Accelerating Genome Analysis

A Primer on an Ongoing Journey

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Technion Invited Lecture





Carnegie Mellon

Overview

- System design for bioinformatics is a critical problem
 - It has large scientific, medical, societal, personal implications
- This talk is about accelerating a key step in bioinformatics: genome sequence analysis
 - In particular, read mapping
- Many bottlenecks exist in accessing and manipulating huge amounts of genomic data during analysis
- We will cover various recent ideas to accelerate read mapping
 - My personal journey since September 2006

Our Dream (circa 2007)

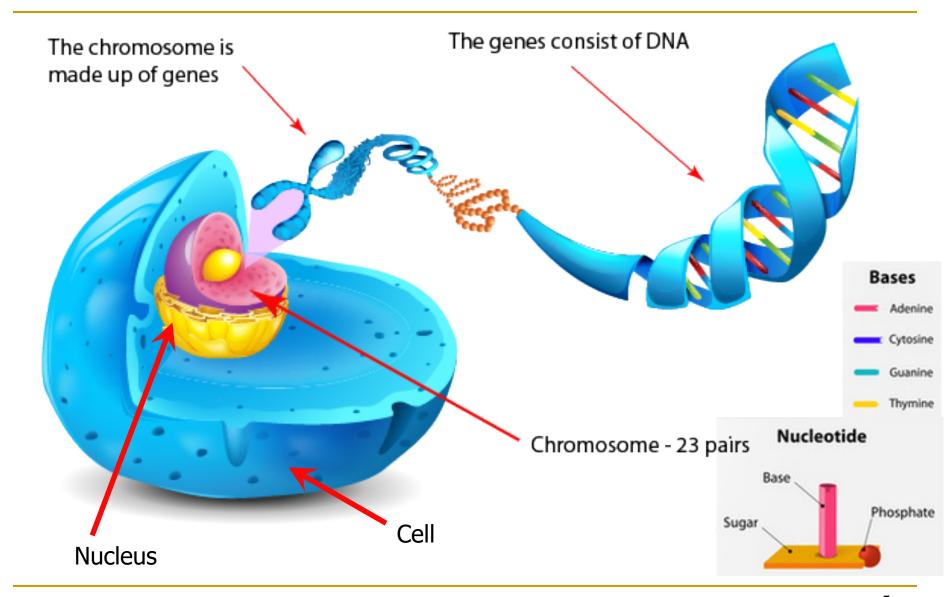
- An embedded device that can perform comprehensive genome analysis in real time (within a minute)
 - Which of these DNAs does this DNA segment match with?
 - What is the likely genetic disposition of this patient to this drug?
 - What disease/condition might this particular DNA/RNA piece associated with?

u . . .

Agenda

- The Problem: DNA Read Mapping
 - State-of-the-art Read Mapper Design
- Algorithmic Acceleration
 - Exploiting Structure of the Genome
 - Exploiting SIMD Instructions
- Hardware Acceleration
 - Specialized Architectures
 - Processing in Memory
- Future Opportunities: New Sequencing Technologies

What Is a Genome Made Of?



The Central Dogma of Molecular Biology



Transcription

UACGU



Protein Phenotypes

Translation



Genotypes

Replication

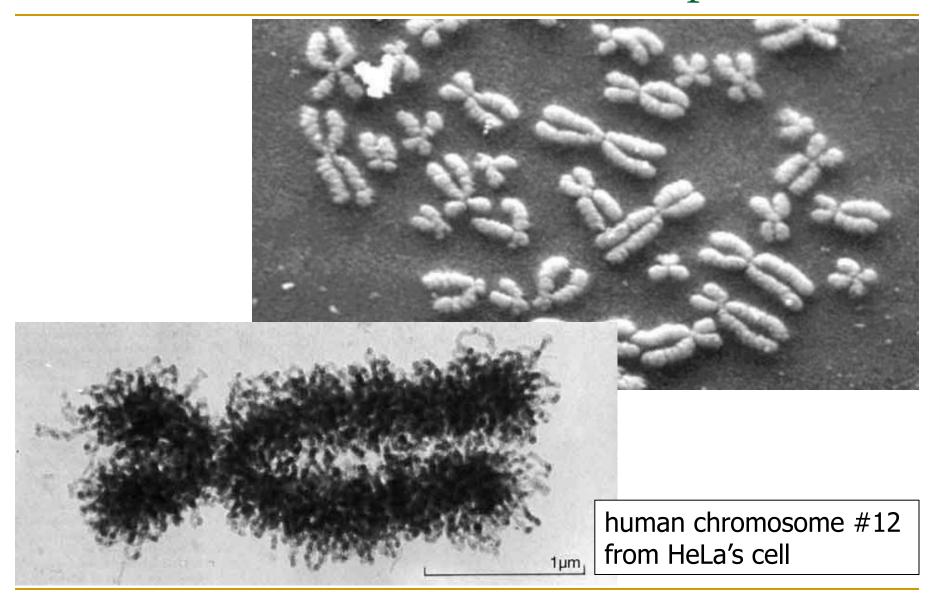








DNA Under Electron Microscope



DNA Sequencing

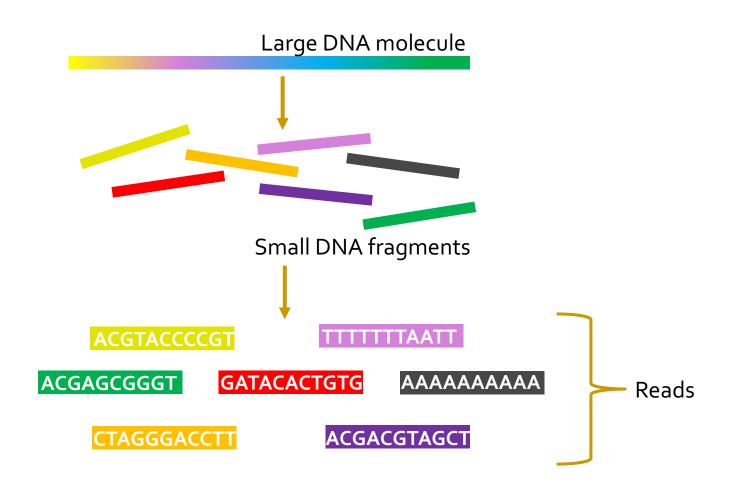
Goal:

Find the complete sequence of A, C, G, T's in DNA.

Challenge:

- There is no machine that takes long DNA as an input, and gives the complete sequence as output
- All sequencing machines chop DNA into pieces and identify relatively small pieces (but not how they fit together)

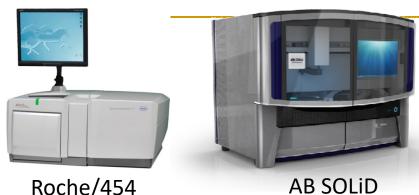
Genome Sequencing



Untangling Yarn Balls & DNA Sequencing



Genome Sequencers



Roche/454



Illumina HiSeq2000



Pacific Biosciences RS



Ion Torrent PGM



Ion Torrent Proton



Illumina MiSeq



Complete Genomics



Oxford Nanopore MinION



Illumina NovaSeq 6000



Oxford Nanopore GridION

... and more! All produce data with different properties.

High-Throughput Sequencers



Illumina MiSeq



Illumina NovaSeq 6000



Pacific Biosciences Sequel II



Pacific Biosciences RS II





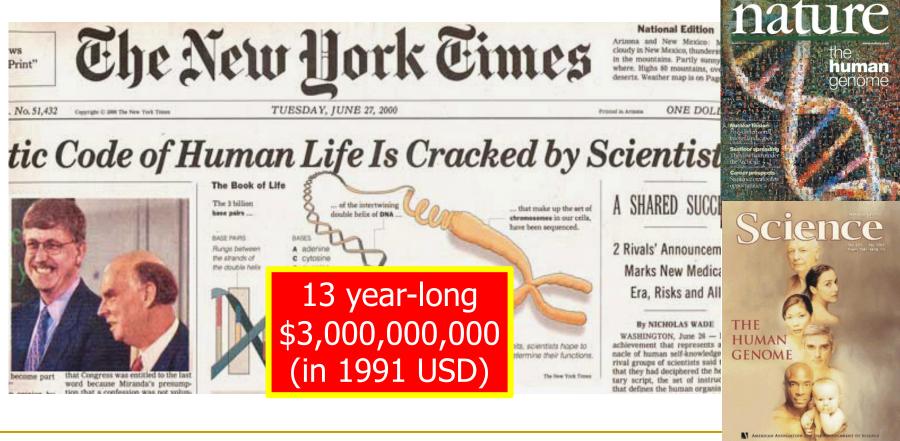
Oxford Nanopore MinION



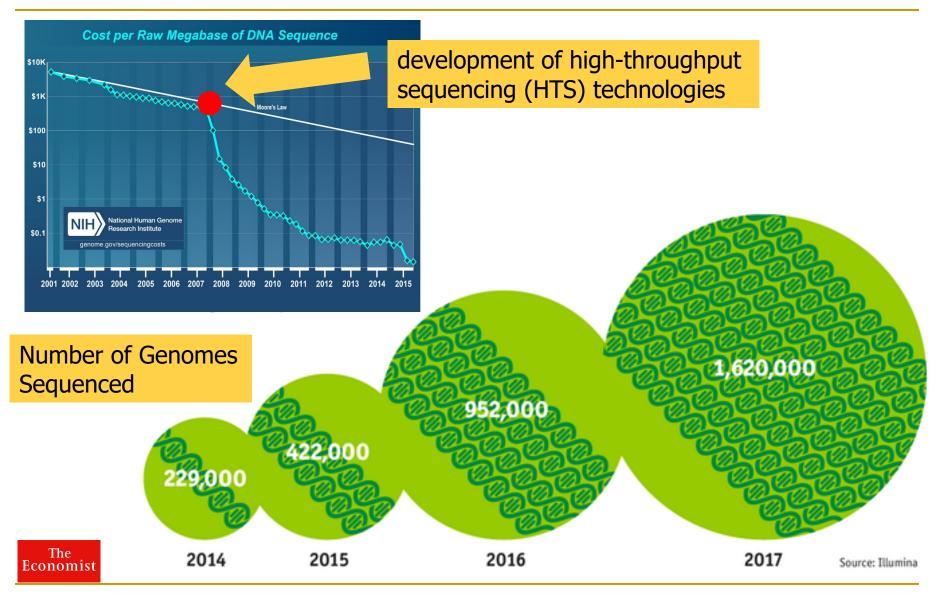
... and more! All produce data with different properties.

The Genomic Era

 1990-2003: The Human Genome Project (HGP) provides a complete and accurate sequence of all **DNA base pairs** that make up the human genome and finds 20,000 to 25,000 human genes.

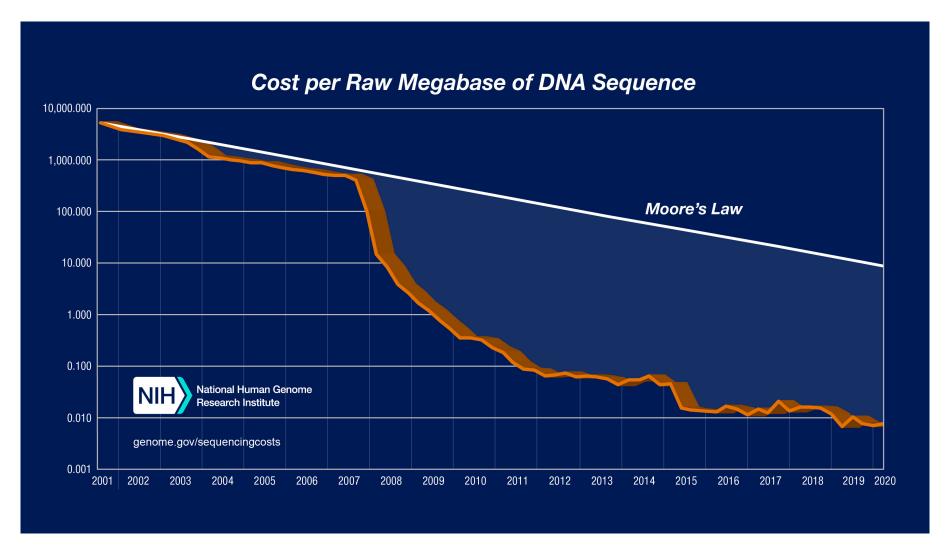


The Genomic Era (continued)





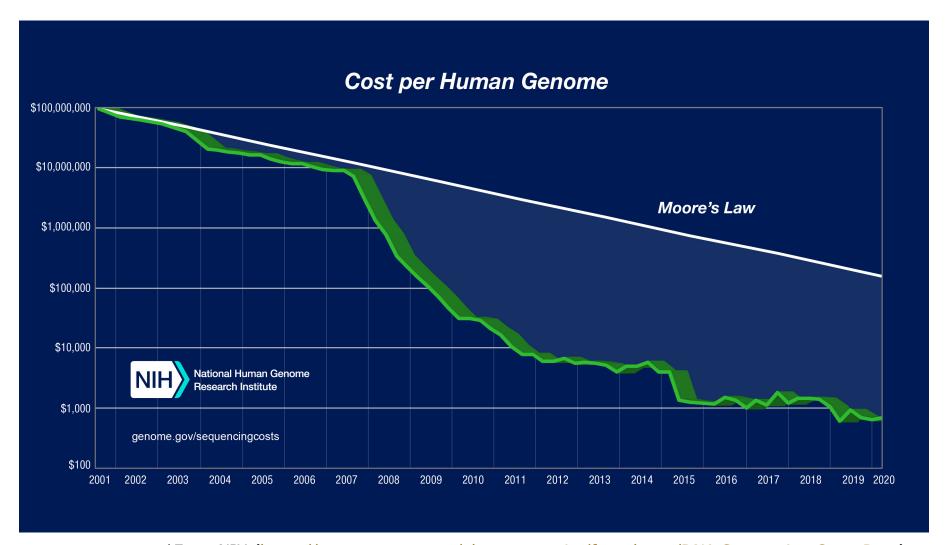
Cost of Sequencing



*From NIH (https://www.genome.gov/about-genomics/fact-sheets/DNA-Sequencing-Costs-Data)



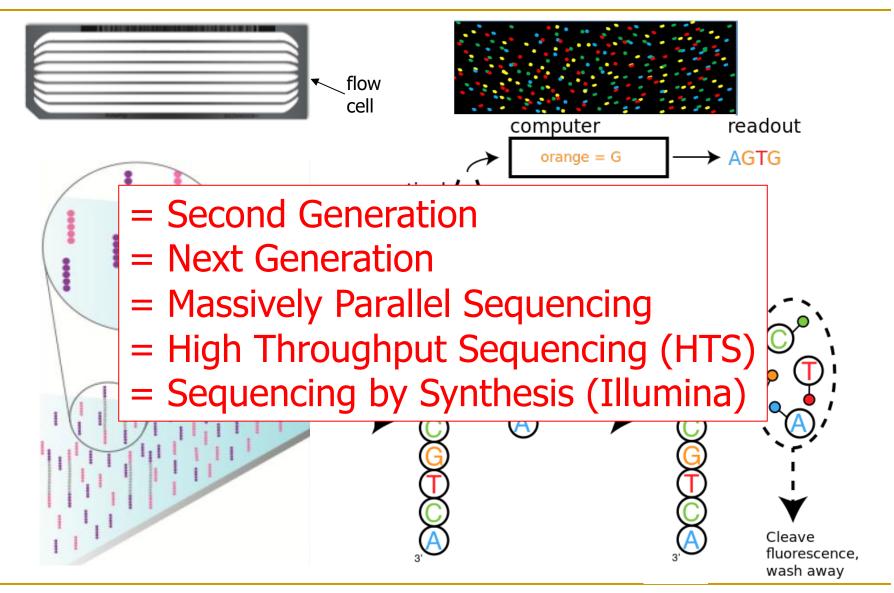
Cost of Sequencing (cont.)



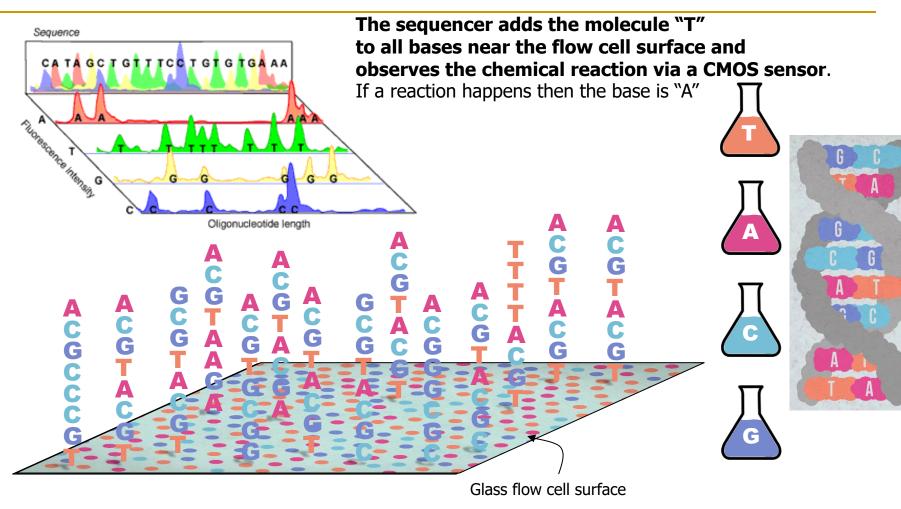
*From NIH (https://www.genome.gov/about-genomics/fact-sheets/DNA-Sequencing-Costs-Data)



High-Throughput Sequencing (HTS)



High-Throughput Sequencing (HTS)



As a workaround, HTS technologies sequence random short DNA fragments (75-300 basepairs long) of copies of the original molecule.

High-Throughput Sequencing

- Massively parallel sequencing technology
 - Illumina, Roche 454, Ion Torrent, SOLID...
- Small DNA fragments are first amplified and then sequenced in parallel, leading to
 - High throughput
 - High speed
 - Low cost
 - Short reads
- Sequencing is done by either reading optical signals as each base is added, or by detecting hydrogen ions instead of light, leading to:
 - Low error rates (relatively)
 - Reads lack information about their order and which part of genome they are originated from



Short Read Read Alignment Reference Genome

Read Mapping

1 Sequencing

Genome Analysis

reference: TTTATCGCTTCCATGACGCAG

read1: ATCGCATCC read2: TATCGCATC

read3: CATCCATGA

read4: CGCTTCCAT

read5: CCATGACGC

read6: TTCCATGAC



Variant Calling

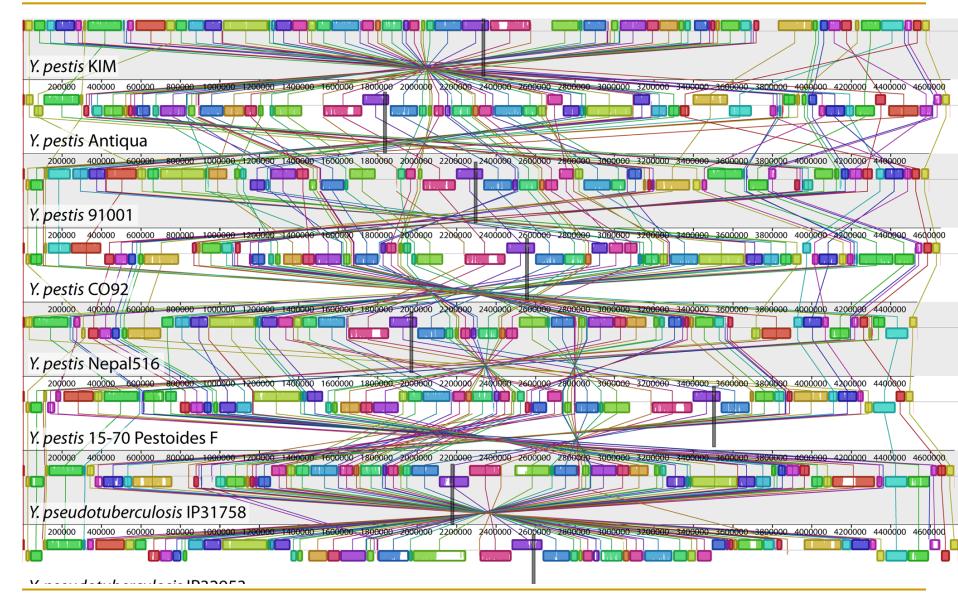
Scientific Discovery 4

Multiple sequence alignment

```
PHDHt.m.
                                                                   MMMMMMMMMMMMMMM-
16082665
            T acid
                         ----MASDRKSEGFOSGAGLIRYFEEEEIKGPALDPKLVVYMGIAVAIIVEIAKIFWPF
                                                                                                    (55)
                         ----masdkksegFOSGAGLIRYFERERIKGPALDPKLVVYIGIAVAIMVELAKIFWPP-
13541150
            T volc
                                                                                                    (55)
                         -MTSMAKDNONENFQSGAGLIRYFNDEDIKGPAIDPKLIIYIGIAMGVIVELAKVFWP
RFAC01077
            F acid
                                                                                                    (58)
                         ----mssgonsggLMSSAGLVRYFDSEDSNALGIDPRSVVAVGAFFGLVVLLAOFFA
15791336
                                                                                                    (53)
            H NRC1
                         MAKAPKGKAKTPPLMSSAGIMBYFER-EKTOINVSPKTILAAGIVTGVLIIILNAYYGLWP-
RAG22196
            A fula
                                                                                                    (68)
                         -----MAKEKTTLPPTGAGLMRFFDE-DTRAIKITPKGAVALTLILIIFEIILHVVGPRIFG
RP001000
                                                                                                    (56)
            P abys
                           ---makekttlpptgAg<mark>lmrff</mark>dE-<mark>dtraikitpkgAialvliliifeillhvv</mark>gpr<mark>i</mark>fg
RPH01741
                                                                                                    (56)
            P hori
                           --makkdkktlppsgag<mark>lvryf</mark>ee<mark>-d</mark>tkg<mark>fkl</mark>tpe<mark>qvvvmsiilavfclvl</mark>rfsg
AE000914
            M ther
                                                                                                    (52)
                         ----MSKRESTGLATSAGLIRYMDE-TFSK<mark>IRV</mark>KPEHVIGVTVAFVIIEAILTYGRF
RMJ09857
            M jann
                                                                                                    (53)
                         -MPSSKKKKETUPLASMAGLIRYYED-PNEKIMISPKLLIIISIIMVAGVIVASILIP
                                                                                                    (58)
15920503
            S toko
AE006662
            S solf
                         -MPSSKKKKETVPVMSMAGLIRYYEE-PNEKVMISPKIVIGASLALTIIVIVITKLF
                                                                                                    (55)
                         --MARRKYEGINPFVAAGLIKFSEEGELEKIKLTPRAAVVISLAIIGLLIAINLLLPPL--
RPK02491
                                                                                                    (58)
            P aero
RAP00437
                         -MSVRRRERRATPVTAAGLLSFYEE-YEGKIKISPTIVVGAAILVSAVVAAA: IFLPAVP-
                                                                                                   (59)
            A pern
                              -----SAGTGGMWRFYTE-DSPGLWVGPVPVLVMSLLFIASVFMLHIWGKYTRS
5803165
                                                                                                    (96)
            H sapi
13324684
            M musc
                                    -SAGTGGMWRFYTR-DSPGLWVGPVPVLVMSLLFIAAVFMLUIWGKYTRS
                                                                                                   (96)
                               -----GAGTGGMWRFYTD-DSPGINVGPVPVLVMSLLFIASVFMLHIWGKYNRS
6002114
            D mela
                                                                                                  (100)
                                 ----ggnngg<mark>lwrfy</mark>t<mark>e-d</mark>stg<mark>lwigevpvlvmslvfiasvfvlhiwgkft</mark>rs
14574310
            C elea
                                                                                                   (81)
                                 ----GGSSSTMLKLYTD-ESQGLK
            Y lipo
10697176
                                                              DPVVVMVLSLGFIFSVVALEILAKVSTK
                                                                                                   (91)
                                 -----GGSSSS<mark>ILKLYTD-D</mark>ANGFRVDSLVVLFLSVGFIFSVIALHLLTKFTHI
6320857
                                                                                                    (88)
6320932
            S cere
                                    -TNSNNS<mark>ILKIYSD-D</mark>ATG<mark>LRV</mark>DPLVVLFLAVGFIFSVVAL<mark>H</mark>VISK<mark>VA</mark>GK
                                                                                                   (82)
```

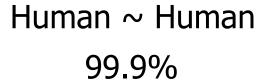
Example Question: If I give you a bunch of sequences, tell me where they are the same and where they are different.

Genome Sequence Alignment: Example



The Genetic Similarity Between Species







Human ~ Chimpanzee 96%



Human ~ Cat 90%



Human ∼ Cow 80%



Human ∼ Banana 50-60%

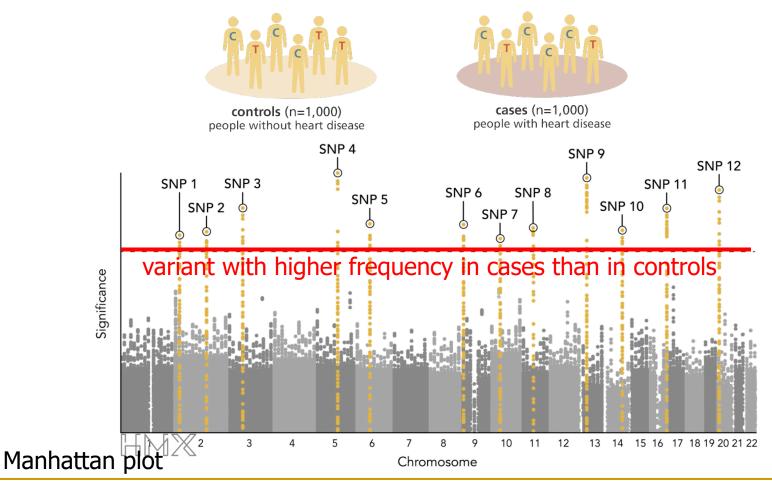
Finding Variations Associated with Traits

| | SNP1 | SNP2 | Blood Pressure |
|----------------|---------------|-----------|-----------------------|
| Individual #1 | ACATGCCGACATT | TCATAGGCC | 180 |
| Individual #2 | ACATGCCGACATT | TCATAAGCC | 175 |
| Individual #3 | ACATGCCGACATT | TCATAGGCC | 170 |
| Individual #4 | ACATGCCGACATT | TCATAAGCC | 165 |
| Individual #5 | ACATGCCGACATT | TCATAGGCC | 160 |
| Individual #6 | ACATGCCGACATT | TCATAGGCC | 145 |
| Individual #7 | ACATGCCGACATT | TCATAAGCC | 140 |
| Individual #8 | ACATGCCGACATT | TCATAAGCC | 130 |
| Individual #9 | ACATGTCGACATT | TCATAGGCC | 120 |
| Individual #10 | ACATGTCGACATT | TCATAAGCC | 120 |
| Individual #11 | ACATGTCGACATT | TCATAGGCC | 115 |
| Individual #12 | ACATGTCGACATT | TCATAAGCC | 110 |
| Individual #13 | ACATGTCGACATT | TCATAGGCC | 110 |
| Individual #14 | ACATGTCGACATT | TCATAAGCC | 110 |
| Individual #15 | ACATGTCGACATT | TCATAGGCC | 105 |
| Individual #16 | ACATGTCGACATT | TCATAAGCC | 100 |

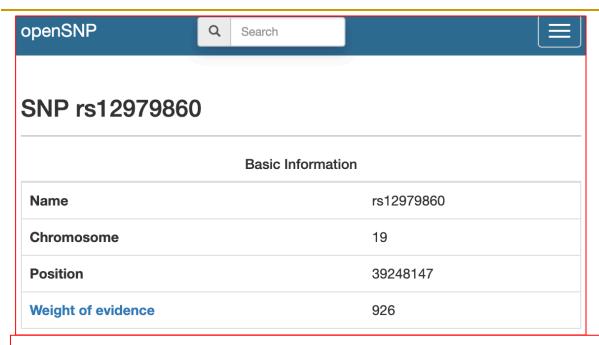
SNP: single nucleotide polymorphism

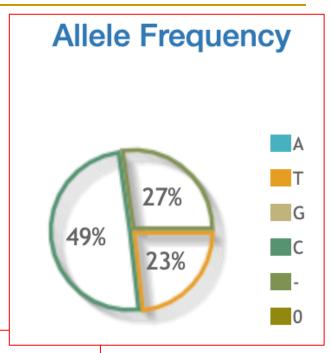
Genome-Wide Association Studies (GWAS)

 Enables detection of genetic variants associated with phenotypes using two groups of people.



