (Inha University), Jung Ho Ahn (Seoul National

University)

RowHammer in 2023 (V)





[28 June, 14:30-16:00] RT-3: Memory 1 (Session Chair: TBD)

Compiler-Implemented Differential Checksums: Effective Detection and Correction of Transient and Permanent Memory Errors (REG) *C. Borchert; H. Schirmeier; O. Spinczyk*

PT-Guard: Integrity-Protected Page Tables to Defend Against Breakthrough Rowhammer Attacks (REG)

A. Saxena; G. Saileshwar; J. Juffinger; A. Kogler; D. Gruss; M. Qureshi

Don't Knock! Rowhammer at the Backdoor of DNN Models (REG)

M. Tol; S. Islam; A. Adiletta; B. Sunar; Z. Zhang

[29 June, 16:00-17:30] DS23-4: Hardware Resilience and Human Factors (Session Chair: TBD)

An Experimental Analysis of RowHammer in HBM2 DRAM Chips

Ataberk Olgun, Majd Osseiran, Abdullah Giray Yaglikci, Yahya Can Tugrul, Juan Gomez Luna, Haocong Luo, Behzad Salami, Steve Rhyner and Onur Mutlu

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Future Memory Reliability/Security Challenges

Future of Main Memory Security

- DRAM is becoming less reliable \rightarrow more vulnerable
- Due to difficulties in DRAM scaling, other problems may also appear (or they may be going unnoticed)
- Some errors may already be slipping into the field
 - Read disturb errors (Rowhammer)
 - Retention errors
 - Read errors, write errors
 - ...

These errors can also pose security vulnerabilities

Future of Main Memory Security

DRAM

- Flash memory
- Emerging Technologies
 - Phase Change Memory
 - STT-MRAM
 - RRAM, memristors
 - ...



Main Memory Needs **Intelligent Controllers** for Security, Safety, Reliability, Scaling



Intelligent Memory Controllers **Can Avoid Many Failures** & Enable Better Scaling

Architecting Future Memory for Security

Understand: Methods for vulnerability modeling & discovery

- Modeling and prediction based on real (device) data and analysis
- Understanding vulnerabilities
- Developing reliable metrics

Architect: Principled architectures with security as key concern

- Good partitioning of duties across the stack
- Cannot give up performance and efficiency
- Patch-ability in the field

Design & Test: Principled design, automation, (online) testing

- Design for security
- High coverage and good interaction with system reliability methods

An Early Proposal for Intelligent Controllers [IMW'13]

Onur Mutlu, <u>"Memory Scaling: A Systems Architecture Perspective"</u> *Proceedings of the <u>5th International Memory</u>* <u>Workshop</u> (IMW), Monterey, CA, May 2013. <u>Slides</u> (pptx) (pdf) <u>EETimes Reprint</u>

Memory Scaling: A Systems Architecture Perspective

Onur Mutlu Carnegie Mellon University onur@cmu.edu http://users.ece.cmu.edu/~omutlu/

https://people.inf.ethz.ch/omutlu/pub/memory-scaling_memcon13.pdf

Industry Is Writing Papers About It, Too

DRAM Process Scaling Challenges

Refresh

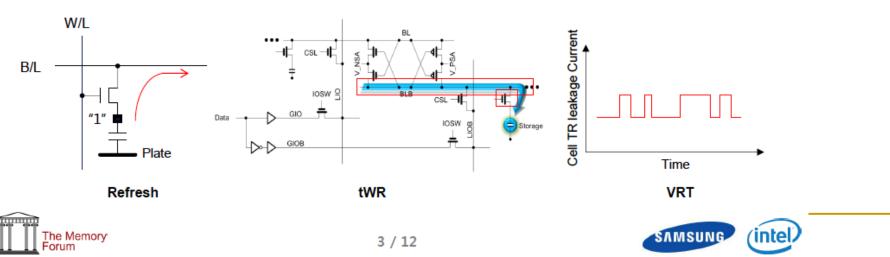
- · Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance
- · Leakage current of cell access transistors increasing

✤ tWR

- · Contact resistance between the cell capacitor and access transistor increasing
- · On-current of the cell access transistor decreasing
- · Bit-line resistance increasing