SNPs and Personalized Medicine





Links to SNPedia

| Title | Summary | |
|----------------|---|--|
| rs12979860 T/T | ~20-25% of such hepatitis c patients respond to treatment | |
| rs12979860 C/C | ~80% of such hepatitis c patients respond to treatment | |
| rs12979860 C/T | ~20-40% of such hepatitis c patients respond to treatment | |

Much Larger Structural Variations



AUTISM

Weiss, *N Eng J Med* 2008 Deletion of 593 kb



SCHIZOPHRENIA

McCarthy, *Nat Genet* 2009 Duplication of 593 kb



OBESITY

Walters, *Nature* 2010 Deletion of 593 kb



UNDERWEIGHT

Jacquemont, *Nature* 2011 Duplication of 593 kb



Deletion in the short arm of chromosome 16 (16p11.2)



Duplication in the short arm of chromosome 16 (16p11.2)

Recommended Reading

nature reviews genetics

Explore our content > Journal information >

nature > nature reviews genetics > review articles > article

Review Article | Published: 15 November 2019

Structural variation in the sequencing era

Steve S. Ho, Alexander E. Urban & Ryan E. Mills ⊠

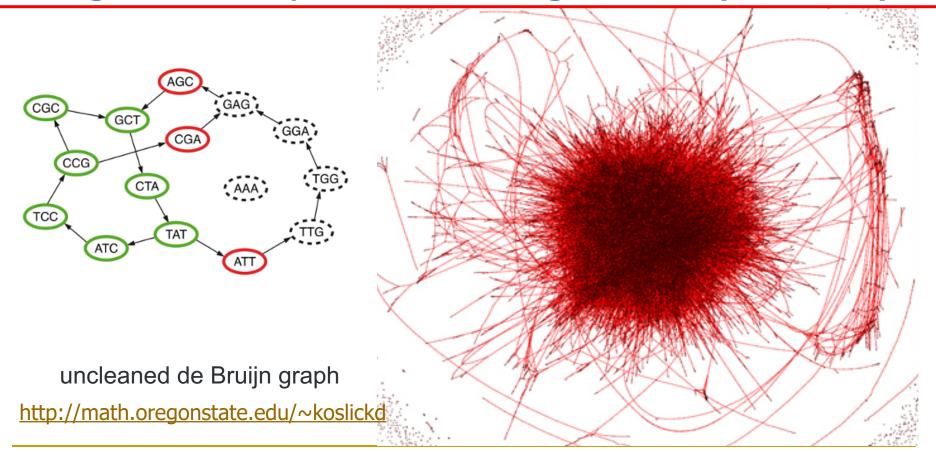
Nature Reviews Genetics 21, 171–189(2020) | Cite this article

15k Accesses | 16 Citations | 309 Altmetric | Metrics

Ho+, "Structural variation in the sequencing era", Nature Reviews Genetics, 2020

Metagenomics, genome assembly, de novo sequencing

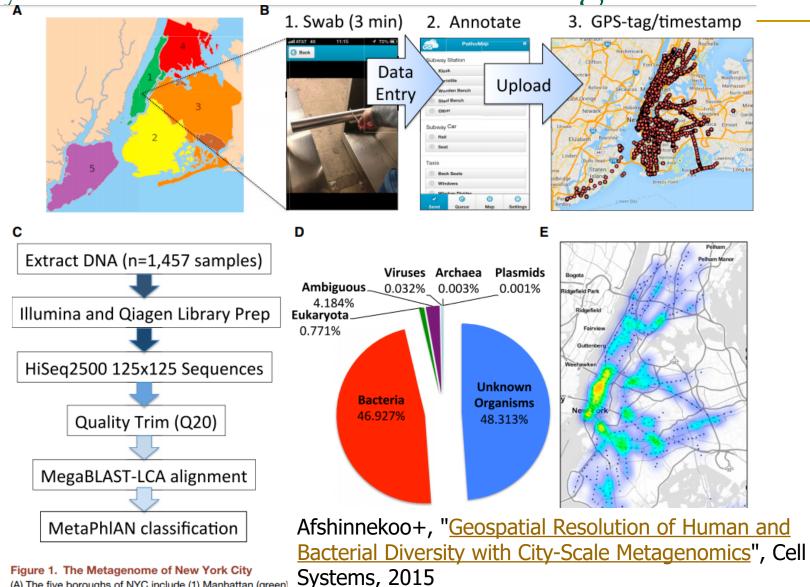
Question 2: Given a bunch of short sequences, Can you identify the approximate species cluster for genomically unknown organisms (bacteria)?



Population-Scale Microbiome Profiling



City-Scale Microbiome Profiling



(A) The five boroughs of NYC include (1) Manhattan (green)

(B) The collection from the 466 subway stations of NYC across the 24 subway lines involved three main steps: (1) collection with Copan Elution swabs, (2) data entry into the database, and (3) uploading of the data. An image is shown of the current collection database, taken from http://pathomap.giscloud.com. (C) Workflow for sample DNA extraction, library preparation, sequencing, quality trimming of the FASTQ files, and alignment with MegaBLAST and MetaPhlAn to discern taxa present



Another Question: Example from 2020

200 Oxford Nanopore sequencers have left UK for China, to support rapid, near-sample coronavirus sequencing for outbreak surveillance

Fri 31st January 2020

Following extensive support of, and collaboration with, public health professionals in China, Oxford Nanopore has shipped an additional 200 MinION sequencers and related consumables to China. These will be used to support the ongoing surveillance of the current coronavirus outbreak, adding to a large number of the devices already installed in the country.



Each MinION sequencer is approximately the size of a stapler, and can provide rapid sequence information about the coronavirus.



700Kg of Oxford Nanopore sequencers and consumables are on their way for use by Chinese scientists in understanding the current coronavirus outbreak.



Example: Scalable SARS-CoV-2 Testing







HOME | ABOL

Search

Comments (I)

Swab-Seq: A high-throughput platform for massively scaled up SARS-CoV-2 testing

Doshua S. Bloom, Eric M. Jones, De Molly Gasperini, De Nathan B. Lubock, Laila Sathe, Chetan Munugala, De A. Sina Booeshaghi, De Oliver F. Brandenberg, De Longhua Guo, De James Boocock, De Scott W. Simpkins, Isabella Lin, Nathan LaPierre, Duke Hong, Yi Zhang, Gabriel Oland, Bianca Judy Choe, Sukantha Chandrasekaran, Evann E. Hilt, De Manish J. Butte, De Robert Damoiseaux, De Aaron R. Cooper, De Yi Yin, De Lior Pachter, De Omai B. Garner, De Jonathan Flint, De Eleazar Eskin, De Chongyuan Luo, De Sriram Kosuri, De Leonid Kruglyak, De Valerie A. Arboleda

doi: https://doi.org/10.1101/2020.08.04.20167874

Bloom+, "Swab-Seq: A high-throughput platform for massively scaled up SARS-CoV-2 testing", medRxiv, 2020

Example: Rapid Surveillance of Ebola Outbreak

Figure 1: Deployment of the portable genome surveillance system in Guinea.

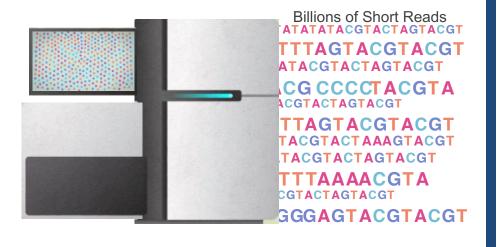


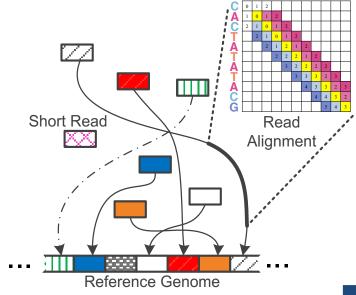






Quick+, "Real-time, portable genome sequencing for Ebola surveillance", Nature, 2016

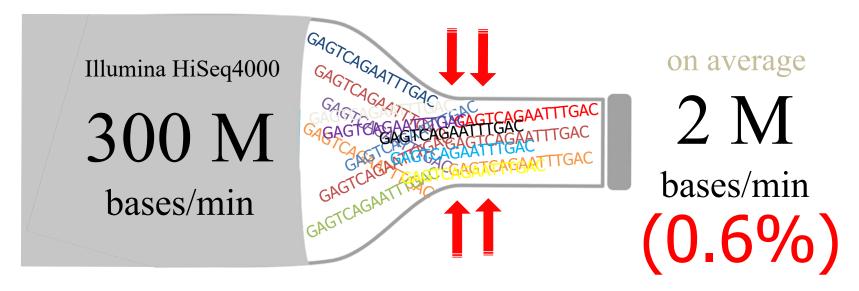




Read Mapping

Sequencing

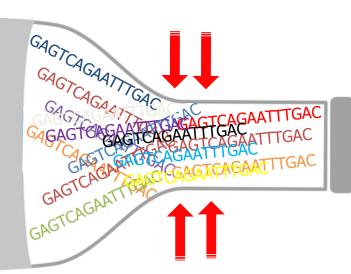
Bottlenecked in Mapping!!



The Read Mapping Bottleneck

300 Million bases/minute

Read Sequencing **



→ Million bases/minute

Read Mapping*

150x slower

* BWA-MEM

** HiSeqX10, MinION

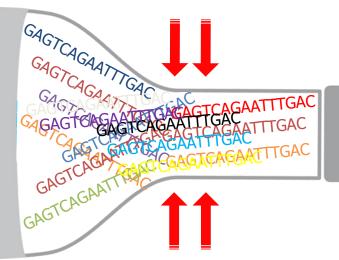
The Read Mapping Bottleneck

48 Human whole genomes

at 30× coverage

in about 2 days

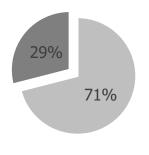
Illumina NovaSeq 6000



1 Human genome

32 CPU hours

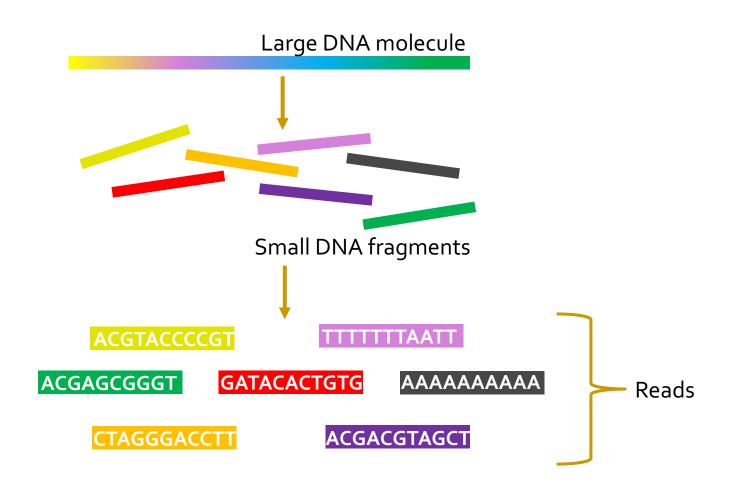
on a 48-core processor



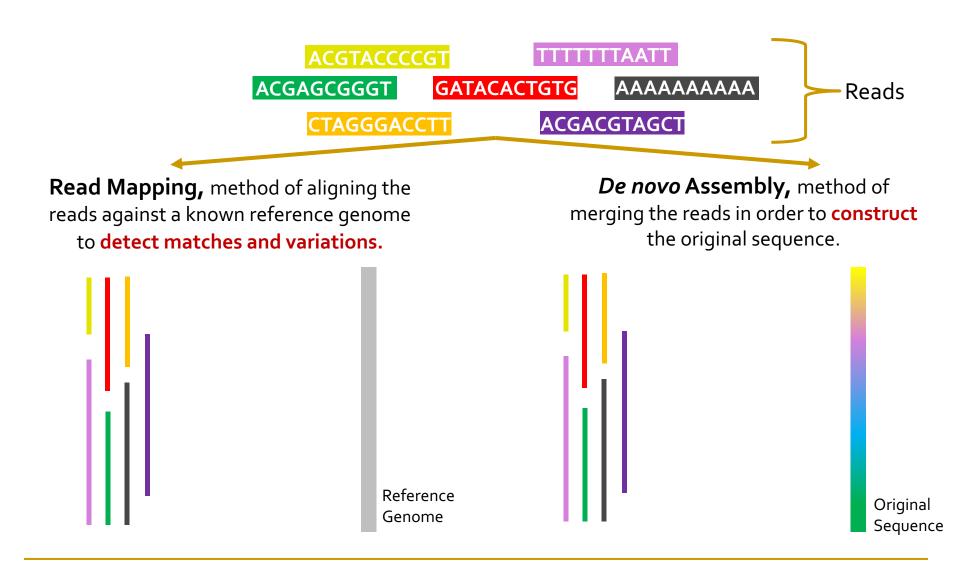
■ Read Mapping ■ Others

Need to construct the entire genome from many reads

Genome Sequencing



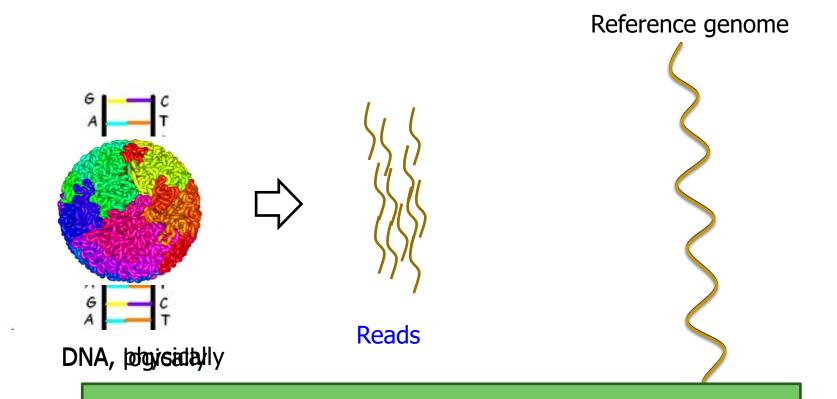
Genome Sequence Analysis





Read Mapping

 Map many short DNA fragments (reads) to a known reference genome with some differences allowed

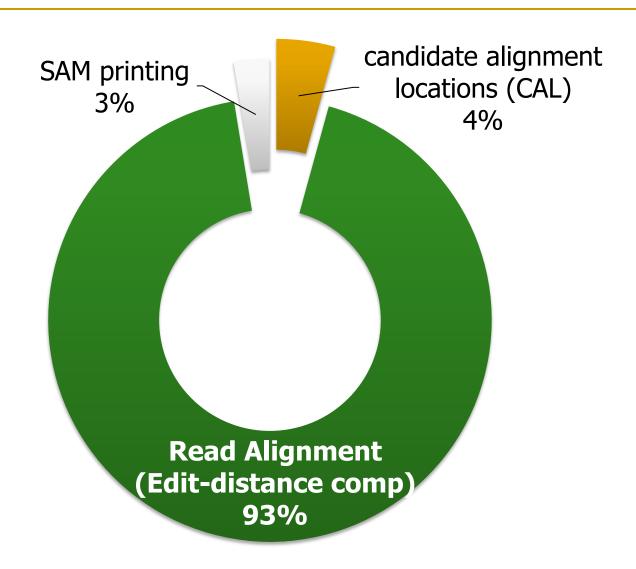


Mapping short reads to reference genome is challenging (billions of 50-300 base pair reads)

Read Mapping for Metagenomic Analysis

Reads from different unknown donors at sequencing time are mapped to many known reference genomes genetic material recovered directly from environmental Reads Reference samples "text format" Database

Read Mapping Execution Time Breakdown



Read Alignment/Verification

 Edit distance is defined as the minimum number of edits (i.e. insertions, deletions, or substitutions) needed to make the read exactly match the reference segment.

NETHERLANDS x SWITZERLAND





Challenges in Read Mapping

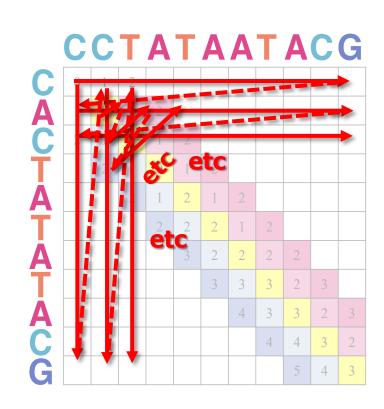
- Need to find many mappings of each read
 - A short read may map to many locations, especially with High-Throughput DNA Sequencing technologies
 - How can we find all mappings efficiently?
- Need to tolerate small variances/errors in each read
 - Each individual is different: Subject's DNA may slightly differ from the reference (Mismatches, insertions, deletions)
 - How can we efficiently map each read with up to e errors present?
- Need to map each read very fast (i.e., performance is important)
 - □ Human DNA is 3.2 billion base pairs long → Millions to billions of reads (State-of-the-art mappers take weeks to map a human's DNA)
 - How can we design a much higher performance read mapper?

Why Is Read Alignment Slow?

Quadratic-time dynamicprogramming algorithm(s)

 Data dependencies limit the computation parallelism

 Entire matrix computed even though strings may be dissimilar



Read Alignment

Example: Dynamic Programming Table

NETHERLANDS x SWITZERLAND

immediate left, upper left, upper entries of its own

| | | N | Е | Т | Н | Ε | R | L | Α | N | D | S |
|---|----|---|---|----------|----------|---|---|---|---|---|----|----|
| | | | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| S | | 1 | | | | | | | | | | |
| W | 2 | | | K | \ | | | | | | | |
| Ι | 3 | | | * | 7 | | | | | | | |
| Т | 4 | | | | | | | | | | | |
| Z | 5 | | | | | | | | | | | |
| Е | 6 | | | | | | | | | | | |
| R | 7 | | | | | | | | | | | |
| L | 8 | | | | | | | | | | | |
| Α | 9 | | | | | | | | | | | |
| N | 10 | | | | | | | | | | | |
| D | 11 | | | | | | | | | | | |

Example: Dynamic Programming Table

NETHERLANDS x SWITZERLAND

| | | N | Е | Т | Н | Е | R | L | Α | N | D | S |
|---|----|----|----|----|----|---|---|---|---|---|----|----|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| S | 1 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 10 |
| W | 2 | 2 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| Ι | 3 | 3 | 3 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| Т | 4 | 4 | 4 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| Z | 5 | 5 | 5 | 4 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| Е | 6 | 6 | 5 | 5 | 5 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| R | 7 | 7 | 6 | 6 | 6 | 5 | 4 | 5 | 6 | 7 | 8 | 9 |
| L | 8 | 8 | 7 | 7 | 7 | 6 | 5 | 4 | 5 | 6 | 7 | 8 |
| Α | 9 | 9 | 8 | 8 | 8 | 7 | 6 | 5 | 4 | 5 | 6 | 7 |
| N | 10 | 9 | 9 | 9 | 9 | 8 | 7 | 6 | 5 | 4 | 5 | 6 |
| D | 11 | 10 | 10 | 10 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 5 |

- Matrix-filling is O(mn) time and space.
- Backtrace is O(m + n) time.

Example: Dynamic Programming

Quadratic-time dynamicprogramming algorithm WHY?!

Enumerate all possible prefixes

NETHERLANDS x SWITZERLAND

- NETHERLANDS x S
 NETHERLANDS x SW
- NETHERLANDS x SWI

NETERLANDS x SWIT

NETHERLANDS x SWITZ

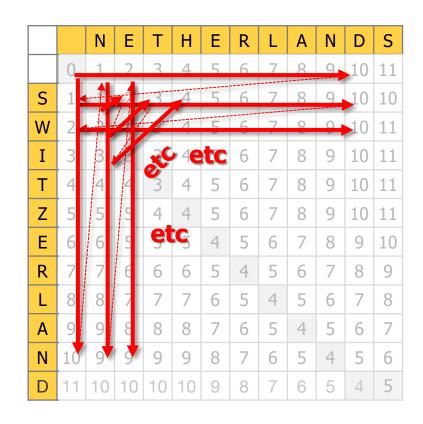
NETHERLANDS x SWITZE

NETHERLANDS x SWITZER

NETHERLANDS x SWITZERL

- NETHERLANDS x SWITZERLA
 - ر NETHERLANDS x SWITZERLAN

NETHERLANDS x SWITZERLAND



Read Mapping Survey in 111 Pages

In-depth analysis of 107 read mapping techniques (1988-2020)

arXiv.org > q-bio > arXiv:2003.00110

Search...

Help | Advanced

Quantitative Biology > Genomics

[Submitted on 28 Feb 2020 (v1), last revised 9 Jul 2020 (this version, v3)]

Technology dictates algorithms: Recent developments in read alignment

Mohammed Alser, Jeremy Rotman, Kodi Taraszka, Huwenbo Shi, Pelin Icer Baykal, Harry Taegyun Yang, Victor Xue, Sergey Knyazev, Benjamin D. Singer, Brunilda Balliu, David Koslicki, Pavel Skums, Alex Zelikovsky, Can Alkan, Onur Mutlu, Serghei Mangul

Alser+, "<u>Technology dictates algorithms: Recent developments in read alignment</u>", arXiv, 2020

GitHub: https://github.com/Mangul-Lab-USC/review_technology_dictates_algorithms

Agenda

- The Problem: DNA Read Mapping
 - State-of-the-art Read Mapper Design
- Algorithmic Acceleration
 - Exploiting Structure of the Genome
 - Exploiting SIMD Instructions
- Hardware Acceleration
 - Specialized Architectures
 - Processing in Memory
- Future Opportunities: New Sequencing Technologies

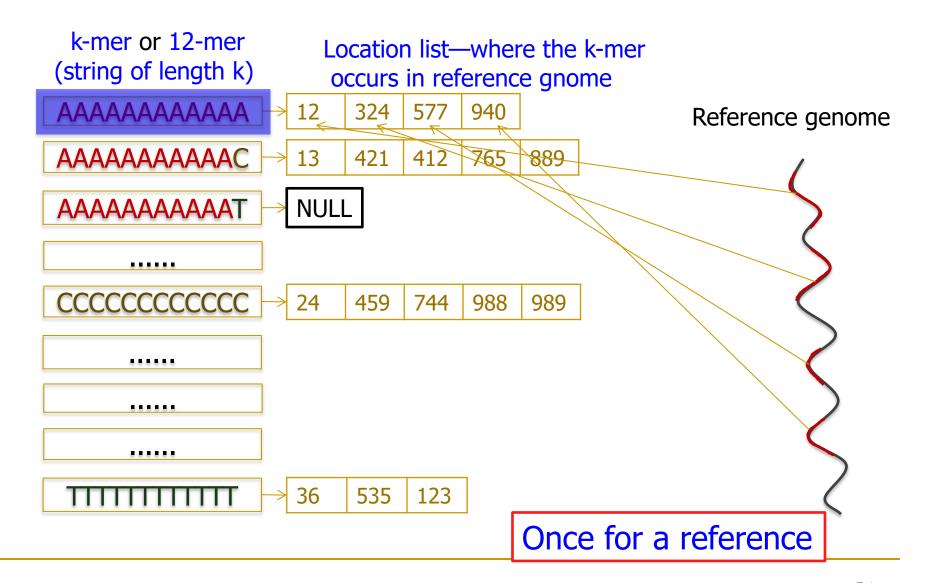
Read Mapping Algorithms: Two Styles

- Hash based seed-and-extend (hash table, suffix array, suffix tree)
 - Index the "k-mers" in the genome into a hash table (pre-processing)
 - When searching a read, find the location of a k-mer in the read; then extend through alignment
 - More sensitive (can find all mapping locations), but slow
 - Requires large memory; this can be reduced with cost to run time
- Burrows-Wheeler Transform & Ferragina-Manzini Index based aligners
 - BWT is a compression method used to compress the genome index
 - Perfect matches can be found very quickly, memory lookup costs increase for imperfect matches
 - Reduced sensitivity

Hash Table Based Read Mappers

- Key Idea
 - Preprocess the reference into a Hash Table
 - Use Hash Table to map reads

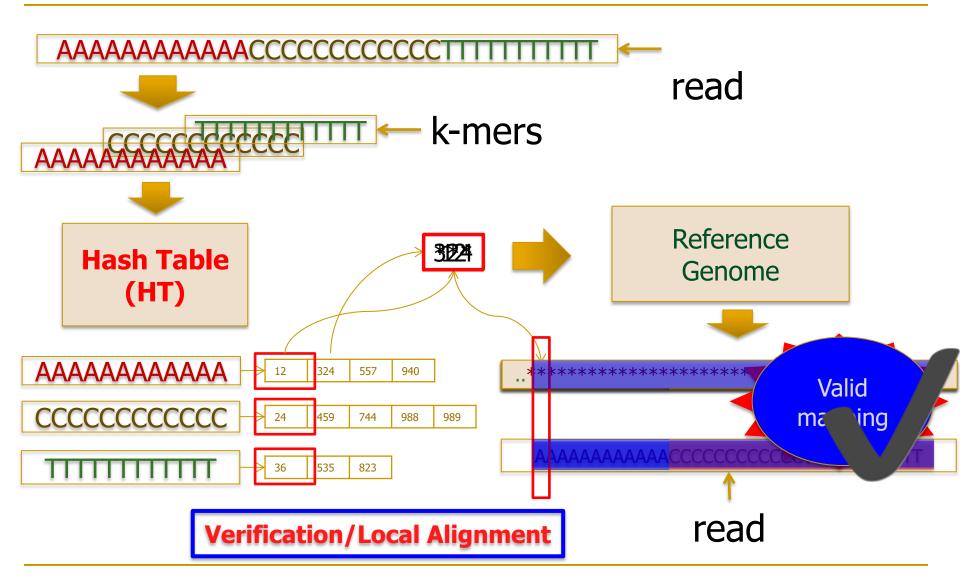
Hash Table-Based Mappers [Alkan+ Nature Gen'09]



Hash Table Based Read Mappers

- Key Idea
 - Preprocess the reference into a Hash Table
 - Use Hash Table to map reads

Hash Table-Based Mappers [Alkan+ Nature Gen'09]



Our First Step: Comprehensive Mapping

- + Guaranteed to find all mappings → sensitive
- + Can tolerate up to e errors

nature genetics

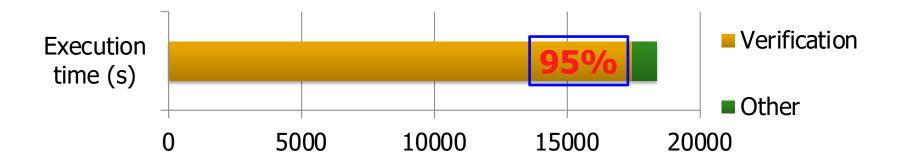
http://mrfast.sourceforge.net/

Personalized copy number and segmental duplication maps using next-generation sequencing

Can Alkan^{1,2}, Jeffrey M Kidd¹, Tomas Marques-Bonet^{1,3}, Gozde Aksay¹, Francesca Antonacci¹, Fereydoun Hormozdiari⁴, Jacob O Kitzman¹, Carl Baker¹, Maika Malig¹, Onur Mutlu⁵, S Cenk Sahinalp⁴, Richard A Gibbs⁶ & Evan E Eichler^{1,2}

Problem and Goal

- Poor performance of existing read mappers: Very slow
 - Verification/alignment takes too long to execute
 - Verification requires a memory access for reference genome + many base-pair-wise comparisons between the reference and the read (edit distance computation)



 Goal: Speed up the mapper by reducing the cost of verification

Overarching Key Idea

Filter fast before you align

Minimize costly edit distance computations

("approximate string comparisons")

Accelerating Genome Analysis: Overview

 Mohammed Alser, Zulal Bingol, Damla Senol Cali, Jeremie Kim, Saugata Ghose, Can Alkan, and Onur Mutlu,

"Accelerating Genome Analysis: A Primer on an Ongoing Journey"

IEEE Micro (IEEE MICRO), Vol. 40, No. 5, pages 65-75, September/October 2020.

[Slides (pptx)(pdf)]

[Talk Video (1 hour 2 minutes)]

Accelerating Genome Analysis: A Primer on an Ongoing Journey

Mohammed Alser

ETH Zürich

Zülal Bingöl

Bilkent University

Damla Senol Cali

Carnegie Mellon University

Jeremie Kim

ETH Zurich and Carnegie Mellon University

Saugata Ghose

University of Illinois at Urbana–Champaign and Carnegie Mellon University

Can Alkan

Bilkent University

Onur Mutlu

ETH Zurich, Carnegie Mellon University, and Bilkent University

Agenda

- The Problem: DNA Read Mapping
 - State-of-the-art Read Mapper Design
- Algorithmic Acceleration
 - Exploiting Structure of the Genome
 - Exploiting SIMD Instructions
- Hardware Acceleration
 - Specialized Architectures
 - Processing in Memory
- Future Opportunities: New Sequencing Technologies

Our First Filter: Pure Software Approach

- Download the source code and try for yourself
 - Download link to FastHASH

Xin et al. BMC Genomics 2013, **14**(Suppl 1):S13 http://www.biomedcentral.com/1471-2164/14/S1/S13



PROCEEDINGS

Open Access

Accelerating read mapping with FastHASH

Hongyi Xin¹, Donghyuk Lee¹, Farhad Hormozdiari², Samihan Yedkar¹, Onur Mutlu^{1*}, Can Alkan^{3*}

From The Eleventh Asia Pacific Bioinformatics Conference (APBC 2013) Vancouver, Canada. 21-24 January 2013

Reducing the Cost of Verification

- We observe that most verification (edit distance computation) calculations are unnecessary
 - 1 out of 1000 potential locations passes the verification process
- We observe that we can get rid of unnecessary verification calculations by
 - Detecting and rejecting early invalid mappings (filtering)
 - Reducing the number of potential mappings to examine

Key Observations [Xin+, BMC Genomics 2013]

Observation 1

- Adjacent k-mers in the read should also be adjacent in the reference genome
- Read mapper can quickly reject mappings that do **not** satisfy this property

Observation 2

- Some k-mers are cheaper to verify than others because they have shorter location lists (they occur less frequently in the reference genome)
 - Mapper needs to examine only e+1 k-mers' locations to tolerate e errors
- Read mapper can choose the cheapest e+1 k-mers and verify their locations

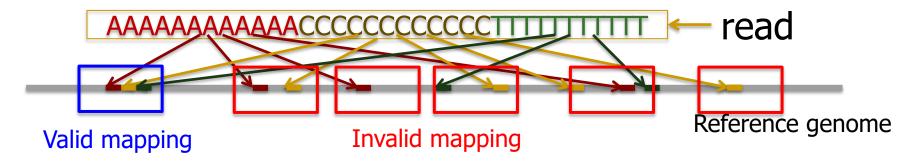
FastHASH Mechanisms [Xin+, BMC Genomics 2013]

 Adjacency Filtering (AF): Rejects obviously invalid mapping locations at early stage to avoid unnecessary verifications

Cheap K-mer Selection (CKS): Reduces the absolute number of potential mapping locations to verify

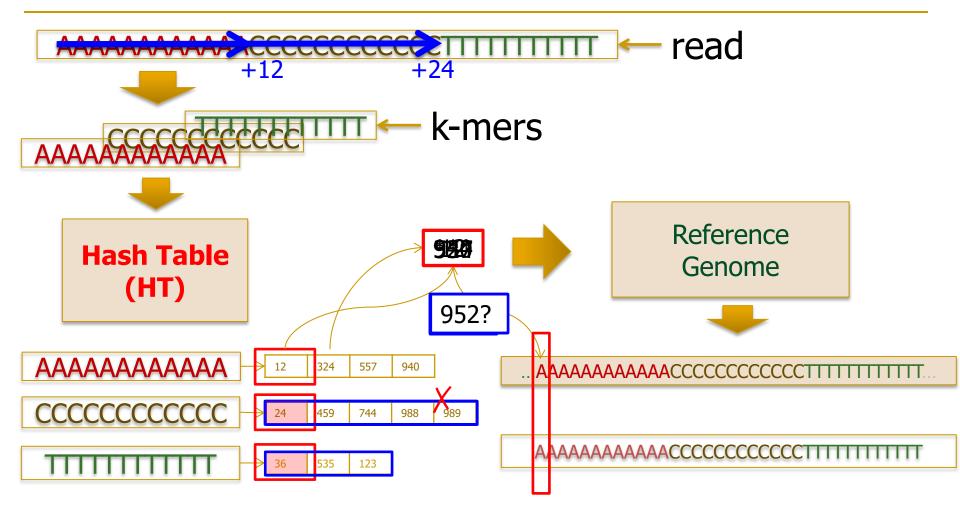
Adjacency Filtering (AF)

- Goal: detect and filter out invalid mappings at early stage
- Key Insight: For a valid mapping, adjacent k-mers in the read are also adjacent in the reference genome



- Key Idea: search for adjacent locations in the k-mers' location lists
 - If more than e k-mers fail → there must be more than e errors → invalid mapping

Adjacency Filtering (AF)



FastHASH Mechanisms [Xin+, BMC Genomics 2013]

 Adjacency Filtering (AF): Rejects obviously invalid mapping locations at early stage to avoid unnecessary verifications

Cheap K-mer Selection (CKS): Reduces the absolute number of potential mapping locations to verify

Cheap K-mer Selection (CKS)

Goal: Reduce the number of potential mappings to examine

Key insight:

 K-mers have different cost to examine: Some k-mers are cheaper as they have fewer locations than others (occur less frequently in reference genome)

Key idea:

- Sort the k-mers based on their number of locations
- Select the k-mers with the fewest number locations to verify

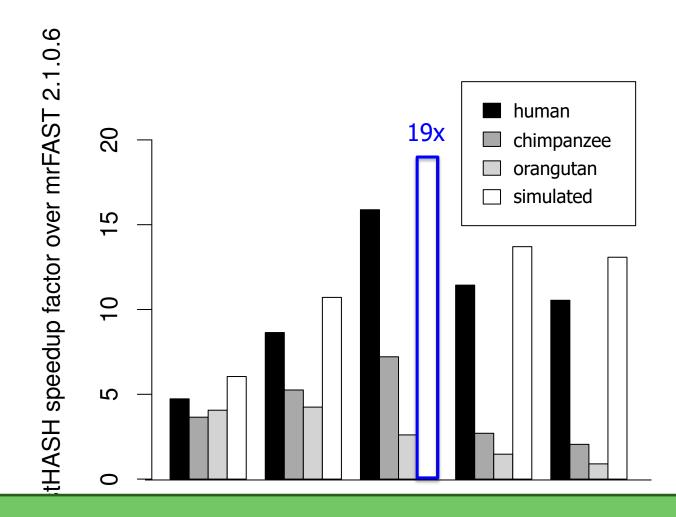
Cheap K-mer Selection

read e=2 (examine 3 k-mers) AAGCTCAATTIC CCTCCTTAATTI TOCTCTTAAGAA GGGTATGGCTAG AAGGTTGAGAGC CTTAGGCTTACC 326 338 326 376 388 314 1231 Location 151 1451 4414 2 loc. 2 loc. 9219 Number of Locations 4 loc. 1K loc. 2K loc. 1K loc. Expensive 3 kmers Previous work needs FastHASH verifies only: to verify: 8 locations 3004 locations

Methodology

- Implemented FastHASH on top of state-of-the-art mapper: mrFAST
 - New version mrFAST-2.5.0.0 over mrFAST-2.1.0.6
- Tested with real read sets generated from Illumina platform
 - 1M reads of a human (160 base pairs)
 - 500K reads of a chimpanzee (101 base pairs)
 - 500K reads of a orangutan (70 base pairs)
- Tested with simulated reads generated from reference genome
 - 1M simulated reads of human (180 base pairs)
- Evaluation system
 - Intel Core i7 Sandy Bridge machine
 - 16 GB of main memory

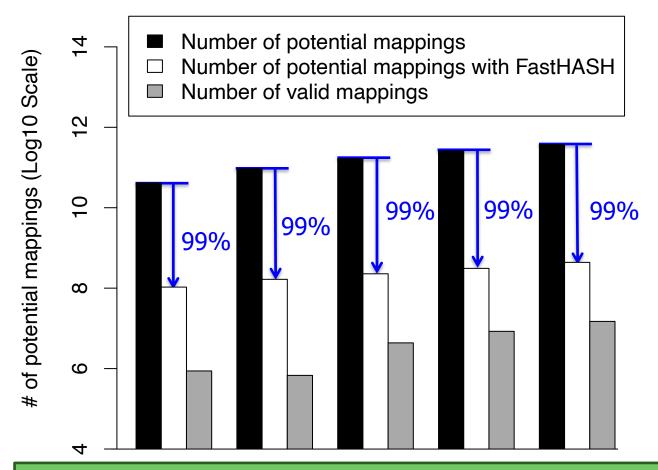
FastHASH Speedup: Entire Read Mapper



With FastHASH, new mrFAST obtains up to 19x speedup over previous version, without losing valid mappings

Analysis

Reduction of potential mappings with FastHASH



FastHASH filters out over 99% of the potential mappings without sacrificing any valid mappings

FastHASH Conclusion

- Problem: Existing read mappers perform poorly in mapping millions of short reads to the reference genome, in the presence of errors
- Observation: Most of the verification calculations are unnecessary → filter them out
- Key Idea: Exploit the structure of the genome to
 - Reject invalid mappings early (Adjacency Filtering)
 - Reduce the number of possible mappings to examine (Cheap K-mer Selection)
- Key Result: FastHASH obtains up to 19x speedup over the state-of-the-art mapper without losing valid mappings

More on FastHASH

- Download source code and try for yourself
 - Download link to FastHASH

Xin et al. BMC Genomics 2013, **14**(Suppl 1):S13 http://www.biomedcentral.com/1471-2164/14/S1/S13



PROCEEDINGS

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Hongyi Xin¹, Donghyuk Lee¹, Farhad Hormozdiari², Samihan Yedkar¹, Onur Mutlu^{1*}, Can Alkan^{3*}

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- Hardware Acceleration
 - Specialized Architectures
 - Processing in Memory
- Future Opportunities: New Sequencing Technologies

Shifted Hamming Distance: SIMD Acceleration

https://github.com/CMU-SAFARI/Shifted-Hamming-Distance

Bioinformatics, 31(10), 2015, 1553–1560

doi: 10.1093/bioinformatics/btu856

Advance Access Publication Date: 10 January 2015

Original Paper



Sequence analysis

Shifted Hamming distance: a fast and accurate SIMD-friendly filter to accelerate alignment verification in read mapping

Hongyi Xin^{1,*}, John Greth², John Emmons², Gennady Pekhimenko¹, Carl Kingsford³, Can Alkan^{4,*} and Onur Mutlu^{2,*}

Xin+, "Shifted Hamming Distance: A Fast and Accurate SIMD-friendly Filter to Accelerate Alignment Verification in Read Mapping", Bioinformatics 2015.

Shifted Hamming Distance

Key observation:

- If two strings differ by E edits, then every bp match can be aligned in at most 2E shifts (of one of the strings).
 - Insight: Shifting a string by one "corrects" for one "error"

Key idea:

- Compute "Shifted Hamming Distance": AND of 2E Hamming Distances of two strings, to filter out invalid mappings
 - Uses bit-parallel operations that nicely map to SIMD instructions

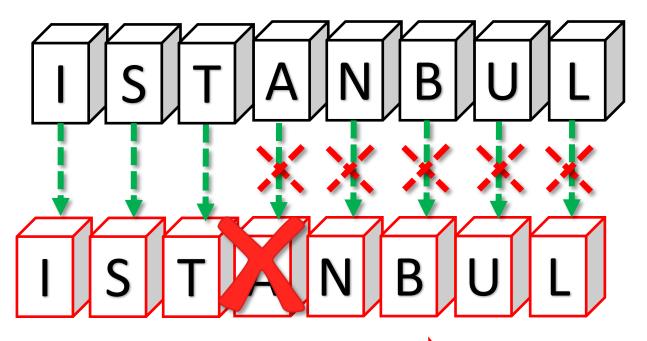
Key result:

- SHD is 3x faster than SeqAn (the best implementation of Gene Myers' bit-vector algorithm), with only a 7% false positive rate
- The fastest CPU-based filtering (pre-alignment) mechanism

Hamming Distance ($\Sigma \oplus$)

3 matches 5 mismatches

<u>Edit = 1 Deletion</u>

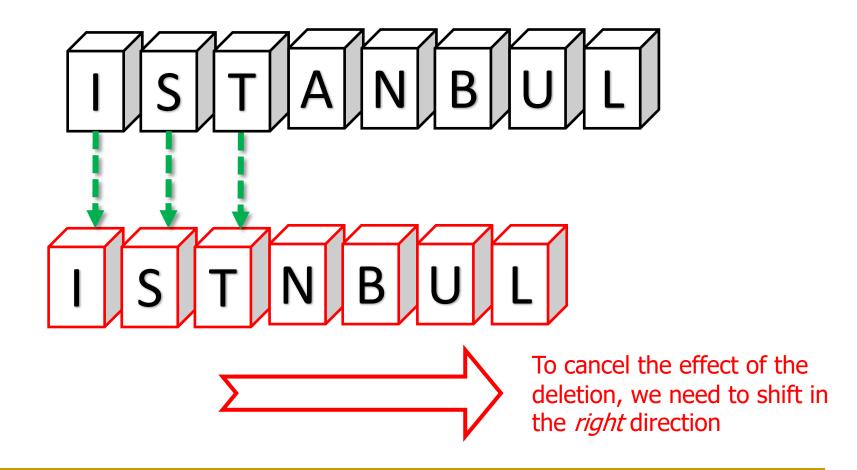




To cancel the effect of a deletion, we need to shift in the *right* direction

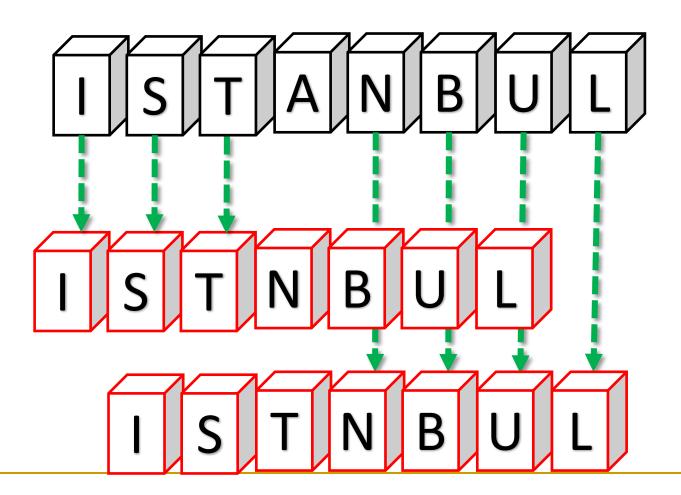
Insight: Shifting a String Helps Similarity Search

3 matches 5 mismatches

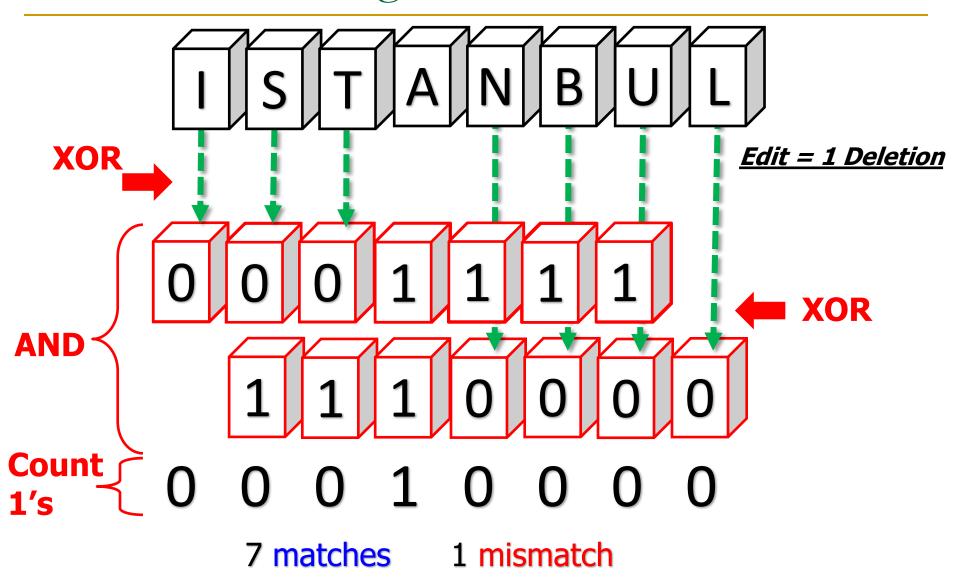


Insight: Shifting a String Helps Similarity Search

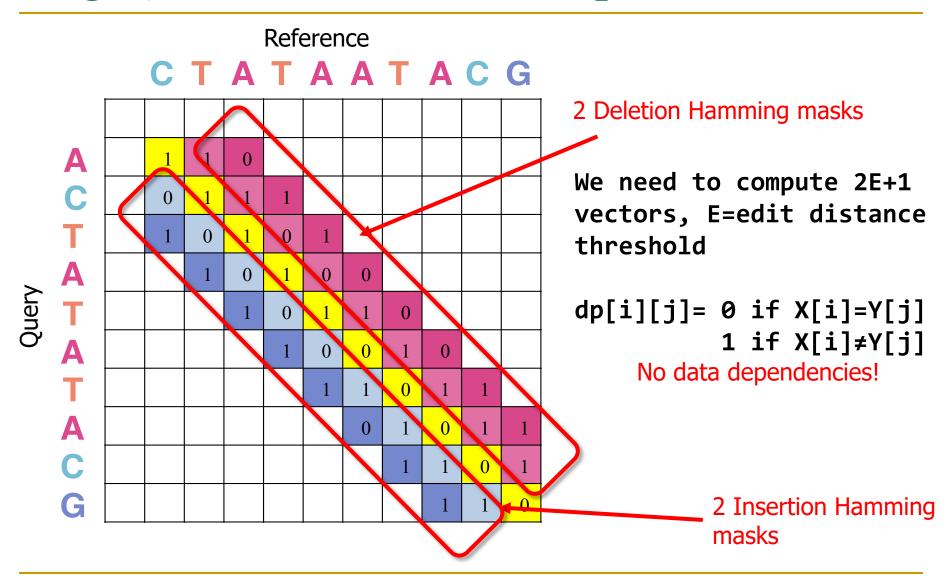
7 matches 1 mismatch



Shifted Hamming Distance



Highly Parallel Matrix Computation



Key Idea of SHD Filtering

Generate 2E+1 masks

Needleman-Wunsch

Alignment

Amend random zeros: $101 \rightarrow 111 \& 1001 \rightarrow 1111$

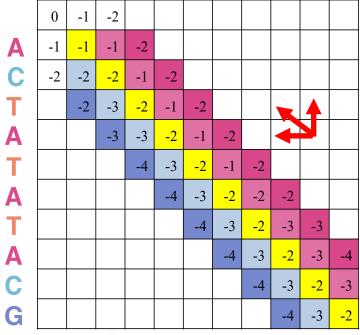
AND all masks, ACCEPT iff number of `1' ≤ Threshold

```
:GAGAGAGATATTTAGTGTTGCAGCACTACAACACAAAAGAGGACCAACTTACGTGTCTAAAAAGGGGGAACATTGTTGGGCCGGA
Reference
  *GAGAGAGATAGTTAGTGTTGCAGCCACTACAACACAAAAGAGGACCAACTTACGTGTCTAAAAGGGGAGACATTGTTGGGCCGG
01111010111001000
     --- Masks after amendment ---
```

Alignment vs. Pre-alignment (Filtering)

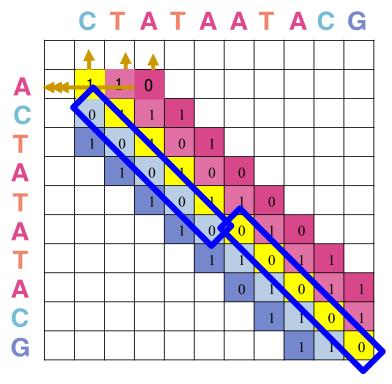


CTATAATACG



|dp[i][j-1] -1 // Inser.

Neighborhood Map



dp[i][i]=|0 if X[i]=Y[i]

Our goal is to track the diagonally consecutive matches in the neighborhood map.

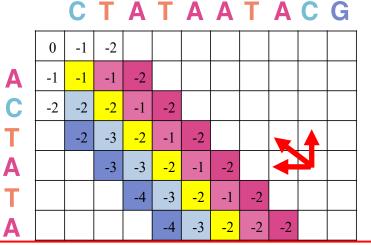
pre-computed cells!

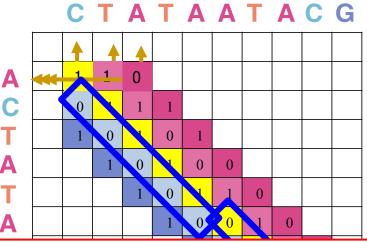
<u>No data dependencies!</u>

Alignment Matrix vs. Neighborhood Map

Needleman-Wunsch

Neighborhood Map

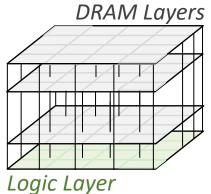




Independent vectors can be processed in parallel using hardware technologies







New Bottleneck: Filtering (Pre-Alignment)

Sequencing generates many reads, each of which potentially mapping to many locations

 \rightarrow

Filtering (Pre-alignment) eliminates the need to verify/align read to invalid mapping locations

 \rightarrow

Alignment/verification (costly edit distance computation) is performed **only** on reads that pass the filter

 New bottleneck in read mapping becomes the "filtering (pre-alignment)" step

More on Shifted Hamming Distance

https://github.com/CMU-SAFARI/Shifted-Hamming-Distance

Bioinformatics, 31(10), 2015, 1553-1560

doi: 10.1093/bioinformatics/btu856

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Sequence analysis

Shifted Hamming distance: a fast and accurate SIMD-friendly filter to accelerate alignment verification in read mapping

Hongyi Xin^{1,*}, John Greth², John Emmons², Gennady Pekhimenko¹, Carl Kingsford³, Can Alkan^{4,*} and Onur Mutlu^{2,*}

Xin+, "Shifted Hamming Distance: A Fast and Accurate SIMD-friendly Filter to Accelerate Alignment Verification in Read Mapping", Bioinformatics 2015.

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Location Filtering

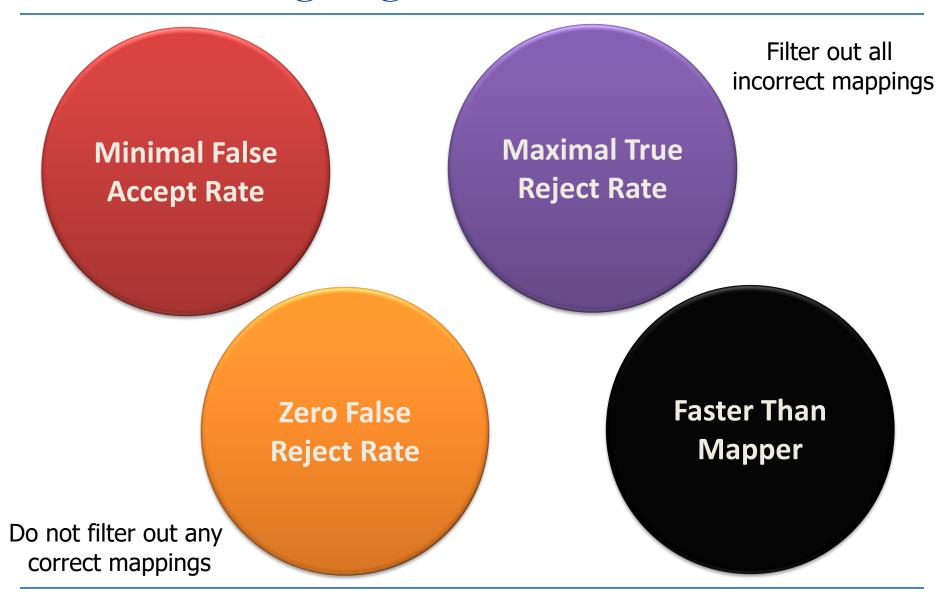
- Alignment is expensive
 - We need to align millions to billions of reads

Our goal is to accelerate read mapping by improving the filtering step

out moments quickly

Both methods are used by mappers today, but filtering has replaced alignment as the bottleneck [Xin+, BMC Genomics 2013]

Ideal Filtering Algorithm



Alignment vs. Pre-alignment (Filtering)



CTATAATACG

SHD

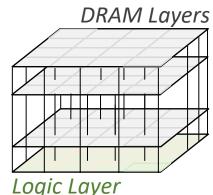
CTATAATACG

| 1 | 1 | 0 | | | | | | | |
|---|---|---------|-----------------------|---|---|---|---|---|---|
| 0 | 1 | 1 | 1 | | | | | | |
| 1 | 0 | 1 | 0 | 1 | | | | | |
| | 1 | 0 | 1 | 0 | 0 | | | | |
| | | 1 | 0 | 1 | 1 | 0 | | | |
| | | | 1 | 0 | 0 | 1 | 0 | | |
| | 0 | 0 1 1 0 | 0 1 1 1 0 1 1 0 | 0 1 1 1 1 0 1 0 1 0 1 | 0 1 1 1 1 0 1 0 1 1 0 1 0 1 0 1 0 1 0 1 | 0 1 1 1 1 0 1 0 1 1 0 1 0 0 1 0 1 0 1 | 0 1 1 1 1 0 1 0 1 1 0 1 0 0 1 0 1 1 0 | 0 1 1 1 1 0 1 0 1 1 0 1 0 0 1 0 1 1 0 | 0 1 1 1 1 0 1 0 1 1 0 1 0 0 1 0 1 1 0 |

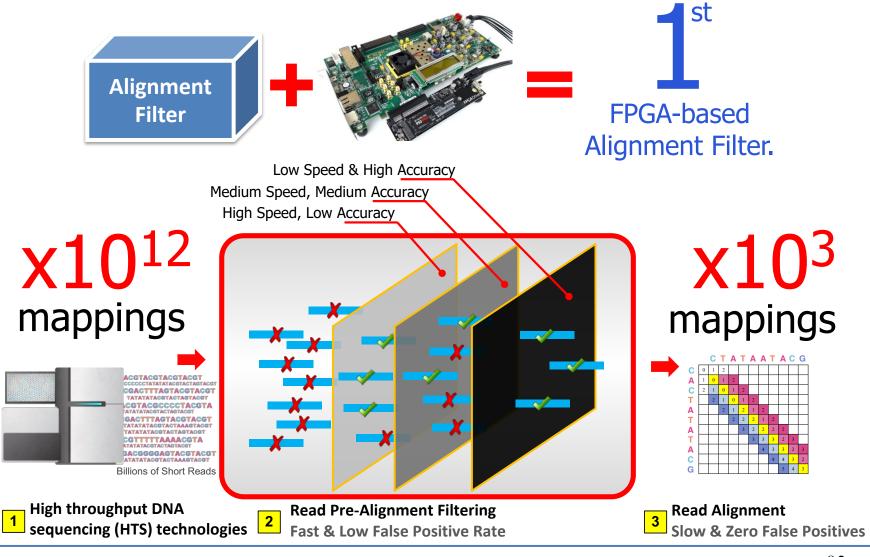
Independent vectors can be processed in parallel using hardware technologies







GateKeeper: FPGA-Based Alignment Filtering



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 Mohammed Alser, Hasan Hassan, Hongyi Xin, Oguz Ergin, Onur Mutlu, and Can Alkan

"GateKeeper: A New Hardware Architecture for Accelerating Pre-Alignment in DNA Short Read Mapping" Bioinformatics, [published online, May 31], 2017.