VRT

· Occurring more frequently with cell capacitance decreasing



197

Industry Is Writing Papers About It, Too

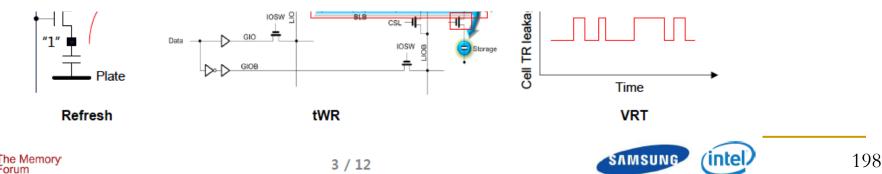
DRAM Process Scaling Challenges

* Refresh

Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance
THE MEMORY FORUM 2014

Co-Architecting Controllers and DRAM to Enhance DRAM Process Scaling

Uksong Kang, Hak-soo Yu, Churoo Park, *Hongzhong Zheng, **John Halbert, **Kuljit Bains, SeongJin Jang, and Joo Sun Choi



Samsung Electronics, Hwasung, Korea / *Samsung Electronics, San Jose / **Intel

Final Thoughts on RowHammer

Aside: Byzantine Failures

- This class of failures is known as Byzantine failures
- Characterized by
 - Undetected erroneous computation
 - Opposite of "fail fast (with an error or no result)"
- "erroneous" can be "malicious" (intent is the only distinction)
- Very difficult to detect and confine Byzantine failures
- Do all you can to avoid them
- Lamport et al., "The Byzantine Generals Problem," ACM TOPLAS 1982.

Aside: Byzantine Generals Problem

The Byzantine Generals Problem

LESLIE LAMPORT, ROBERT SHOSTAK, and MARSHALL PEASE SRI International

Reliable computer systems must handle malfunctioning components that give conflicting information to different parts of the system. This situation can be expressed abstractly in terms of a group of generals of the Byzantine army camped with their troops around an enemy city. Communicating only by messenger, the generals must agree upon a common battle plan. However, one or more of them may be traitors who will try to confuse the others. The problem is to find an algorithm to ensure that the loyal generals will reach agreement. It is shown that, using only oral messages, this problem is solvable if and only if more than two-thirds of the generals are loyal; so a single traitor can confound two loyal generals. With unforgeable written messages, the problem is solvable for any number of generals and possible traitors. Applications of the solutions to reliable computer systems are then discussed.

Categories and Subject Descriptors: C.2.4. [Computer-Communication Networks]: Distributed Systems—network operating systems; D.4.4 [Operating Systems]: Communications Management network communication; D.4.5 [Operating Systems]: Reliability—fault tolerance

General Terms: Algorithms, Reliability

Additional Key Words and Phrases: Interactive consistency

ACM TOPLAS 1982

https://dl.acm.org/citation.cfm?id=357176

Using Memory Errors to Attack a Virtual Machine

Sudhakar Govindavajhala * Andrew W. Appel Princeton University {sudhakar,appel}@cs.princeton.edu

We present an experimental study showing that soft memory errors can lead to serious security vulnerabilities in Java and .NET virtual machines, or in any system that relies on type-checking of untrusted programs as a protection mechanism. Our attack works by sending to the JVM a Java program that is designed so that almost any memory error in its address space will allow it to take control of the JVM. All conventional Java and .NET virtual machines are vulnerable to this attack. The technique of the attack is broadly applicable against other language-based security schemes such as proof-carrying code.

We measured the attack on two commercial Java Virtual Machines: Sun's and IBM's. We show that a singlebit error in the Java program's data space can be exploited to execute arbitrary code with a probability of about 70%, and multiple-bit errors with a lower probability.

Our attack is particularly relevant against smart cards or tamper-resistant computers, where the user has physical access (to the outside of the computer) and can use various means to induce faults; we have successfully used heat. Fortunately, there are some straightforward defenses against this attack.

7 Physical fault injection

If the attacker has physical access to the outside of the machine, as in the case of a smart card or other tamperresistant computer, the attacker can induce memory errors. We considered attacks on boxes in form factors ranging from a credit card to a palmtop to a desktop PC.

We considered several ways in which the attacker could induce errors.⁴

IEEE S&P 2003

Before RowHammer (II)

Using Memory Errors to Attack a Virtual Machine

Sudhakar Govindavajhala * Andrew W. Appel Princeton University {sudhakar,appel}@cs.princeton.edu



Figure 3. Experimental setup to induce memory errors, showing a PC built from surplus components, clip-on gooseneck lamp, 50-watt spotlight bulb, and digital thermometer. Not shown is the variable AC power supply for the lamp.

IEEE S&P 2003

https://www.cs.princeton.edu/~appel/papers/memerr.pdf

After RowHammer

A simple memory error can be induced by software



RowHammer: Retrospective

- New mindset that has enabled a renewed interest in HW security attack research:
 - Real (memory) chips are vulnerable, in a simple and widespread manner
 → this causes real security problems
 - Hardware reliability \rightarrow security connection is now mainstream discourse
- Many new RowHammer attacks...
 - Tens of papers in top security & architecture venues
 - More to come as RowHammer is getting worse (DDR4 & beyond)
- Many new RowHammer solutions...
 - Apple security release; Memtest86 updated
 - Many solution proposals in top venues (latest in HPCA/S&P 2023)
 - Principled system-DRAM co-design (in original RowHammer paper)
 - More to come...