Source Code

[Online link at Bioinformatics Journal]

GateKeeper: a new hardware architecture for accelerating pre-alignment in DNA short read mapping

Mohammed Alser ™, Hasan Hassan, Hongyi Xin, Oğuz Ergin, Onur Mutlu ™, Can Alkan ™

Bioinformatics, Volume 33, Issue 21, 1 November 2017, Pages 3355–3363,

https://doi.org/10.1093/bioinformatics/btx342

Published: 31 May 2017 Article history ▼

SAFARI

GateKeeper Walkthrough

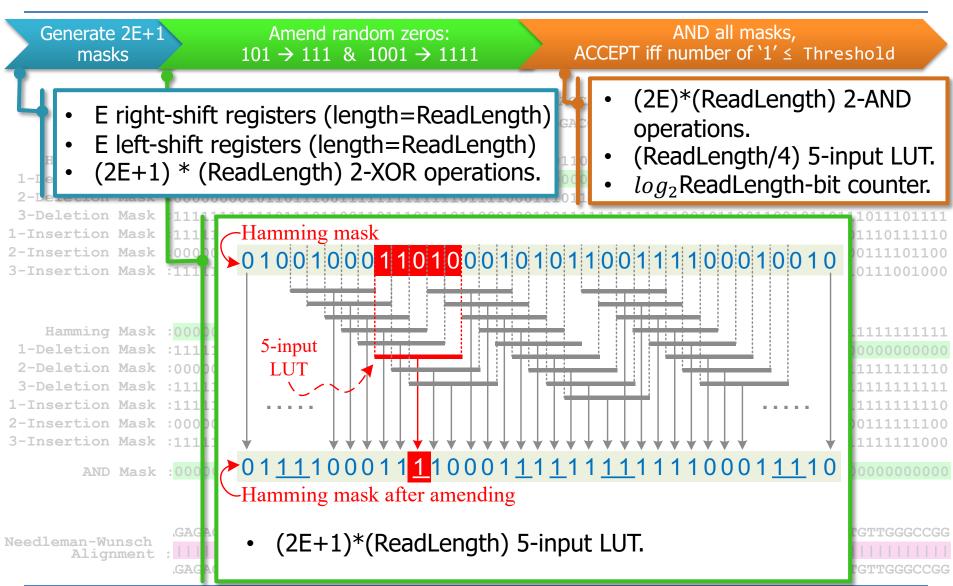
Generate 2E+1

Amend random zeros: $101 \rightarrow 111 \& 1001 \rightarrow 1111$

AND all masks, ACCEPT iff number of `1' ≤ Threshold

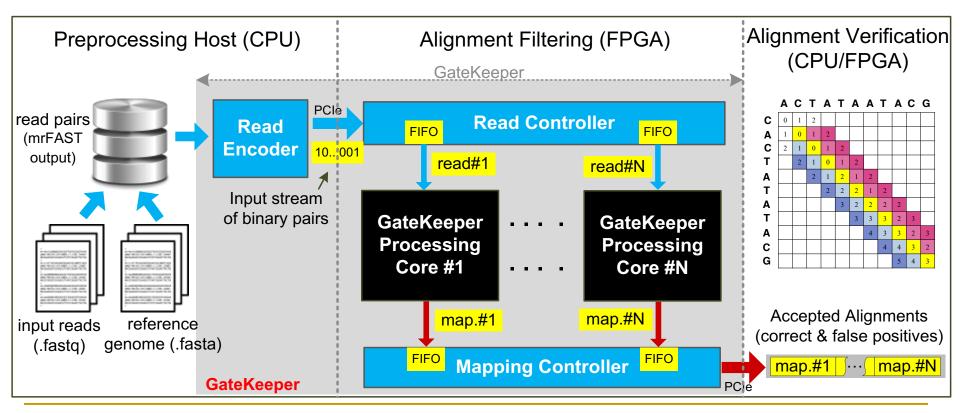
```
:GAGAGAGATATTTAGTGTTGCAGCACTACAACACAAAAGAGGACCAACTTACGTGTCTAAAAGGGGGAACATTGTTGGGCCGGA
 Reference
   GAGAGAGATAGTTAGTGTTGCAGCCACTACAACACAAAAGAGGACCAACTTACGTGTCTAAAAGGGGAGACATTGTTGGGCCGG
--- Masks after amendment ---
GAGAGAGATATTTAGTGTTGCAG-CACTACAACACAAAAGAGGACCAACTTACGTGTCTAAAAAGGGGGAACATTGTTGGGCCGG
Needleman-Wunsch
 Alignment
   GAGAGAGATAGTTAGTGTTGCAGCCACTACAACACAAAAGAGGACCAACTTACGTGTCTAAAAAGGGGAGACATTGTTGGGCCGG
```

GateKeeper Walkthrough (cont'd)

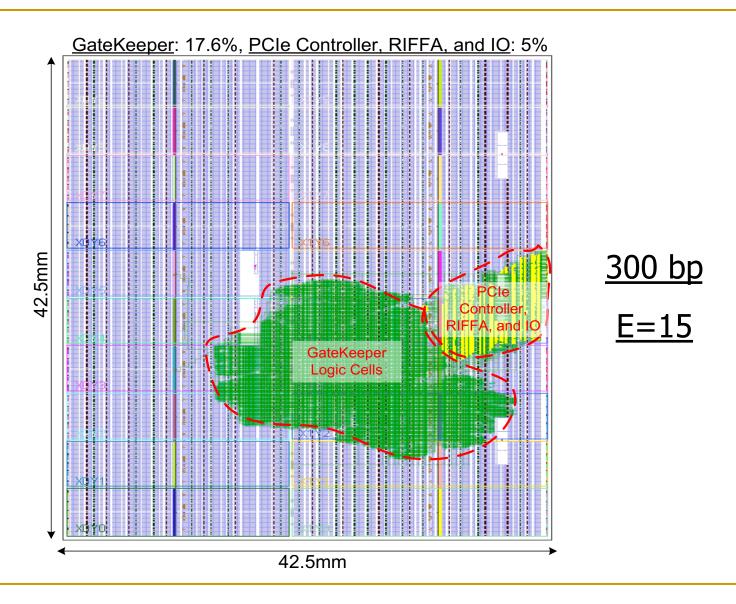


GateKeeper Accelerator Architecture

- Maximum data throughput =~13.3 billion bases/sec
- Can examine 8 (300 bp) or 16 (100 bp) mappings concurrently at 250 MHz
- Occupies 50% (100 bp) to 91% (300 bp) of the FPGA slice LUTs and registers



FPGA Chip Layout



GateKeeper vs. SHD

GateKeeper

- FPGA (Xilinx VC709)
- Multi-core (parallel)
- Examines a single mapping @ 125 MHz
- Limited to PCIe Gen3(4x) transfer rate (128 bits @ 250MHz)
- Amending requires:
 - (2E+1) 5-input LUT.

SHD

- Intel SIMD
- Single-core (sequential)
- Examines a single mapping @ ~2MHz
- Limited to a read length of 128 bp (SSE register size)
- Amending requires:
 - \bullet 4(2E+1) bitwise OR.
 - 4(2E+1) packed shuffle.
 - □ 3(2E+1) shift.

GateKeeper: Speed & Accuracy Results

90x-130x faster filter

than SHD (Xin et al., 2015) and the Adjacency Filter (Xin et al., 2013)

4x lower false accept rate

than the Adjacency Filter (Xin et al., 2013)

10x speedup in read mapping

with the addition of GateKeeper to the mrFAST mapper (Alkan et al., 2009)

Freely available online

github.com/BilkentCompGen/GateKeeper

GateKeeper Conclusions

- FPGA-based pre-alignment greatly speeds up read mapping
 - 10x speedup of a state-of-the-art mapper (mrFAST)

- FPGA-based pre-alignment can be integrated with the sequencer
 - It can help to hide the complexity and details of the FPGA
 - Enables real-time filtering while sequencing

More on GateKeeper

 Mohammed Alser, Hasan Hassan, Hongyi Xin, Oguz Ergin, Onur Mutlu, and Can Alkan

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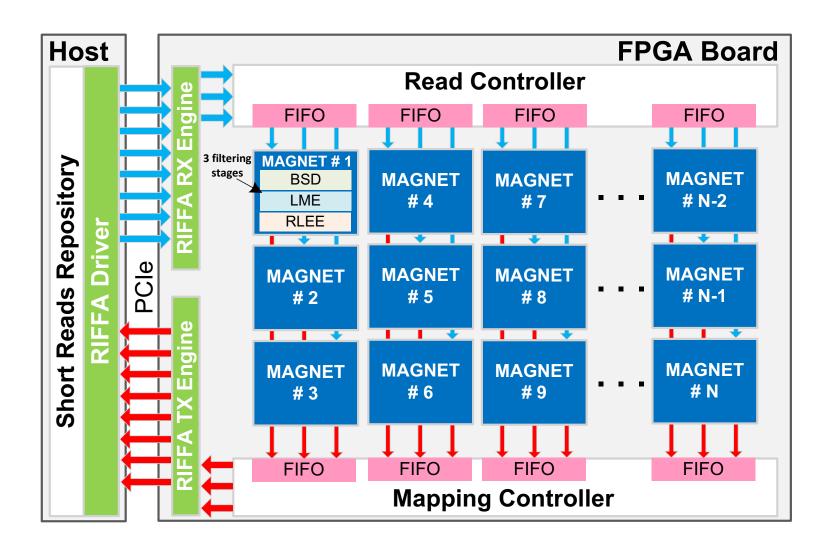
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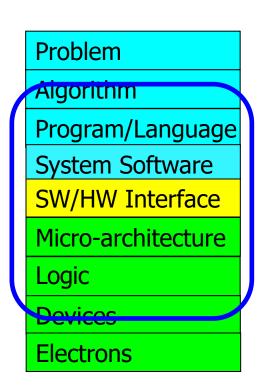
MAGNET Accelerator [Alser+, TIR 2017]



Faster, More Accurate, More Scalable Pre-Alignment Filtering

Algorithm-Arch-Device Co-Design is Critical

Computer Architecture (expanded view)



Shouji (障子) [Alser+, Bioinformatics 2019]

Mohammed Alser, Hasan Hassan, Akash Kumar, Onur Mutlu, and Can Alkan, "Shouji: A Fast and Efficient Pre-Alignment Filter for Sequence Alignment" Bioinformatics, [published online, March 28], 2019.

Source Code

Online link at Bioinformatics Journal

Bioinformatics, 2019, 1–9 doi: 10.1093/bioinformatics/btz234 Advance Access Publication Date: 28 March 2019 Original Paper



Sequence alignment

Shouji: a fast and efficient pre-alignment filter for sequence alignment

Mohammed Alser^{1,2,3,*}, Hasan Hassan¹, Akash Kumar², Onur Mutlu^{1,3,*} and Can Alkan^{3,*}

¹Computer Science Department, ETH Zürich, Zürich 8092, Switzerland, ²Chair for Processor Design, Center For Advancing Electronics Dresden, Institute of Computer Engineering, Technische Universität Dresden, 01062 Dresden, Germany and ³Computer Engineering Department, Bilkent University, 06800 Ankara, Turkey

*To whom correspondence should be addressed.

Associate Editor: Inanc Birol

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Shouji

Key observation:

- Correct alignment always includes long identical subsequences.
- Processing the entire mapping at once is ineffective for hardware design.

Key idea:

 Use an **overlapping** sliding window approach to quickly and accurately find all long segments of consecutive zeros.

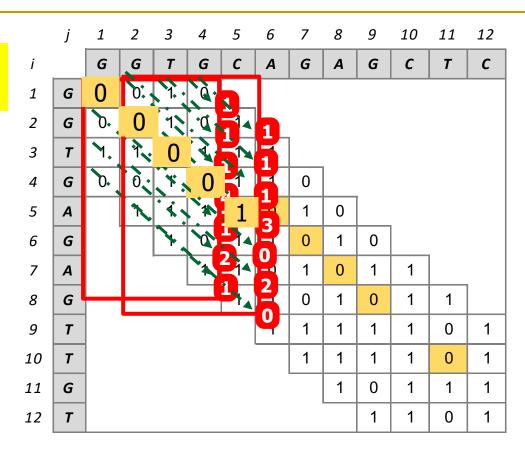
Key result:

- Shouji accelerates the best-performing CPU read aligner Edlib (Bioinformatics 2017) by up to 18.8x using 16 filtering units that work in parallel.
- Shouji on FPGA is up to 10,000x faster than on CPU.
- Shouji is 2.4x to 467x more accurate than GateKeeper (Bioinformatics 2017) and SHD (Bioinformatics 2015).

Shouji Walkthrough

Build the Neighborhood Map

Find all common subsequences (diagonal segments of consecutive zeros) shared between two given sequences.



Store longest subsequence in Shouji Bit-vector

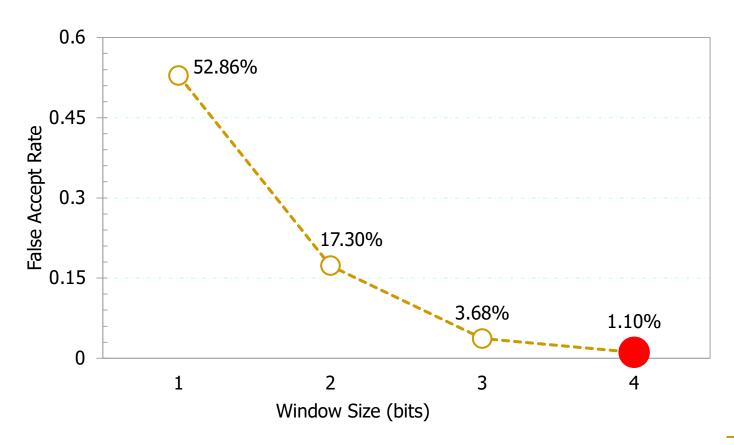


ACCEPT iff number of '1's ≤ Threshold

<u>Shouji: a fast and efficient pre-alignment filter for sequence alignment</u>, *Bioinformatics* 2019, https://doi.org/10.1093/bioinformatics/btz234

Effect of Sliding Window Size

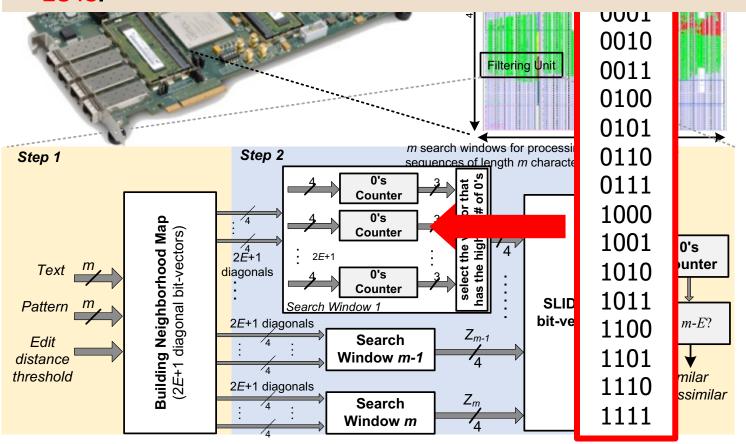
- Large enough window to accurately capture longer streaks of matches → lower false positives
- Small enough window to perform fast computation





Hardware Implementation

 Counting is performed concurrently for all bit-vectors and all sliding windows in a single clock cycle using multiple 4-input LUTs.



More on Shouji (障子) [Alser+, Bioinformatics 2019]

Mohammed Alser, Hasan Hassan, Akash Kumar, Onur Mutlu, and Can Alkan, "Shouji: A Fast and Efficient Pre-Alignment Filter for Sequence Alignment" Bioinformatics, [published online, March 28], 2019.

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Associate Editor: Inanc Birol

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SneakySnake [Alser+, Bioinformatics 2020]

Mohammed Alser, Taha Shahroodi, Juan-Gomez Luna, Can Alkan, and Onur Mutlu, "SneakySnake: A Fast and Accurate Universal Genome Pre-Alignment Filter for CPUs, GPUs, and FPGAs"

Bioinformatics, to appear in 2020.

[Source Code]

Online link at Bioinformatics Journal

Bioinformatics

doi.10.1093/bioinformatics/xxxxxx

Advance Access Publication Date: Day Month Year

Manuscript Category



Subject Section

SneakySnake: A Fast and Accurate Universal Genome Pre-Alignment Filter for CPUs, GPUs, and FPGAs

Mohammed Alser^{1,2,*}, Taha Shahroodi¹, Juan Gómez-Luna^{1,2}, Can Alkan^{4,*}, and Onur Mutlu^{1,2,3,4,*}

¹Department of Computer Science, ETH Zurich, Zurich 8006, Switzerland

²Department of Information Technology and Electrical Engineering, ETH Zurich, Zurich 8006, Switzerland

³Department of Electrical and Computer Engineering, Carnegie Mellon University, Pittsburgh 15213, PA, USA

⁴Department of Computer Engineering, Bilkent University, Ankara 06800, Turkey

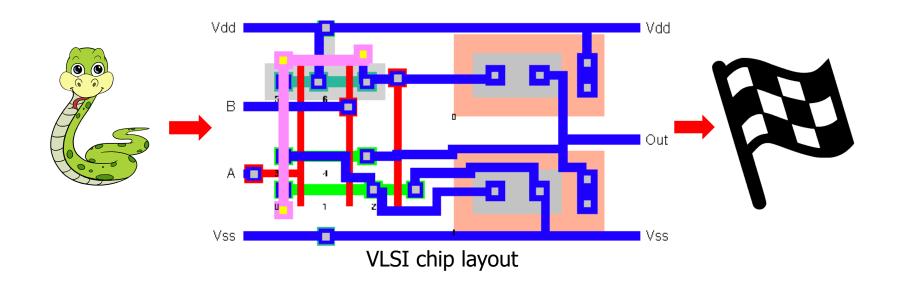
SneakySnake

Key observation:

Correct alignment is a sequence of non-overlapping long matches.

Key idea:

 Reduce the approximate string matching problem to the Single Net Routing problem in VLSI chip layout.



SneakySnake

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Key idea:

 Reduce the approximate string matching problem to the Single Net Routing problem in VLSI chip layout.

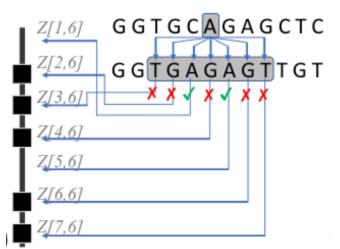
Key result:

- SneakySnake is up to four orders of magnitude more accurate than Shouji (Bioinformatics'19) and GateKeeper (Bioinformatics'17).
- SneakySnake greatly accelerates state-of-the-art CPU sequence aligners, Edlib (Bioinformatics'17) and Parasail (BMC Bioinformatics'16)
 - \Box by up to 37.7× and 43.9× (>12× on average), on CPUs
 - \Box by up to 413× and 689× (>400× on average) with FPGA acceleration

Building Neighborhood Map

Finding the Optimal Routing Path

Examining the Snake Survival



| column | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
|--------------------------------|---|---|---|---|---|---|---|---|---|----|----|----|
| 3 rd Upper Diagonal | 1 | 1 | 1 | 0 | 1 | 1 | 0 | 0 | 0 | 1 | 1 | 1 |
| 2 nd Upper Diagonal | 1 | 1 | 1 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 0 | 1 |
| 1 st Upper Diagonal | 1 | 0 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 1 |
| Main Diagonal | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| 1 st Lower Diagonal | 0 | 1 | 1 | 1 | 1 | 0 | 0 | 1 | 1 | 1 | 0 | 1 |
| 2 nd Lower Diagonal | 1 | 0 | 1 | 0 | 1 | 1 | 1 | 1 | 0 | 1 | 1 | 1 |
| 3 rd Lower Diagonal | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |

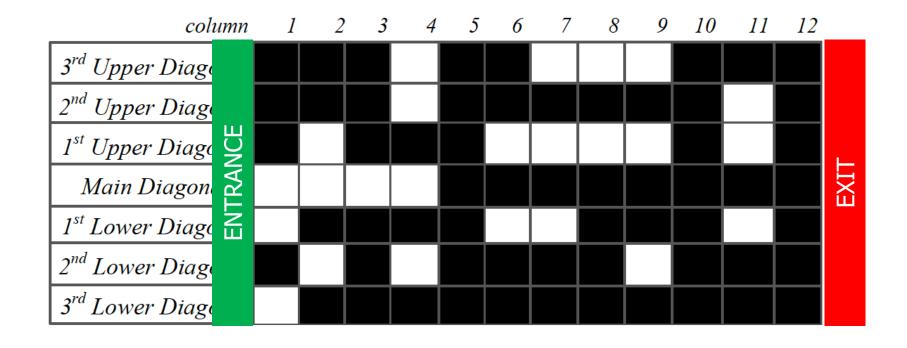
SAFARI

Building Neighborhood Map

Finding the Optimal Routing Path

Examining the Snake Survival

$$E = 3$$

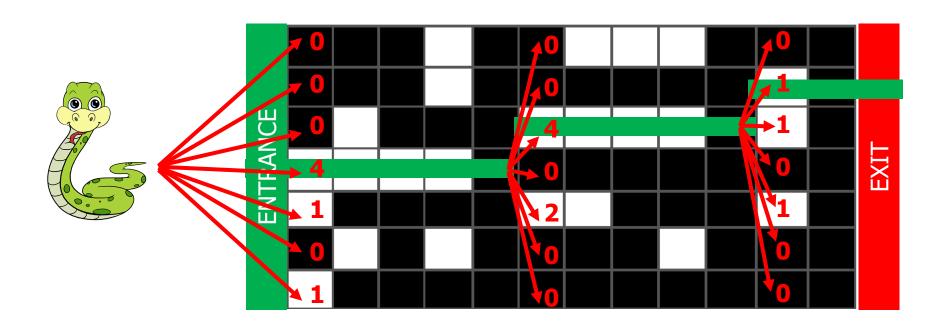


Building Neighborhood Map

Finding the Optimal Routing Path

Examining the Snake Survival





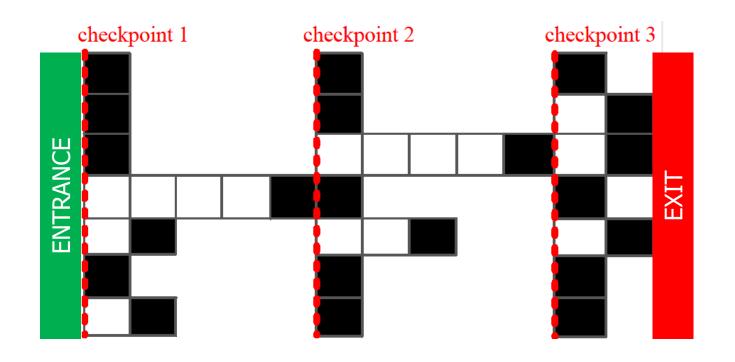
Building Neighborhood Map

Finding the Routing Travel Path

Examining the Snake Survival

This is what you actually need to build and it can be done on-the-fly!





FPGA Resource Analysis

 FPGA resource usage for a single filtering unit of GateKeeper, Shouji, and Snake-on-Chip for a sequence length of 100 and under different edit distance thresholds (E).

| | <i>E</i> (bp) | Slice LUT | Slice Register | No. of Filtering Units |
|---------------|---------------|-----------|----------------|------------------------|
| GateKeeper | 2 | 0.39% | 0.01% | 16 |
| | 5 | 0.71% | 0.01% | 16 |
| Shouji | 2 | 0.69% | 0.08% | 16 |
| | 5 | 1.72% | 0.16% | 16 |
| Snake-on-Chip | 2 | 0.68% | 0.16% | 16 |
| | 5 | 1.42% | 0.34% | 16 |
| | | | | |

Long Read Mapping (SneakySnake vs Parasail)

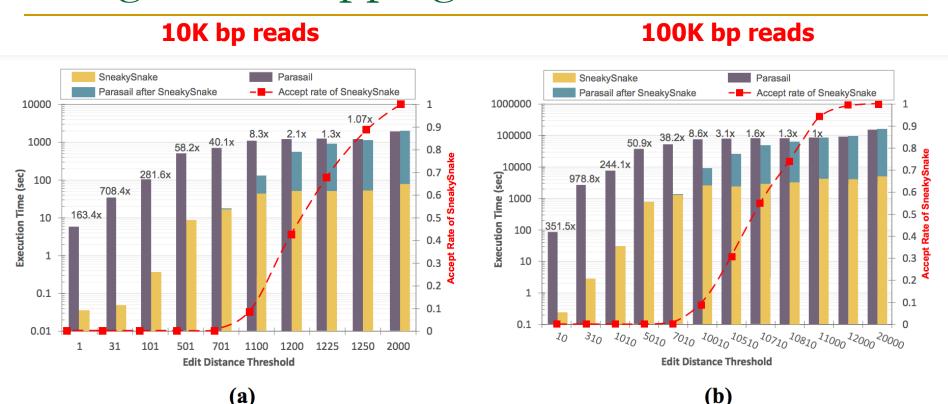


Fig. 10: The execution time of SneakySnake, Parasail, and SneakySnake integrated with Parasail using long sequences, (a) 10Kbp and (b) 100Kbp, and 40 CPU threads. The left y-axes of (a) and (b) are on a logarithmic scale. For each edit distance threshold value, we provide in the right y-axes of (a) and (b) the rate of accepted pairs (out of 100,000 pairs for 10Kbp and out of 74,687 pairs for 100Kbp) by SneakySnake that are passed to Parasail. We present the end-to-end speedup values obtained by integrating SneakySnake with Parasail.

Long Read Mapping (SneakySnake vs KSW2)



Fig. 11: The execution time of SneakySnake, KSW2, and SneakySnake integrated with KSW2 using long sequences, (a) 10Kbp and (b) 100Kbp, and a single CPU thread. The left y-axes of (a) and (b) are on a logarithmic scale. For each edit distance threshold value, we provide in the right y-axes of (a) and (b) the rate of accepted pairs (out of 100,000 pairs for 10Kbp and out of 74,687 pairs for 100Kbp) by SneakySnake that are passed to KSW2. We present the end-to-end speedup values obtained by integrating SneakySnake with KSW2.

More on SneakySnake [Alser+, Bioinformatics 2020]

Mohammed Alser, Taha Shahroodi, Juan-Gomez Luna, Can Alkan, and Onur Mutlu, "SneakySnake: A Fast and Accurate Universal Genome Pre-Alignment Filter for CPUs, GPUs, and FPGAs"

Bioinformatics, to appear in 2020.

[Source Code]

[Online link at Bioinformatics Journal]

Bioinformatics

doi.10.1093/bioinformatics/xxxxxx

Advance Access Publication Date: Day Month Year

Manuscript Category



Subject Section

SneakySnake: A Fast and Accurate Universal Genome Pre-Alignment Filter for CPUs, GPUs, and FPGAs

Mohammed Alser ^{1,2,*}, Taha Shahroodi ¹, Juan Gómez-Luna ^{1,2}, Can Alkan ^{4,*}, and Onur Mutlu ^{1,2,3,4,*}

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³Department of Electrical and Computer Engineering, Carnegie Mellon University, Pittsburgh 15213, PA, USA

GenASM Framework [MICRO 2020]

Damla Senol Cali, Gurpreet S. Kalsi, Zulal Bingol, Can Firtina, Lavanya Subramanian, Jeremie S. Kim, Rachata Ausavarungnirun, Mohammed Alser, Juan Gomez-Luna, Amirali Boroumand, Anant Nori, Allison Scibisz, Sreenivas Subramoney, Can Alkan, Saugata Ghose, and Onur Mutlu, "GenASM: A High-Performance, Low-Power Approximate String Matching Acceleration Framework for Genome Sequence Analysis"
Proceedings of the 53rd International Symposium on Microarchitecture (MICRO), Virtual, October 2020.

[<u>Lighting Talk Video</u> (1.5 minutes)]
[<u>Lightning Talk Slides (pptx) (pdf)</u>]
[<u>Talk Video</u> (18 minutes)]
[<u>Slides (pptx) (pdf)</u>]

GenASM: A High-Performance, Low-Power Approximate String Matching Acceleration Framework for Genome Sequence Analysis

Damla Senol Cali^{†™} Gurpreet S. Kalsi[™] Zülal Bingöl[▽] Can Firtina[⋄] Lavanya Subramanian[‡] Jeremie S. Kim^{⋄†} Rachata Ausavarungnirun[⊙] Mohammed Alser[⋄] Juan Gomez-Luna[⋄] Amirali Boroumand[†] Anant Nori[™] Allison Scibisz[†] Sreenivas Subramoney[™] Can Alkan[▽] Saugata Ghose^{*†} Onur Mutlu^{⋄†▽}

† Carnegie Mellon University [™] Processor Architecture Research Lab, Intel Labs [▽] Bilkent University [⋄] ETH Zürich

‡ Facebook [⊙] King Mongkut's University of Technology North Bangkok ^{*} University of Illinois at Urbana–Champaign

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Problem & Our Goal

- ☐ Multiple steps of read mapping require *approximate string matching*
 - ASM enables read mapping to account for sequencing errors and genetic variations in the reads
- □ ASM makes up a significant portion of read mapping (more than 70%)
- One of the major bottlenecks of genome sequence analysis

Our Goal:

Accelerate approximate string matching by designing a fast and flexible framework, which can be used to accelerate *multiple steps* of the genome sequence analysis pipeline

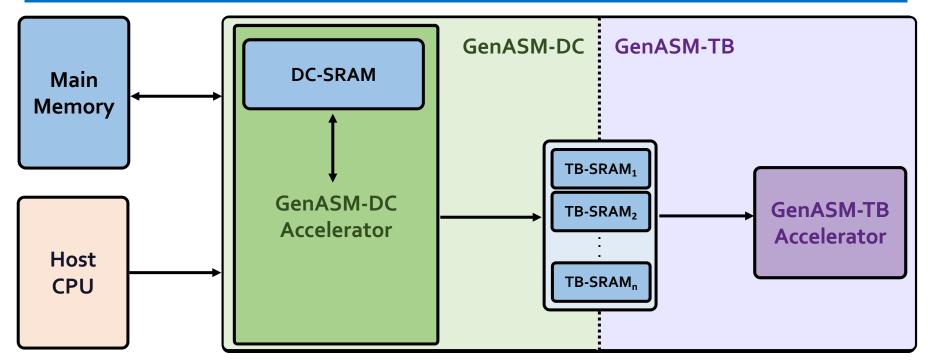
GenASM: ASM Framework for GSA

Our Goal:

Accelerate approximate string matching by designing a fast and flexible framework, which can accelerate *multiple steps* of genome sequence analysis

- GenASM: First ASM acceleration framework for GSA
 - Based on the Bitαp algorithm
 - Uses fast and simple bitwise operations to perform ASM
 - Modified and extended ASM algorithm
 - Highly-parallel Bitap with long read support
 - Bitvector-based novel algorithm to perform traceback
 - Co-design of our modified scalable and memory-efficient algorithms with low-power and area-efficient hardware accelerators

GenASM: Hardware Design



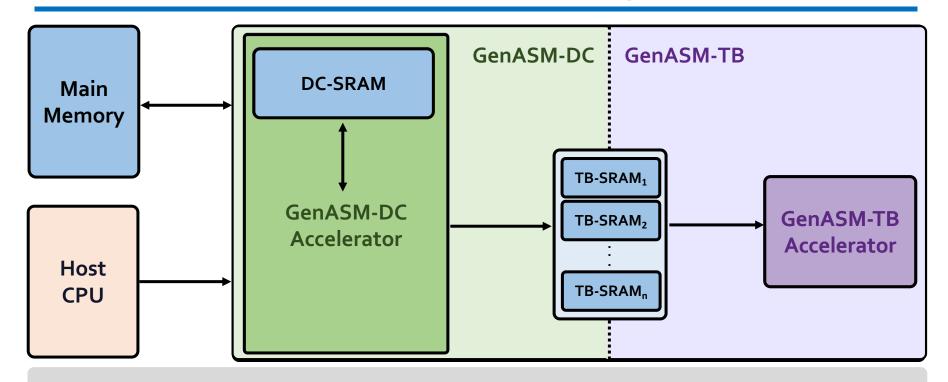
GenASM-DC:

generates bitvectors and performs edit Distance Calculation

GenASM-TB:

performs TraceBack and assembles the optimal alignment

GenASM: Hardware Design



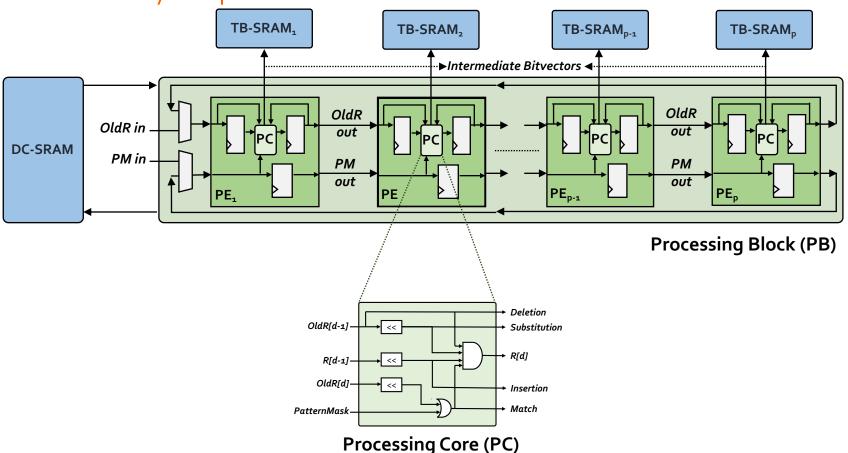
Our specialized compute units and on-chip SRAMs help us to:

- → Match the rate of computation with memory capacity and bandwidth
 - → Achieve high performance and power efficiency
 - → Scale linearly in performance with

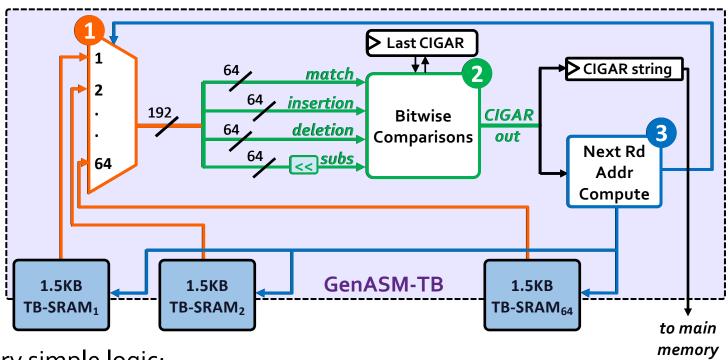
the number of parallel compute units that we add to the system

GenASM-DC: Hardware Design

- Linear cyclic systolic array based accelerator
 - Designed to maximize parallelism and minimize memory bandwidth and memory footprint



GenASM-TB: Hardware Design



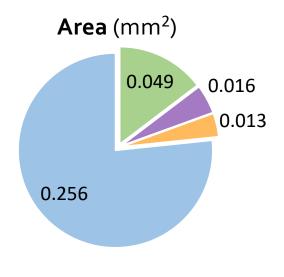
- Very simple logic:
 - 1 Reads the bitvectors from one of the TB-SRAMs using the computed address
 - 2 Performs the required bitwise comparisons to find the traceback output for the current position
 - 3 Computes the next TB-SRAM address to read the new set of bitvectors

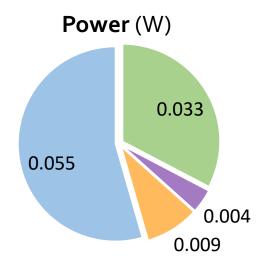
Key Results – Area and Power

- Based on our **synthesis** of **GenASM-DC** and **GenASM-TB** accelerator datapaths using the Synopsys Design Compiler with a **28nm** LP process:
 - Both GenASM-DC and GenASM-TB operate @ 1GHz



- GenASM-TB
- DC-SRAM (8 KB)
- TB-SRAMs (64 x 1.5 KB)





Total (1 vault): 0.334 mm²

Total (32 vaults): 10.69 mm²

% of a Xeon CPU core: 1%

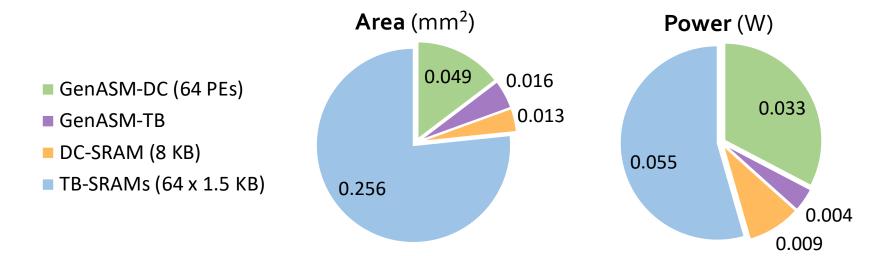
0.101 W

3.23 W

1%

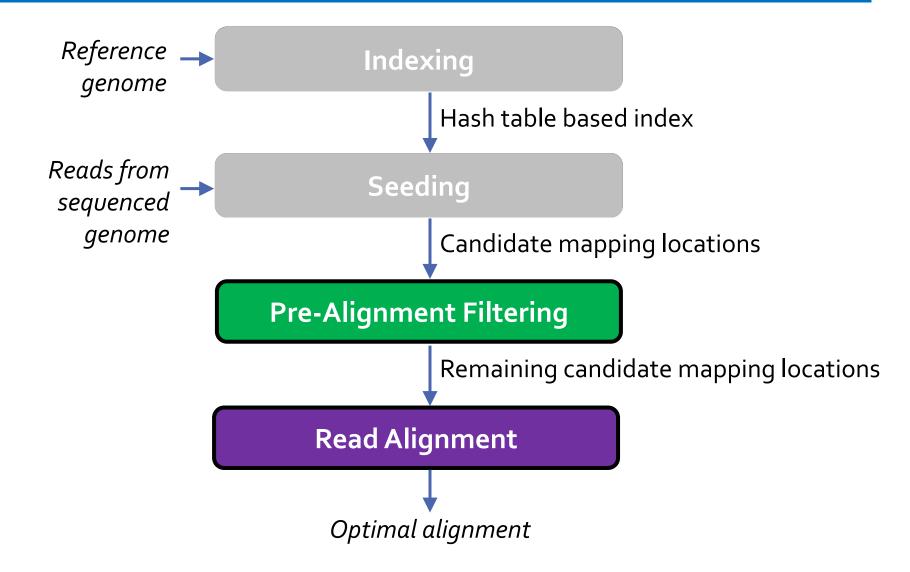
Key Results – Area and Power

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 - Both GenASM-DC and GenASM-TB operate @ 1GHz



GenASM has low area and power overheads

Use Cases of GenASM



Use Cases of GenASM (cont'd.)

(1) Read Alignment Step of Read Mapping

 Find the optimal alignment of how reads map to candidate reference regions

(2) Pre-Alignment Filtering for Short Reads

 Quickly identify and filter out the unlikely candidate reference regions for each read

(3) Edit Distance Calculation

- Measure the similarity or distance between two sequences
- We also discuss other possible use cases of GenASM in our paper:
 - Read-to-read overlap finding, hash-table based indexing, whole genome alignment, generic text search

Key Results

(1) Read Alignment

- 116× speedup, 37× less power than Minimap2 (state-of-the-art SW)
- 111× speedup, 33× less power than BWA-MEM (state-of-the-art SW)
- □ 3.9× better throughput, 2.7× less power than Darwin (state-of-the-art HW)
- 1.9× better throughput, 82% less logic power than GenAx (state-of-the-art HW)

(2) Pre-Alignment Filtering

□ 3.7× speedup, 1.7× less power than Shouji (state-of-the-art HW)

(3) Edit Distance Calculation

- □ 22-12501× speedup, 548-582× less power than Edlib (state-of-the-art SW)
- 9.3–400× speedup, 67× less power than ASAP (state-of-the-art HW)

More on GenASM Framework [MICRO 2020]

Damla Senol Cali, Gurpreet S. Kalsi, Zulal Bingol, Can Firtina, Lavanya Subramanian, Jeremie S. Kim, Rachata Ausavarungnirun, Mohammed Alser, Juan Gomez-Luna, Amirali Boroumand, Anant Nori, Allison Scibisz, Sreenivas Subramoney, Can Alkan, Saugata Ghose, and Onur Mutlu, "GenASM: A High-Performance, Low-Power Approximate String Matching Acceleration Framework for Genome Sequence Analysis"
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GenASM: A High-Performance, Low-Power Approximate String Matching Acceleration Framework for Genome Sequence Analysis

Damla Senol Cali^{†™} Gurpreet S. Kalsi[™] Zülal Bingöl[▽] Can Firtina[⋄] Lavanya Subramanian[‡] Jeremie S. Kim^{⋄†} Rachata Ausavarungnirun[⊙] Mohammed Alser[⋄] Juan Gomez-Luna[⋄] Amirali Boroumand[†] Anant Nori[™] Allison Scibisz[†] Sreenivas Subramoney[™] Can Alkan[▽] Saugata Ghose^{*†} Onur Mutlu^{⋄†▽}

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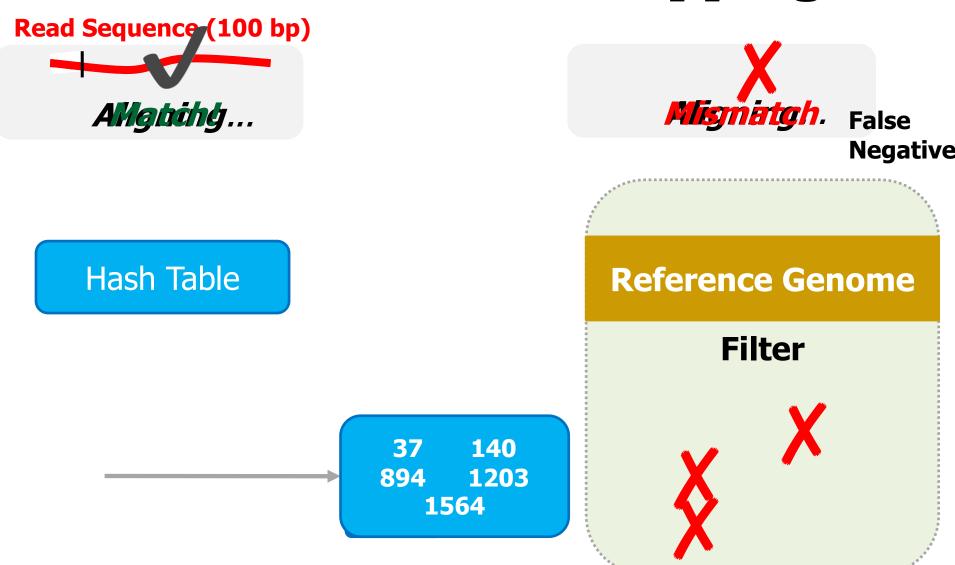
Agenda

- The Problem: DNA Read Mapping
 - State-of-the-art Read Mapper Design
- Algorithmic Acceleration
 - Exploiting Structure of the Genome
 - Exploiting SIMD Instructions
- Hardware Acceleration
 - Specialized Architectures
 - Processing in Memory
- Future Opportunities: New Sequencing Technologies

Read Mapping & Filtering

- Problem: Heavily bottlenecked by Data Movement
- GateKeeper, Shouji, SneakySnake performance limited by DRAM bandwidth [Alser+, Bioinformatics 2017,2019,2020]
- Ditto for SHD [Xin+, Bioinformatics 2015]
- Solution: Processing-in-memory can alleviate the bottleneck
- We need to design mapping & filtering algorithms to fit processing-in-memory

Hash Tables in Read Mapping



Read Mapping & Filtering in Memory

We need to design mapping & filtering algorithms that fit processing-in-memory

More on GRIM-Filter

 Jeremie S. Kim, Damla Senol Cali, Hongyi Xin, Donghyuk Lee, Saugata Ghose, Mohammed Alser, Hasan Hassan, Oguz Ergin, Can Alkan, and <u>Onur Mutlu</u>,

"GRIM-Filter: Fast Seed Location Filtering in DNA Read Mapping Using Processing-in-Memory Technologies"

BMC Genomics, 2018.

Proceedings of the <u>16th Asia Pacific Bioinformatics Conference</u> (**APBC**), Yokohama, Japan, January 2018.

[Slides (pptx) (pdf)]

[Source Code]

[arxiv.org Version (pdf)]

[Talk Video at AACBB 2019]

GRIM-Filter: Fast seed location filtering in DNA read mapping using processing-in-memory technologies

Jeremie S. Kim^{1,6*}, Damla Senol Cali¹, Hongyi Xin², Donghyuk Lee³, Saugata Ghose¹, Mohammed Alser⁴, Hasan Hassan⁶, Oguz Ergin⁵, Can Alkan^{4*} and Onur Mutlu^{6,1*}

From The Sixteenth Asia Pacific Bioinformatics Conference 2018 Yokohama, Japan. 15-17 January 2018



Our Proposal: GRIM-Filter

- 1. Data Structures: Bins & Bitvectors
- 2. Checking a Bin
- 3. Integrating GRIM-Filter into a Mapper

GRIM-Filter: Bins

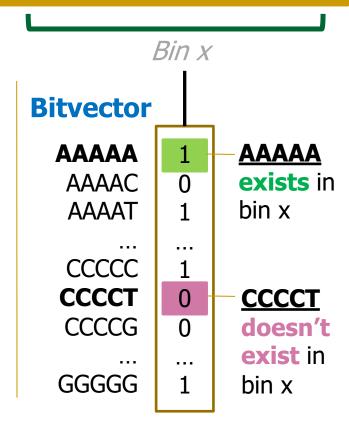
We partition the genome into large sequences (bins).

Bin x - 3

Bin x - 1

Bin x - 2

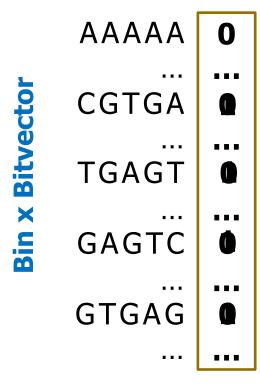
- Represent each bin with a bitvector that holds the occurrence of all permutations of a small string (token) in the bin
- To account for matches that straddle bins, we employ overlapping bins
 - A read will now always completely fall within a single bin



GRIM-Filter: Bitvectors

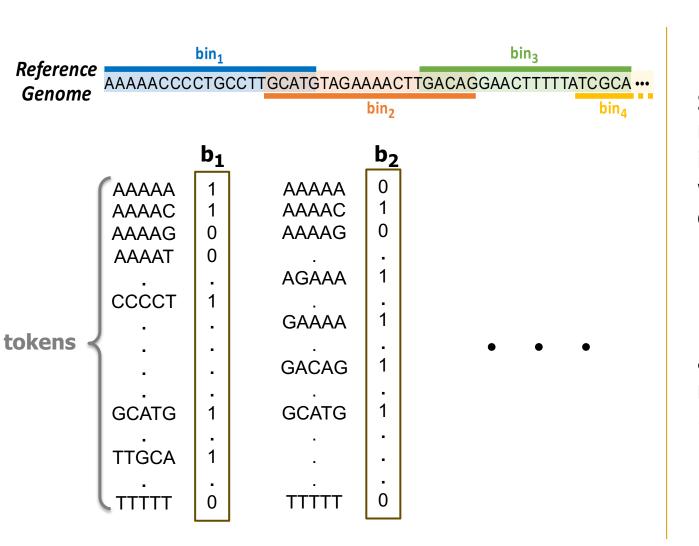


Bin x





GRIM-Filter: Bitvectors



Storing all bitvectors requires $4^n * t$ bits in memory, where t = number of bins.

For **bin size** ~200, and **n** = 5, **memory footprint** ~3.8 GB

Our Proposal: GRIM-Filter

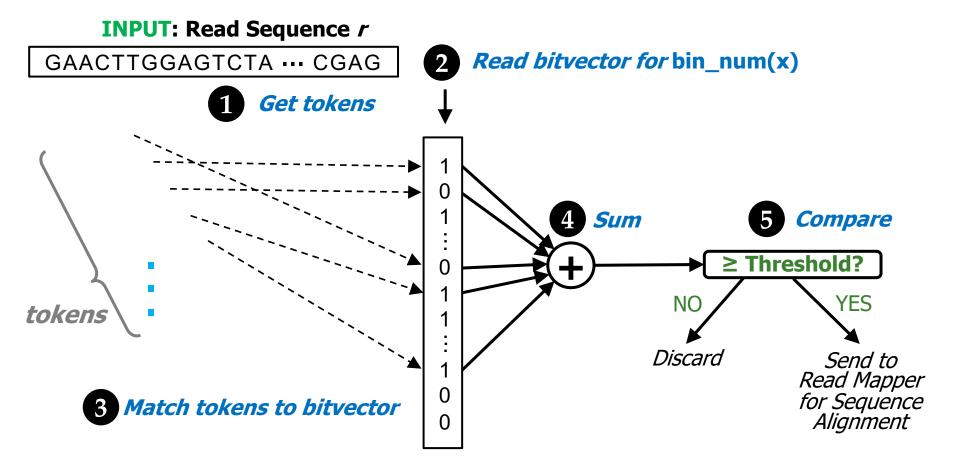
1. Data Structures: Bins & Bitvectors

- 2. Checking a Bin
- 3. Integrating GRIM-Filter into a Mapper



GRIM-Filter: Checking a Bin

How GRIM-Filter determines whether to **discard** potential match locations in a given bin **prior** to alignment

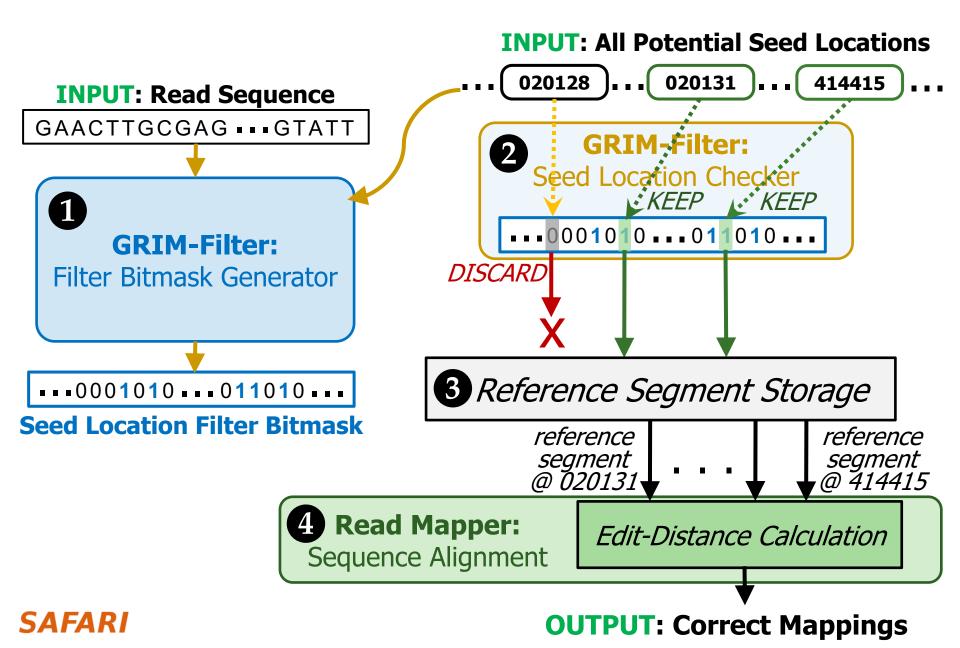




Our Proposal: GRIM-Filter

- 1. Data Structures: Bins & Bitvectors
- Checking a Bin
- 3. Integrating GRIM-Filter into a Mapper

Integrating GRIM-Filter into a Read Mapper



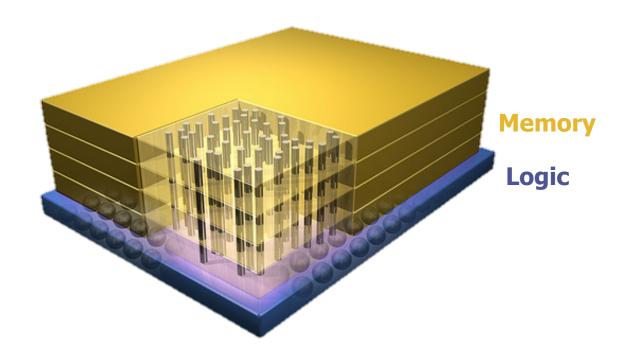
Key Properties of GRIM-Filter

1. Simple Operations:

- To check a given bin, find the sum of all bits corresponding to each token in the read
- Compare against threshold to determine whether to align
- 2. Highly Parallel: Each bin is operated on independently and there are many many bins
- 3. Memory Bound: Given the frequent accesses to the large bitvectors, we find that GRIM-Filter is memory bound

These properties together make GRIM-Filter a good algorithm to be run in 3D-Stacked DRAM

Opportunity: 3D-Stacked Logic+Memory



Other "True 3D" technologies under development

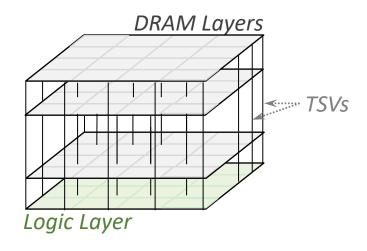
DRAM Landscape (circa 2015)

| Segment | DRAM Standards & Architectures |
|-------------|--|
| Commodity | DDR3 (2007) [14]; DDR4 (2012) [18] |
| Low-Power | LPDDR3 (2012) [17]; LPDDR4 (2014) [20] |
| Graphics | GDDR5 (2009) [15] |
| Performance | eDRAM [28], [32]; RLDRAM3 (2011) [29] |
| 3D-Stacked | WIO (2011) [16]; WIO2 (2014) [21]; MCDRAM (2015) [13]; HBM (2013) [19]; HMC1.0 (2013) [10]; HMC1.1 (2014) [11] |
| Academic | SBA/SSA (2010) [38]; Staged Reads (2012) [8]; RAIDR (2012) [27]; SALP (2012) [24]; TL-DRAM (2013) [26]; RowClone (2013) [37]; Half-DRAM (2014) [39]; Row-Buffer Decoupling (2014) [33]; SARP (2014) [6]; AL-DRAM (2015) [25] |

Table 1. Landscape of DRAM-based memory

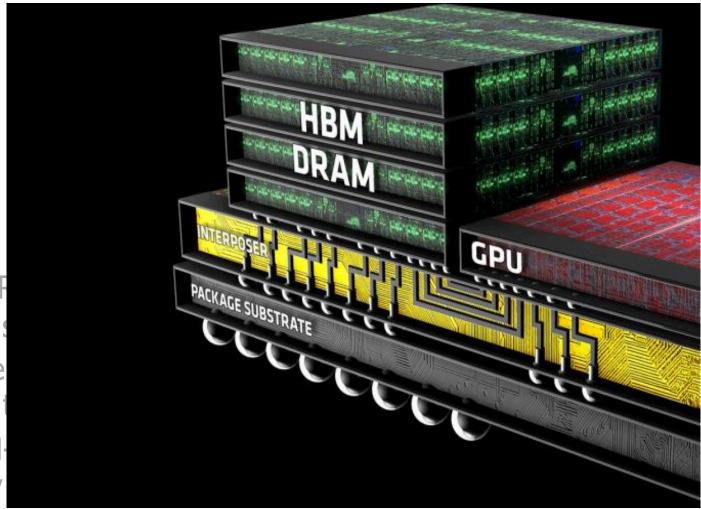
Kim+, "Ramulator: A Flexible and Extensible DRAM Simulator", IEEE CAL 2015.