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Perhaps Most Importantly...

- RowHammer enabled a shift of mindset in mainstream security researchers
 - □ General-purpose hardware is fallible, in a widespread manner
 - Its problems are exploitable
- This mindset has enabled many systems security researchers to examine hardware in more depth
 - And understand HW's inner workings and vulnerabilities
- It is no coincidence that two of the groups that discovered Meltdown and Spectre heavily worked on RowHammer attacks before
 - More to come...

Conclusion

Summary: RowHammer

- Memory reliability is reducing
- Reliability issues open up security vulnerabilities
 - Very hard to defend against

Rowhammer is a prime example

- First example of how a simple hardware failure mechanism can create a widespread system security vulnerability
- Its implications on system security research are tremendous & exciting
- Bad news: RowHammer is getting worse

Good news: We have a lot more to do

- □ We are now fully aware hardware is easily fallible
- We are developing both attacks and solutions
- We are developing principled models, methodologies, solutions

A RowHammer Survey Across the Stack

Onur Mutlu and Jeremie Kim,
 "RowHammer: A Retrospective"
 IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems (TCAD) Special Issue on Top Picks in Hardware and Embedded Security, 2019.
 [Preliminary arXiv version]
 [Slides from COSADE 2019 (pptx)]
 [Slides from VLSI-SOC 2020 (pptx) (pdf)]
 [Talk Video (1 hr 15 minutes, with Q&A)]

RowHammer: A Retrospective

Onur Mutlu§‡Jeremie S. Kim‡§§ETH Zürich‡Carnegie Mellon University

A RowHammer Survey: Recent Update

Onur Mutlu, Ataberk Olgun, and A. Giray Yaglikci, "Fundamentally Understanding and Solving RowHammer" Invited Special Session Paper at the <u>28th Asia and South Pacific Design</u> <u>Automation Conference (ASP-DAC)</u>, Tokyo, Japan, January 2023. [arXiv version] [Slides (pptx) (pdf)] [Talk Video (26 minutes)]

Fundamentally Understanding and Solving RowHammer

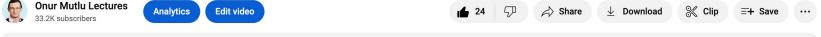
Onur Mutlu onur.mutlu@safari.ethz.ch ETH Zürich Zürich, Switzerland Ataberk Olgun ataberk.olgun@safari.ethz.ch ETH Zürich Zürich, Switzerland A. Giray Yağlıkcı giray.yaglikci@safari.ethz.ch ETH Zürich Zürich, Switzerland

https://arxiv.org/pdf/2211.07613.pdf

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Latest RowHammer Lecture





408 views Streamed 2 months ago Livestream - Seminar in Computer Architecture - ETH Zürich (Spring 2023) Seminar in Computer Architecture, ETH Zürich, Spring 2023 (https://safari.ethz.ch/architecture_s...)

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https://www.youtube.com/watch?v=e6G_Vbrqr_c

The Story of RowHammer Tutorial ...

Onur Mutlu, <u>"Security Aspects of DRAM: The Story of RowHammer"</u> *Invited Tutorial at <u>14th IEEE Electron Devices Society International Memory</u> <u>Workshop</u> (IMW), Dresden, Germany, May 2022. [<u>Slides (pptx)(pdf)</u>] [<u>Tutorial Video</u> (57 minutes)]*



The Story of RowHammer -- Invited Tutorial at IMW 2022 (Intl. Memory Workshop) - Onur Mutlu

598 views • Premiered Jul 27, 2022

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https://www.youtube.com/watch?v=37hWglkQRG0

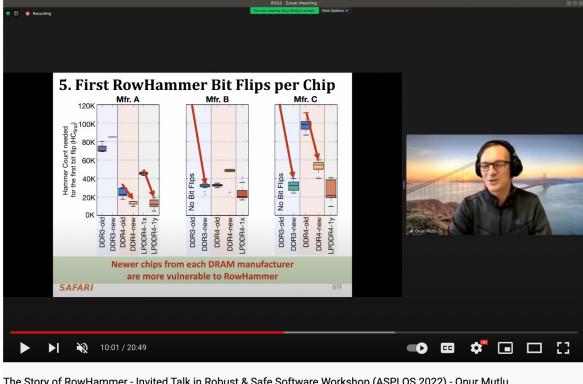
ANALYTICS

The Story of RowHammer in 20 Minutes

Onur Mutlu,

"The Story of RowHammer"