3D-Stacked Memory



- 3D-Stacked DRAM architecture has extremely high bandwidth as well as a stacked customizable logic layer
 - Logic Layer enables Processing-in-Memory, via highbandwidth low-latency access to DRAM layers
 - Embed GRIM-Filter operations into DRAM logic layer and appropriately distribute bitvectors throughout memory

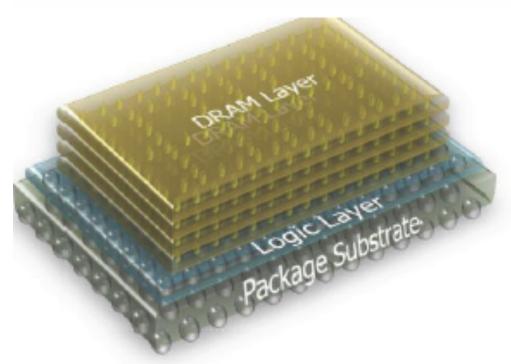
3D-Stacked Memory



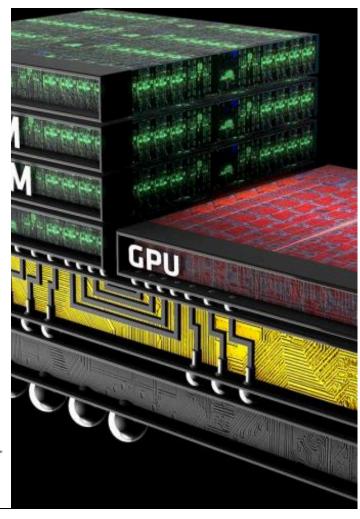
- 3D-Stacked DF bandwidth as
 - Logic Layer e computation t
 - Embed GRIMappropriately

3D-Stacked Memory

Micron's HMC

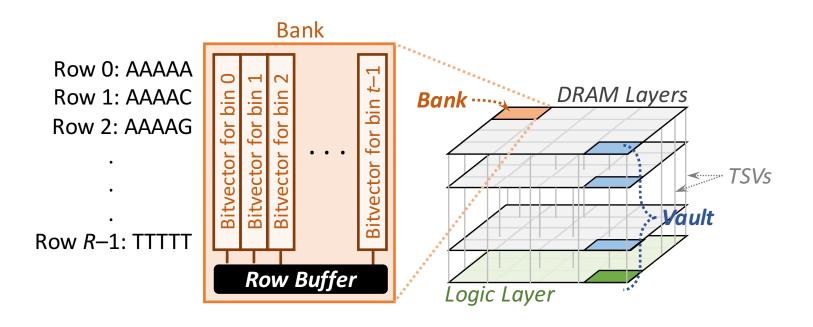


Micron has working demonstration components



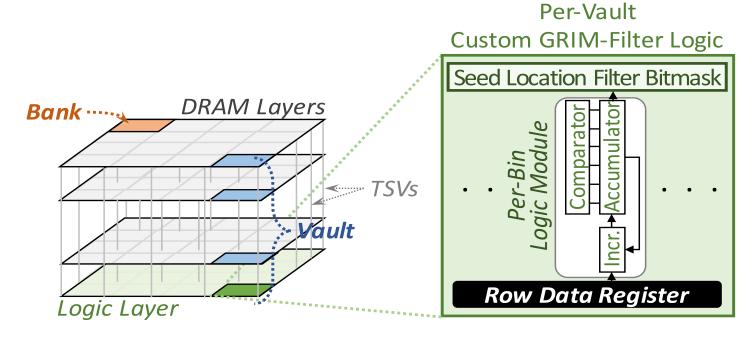
http://images.anandtech.com/doci/9266/HBMCar_678x452.jpg

GRIM-Filter in 3D-Stacked DRAM



- Each DRAM layer is organized as an array of banks
 - □ A bank is an array of cells with a row buffer to transfer data
- The layout of bitvectors in a bank enables filtering many bins in parallel

GRIM-Filter in 3D-Stacked DRAM



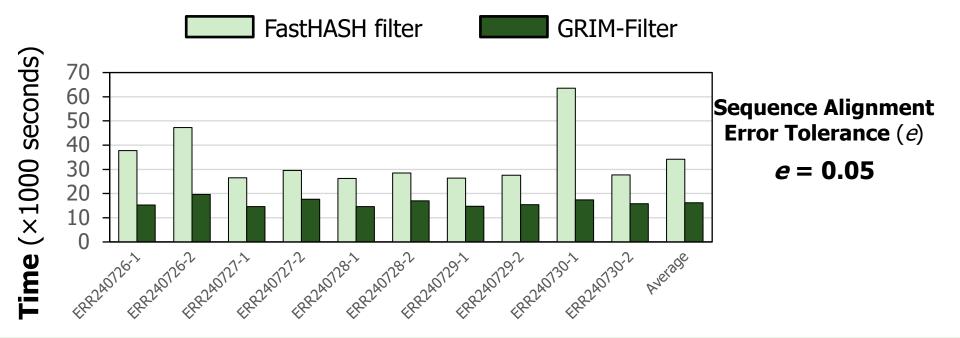
- Customized logic for accumulation and comparison per genome segment
 - Low area overhead, simple implementation
 - For HBM2, we use 4096 incrementer LUTs, 7-bit counters, and comparators in logic layer

Methodology

- Performance simulated using an in-house 3D-Stacked DRAM simulator
- Evaluate 10 real read data sets (From the 1000 Genomes Project)
 - Each data set consists of 4 million reads of length 100
- Evaluate two key metrics
 - Performance
 - False negative rate
 - The fraction of locations that pass the filter but result in a mismatch
- Compare against a state-of-the-art filter, FastHASH [xin+, BMC Genomics 2013] when using mrFAST, but GRIM-Filter can be used with ANY read mapper

GRIM-Filter Performance

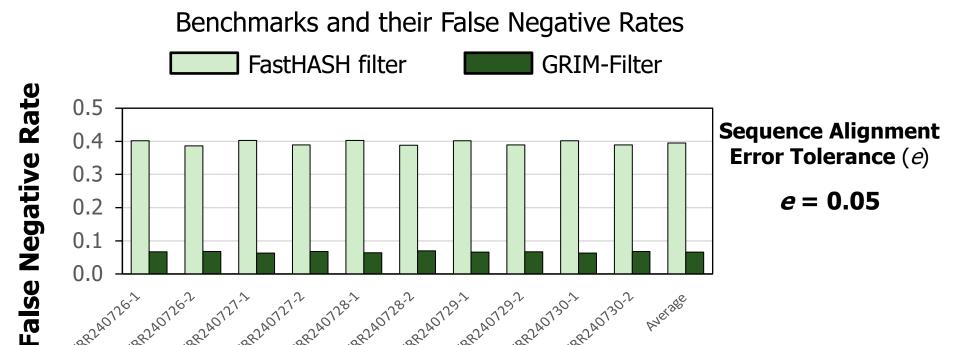




1.8x-3.7x performance benefit across real data sets
2.1x average performance benefit

GRIM-Filter gets performance due to its hardware-software co-design

GRIM-Filter False Negative Rate



5.6x-6.4x False Negative reduction across real data sets 6.0x average reduction in False Negative Rate

GRIM-Filter utilizes more information available in the read to filter

More on GRIM-Filter

 Jeremie S. Kim, Damla Senol Cali, Hongyi Xin, Donghyuk Lee, Saugata Ghose, Mohammed Alser, Hasan Hassan, Oguz Ergin, Can Alkan, and <u>Onur Mutlu</u>,

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From The Sixteenth Asia Pacific Bioinformatics Conference 2018 Yokohama, Japan. 15-17 January 2018



Aside: In-Memory Graph Processing

Large graphs are everywhere (circa 2015)



36 Million Wikipedia Pages



1.4 Billion Facebook Users

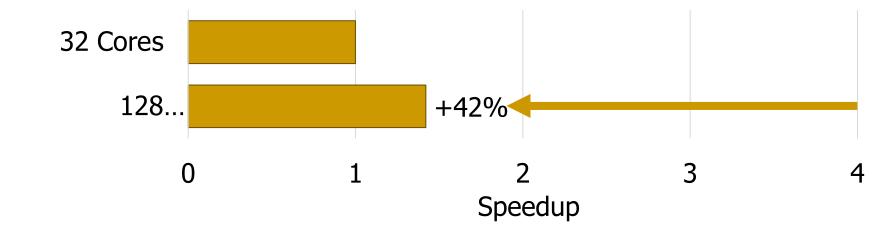


300 Million Twitter Users



30 Billion Instagram Photos

Scalable large-scale graph processing is challenging

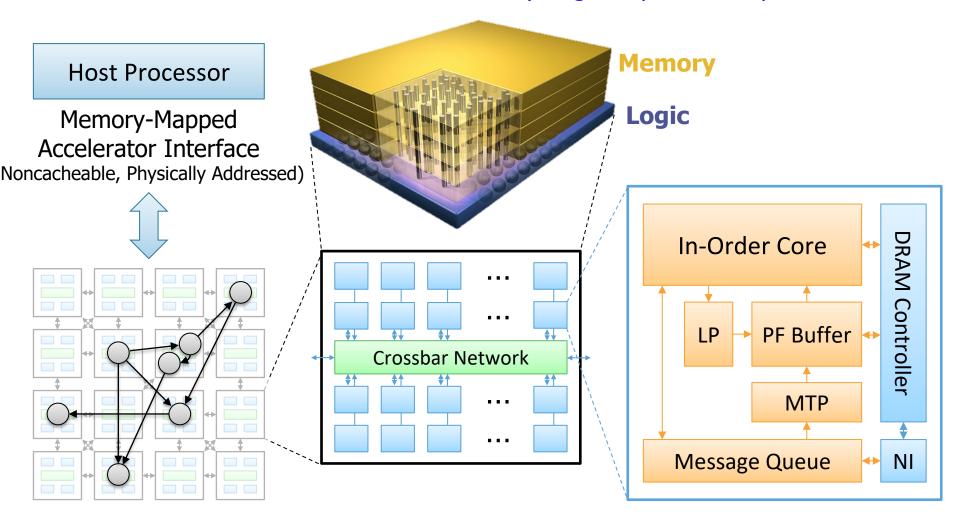


Key Bottlenecks in Graph Processing

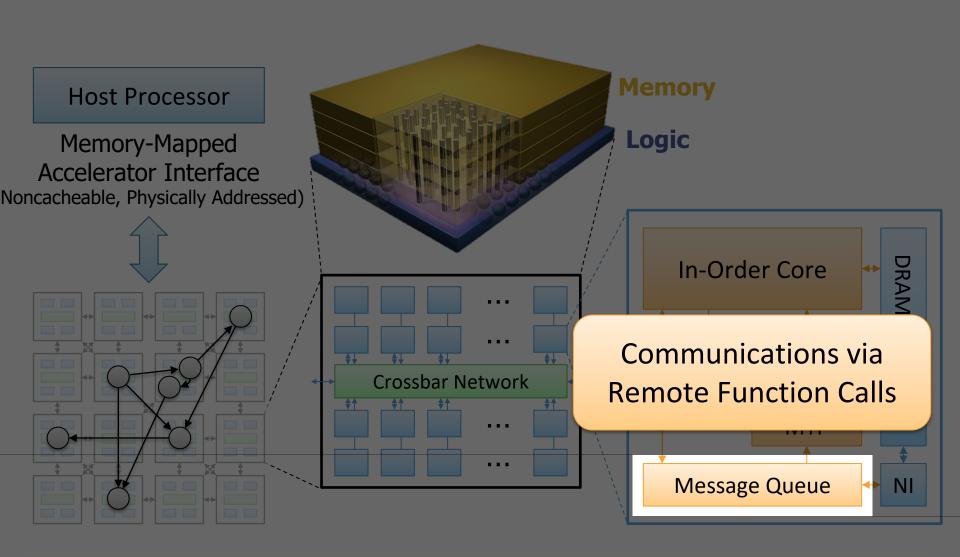
```
for (v: graph.vertices) {
     for (w: v.successors) {
       w.next rank += weight * v.rank;
                       1. Frequent random memory accesses
                                   &w
            V
 w.rank
w.next rank
                              weight * v.rank
 w.edges
            W
                              2. Little amount of computation
```

Tesseract System for Graph Processing

Interconnected set of 3D-stacked memory+logic chips with simple cores

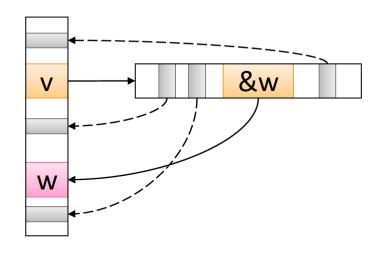


Tesseract System for Graph Processing



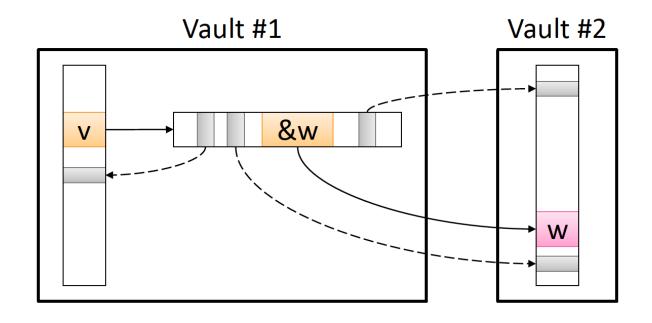
Communications In Tesseract (I)

```
for (v: graph.vertices) {
   for (w: v.successors) {
      w.next_rank += weight * v.rank;
   }
}
```



Communications In Tesseract (II)

```
for (v: graph.vertices) {
   for (w: v.successors) {
      w.next_rank += weight * v.rank;
   }
}
```



Communications In Tesseract (III)

```
for (v: graph.vertices) {
                              Non-blocking Remote Function Call
  for (w: v.successors) {
    put(w.id, function() { w.next_rank += weight * v.rank; });
                                 Can be delayed
                                 until the nearest barrier
barrier();
                  Vault #1
                                               Vault #2
                                         put
                           &w
         V
                put
                                         put
                                                  W
                                         put
```

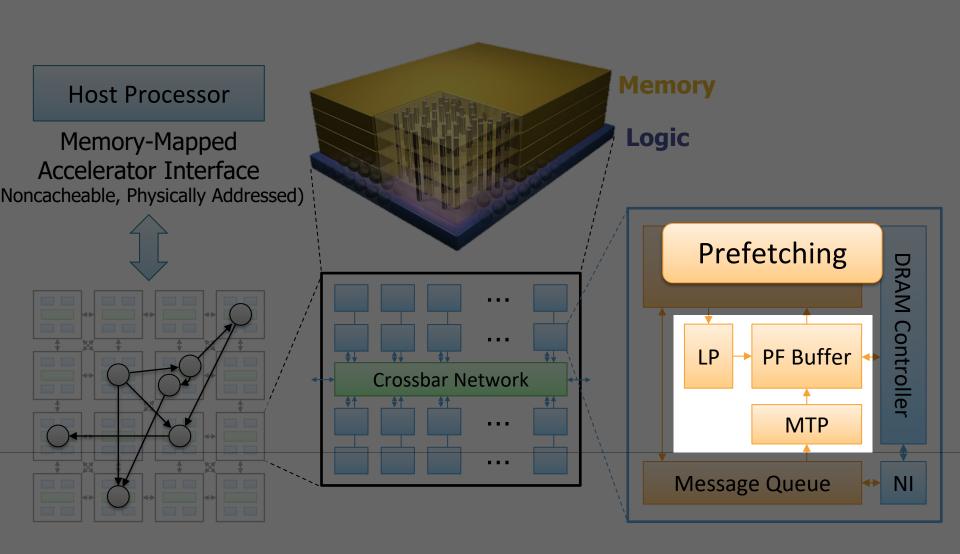
Remote Function Call (Non-Blocking)

- 1. Send function address & args to the remote core
- 2. Store the incoming message to the message queue
- Flush the message queue when it is full or a synchronization barrier is reached

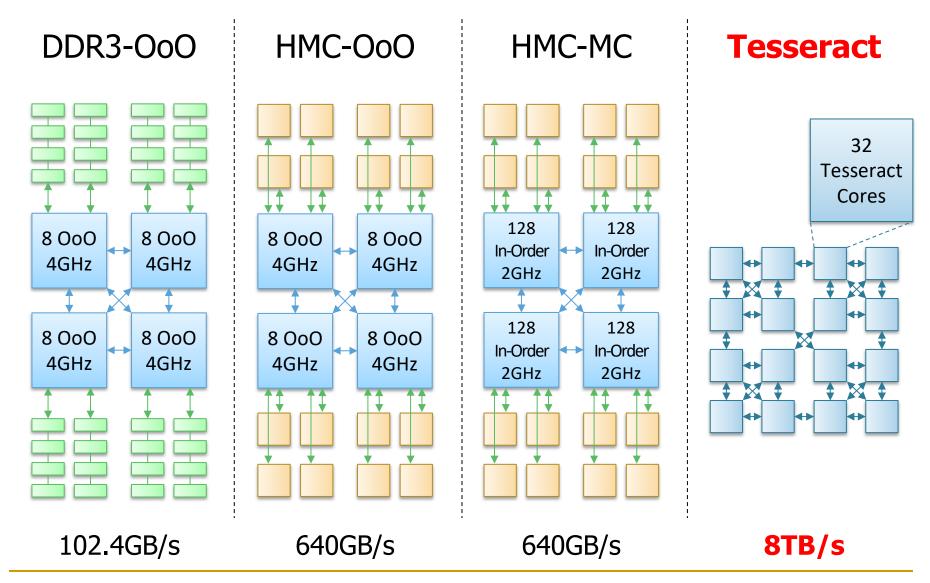


put(w.id, function() { w.next_rank += value; })

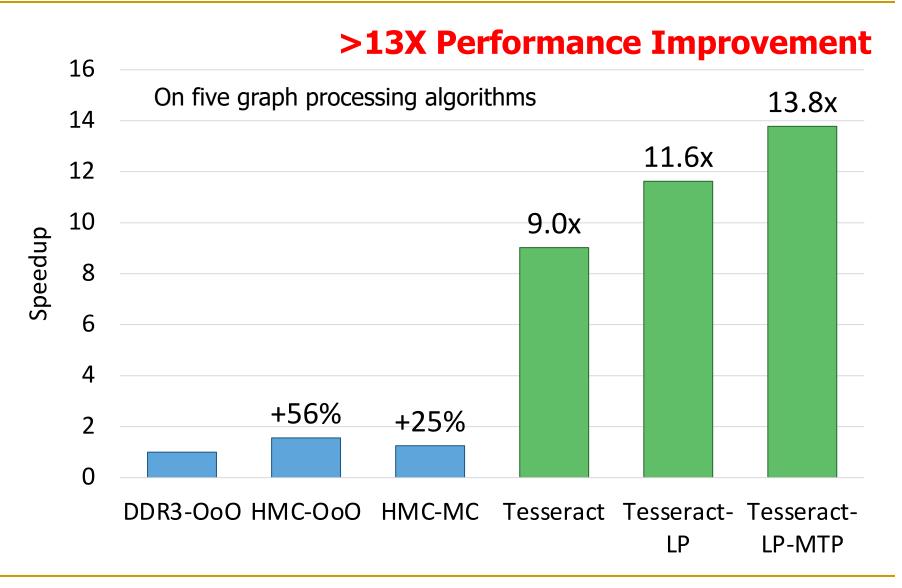
Tesseract System for Graph Processing



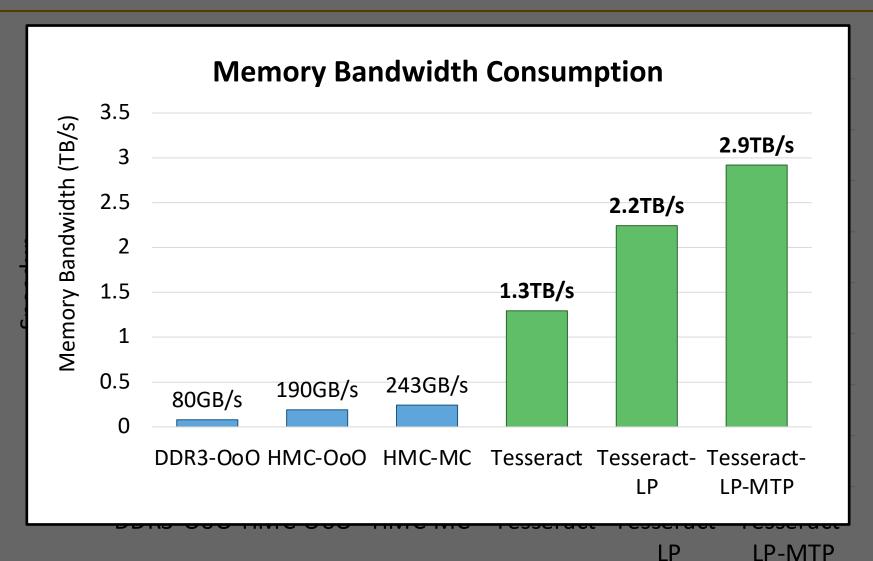
Evaluated Systems



Tesseract Graph Processing Performance

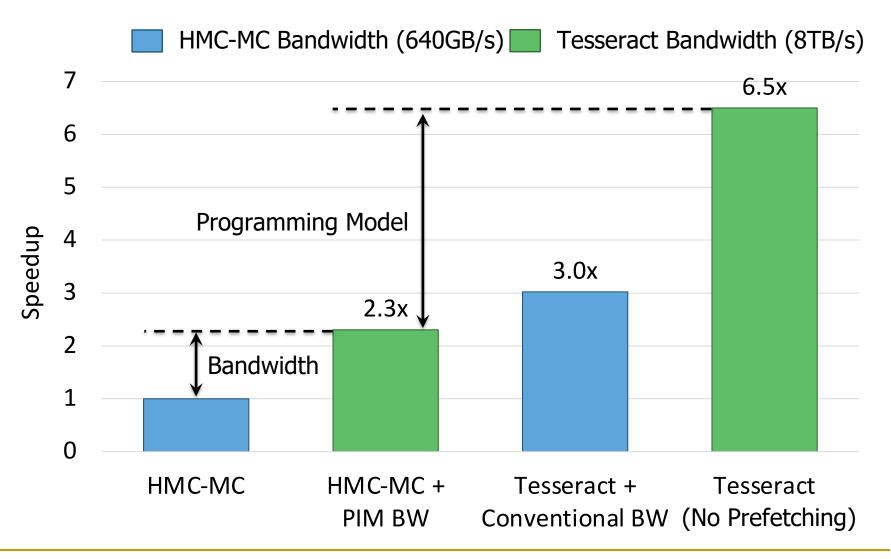


Tesseract Graph Processing Performance

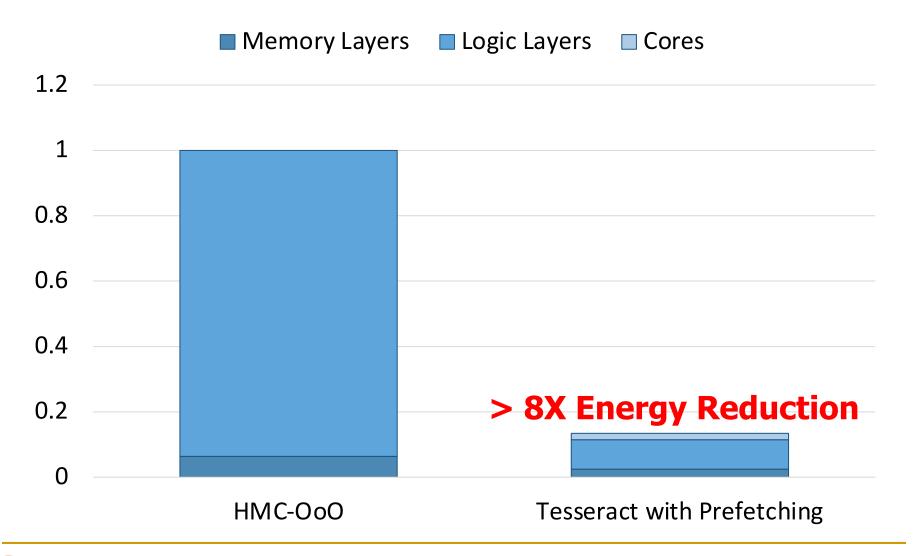


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Effect of Bandwidth & Programming Model



Tesseract Graph Processing System Energy



SAFARI Ahn+, "A Scalable Processing-in-Memory Accelerator for Parallel Graph Processing" ISCA 2015.

More on Tesseract

Junwhan Ahn, Sungpack Hong, Sungjoo Yoo, Onur Mutlu, and Kiyoung Choi,

"A Scalable Processing-in-Memory Accelerator for Parallel Graph Processing"

Proceedings of the <u>42nd International Symposium on</u> <u>Computer Architecture</u> (**ISCA**), Portland, OR, June 2015. [Slides (pdf)] [Lightning Session Slides (pdf)]

A Scalable Processing-in-Memory Accelerator for Parallel Graph Processing

Junwhan Ahn Sungpack Hong[§] Sungjoo Yoo Onur Mutlu[†] Kiyoung Choi junwhan@snu.ac.kr, sungpack.hong@oracle.com, sungjoo.yoo@gmail.com, onur@cmu.edu, kchoi@snu.ac.kr Seoul National University [§]Oracle Labs [†]Carnegie Mellon University

PIM Review and Open Problems

A Modern Primer on Processing in Memory

Onur Mutlu^{a,b}, Saugata Ghose^{b,c}, Juan Gómez-Luna^a, Rachata Ausavarungnirun^d

SAFARI Research Group

^aETH Zürich

^bCarnegie Mellon University

^cUniversity of Illinois at Urbana-Champaign

^dKing Mongkut's University of Technology North Bangkok

Onur Mutlu, Saugata Ghose, Juan Gomez-Luna, and Rachata Ausavarungnirun,

"A Modern Primer on Processing in Memory"

Invited Book Chapter in <u>Emerging Computing: From Devices to Systems -</u>

Looking Beyond Moore and Von Neumann, Springer, to be published in 2021.

PIM Review and Open Problems (II)

A Workload and Programming Ease Driven Perspective of Processing-in-Memory

Saugata Ghose[†] Amirali Boroumand[†] Jeremie S. Kim[†]§ Juan Gómez-Luna[§] Onur Mutlu^{§†}

†Carnegie Mellon University §ETH Zürich

Saugata Ghose, Amirali Boroumand, Jeremie S. Kim, Juan Gomez-Luna, and Onur Mutlu, "Processing-in-Memory: A Workload-Driven Perspective"

Invited Article in IBM Journal of Research & Development, Special Issue on Hardware for Artificial Intelligence, to appear in November 2019.

[Preliminary arXiv version]

More on Processing-in-Memory

Onur Mutlu, "Memory-Centric Computing Systems" Invited Tutorial at 66th International Electron Devices Meeting (IEDM), Virtual, 12 December 2020. [Slides (pptx) (pdf)] [Executive Summary Slides (pptx) (pdf)] [Tutorial Video (1 hour 51 minutes)] [Executive Summary Video (2 minutes)] [Abstract and Bio] [Related Keynote Paper from VLSI-DAT 2020]

[Related Review Paper on Processing in Memory]

https://www.youtube.com/watch?v=H3sEaINPBOE

Agenda

- The Problem: DNA Read Mapping
 - State-of-the-art Read Mapper Design
- Algorithmic Acceleration
 - Exploiting Structure of the Genome
 - Exploiting SIMD Instructions
- Hardware Acceleration
 - Specialized Architectures
 - Processing in Memory
- Future Opportunities: New Sequencing Technologies

New Genome Sequencing Technologies

Nanopore sequencing technology and tools for genome assembly: computational analysis of the current state, bottlenecks and future directions

Damla Senol Cali ™, Jeremie S Kim, Saugata Ghose, Can Alkan, Onur Mutlu

Briefings in Bioinformatics, bby017, https://doi.org/10.1093/bib/bby017

Published: 02 April 2018 Article history ▼



Oxford Nanopore MinION

Senol Cali+, "Nanopore Sequencing Technology and Tools for Genome Assembly: Computational Analysis of the Current State, Bottlenecks and Future Directions," Briefings in Bioinformatics, 2018.

[Preliminary arxiv.org version]

Recall: High-Throughput Sequencing

- Massively parallel sequencing technology
 - Illumina, Roche 454, Ion Torrent, SOLID...
- Small DNA fragments are first amplified and then sequenced in parallel, leading to
 - High throughput
 - High speed
 - Low cost
 - Short reads
 - Amplification step limits the read length since too short or too long fragments are not amplified well.
- Sequencing is done by either reading optical signals as each base is added, or by detecting hydrogen ions instead of light, leading to:
 - Low error rates (relatively)
 - Reads lack information about their order and which part of genome they are originated from

Nanopore Sequencing Technology

Nanopore sequencing is an emerging and a promising single-molecule DNA sequencing technology

- First nanopore sequencing device, MinION, made commercially available by Oxford Nanopore
 Technologies (ONT) in May 2014.
 - Inexpensive
 - Long read length (> 882K bp)
 - Portable: Pocket-sized
 - Produces data in real-time

Nanopore Sequencing Technology



an emerging and a promising named na

read length → Longer read length

- First nanopore sequencing device, MinION, made commercially available by Oxford Nanopore
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 - Inexpensive
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 - Portable: Pocket-sized
 - Produces data in real-time



Oxford Nanopore Sequencers NANOPORE











MinION Mk1B

MinION Mk1C

GridION Mk1

PromethION 24/48

| | MinION Mk1B | MinION Mk1C | GridION Mk1 | PromethION 24 | PromethION 48 | |
|---------------------------------|----------------|----------------|-------------|---------------|---------------|--|
| Read length | > 2Mb | > 2Mb | > 2Mb | > 2Mb | > 2Mb | |
| Yield per flow cell | 50 Gb | 50 Gb | 50 Gb | 220 Gb | 220 Gb | |
| Number of flow cells per device | 1 | 1 | 5 | 24 | 48 | |
| Yield per device | <50 Gb | <50 Gb | <250 Gb | <5.2 Tb | <10.5 Tb | |
| Starting price | \$1,000 | \$4,990 | \$49,995 | \$195,455 | \$327,455 | |

Illumina Sequencers

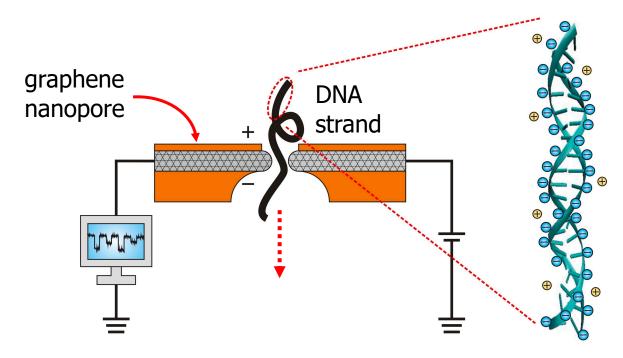




| Run time | 9.5–19 hrs | 4–24 hrs | 4–55 hrs | 12–30 hrs | 24-48 hrs | 13-44 hrs |
|-----------------------|------------|------------|------------|-------------|------------|------------|
| Max. reads per run | 4 million | 25 million | 25 million | 400 million | 1 billion | 20 billion |
| Max. read length | 2 × 150 bp | 2 × 150 bp | 2 × 300 bp | 2 × 150 bp | 2 × 150 bp | 2 x 250 |
| Max. output | 1.2 Gb | 7.5 Gb | 15 Gb | 120 Gb | 300 Gb | 6000 Gb |
| Estimated price | \$19,900 | \$49,500 | \$128,000 | \$275,000 | \$335,000 | \$985,000 |

SAFARI https://www.illumina.com/systems/sequencing-platforms.html

How Does Nanopore Sequencing Work?