Invited Talk at the Workshop on Robust and Safe Software 2.0 (RSS2), held with the 27th International Conference on Architectural Support for Programming Languages and *Operating Systems (ASPLOS)*, Virtual, 28 February 2022. [Slides (pptx) (pdf)]



The Story of RowHammer - Invited Talk in Robust & Safe Software Workshop (ASPLOS 2022) - Onur Mutlu

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402 views · Premiered Apr 27, 2022



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Funding Acknowledgments

- Alibaba, AMD, ASML, Google, Facebook, Hi-Silicon, HP Labs, Huawei, IBM, Intel, Microsoft, Nvidia, Oracle, Qualcomm, Rambus, Samsung, Seagate, VMware, Xilinx
- Microsoft Swiss JRC
- NSF
- NIH
- GSRC
- SRC
- CyLab
- EFCL

Thank you!

Acknowledgments

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Computer architecture, HW/SW, systems, bioinformatics, security, memory

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Detailed Lectures on RowHammer

- Computer Architecture, Fall 2021, Lecture 5
 - RowHammer (ETH Zürich, Fall 2021)
 - https://www.youtube.com/watch?v=7wVKnPj3NVw&list=P L5Q2soXY2Zi-Mnk1PxjEIG32HAGILkTOF&index=5
- Computer Architecture, Fall 2021, Lecture 6
 - RowHammer and Secure & Reliable Memory (ETH Zürich, Fall 2021)
 - https://www.youtube.com/watch?v=HNd4skQrt6I&list=PL 5Q2soXY2Zi-Mnk1PxjEIG32HAGILkTOF&index=6

https://www.youtube.com/onurmutlulectures

First RowHammer Analysis [ISCA 2014]

Yoongu Kim, Ross Daly, Jeremie Kim, Chris Fallin, Ji Hye Lee, Donghyuk Lee, Chris Wilkerson, Konrad Lai, and Onur Mutlu,
 "Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors"
 Proceedings of the 41st International Symposium on Computer Architecture (ISCA), Minneapolis, MN, June 2014.
 [Slides (pptx) (pdf)] [Lightning Session Slides (pptx) (pdf)] [Source Code and Data] [Lecture Video (1 hr 49 mins), 25 September 2020]
 One of the 7 papers of 2012-2017 selected as Top Picks in Hardware and Embedded Security for IEEE TCAD (link).
 Selected to the ISCA-50 25-Year Retrospective Issue covering 1996-2020 in 2023 (Retrospective (pdf) Full Issue).

Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors

Yoongu Kim¹ Ross Daly^{*} Jeremie Kim¹ Chris Fallin^{*} Ji Hye Lee¹ Donghyuk Lee¹ Chris Wilkerson² Konrad Lai Onur Mutlu¹

¹Carnegie Mellon University ²Intel Labs

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Retrospective on RowHammer & Future

Onur Mutlu, "The RowHammer Problem and Other Issues We May Face as Memory Becomes Denser" Invited Paper in Proceedings of the Design, Automation, and Test in Europe Conference (DATE), Lausanne, Switzerland, March 2017. [Slides (pptx) (pdf)]

The RowHammer Problem and Other Issues We May Face as Memory Becomes Denser

Onur Mutlu ETH Zürich onur.mutlu@inf.ethz.ch https://people.inf.ethz.ch/omutlu

SAFARI https://people.inf.ethz.ch/omutlu/pub/rowhammer-and-other-memory-issues date17.pdf 222

A More Recent RowHammer Retrospective

Onur Mutlu and Jeremie Kim,
 "RowHammer: A Retrospective"
 IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems (TCAD) Special Issue on Top Picks in Hardware and Embedded Security, 2019.
 [Preliminary arXiv version]
 [Slides from COSADE 2019 (pptx)]
 [Slides from VLSI-SOC 2020 (pptx) (pdf)]
 [Talk Video (1 hr 15 minutes, with Q&A)]

RowHammer: A Retrospective

Onur Mutlu§‡Jeremie S. Kim‡§§ETH Zürich‡Carnegie Mellon University

A RowHammer Survey: Recent Update

 Onur Mutlu, Ataberk Olgun, and A. Giray Yaglikci, "Fundamentally Understanding and Solving RowHammer" Invited Special Session Paper at the <u>28th Asia and South Pacific Design</u> Automation Conference (ASP-DAC), Tokyo, Japan, January 2023. [arXiv version] [Slides (pptx) (pdf)] [Talk Video (26 minutes)]

Fundamentally Understanding and Solving RowHammer

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https://arxiv.org/pdf/2211.07613.pdf