- Nanopore is a nano-scale hole (<20nm).</p>
- In nanopore sequencers, an ionic current passes through the nanopores
- When the DNA strand passes through the nanopore, the sequencer measures the the change in current
- This change is used to identify the bases in the strand with the help of different electrochemical structures of the different bases

Advantages of Nanopore Sequencing

Nanopores:

- Do not require any labeling of the DNA or nucleotide for detection during sequencing
- Rely on the electronic or chemical structure of the different nucleotides for identification
- Allow sequencing very long reads, and
- Provide portability, low cost, and high throughput.

Challenges of Nanopore Sequencing

- One major drawback: high error rates
- Nanopore sequence analysis tools have a critical role to:
 - overcome high error rates
 - take better advantage of the technology
- Faster tools are critically needed to:
 - Take better advantage of the real-time data production capability of nanopore sequencing
 - Enable fast, real-time data analysis

Nanopore Genome Assembly Pipeline

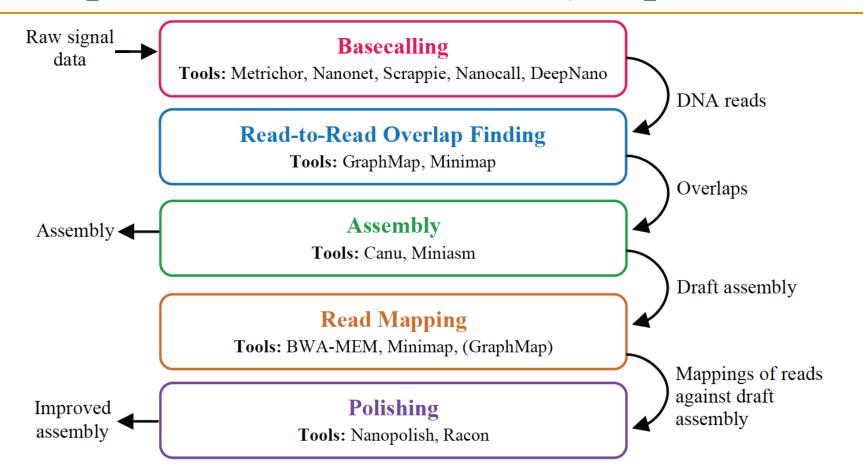


Figure 1. The analyzed genome assembly pipeline using nanopore sequence data, with its five steps and the associated tools for each

step.

Senol Cali+, "Nanopore Sequencing Technology and Tools for Genome Assembly" Briefings in Bioinformatics, 2018.

Nanopore Genome Assembly Tools (I)

Table 12. Accuracy analysis results for the full pipeline with a focus on the last two steps.

| | | | | | | | | | Number of | Number of | Identity | Coverage | Number of | Number of |
|----|-----------|---|----------|------------|----------|----------------|---|------------|-----------|-----------|----------|----------|------------|-----------|
| | | | | | | | | | Bases | Contigs | (%) | (%) | Mismatches | Indels |
| 1 | Metrichor | + | _ | + | Canu + | BWA-MEM | + | Nanopolish | 4,683,072 | 1 | 99.48 | 99.93 | 8,198 | 15,581 |
| 2 | Metrichor | + | Minimap | + | Miniasm+ | BWA-MEM | + | Nanopolish | 4,540,352 | 1 | 92.33 | 96.31 | 162,884 | 182,965 |
| 3 | Metrichor | + | GraphMaj | p + | Miniasm+ | BWA-MEM | + | Nanopolish | 4,637,916 | 2 | 92.38 | 95.80 | 159,206 | 180,603 |
| 4 | Metrichor | + | _ | + | Canu + | BWA-MEM | + | Racon | 4,650,502 | 1 | 98.46 | 100.00 | 18,036 | 51,842 |
| 5 | Metrichor | + | _ | + | Canu + | Minimap | + | Racon | 4,648,710 | 1 | 98.45 | 100.00 | 17,906 | 52,168 |
| 6 | Metrichor | + | Minimap | + | Miniasm+ | BWA-MEM | + | Racon | 4,598,267 | 1 | 97.70 | 99.91 | 24,014 | 82,906 |
| 7 | Metrichor | + | Minimap | + | Miniasm+ | Minimap | + | Racon | 4,600,109 | 1 | 97.78 | 100.00 | 23,339 | 79,721 |
| 8 | Nanonet | + | _ | + | Canu + | BWA-MEM | + | Racon | 4,622,285 | 1 | 98.48 | 100.00 | 16,872 | 52,509 |
| 9 | Nanonet | + | _ | + | Canu + | Minimap | + | Racon | 4,620,597 | 1 | 98.49 | 100.00 | 16,874 | 52,232 |
| 10 | Nanonet | + | Minimap | + | Miniasm+ | BWA-MEM | + | Racon | 4,593,402 | 1 | 98.01 | 99.97 | 20,322 | 72,284 |
| 11 | Nanonet | + | Minimap | + | Miniasm+ | Minimap | + | Racon | 4,592,907 | 1 | 98.04 | 100.00 | 20,170 | 70,705 |
| 12 | Scrappie | + | _ | + | Canu + | BWA-MEM | + | Racon | 4,673,871 | 1 | 98.40 | 99.98 | 13,583 | 60,612 |
| 13 | Scrappie | + | _ | + | Canu + | Minimap | + | Racon | 4,673,606 | 1 | 98.40 | 99.98 | 13,798 | 60,423 |
| 14 | Scrappie | + | Minimap | + | Miniasm+ | BWA-MEM | + | Racon | 5,157,041 | 8 | 97.87 | 99.80 | 18,085 | 78,492 |
| 15 | Scrappie | + | Minimap | + | Miniasm+ | Minimap | + | Racon | 5,156,375 | 8 | 97.87 | 99.94 | 17,922 | 77,807 |
| 16 | Nanocall | + | _ | + | Canu + | BWA-MEM | + | Racon | 1,383,851 | 86 | 93.49 | 28.82 | 19,057 | 65,244 |
| 17 | Nanocall | + | _ | + | Canu + | Minimap | + | Racon | 1,367,834 | 86 | 94.43 | 28.74 | 15,610 | 55,275 |
| 18 | Nanocall | + | Minimap | + | Miniasm+ | BWA-MEM | + | Racon | 4,707,961 | 5 | 90.75 | 97.11 | 91,502 | 347,005 |
| 19 | Nanocall | + | Minimap | + | Miniasm+ | Minimap | + | Racon | 4,673,069 | 5 | 92.23 | 97.10 | 72,646 | 291,918 |
| 20 | DeepNano | + | _ | + | Canu + | BWA-MEM | + | Racon | 7,429,290 | 106 | 96.46 | 99.24 | 27,811 | 102,682 |
| 21 | DeepNano | + | _ | + | Canu + | Minimap | + | Racon | 7,404,454 | 106 | 96.03 | 99.21 | 34,023 | 110,640 |
| 22 | DeepNano | + | Minimap | + | Miniasm+ | BWA-MEM | + | Racon | 4,566,253 | 1 | 96.76 | 99.86 | 25,791 | 125,386 |
| 23 | DeepNano | + | Minimap | + | Miniasm+ | Minimap | + | Racon | 4,571,810 | 1 | 96.90 | 99.97 | 24,994 | 119,519 |

Senol Cali+, "Nanopore Sequencing Technology and Tools for Genome Assembly" Briefings in Bioinformatics, 2018.

Nanopore Genome Assembly Tools (II)

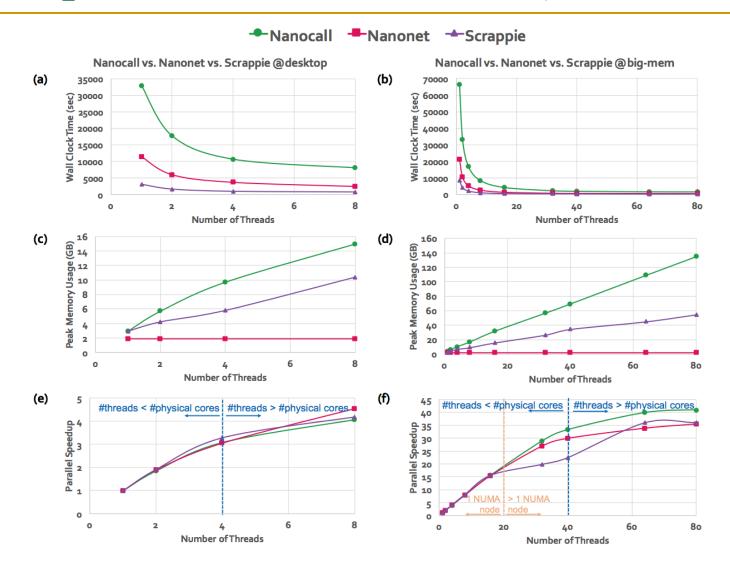
Table 13. Performance analysis results for the full pipeline with a focus on the last two steps.

| | | | | | | | | | | Step 4: Read Mapper | | | Step 5: Polisher | | |
|----|-----------|---|----------|---|-----------|---|----------------|---|------------|----------------------------------|------------------|-------------------------|----------------------------------|------------------|-------------------------|
| | | | | | | | | | | Wall Clock Time (h:m:s) | CPU Time (h:m:s) | Memory Usage (GB) | Wall Clock Time (h:m:s) | CPU Time (h:m:s) | Memory Usage (GB) |
| 1 | Metrichor | + | _ | + | Canu + | + | BWA-MEM | + | Nanopolish | 24:43 | 15:47:21 | 5.26 | 5:51:00 | 191:18:52 | 13.38 |
| 2 | Metrichor | + | Minimap | + | Miniasm + | H | BWA-MEM | + | Nanopolish | 12:33 | 7:50:54 | 3.75 | 122:52:00 | 4458:36:10 | 31.36 |
| 3 | Metrichor | + | GraphMap | + | Miniasm + | + | BWA-MEM | + | Nanopolish | 12:47 | 7:57:58 | 3.60 | 129:46:00 | 4799:03:51 | 31.31 |
| 4 | Metrichor | + | _ | + | Canu + | + | BWA-MEM | + | Racon | 24:20 | 15:43:40 | 6.60 | 14:44 | 9:09:22 | 8.11 |
| 5 | Metrichor | + | _ | + | Canu + | + | Minimap | + | Racon | 3 | 1:35 | 0.26 | 15:12 | 9:45:33 | 14.55 |
| 6 | Metrichor | + | Minimap | + | Miniasm + | H | BWA-MEM | + | Racon | 12:10 | 7:48:10 | 5.19 | 15:43 | 9:33:39 | 9.98 |
| 7 | Metrichor | + | Minimap | + | Miniasm + | H | Minimap | + | Racon | 3 | 1:24 | 0.26 | 20:28 | 8:57:40 | 18.24 |
| 8 | Nanonet | + | _ | + | Canu + | + | BWA-MEM | + | Racon | 9:08 | 5:53:18 | 4.84 | 6:33 | 4:02:10 | 4.47 |
| 9 | Nanonet | + | _ | + | Canu + | + | Minimap | + | Racon | 2 | 54 | 0.26 | 6:45 | 4:17:26 | 7.93 |
| 10 | Nanonet | + | Minimap | + | Miniasm + | H | BWA-MEM | + | Racon | 4:40 | 2:58:02 | 3.88 | 7:08 | 4:19:30 | 5.35 |
| 11 | Nanonet | + | Minimap | + | Miniasm + | H | Minimap | + | Racon | 2 | 46 | 0.26 | 7:01 | 4:18:48 | 9.53 |
| 12 | Scrappie | + | _ | + | Canu + | + | BWA-MEM | + | Racon | 33:41 | 21:11:06 | 8.66 | 13:32 | 8:24:44 | 7.58 |
| 13 | Scrappie | + | _ | + | Canu + | H | Minimap | + | Racon | 3 | 1:39 | 0.27 | 18:45 | 7:43:17 | 13.20 |
| 14 | Scrappie | + | Minimap | + | Miniasm + | + | BWA-MEM | + | Racon | 22:41 | 14:31:00 | 6.08 | 14:37 | 8:53:59 | 9.50 |
| 15 | Scrappie | + | Minimap | + | Miniasm + | + | Minimap | + | Racon | 3 | 1:27 | 0.27 | 15:10 | 9:02:45 | 12.72 |
| 16 | Nanocall | + | _ | + | Canu + | + | BWA-MEM | + | Racon | 4:52 | 3:01:15 | 3.80 | 11:07 | 3:26:52 | 5.63 |
| 17 | Nanocall | + | _ | + | Canu + | H | Minimap | + | Racon | 3 | 1:16 | 0.22 | 7:28 | 2:50:35 | 3.62 |
| 18 | Nanocall | + | Minimap | + | Miniasm + | H | BWA-MEM | + | Racon | 16:06 | 10:27:20 | 5.06 | 18:56 | 11:32:45 | 11.47 |
| 19 | Nanocall | + | Minimap | + | Miniasm + | H | Minimap | + | Racon | 4 | 1:18 | 0.26 | 11:49 | 7:08:59 | 10.98 |
| 20 | DeepNano | + | _ | + | Canu + | + | BWA-MEM | + | Racon | 17:36 | 11:30:20 | 4.43 | 12:48 | 7:13:04 | 8.88 |
| 21 | DeepNano | + | _ | + | Canu + | H | Minimap | + | Racon | 3 | 1:24 | 0.28 | 11:39 | 6:55:01 | 3.73 |
| 22 | DeepNano | + | Minimap | + | Miniasm + | H | BWA-MEM | + | Racon | 8:15 | 5:22:29 | 4.11 | 14:16 | 8:34:32 | 10.30 |
| 23 | DeepNano | + | Minimap | + | Miniasm + | H | Minimap | + | Racon | 3 | 1:10 | 0.26 | 12:29 | 7:55:32 | 17.11 |

Senol Cali+, "Nanopore Sequencing Technology and Tools for Genome Assembly" Briefings in Bioinformatics, 2018.



Nanopore Genome Assembly Tools (III)



Senol Cali+, "Nanopore Sequencing Technology and Tools for Genome Assembly" to appear in Briefings in Bioinformatics, 2018.

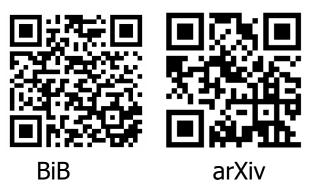
More on Nanopore Sequencing & Tools

Nanopore sequencing technology and tools for genome assembly: computational analysis of the current state, bottlenecks and future directions

Damla Senol Cali ™, Jeremie S Kim, Saugata Ghose, Can Alkan, Onur Mutlu

Briefings in Bioinformatics, bby017, https://doi.org/10.1093/bib/bby017

Published: 02 April 2018 Article history ▼



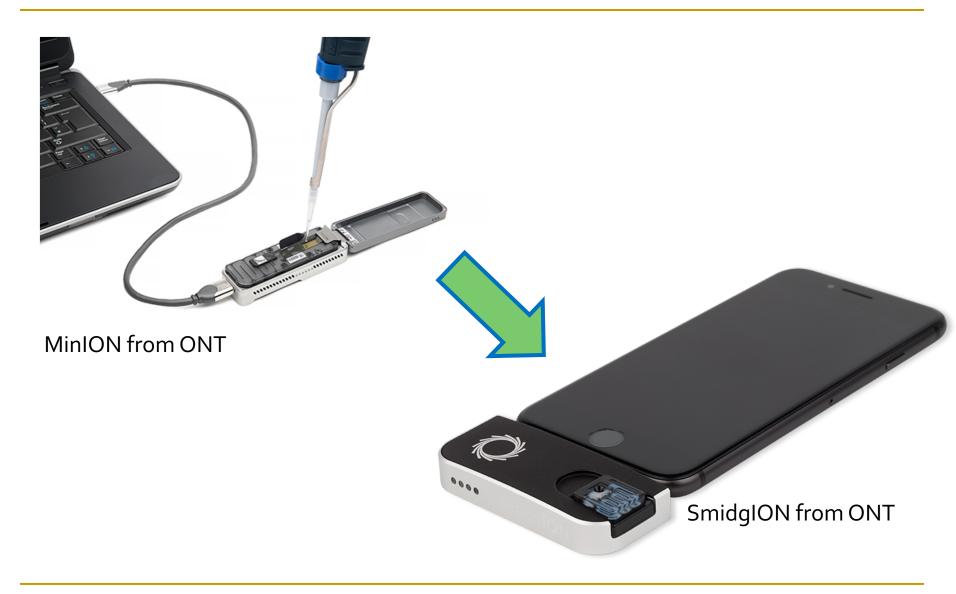
Senol Cali+, "Nanopore Sequencing Technology and Tools for Genome Assembly: Computational Analysis of the Current State, Bottlenecks and Future Directions," Briefings in Bioinformatics, 2018.

[Preliminary arxiv.org version]

Recall Our Dream (from 2007)

- An embedded device that can perform comprehensive genome analysis in real time (within a minute)
- Still a long ways to go
 - Energy efficiency
 - Performance (latency)
 - Security
 - Huge memory bottleneck

Future of Genome Sequencing & Analysis



Why Do We Care? An Example from 2020

200 Oxford Nanopore sequencers have left UK for China, to support rapid, near-sample coronavirus sequencing for outbreak surveillance

Fri 31st January 2020

Following extensive support of, and collaboration with, public health professionals in China, Oxford Nanopore has shipped an additional 200 MinION sequencers and related consumables to China. These will be used to support the ongoing surveillance of the current coronavirus outbreak, adding to a large number of the devices already installed in the country.



Each MinION sequencer is approximately the size of a stapler, and can provide rapid sequence information about the coronavirus.



700Kg of Oxford Nanopore sequencers and consumables are on their way for use by Chinese scientists in understanding the current coronavirus outbreak.



Sequencing of COVID-19

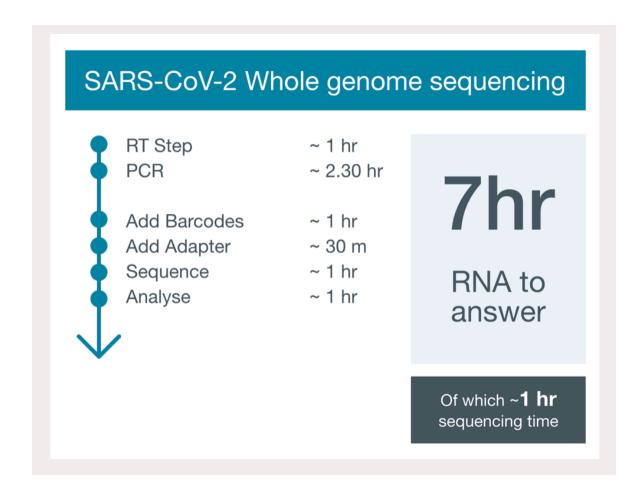
Whole genome sequencing (WGS) and sequence data analysis are important

- To detect the virus from a human sample such as saliva,
 Bronchoalveolar fluid etc.
- To understand the sources and modes of transmission of the virus
- To discover the genomic characteristics of the virus, and compare with better-known viruses (e.g., 02-03 SARS epidemic)
- To design and evaluate the diagnostic tests and deep-dive studies

Two key areas of COVID-19 genomic research

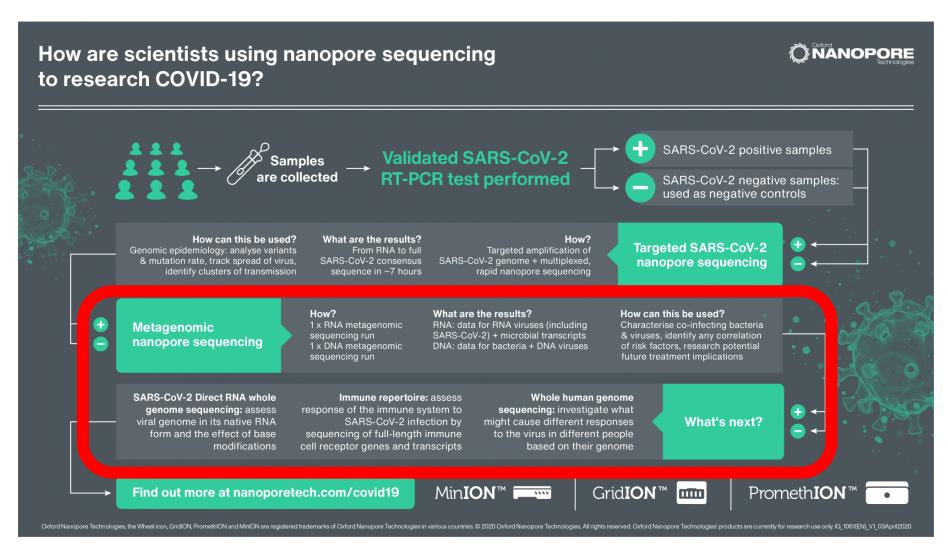
- To sequence the genome of the virus itself, COVID-19, in order to track the mutations in the virus.
- To explore the genes of infected patients. This analysis can be used to understand why some people get more severe symptoms than others, as well as, help with the development of new treatments in the future.

COVID-19 Nanopore Sequencing (I)



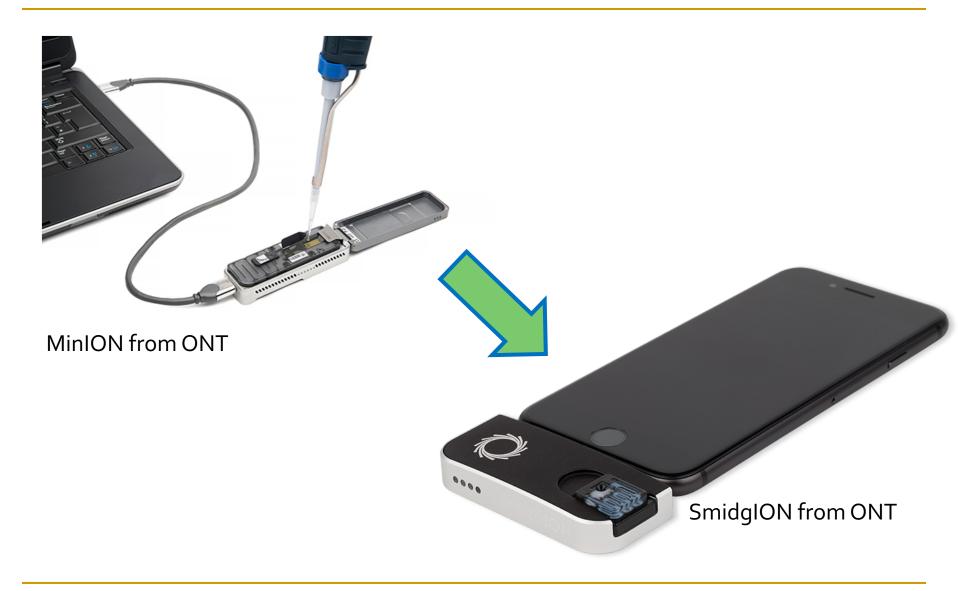
From ONT (https://nanoporetech.com/covid-19/overview)

COVID-19 Nanopore Sequencing (II)



From ONT (https://nanoporetech.com/covid-19/overview)

Future of Genome Sequencing & Analysis



Agenda

- The Problem: DNA Read Mapping
 - State-of-the-art Read Mapper Design
- Algorithmic Acceleration
 - Exploiting Structure of the Genome
 - Exploiting SIMD Instructions
- Hardware Acceleration
 - Specialized Architectures
 - Processing in Memory
- Future Opportunities: New Sequencing Technologies

Conclusion

- System design for bioinformatics is a critical problem
 - It has large scientific, medical, societal, personal implications
- This talk is about accelerating a key step in bioinformatics: genome sequence analysis
 - In particular, read mapping
- We covered various recent ideas to accelerate read mapping
 - My personal journey since September 2006
- Many future opportunities exist
 - Especially with new sequencing technologies
 - Especially with new applications and use cases

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Accelerating Genome Analysis: Overview

 Mohammed Alser, Zulal Bingol, Damla Senol Cali, Jeremie Kim, Saugata Ghose, Can Alkan, and Onur Mutlu,

"Accelerating Genome Analysis: A Primer on an Ongoing Journey"

<u>IEEE Micro</u> (IEEE MICRO), Vol. 40, No. 5, pages 65-75, September/October 2020.

[Slides (pptx)(pdf)]

[Talk Video (1 hour 2 minutes)]

Accelerating Genome Analysis: A Primer on an Ongoing Journey

Mohammed Alser

ETH Zürich

Zülal Bingöl

Bilkent University

Damla Senol Cali

Carnegie Mellon University

Jeremie Kim

ETH Zurich and Carnegie Mellon University

Saugata Ghose

University of Illinois at Urbana–Champaign and Carnegie Mellon University

Can Alkan

Bilkent University

Onur Mutlu

ETH Zurich, Carnegie Mellon University, and Bilkent University

PIM Review and Open Problems

A Modern Primer on Processing in Memory

Onur Mutlu^{a,b}, Saugata Ghose^{b,c}, Juan Gómez-Luna^a, Rachata Ausavarungnirun^d

SAFARI Research Group

^aETH Zürich

^bCarnegie Mellon University

^cUniversity of Illinois at Urbana-Champaign

^dKing Mongkut's University of Technology North Bangkok

Onur Mutlu, Saugata Ghose, Juan Gomez-Luna, and Rachata Ausavarungnirun,

"A Modern Primer on Processing in Memory"

Invited Book Chapter in <u>Emerging Computing: From Devices to Systems -</u>

Looking Beyond Moore and Von Neumann, Springer, to be published in 2021.

PIM Review and Open Problems (II)

A Workload and Programming Ease Driven Perspective of Processing-in-Memory

Saugata Ghose[†] Amirali Boroumand[†] Jeremie S. Kim[†]§ Juan Gómez-Luna[§] Onur Mutlu^{§†}

†Carnegie Mellon University §ETH Zürich

Saugata Ghose, Amirali Boroumand, Jeremie S. Kim, Juan Gomez-Luna, and Onur Mutlu, "Processing-in-Memory: A Workload-Driven Perspective"

Invited Article in IBM Journal of Research & Development, Special Issue on Hardware for Artificial Intelligence, to appear in November 2019.

[Preliminary arXiv version]

More on Memory-Centric System Design

Onur Mutlu,

"Memory-Centric Computing Systems"

Invited Tutorial at <u>66th International Electron Devices</u>

Meeting (IEDM), Virtual, 12 December 2020.

[Slides (pptx) (pdf)]

[Executive Summary Slides (pptx) (pdf)]

[Tutorial Video (1 hour 51 minutes)]

[Executive Summary Video (2 minutes)]

[Abstract and Bio]

[Related Keynote Paper from VLSI-DAT 2020]

[Related Review Paper on Processing in Memory]

https://www.youtube.com/watch?v=H3sEaINPBOE

Detailed Lectures on Genome Analysis

- Computer Architecture, Fall 2020, Lecture 3a
 - Introduction to Genome Sequence Analysis (ETH Zürich, Fall 2020)
 - https://www.youtube.com/watch?v=CrRb32v7SJc&list=PL5Q2soXY2Zi9xidyIgBxUz7xRPS-wisBN&index=5
- Computer Architecture, Fall 2020, Lecture 8
 - □ **Intelligent Genome Analysis** (ETH Zürich, Fall 2020)
 - https://www.youtube.com/watch?v=ygmQpdDTL7o&list=PL5Q2soXY2Zi9xidyIgBxUz7xRPS-wisBN&index=14
- Computer Architecture, Fall 2020, Lecture 9a
 - □ **GenASM: Approx. String Matching Accelerator** (ETH Zürich, Fall 2020)
 - https://www.youtube.com/watch?v=XoLpzmN Pas&list=PL5Q2soXY2Zi9xidyIgBxUz7xRPS-wisBN&index=15
- Accelerating Genomics Project Course, Fall 2020, Lecture 1
 - Accelerating Genomics (ETH Zürich, Fall 2020)
 - https://www.youtube.com/watch?v=rgjl8ZyLsAg&list=PL5Q2soXY2Zi9E2bBVAgCqL gwiDRQDTyId

Acknowledgments

- Can Alkan, Bilkent University
- Many students at ETH, CMU, Bilkent
 - Mohammed Alser, Damla Senol Cali, Jeremie Kim, Hasan Hassan, Donghyuk Lee, Hongyi Xin, ...

Funders:

- NIH and Industrial Partners (Alibaba, AMD, Google, Facebook, HP Labs, Huawei, IBM, Intel, Microsoft, Nvidia, Oracle, Qualcomm, Rambus, Samsung, Seagate, VMware)
- All papers, source code, and more are at:
 - https://people.inf.ethz.ch/omutlu/projects.htm

Funding Acknowledgments

- Alibaba, AMD, ASML, Google, Facebook, Hi-Silicon, HP Labs, Huawei, IBM, Intel, Microsoft, Nvidia, Oracle, Qualcomm, Rambus, Samsung, Seagate, VMware
- NSF
- NIH
- GSRC
- SRC
- CyLab

Acknowledgments



Think BIG, Aim HIGH!

https://safari.ethz.ch

Onur Mutlu's SAFARI Research Group

Computer architecture, HW/SW, systems, bioinformatics, security, memory

https://safari.ethz.ch/safari-newsletter-january-2021/



Think BIG, Aim HIGH!

SAFARI

https://safari.ethz.ch

SAFARI Newsletter April 2020 Edition

https://safari.ethz.ch/safari-newsletter-april-2020/





View in your browser

Think Big, Aim High



Dear SAFARI friends,

SAFARI Newsletter January 2021 Edition

https://safari.ethz.ch/safari-newsletter-january-2021/





Newsletter January 2021

Think Big, Aim High, and Have a Wonderful 2021!



Dear SAFARI friends,

Accelerating Genome Analysis

A Primer on an Ongoing Journey

Onur Mutlu

omutlu@gmail.com

https://people.inf.ethz.ch/omutlu

26 January 2021

Technion Invited Lecture





Carnegie Mellon

Backup Slides for Further Info

Referenced Papers and Talks

All are available at

https://people.inf.ethz.ch/omutlu/projects.htm

http://scholar.google.com/citations?user=7XyGUGkAAAAJ&hl=en

https://www.youtube.com/onurmutlulectures

Research & Teaching: Some Overview Talks

https://www.youtube.com/onurmutlulectures

- Future Computing Architectures
 - https://www.youtube.com/watch?v=kqiZISOcGFM&list=PL5Q2soXY2Zi8D 5MGV6EnXEJHnV2YFBJI&index=1
- Enabling In-Memory Computation
 - https://www.youtube.com/watch?v=njX 14584Jw&list=PL5Q2soXY2Zi8D 5MGV6EnXEJHnV2YFBJl&index=16
- Accelerating Genome Analysis
 - https://www.youtube.com/watch?v=hPnSmfwu2-A&list=PL5Q2soXY2Zi8D_5MGV6EnXEJHnV2YFBJl&index=9
- Rethinking Memory System Design
 - https://www.youtube.com/watch?v=F7xZLNMIY1E&list=PL5Q2soXY2Zi8D_5MGV6EnXEJHnV2YFBJl&index=3
- Intelligent Architectures for Intelligent Machines
 - https://www.youtube.com/watch?v=n8Aj_A0WSg8&list=PL5Q2soXY2Zi8D_5MGV6EnXEJHnV2YFBJl&index=22
- Revisiting RowHammer
 - https://www.youtube.com/watch?v=B58YT9hZM4g&list=PL5Q2soXY2Zi8D_5MGV6EnXEJHnV2YFBJl&index=25

An Interview on Research and Education

- Computing Research and Education (@ ISCA 2019)
 - https://www.youtube.com/watch?v=8ffSEKZhmvo&list=PL5Q2 soXY2Zi_4oP9LdL3cc8G6NIjD2Ydz

- Maurice Wilkes Award Speech (10 minutes)
 - https://www.youtube.com/watch?v=tcQ3zZ3JpuA&list=PL5Q2 soXY2Zi8D_5MGV6EnXEJHnV2YFBJl&index=15

More Thoughts and Suggestions

Onur Mutlu,

"Some Reflections (on DRAM)"

Award Speech for <u>ACM SIGARCH Maurice Wilkes Award</u>, at the **ISCA** Awards Ceremony, Phoenix, AZ, USA, 25 June 2019.

[Slides (pptx) (pdf)]

[Video of Award Acceptance Speech (Youtube; 10 minutes) (Youku; 13 minutes)]

[Video of Interview after Award Acceptance (Youtube; 1 hour 6 minutes)]

1 hour 6 minutes)

[News Article on "ACM SIGARCH Maurice Wilkes Award goes to Prof. Onur Mutlu"]

Onur Mutlu,

"How to Build an Impactful Research Group"

57th Design Automation Conference Early Career Workshop (DAC), Virtual, 19 July 2020.

[Slides (pptx) (pdf)]

Detailed Lectures on PIM (I)

- Computer Architecture, Fall 2020, Lecture 6
 - Computation in Memory (ETH Zürich, Fall 2020)
 - https://www.youtube.com/watch?v=oGcZAGwfEUE&list=PL5Q2soXY2Zi9xidyIgBxUz 7xRPS-wisBN&index=12
- Computer Architecture, Fall 2020, Lecture 7
 - Near-Data Processing (ETH Zürich, Fall 2020)
 - https://www.youtube.com/watch?v=j2GIigqn1Qw&list=PL5Q2soXY2Zi9xidyIgBxUz7xRPS-wisBN&index=13
- Computer Architecture, Fall 2020, Lecture 11a
 - Memory Controllers (ETH Zürich, Fall 2020)
 - https://www.youtube.com/watch?v=TeG773OgiMQ&list=PL5Q2soXY2Zi9xidyIgBxUz 7xRPS-wisBN&index=20
- Computer Architecture, Fall 2020, Lecture 12d
 - Real Processing-in-DRAM with UPMEM (ETH Zürich, Fall 2020)
 - https://www.youtube.com/watch?v=Sscy1Wrr22A&list=PL5Q2soXY2Zi9xidyIgBxUz7xRPS-wisBN&index=25

Detailed Lectures on PIM (II)

- Computer Architecture, Fall 2020, Lecture 15
 - Emerging Memory Technologies (ETH Zürich, Fall 2020)
 - https://www.youtube.com/watch?v=AlE1rD9G_YU&list=PL5Q2soXY2Zi9xidyIgBxUz 7xRPS-wisBN&index=28
- Computer Architecture, Fall 2020, Lecture 16a
 - Opportunities & Challenges of Emerging Memory Technologies
 (ETH Zürich, Fall 2020)
 - https://www.youtube.com/watch?v=pmLszWGmMGQ&list=PL5Q2soXY2Zi9xidyIgBx Uz7xRPS-wisBN&index=29
- Computer Architecture, Fall 2020, Guest Lecture
 - In-Memory Computing: Memory Devices & Applications (ETH Zürich, Fall 2020)
 - https://www.youtube.com/watch?v=wNmqQHiEZNk&list=PL5Q2soXY2Zi9xidyIgBxUz7xRPS-wisBN&index=41

Genome Analysis



No machine can read the *entire* content of a genome



Genome Analysis



No machine can read the *entire* content of a genome



>CCT GACC CATGT GAAG ACTA AAGT GAAA

Why?!

CAAG

TCTT

CATTG

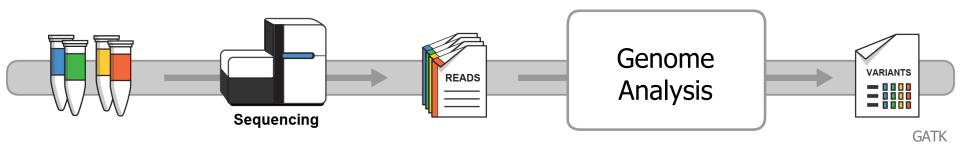
AAAA

ATTT

AAAA

ATGG GAAA

Genome Sequencer is a Chopper



CCCCCTATATATACGTACTAGTACGT

ACGACTTTAGTACGTACGT TATATACGTACTAGTACGT

ACGTACG CCCCTACGTA
TATATATACGTACTAGTACGT

ACGACTTTAGTACGTACGT TATATATACGTACTAGAGTACGT TATATATACGTACTAGTACGT

ACG TTTTTAAAACGTA
TATATATACGTACTACGT

ACGAC GGGGAGTACGTACGT



1x10¹² bases*



44 hours*



<1000 \$

* NovaSeq 6000

High-Throughput Sequencers



Illumina MiSeq



Illumina NovaSeq 6000



Pacific Biosciences Sequel II



Pacific Biosciences RS II





Oxford Nanopore MinION



... and more! All produce data with different properties.

Oxford Nanopore Sequencers NANOPORE











MinION Mk1B

MinION Mk1C

GridION Mk1

PromethION 24/48

| | MinION Mk1B | MinION Mk1C | GridION Mk1 | PromethION 24 | PromethION 48 |
|---------------------------------|----------------|----------------|-------------|---------------|---------------|
| Read length | > 2Mb | > 2Mb | > 2Mb | > 2Mb | > 2Mb |
| Yield per flow cell | 50 Gb | 50 Gb | 50 Gb | 220 Gb | 220 Gb |
| Number of flow cells per device | 1 | 1 | 5 | 24 | 48 |
| Yield per device | <50 Gb | <50 Gb | <250 Gb | <5.2 Tb | <10.5 Tb |
| Starting price | \$1,000 | \$4,990 | \$49,995 | \$195,455 | \$327,455 |

Illumina Sequencers

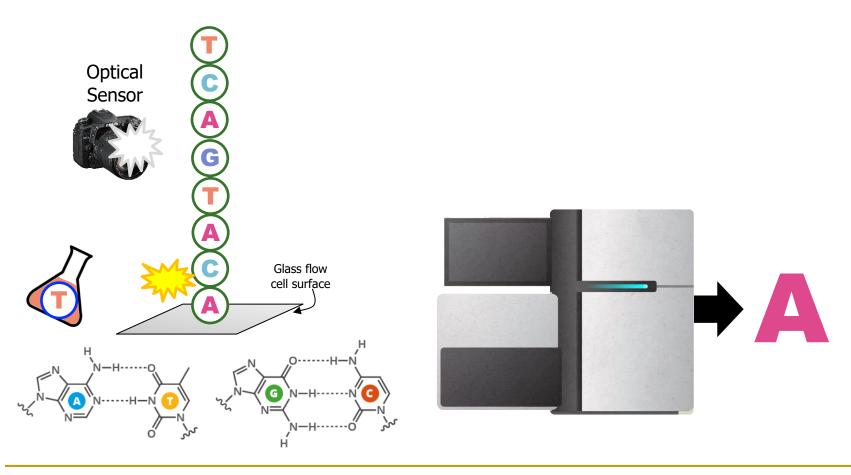




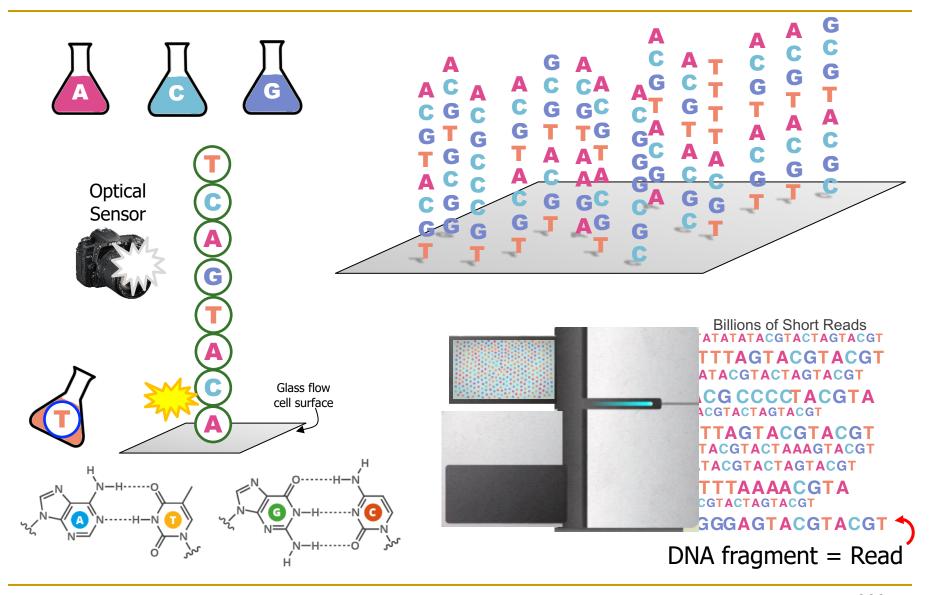
| Run time | 9.5–19 hrs | 4–24 hrs | 4–55 hrs | 12–30 hrs | 24-48 hrs | 13-44 hrs |
|-----------------------|------------|------------|------------|-------------|------------|------------|
| Max. reads per run | 4 million | 25 million | 25 million | 400 million | 1 billion | 20 billion |
| Max. read length | 2 × 150 bp | 2 × 150 bp | 2 × 300 bp | 2 × 150 bp | 2 × 150 bp | 2 x 250 |
| Max. output | 1.2 Gb | 7.5 Gb | 15 Gb | 120 Gb | 300 Gb | 6000 Gb |
| Estimated price | \$19,900 | \$49,500 | \$128,000 | \$275,000 | \$335,000 | \$985,000 |

SAFARI https://www.illumina.com/systems/sequencing-platforms.html

How Does Illumina Machine Work?



How Does Illumina Machine Work?



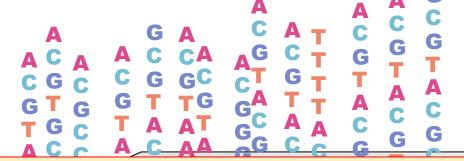
How Does Illumina Machine Work?





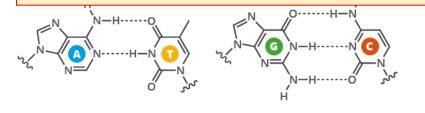






Check Illumina virtual tour:

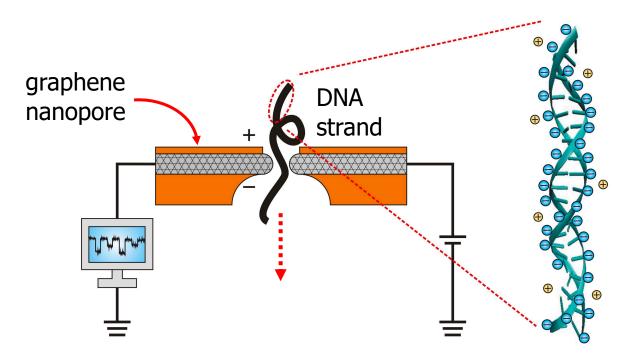
https://emea.illumina.com/systems/sequencing-platforms/iseq/tour.html





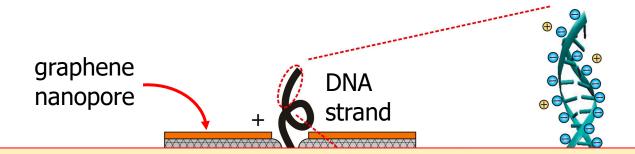
DNA fragment = Read

How Does Nanopore Machine Work?



- Nanopore is a nano-scale hole (<20nm).</p>
- In nanopore sequencers, an ionic current passes through the nanopores
- When the DNA strand passes through the nanopore, the sequencer measures the the change in current
- This change is used to identify the bases in the strand with the help of different electrochemical structures of the different bases

How Does Nanopore Machine Work?



Check Nanopore virtual tour:

https://nanoporetech.com/resource-centre/minion-video

measures the the change in current

 This change is used to identify the bases in the strand with the help of different electrochemical structures of the different bases



Common Disadvantages!

Regardless the sequencing machine,

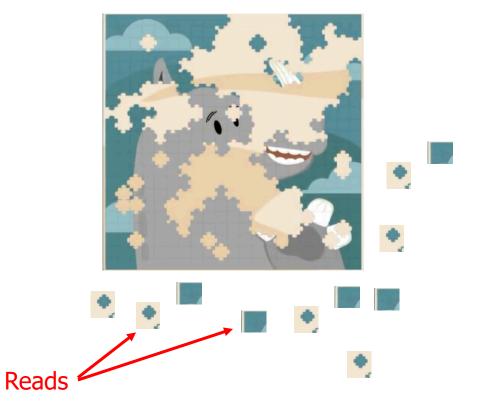
reads still lack information about their order and location

(which part of genome they are originated from)



Solving the Puzzle





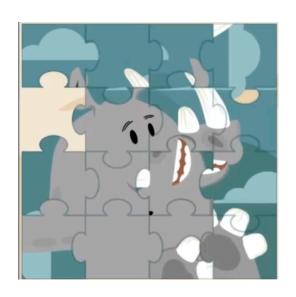
https://www.pacb.com/smrt-science/smrt-sequencing/hifi-reads-for-highly-accurate-long-read-sequencing/

HTS Sequencing Output

Small pieces of a puzzle short reads (Illumina)



Large pieces of a puzzle long reads (ONT & PacBio)



Which sequencing technology is the best?

□ 100-300 bp

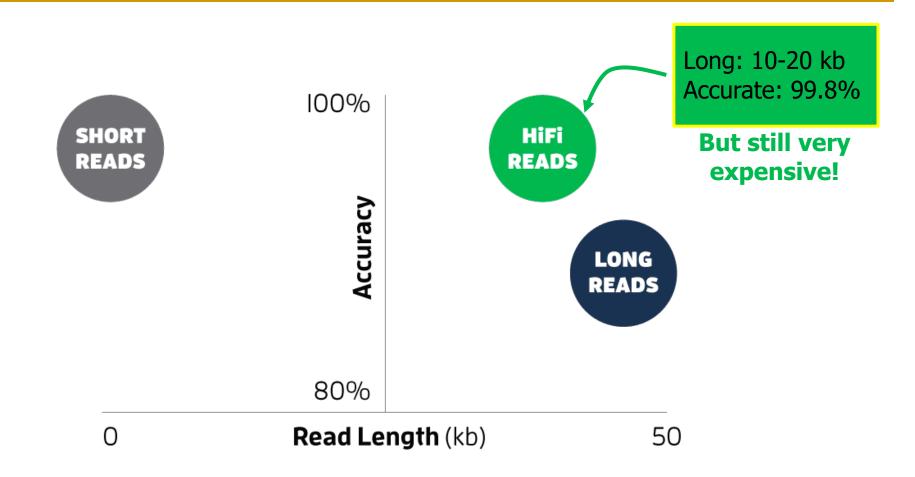
□ 500-2M bp

□ low error rate (~0.1%)

☐ high error rate (~15%)

https://www.pacb.com/smrt-science/smrt-sequencing/hifi-reads-for-highly-accurate-long-read-sequencing/

HiFi Reads (PacBio)



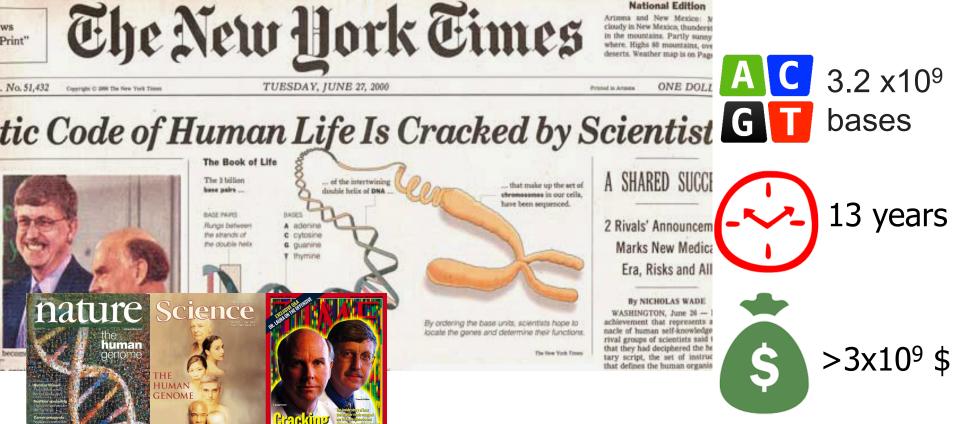
Wenger+, "<u>Accurate circular consensus long-read sequencing improves variant</u> detection and assembly of a human genome", *Nature Biotechnology*, 2019

How Long is DNA?



Cracking the 1st Human Genome Sequence

■ **1990-2003:** The Human Genome Project (HGP) provides a complete and accurate sequence of all **DNA base pairs** that make up the human genome and finds 20,000 to 25,000 human genes.



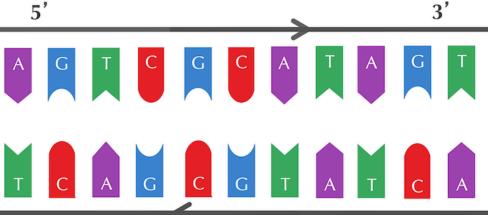
Obtaining the Human Reference Genome

GRCh38.p13

- Description: Genome Reference Consortium Human Build 38 patch release 13 (GRCh38.p13)
- Organism name: <u>Homo sapiens (human)</u>
- Date: 2019/02/28
- **3,099,706,404** bases
- Compressed .fna file (964.9 MB)
- https://www.ncbi.nlm.nih.gov/assembly/GCF 000001405.39

Challenges in Read Mapping

- Need to find many mappings of each read
- Need to tolerate variances/sequencing errors in each read
- Need to map each read very fast (i.e., performance is important, life critical in some cases)
- Need to map reads to both forward and reverse strands



Revisiting the Puzzle



http://www.pacb.com/smrt-science/smrt-sequencing/hifi-reads-for-highly-accurate-long-read-sequencing/



Reference Genome Bias

nature genetics

Letter | Open Access | Published: 19 November 2018

Assembly of a pan-genome from deep sequencing of 910 humans of African descent

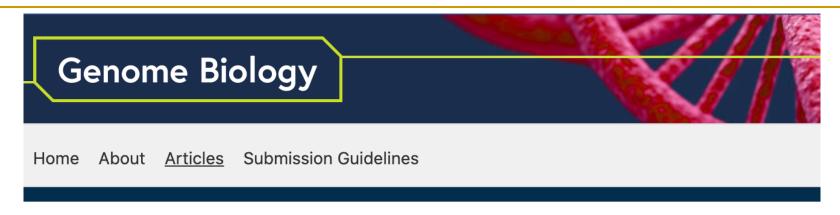
Rachel M. Sherman ⊠, Juliet Forman, [...] Steven L. Salzberg ⊠

Nature Genetics **51**, 30–35(2019) | Cite this article

"African pan-genome contains ~10% more DNA bases than the current human reference genome"



Time to Change the Reference Genome



Opinion | Open Access | Published: 09 August 2019

Is it time to change the reference genome?

Sara Ballouz, Alexander Dobin & Jesse A. Gillis ≥

Genome Biology 20, Article number: 159 (2019) Cite this article

12k Accesses | 11 Citations | 45 Altmetric | Metrics

"Switching to a consensus reference would offer important advantages over the continued use of the current reference with few disadvantages"

SAFARI Ballouz+, "Is it time to change the reference genome?", Genome Biology, 2019^{43}

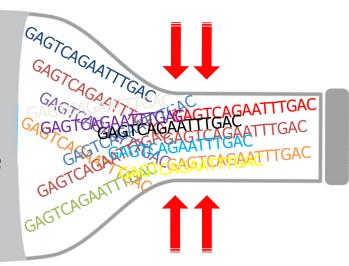
Bottlenecked in Read Mapping!!

48 Human whole genomes

at 30× coverage

in about 2 days

Illumina NovaSeq 6000



1 Human genome

32 CPU hours

on a 48-core processor



■ Read Mapping ■ Others

MAGNET (AACBB 2018, TIR 2017)

- Key observation: the use of AND operation to check if a zero (match) exists in a column introduces filtering inaccuracy.
- Key Idea: count the consecutive zeros in each mask and select the longest in a divide-and-conquer approach.
- MAGNET is 17x to 105x more accurate than GateKeeper and SHD.

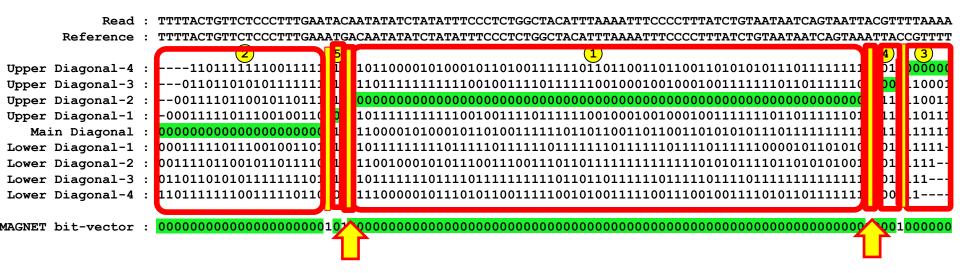
AAAAAAAAAAAAAGAGAGAGAGATAGTTAGTGTTGCAGCCACTACAACACAAAAGAGGACCAACTTACGTGTCTAAAAGGGGAGACATTGTTGGGCC

MAGNET Walkthrough

Build Neighborhood Map

Track the Diagonally Consecutive Matches

ACCEPT iff number of '1' ≤ Threshold



Find the longest segment of consecutive zeros

Exclude the errors from the search space

Divide the problem into two subproblems and repeat



What if we got a new version of the reference genome?

AirLift

- Key observation: Reference genomes are updated frequently.
 Repeating read mapping is a computationally expensive workload.
- Key idea: Update the mapping results of only affected reads depending on how a region in the old reference relates to another region in the new reference.

Key results:

- reduces number of reads that needs to be re-mapped to new reference by up to 99%
- reduces overall runtime to re-map reads by 6.94x, 208x, and 16.4x for large (human), medium (C. elegans), and small (yeast) reference genomes

248

Clustering the Reference Genome Regions

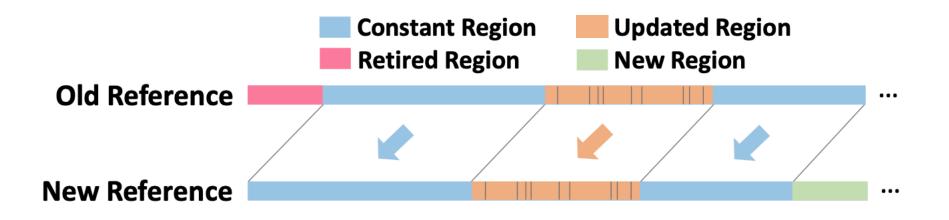


Fig. 2. Reference Genome Regions.

More Details on AirLift

arXiv.org > q-bio > arXiv:1912.08735

Search... Help | Advanc

Quantitative Biology > Genomics

[Submitted on 18 Dec 2019]

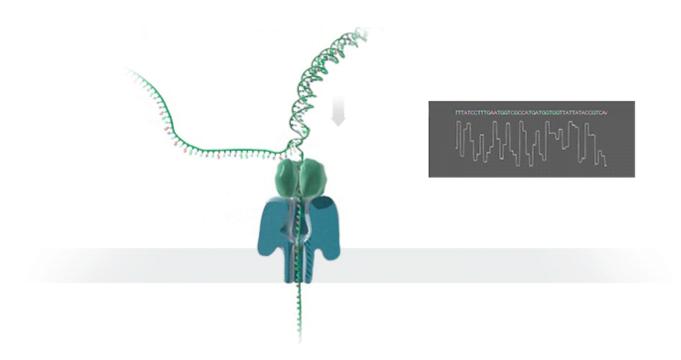
AirLift: A Fast and Comprehensive Technique for Translating Alignments between Reference Genomes

Jeremie S. Kim, Can Firtina, Damla Senol Cali, Mohammed Alser, Nastaran Hajinazar, Can Alkan, Onur Mutlu

GitHub: https://github.com/CMU-SAFARI/AirLift

Kim+, "AirLift: A Fast and Comprehensive Technique for Translating Alignments between Reference Genomes", arXiv, 2020

Nanopore Sequencing



- Nanopore is a nano-scale hole
- In nanopore sequencers, an ionic current passes through the nanopores
- When the DNA strand passes through the nanopore, the sequencer measures the the change in current
- This change is used to identify the bases in the strand with the help of different electrochemical structures of the different bases

The Effect of Pre-Alignment (Theoretically)

